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```
In [5]: import math
```

1 Algorithm Design Technique: Divide and Conquer

- divide into subproblems of same type
- recurse
- combine

1.1 Problem (Maxima)

Given a set P of n points in 2D, we say point p dominates q iff p.x > q.x and p.y > q.y. We say point q is *maximal* if $q \in P$ and no point in P dominates q. Find all maximal points.

```
In [6]: \# O(N^2)
         # brute force
        def find_maximal_1(points):
             big_points = []
             for point1 in points:
                 for point2 in points:
                      if point1 != point2 and (point1[0] < point2[0] or point1[1] < p</pre>
                          break
                 else:
                     big_points.append(point1)
             return big_points
         # 0(N log N)
         # couple of ways e.g. convex hull, 2D range tree (?), sort-then-line-sweep,
        def find_maximal_2(points):
             def f(points):
                 n = len(points)
                 if n <= 1:
                      return points
                 ls = f(points[:n/2])
                 rs = f(points[n/2:])
                 values = [x \text{ for } x \text{ in } 1s \text{ if } x[1] > rs[0][1]] + rs
```

return values

1.2 Closest Pair

Given a set of n points in 2D, find the pair of points s.t. the Euclidean distance, or $(x - y)^2$, is minimized.

```
In [14]: def dist(px, py):
              return (px[0] - py[0]) ** 2 + (px[1] - py[1]) ** 2
          # O(N^2)
          # brute force
          def find_closest_1(points):
              assert len(points) >= 2
              best = float('inf')
              for i in range(len(points)):
                  for j in range(i+1, len(points)):
                       if dist(points[i], points[j]) < best:</pre>
                           best = dist(points[i], points[j])
              return math.sqrt(best)
          \# O(N \log N)
          # Shamos' Algorithm
          # This is because T(n) = 2 T(n/2) + O(n \log n) = O(n^2 \log n)
          # But optimally, we don't need to sort within this algorithm, so T(n) = 2
          def find_closest_2(points):
              sorted_by_ys = sorted(points, key=lambda p:p[1])
              def f(points):
                  if len(points) <= 1:</pre>
                       return float('inf')
                  if len(points) == 2:
                       return math.sqrt(dist(points[0], points[1]))
                  if len(points) == 3:
                       return math.sqrt(min(dist(points[0], points[1]), dist(points[0])
                  x_m = points[len(points)/2][0]
                  p_l = [p \text{ for } p \text{ in } points \text{ if } p[0] \le x_m]
                  p_r = [p \text{ for } p \text{ in } points \text{ if } p[0] > x_m]
                  closest_l = f(p_l)
                  closest_r = f(p_r)
```

delta = min(closest_l, closest_r)

```
\# first idea: look at pairs within some [x - delta, x + delta] ran
                  within_delta_l = [p for p in p_l if x_m - delta <= p[0]]</pre>
                  within_delta_r = [p \text{ for } p \text{ in } p_r \text{ if } p[0] \le x_m + delta]
                  # second idea: if we have sliding window upwards of length delta :
                  # each window contains <= 8 points, because of Pidgeonhole Princip
                  combined = within_delta_l + within_delta_r
                  windows = filter(lambda p: p in combined, sorted_by_ys) # "sorting"
                 best = delta
                  for i in range(len(windows)): # O(n)
                      for j in range(i, min(len(windows), i+8)): # O(1)
                          for k in range(j+1, min(len(windows), i+8)): \# O(1)
                              best = min(best, dist(windows[i], windows[k]))
                  return best
             return f(sorted(points))
In [15]: print find_closest_1([(0, 0), (1, 1), (2, 2), (3, 3)])
         print find_closest_2([(0, 0), (1, 1), (2, 2), (3, 3)])
1.41421356237
1.41421356237
```