

EECS 489

Computer Networks

Fall 2020

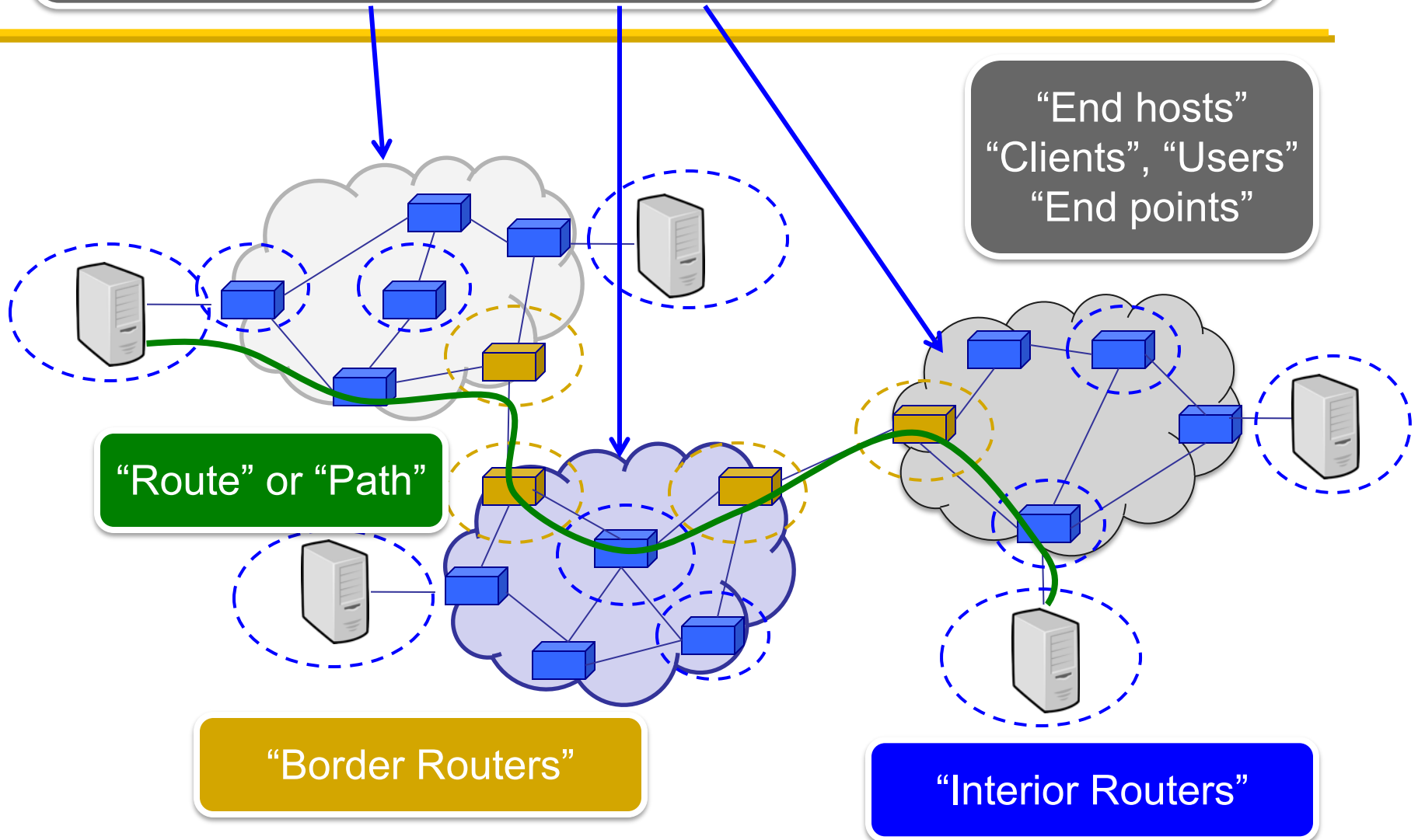
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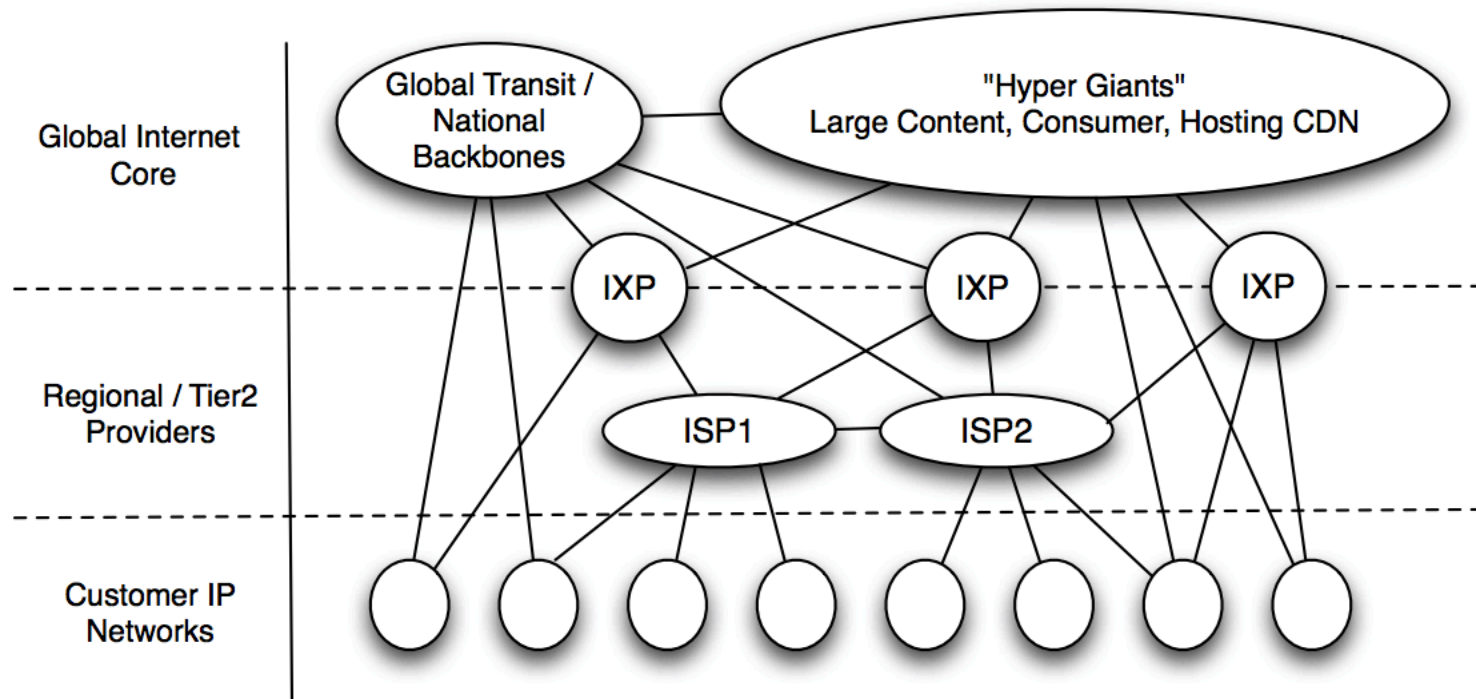
Agenda

- Inter-domain-routing

“Autonomous System (AS)” or “Domain”
Region of a network under a single administrative entity



AS-level Internet



Internet Inter-Domain Traffic, SIGCOMM, 2010

Autonomous systems (AS)

- An AS is a network under a single administrative control
 - Currently over 70,000 ASes
 - Updated daily at <http://www.cidr-report.org/as2.0/>
- ASes are sometimes called “domains”
- Each AS is assigned a unique identifier (ASN)
 - E.g., University of Michigan owns ASNs 177 to 180

“Intra-domain” routing: Within an AS

- Link-State (e.g., OSPF) and Distance-Vector (e.g., RIP)
- Primary focus
 - Finding least-cost paths
 - Fast convergence

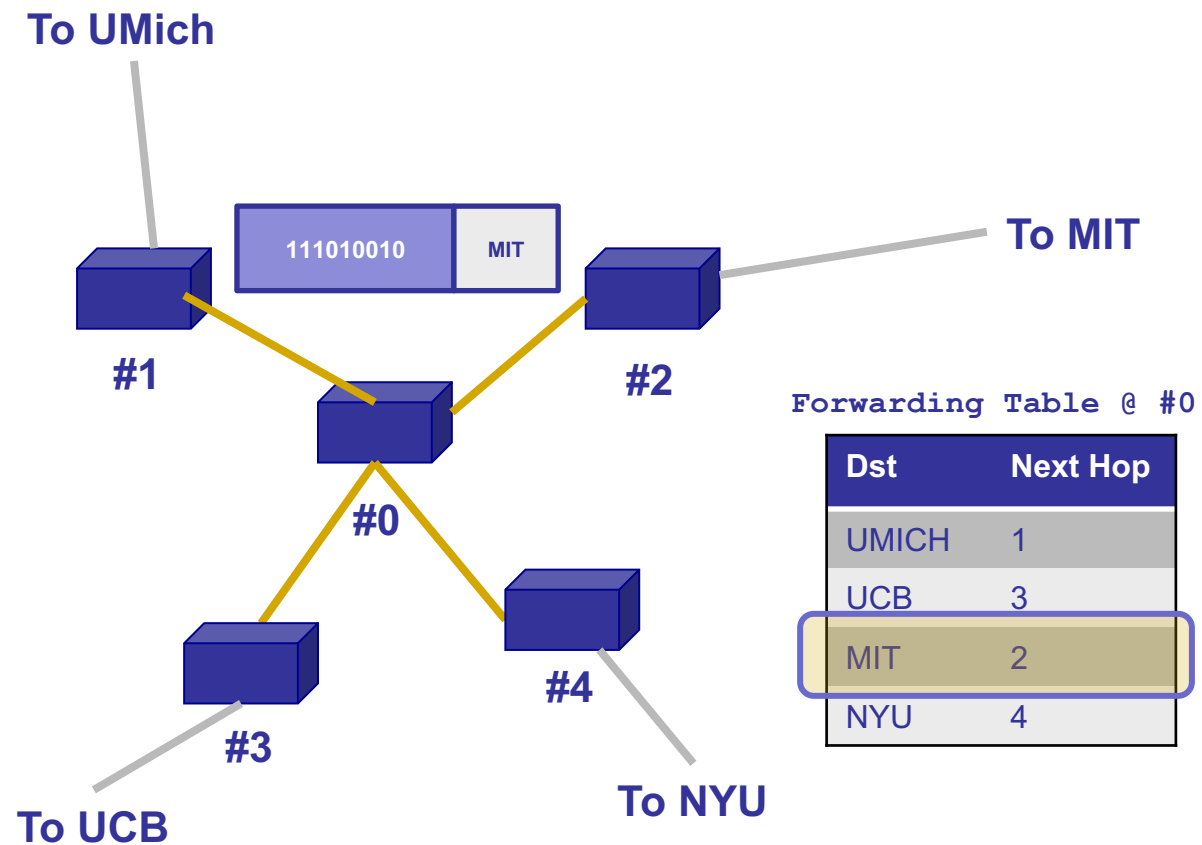
“Inter-domain” routing: Between ASes

- Two key challenges
 - Scaling
 - Administrative structure
 - » Issues of autonomy, policy, privacy

Recall: Addressing (so far)

- Each host has a unique ID
- No particular structure to those IDs

Recall: Forwarding



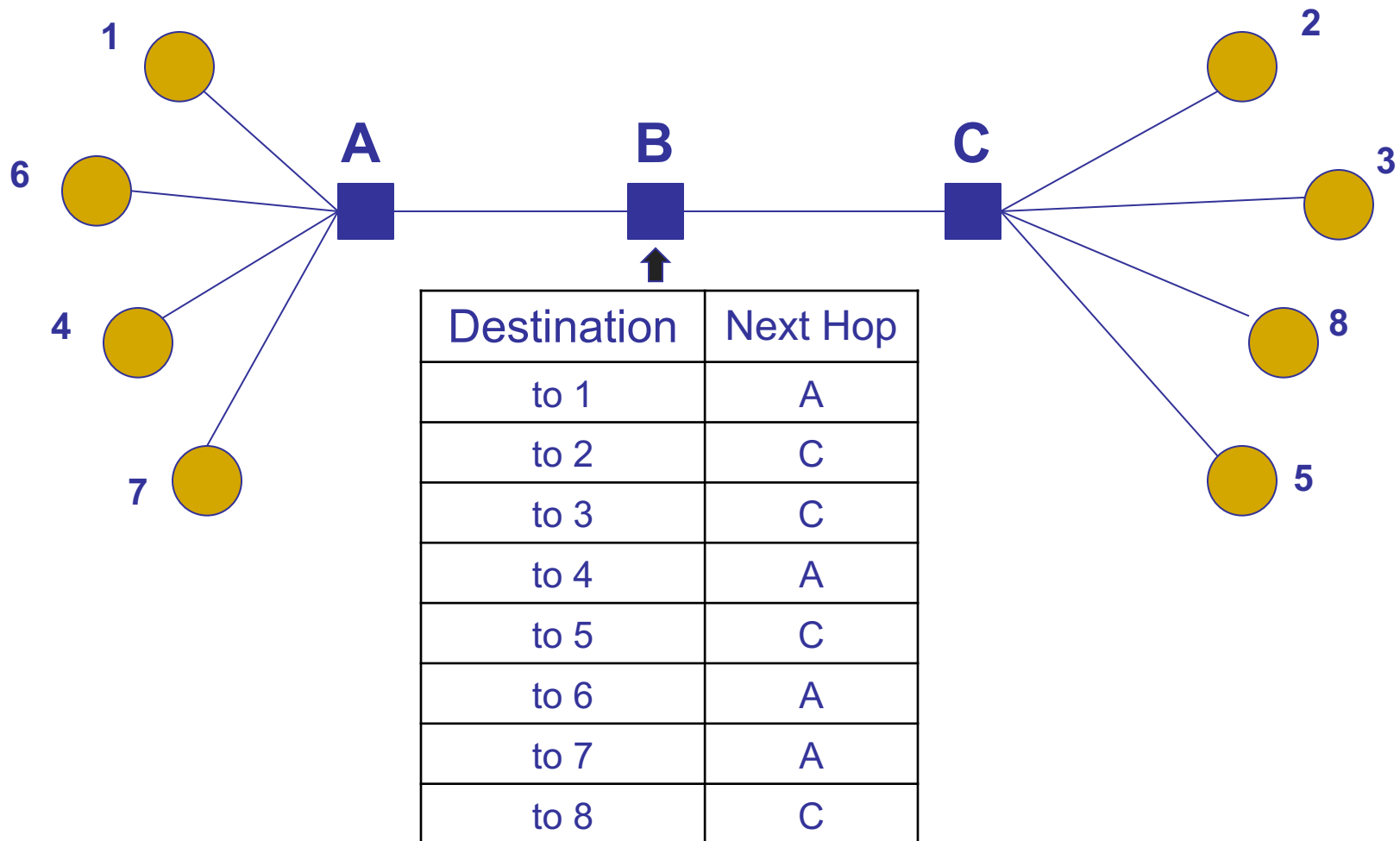
Two key challenges

- Scaling
- Administrative structure
 - Issues of autonomy, policy, privacy

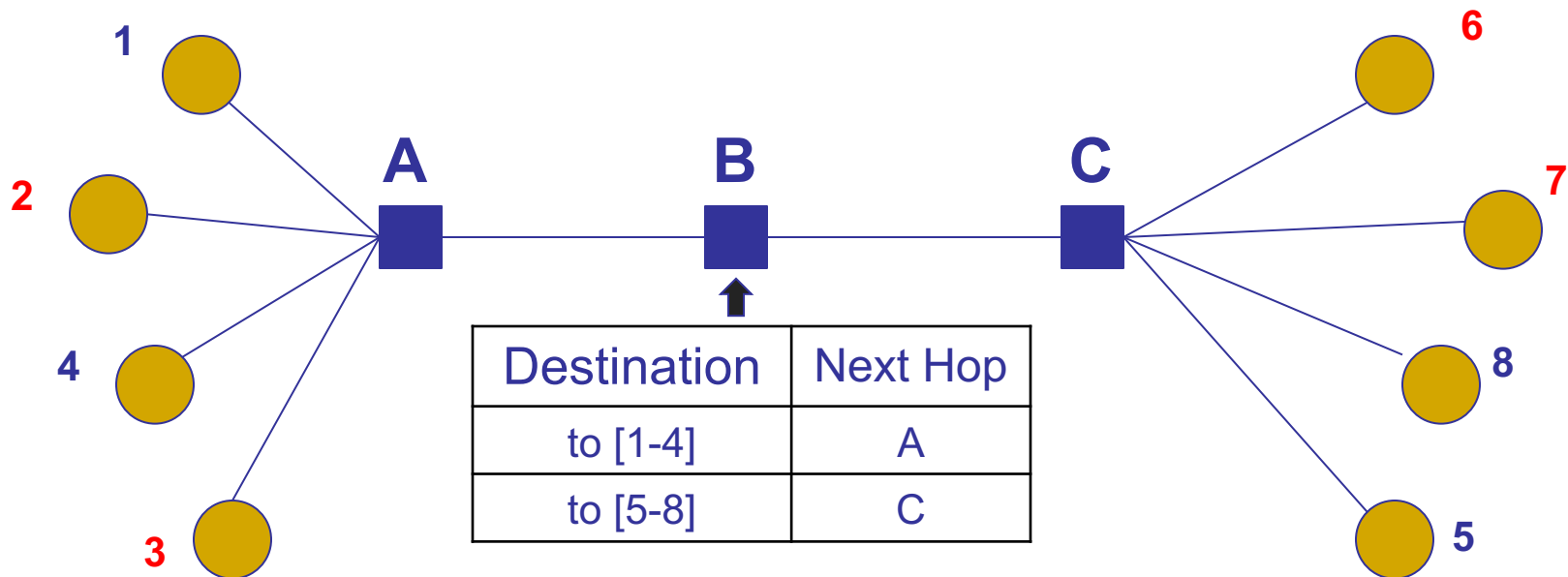
Scaling

- A router must be able to reach any destination
 - Given packet's destination address, lookup next hop
- Naive: Have an entry for each destination
 - There would be over 10^8 entries!
 - **AND** routing updates per destination!
- How can we improve scalability?
 - We have already seen an example: **longest-prefix matching**

A smaller table at node B?



Re-number the end-systems?



- Careful address assignment → can *aggregate* multiple addresses into one range → scalability!
- Akin to reducing the number of destinations

Scaling

- A router must be able to reach any destination
- Naive: Have an entry for each destination
- Better: Have an entry for a range of addresses
 - Can't do this if addresses are assigned randomly!
 - How addresses are allocated will matter!
- Host addressing is key to scaling

Two key challenges

- Scaling
- Administrative structure
 - Issues of autonomy, policy, privacy

Administrative structure shapes inter-domain routing

- ASes want **freedom in picking routes**
 - “My traffic can’t be carried over my competitor’s network”
 - “I don’t want to carry A’s traffic through my network”
 - Not expressible as Internet-wide “least cost”
- ASes want **autonomy**
 - Want to choose their own internal routing protocol
 - Want to choose their own policy
- ASes want **privacy**
 - Choice of network topology, routing policies, etc.

Choice of routing algorithm

- Link-state
 - No privacy – broadcasts all network information
 - Limited autonomy – needs agreement on metric, algo
- Distance-vector is a decent starting point
 - Per-destination updates give some control
 - BUT wasn't designed to implement policy
 - AND is vulnerable to loops
- The “Border Gateway Protocol” (BGP) extends distance-vector ideas to accommodate policy

Agenda

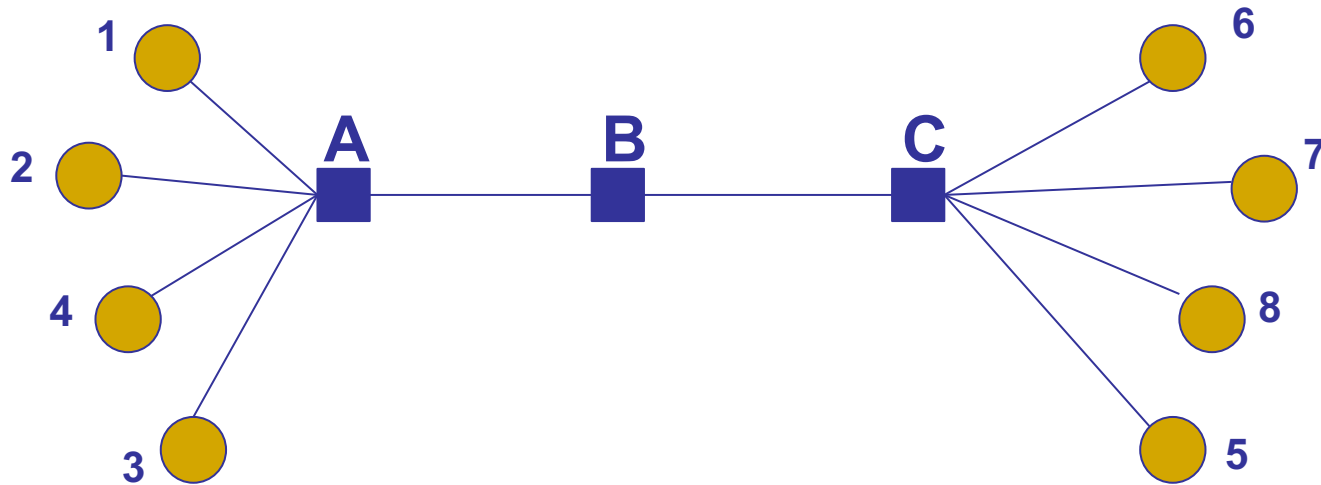
- Inter-domain-routing
 - Addressing (Scalability)
 - BGP (Autonomy, policy, privacy)
 - » Context and basic ideas: today
 - » Details and issues: next lecture

IP ADDRESSING

Goal of addressing: Scalable routing

- State: Small forwarding tables at routers
 - Much less than the number of hosts
- Churn: Limited rate of change in routing tables
- Ability to aggregate addresses is crucial for both

Aggregation works if...



- Groups of destinations reached via the same path
- These groups are assigned contiguous addresses
- These groups are relatively stable
- Few enough groups to make forwarding easy

IP addressing is hierarchical

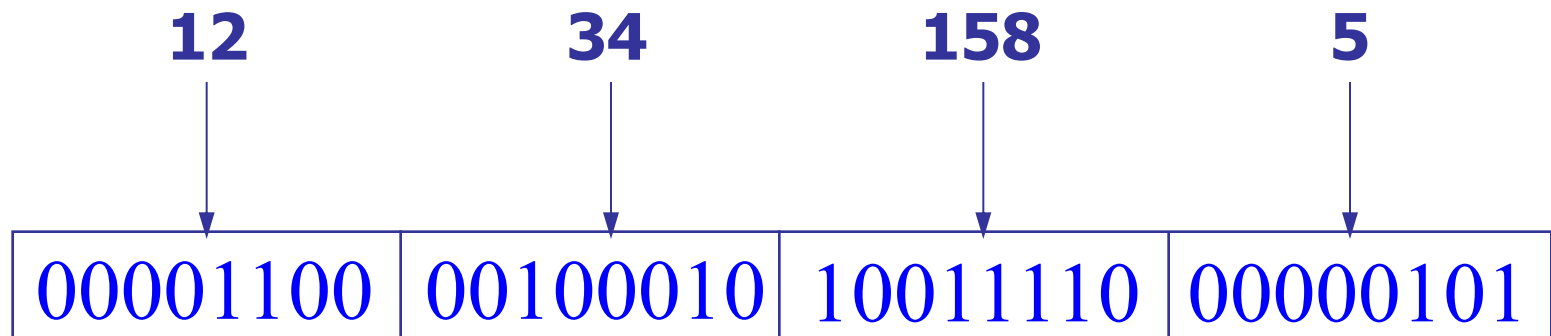
- Hierarchical address structure
- Hierarchical address allocation
- Hierarchical addresses and routing scalability

IP addresses (IPv4)

- Unique 32-bit number associated with a host

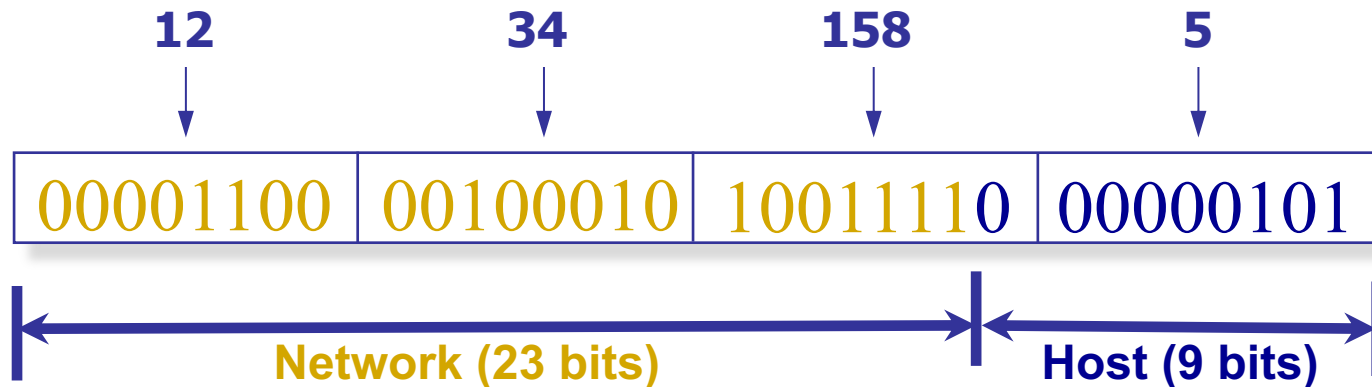
00001100 00100010 10011110 00000101

- Represented with the “dotted-decimal” notation
 - e.g., 12.34.158.5



Hierarchy in IP addressing

- 32 bits are partitioned into a prefix and suffix components
- Prefix is the **network** component; suffix is the **host** component



- Inter-domain routing operates on network prefix

CIDR: Classless inter-domain routing

- Flexible division between network and host addresses
- Offers a better tradeoff between size of the routing table and efficient use of the IP address space

CIDR example

- Suppose a network has 50 computers
 - Allocate 6 bits for host addresses ($2^5 < 50 < 2^6$)
 - Remaining $32 - 6 = 26$ bits as network prefix
- Flexible boundary means the boundary must be explicitly specified with the network address!
 - Informally, “slash 26” \rightarrow 128.23.9/26
 - Formally, prefix represented with a 32-bit mask: 255.255.255.192, where all network prefix bits set to “1” and host suffix bits to “0”
 - Also known as **subnet mask** (a group of machines with the same prefix are in the same subnet)

Before CIDR: Classful addressing

- Three classes
 - 8-bit network prefix (Class A),
 - 16-bit network prefix (Class B), or
 - 24-bit network prefix (Class C)
- Example: an organization needs 500 addresses.
 - A single class C address is not enough (<500 hosts)
 - Instead, a class B address is allocated (~65K hosts)
 - » Huge waste!

IP addressing is hierarchical

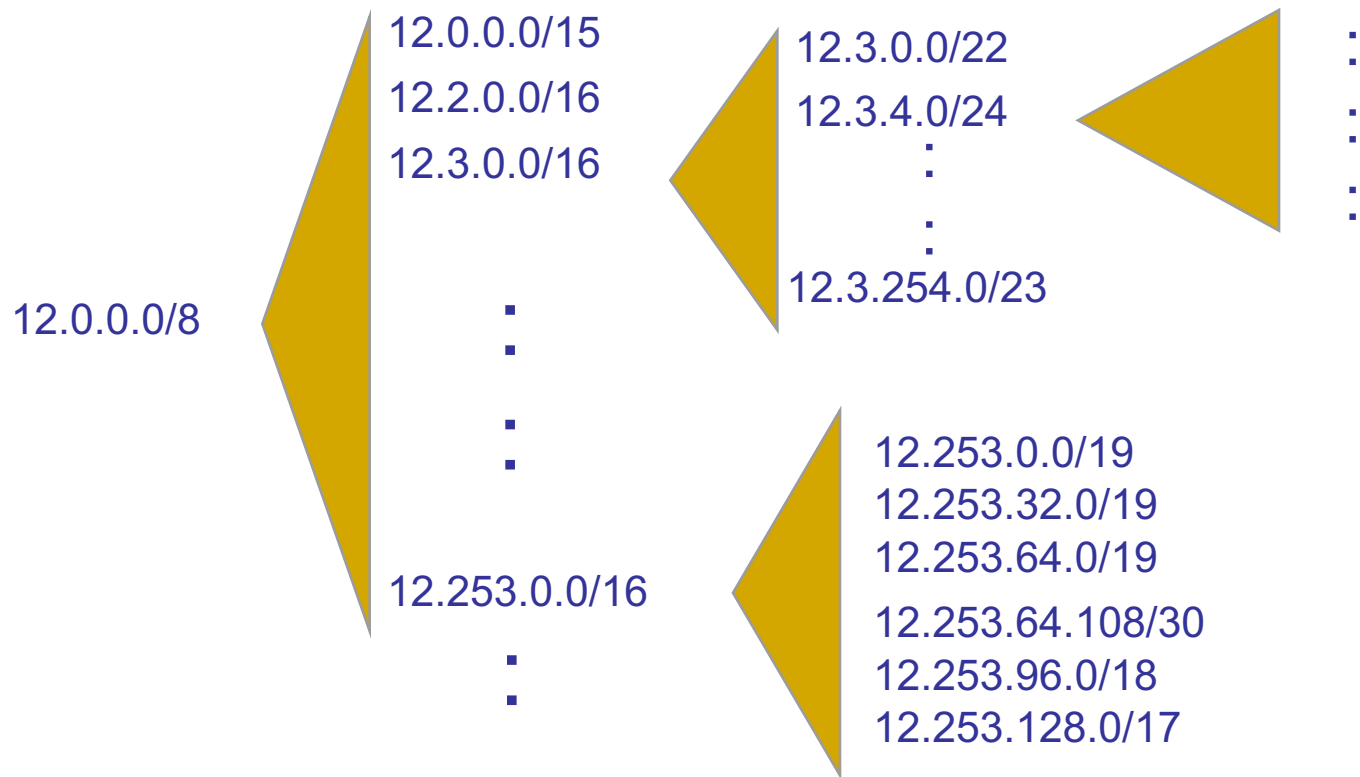
- Hierarchical address structure
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- Hierarchical addresses and routing scalability

Allocation done hierarchically

- Internet Corporation for Assigned Names and Numbers (ICANN) gives large blocks to...
- Regional Internet Registries, such as the American Registry for Internet Names (ARIN), which give blocks to...
- Large institutions (ISPs), which give addresses to...
- Individuals and smaller institutions
- FAKE Example:
 - ICANN → ARIN → AT&T → UMICH → EECS

CIDR: Addresses allocated in contiguous prefix chunks

- Recursively break down chunks as get closer to host



FAKE example in more detail

- ICANN gave ARIN several /8s
- ARIN gave AT&T one /8, **12.0/8**
 - Network Prefix: **00001100**
- AT&T gave UMich a /16, **12.34/16**
 - Network Prefix: **0000110000100010**
- UMich gave EECS a /24, **12.34.56/24**
 - Network Prefix: **000011000010001000111000**
- EECS gave me specific address **12.34.56.78**
 - Address: **00001100001000100011100001001110**

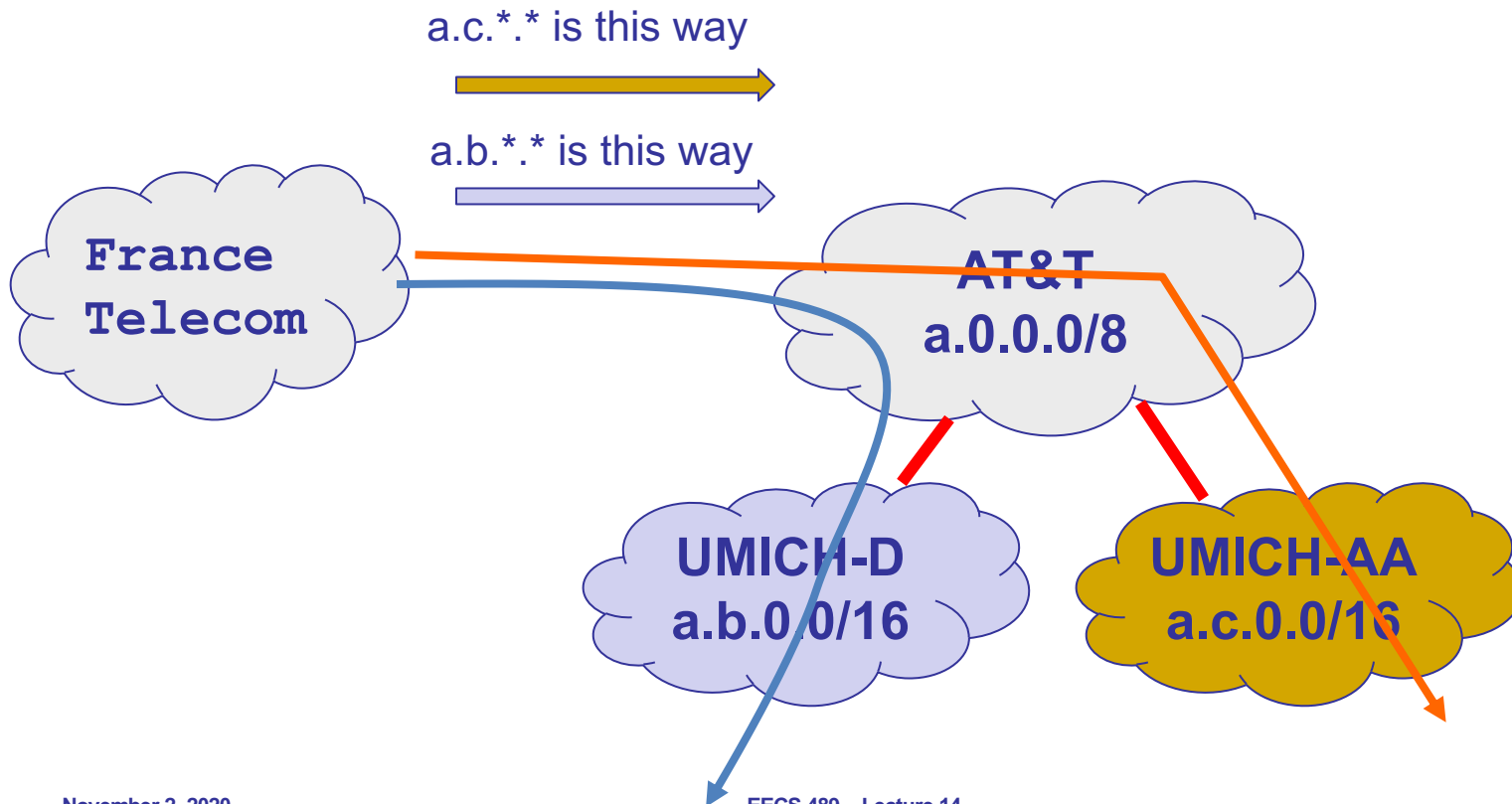
IP addressing is hierarchical

- Hierarchical address structure
- Hierarchical address allocation
- Hierarchical addresses and routing scalability

IP addressing → Scalable routing?

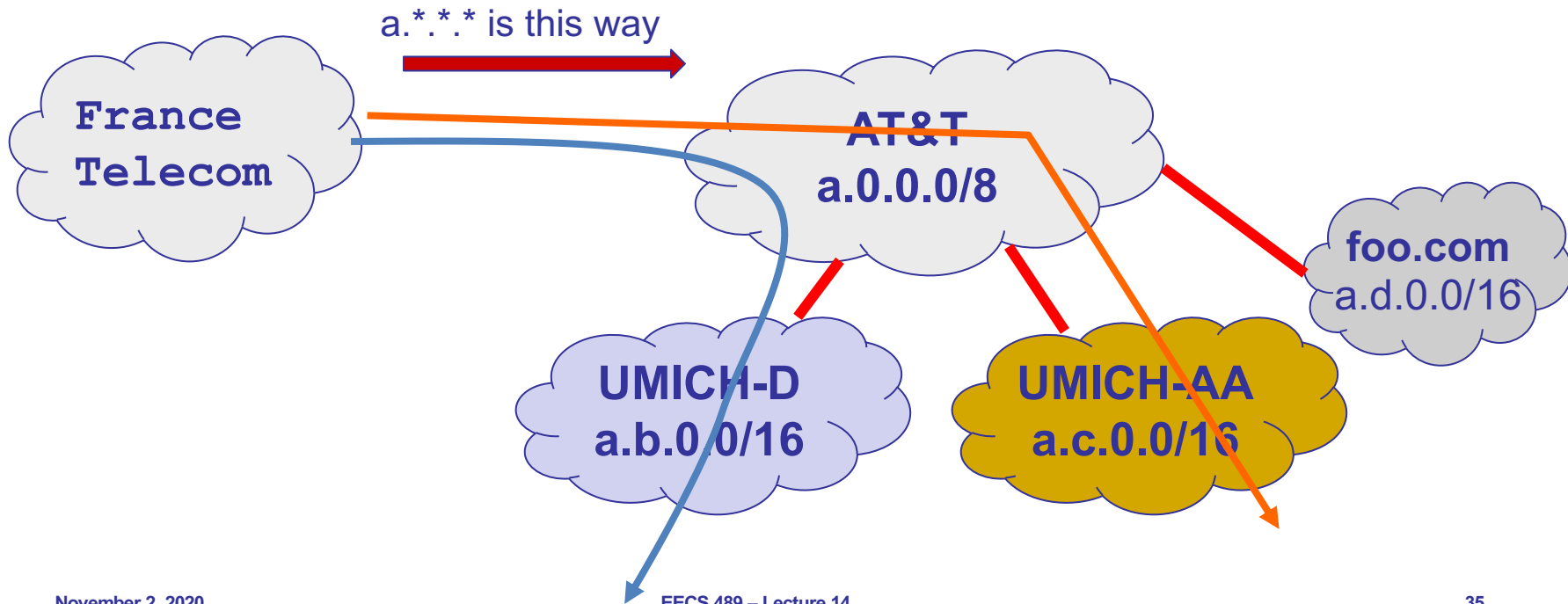
- Hierarchical address allocation only helps routing scalability if allocation matches topological hierarchy

IP addressing → Scalable routing?



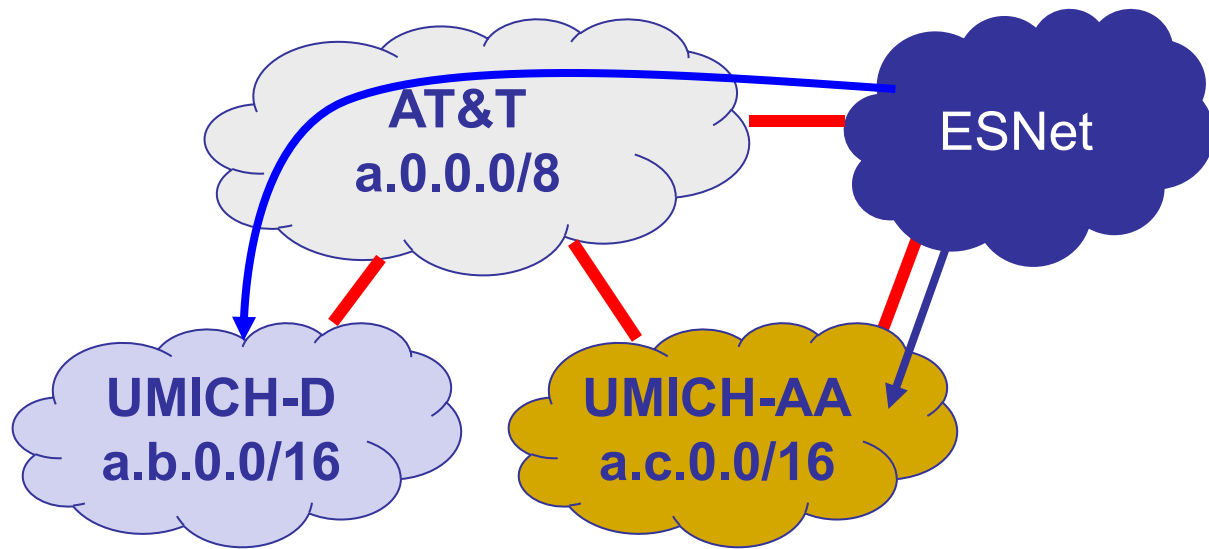
IP addressing → Scalable routing?

Can add new hosts/networks without updating the routing entries at France Telecom



IP addressing → Scalable routing?

ESNet must maintain routing entries for both $a.*.*$ and $a.c.*.*$



IP addressing → Scalable routing?

- Hierarchical address allocation only helps routing scalability if allocation matches topological hierarchy
- May not be able to aggregate addresses for “multi-homed” networks
 - A multi-homed network is connected to more than one ASes for fault-tolerance, load balancing, etc.

5-MINUTE BREAK!

BGP: BORDER GATEWAY PROTOCOL

BGP (Today)

- The role of policy
 - What we mean by it
 - Why we need it
- Overall approach
 - Four non-trivial changes to DV

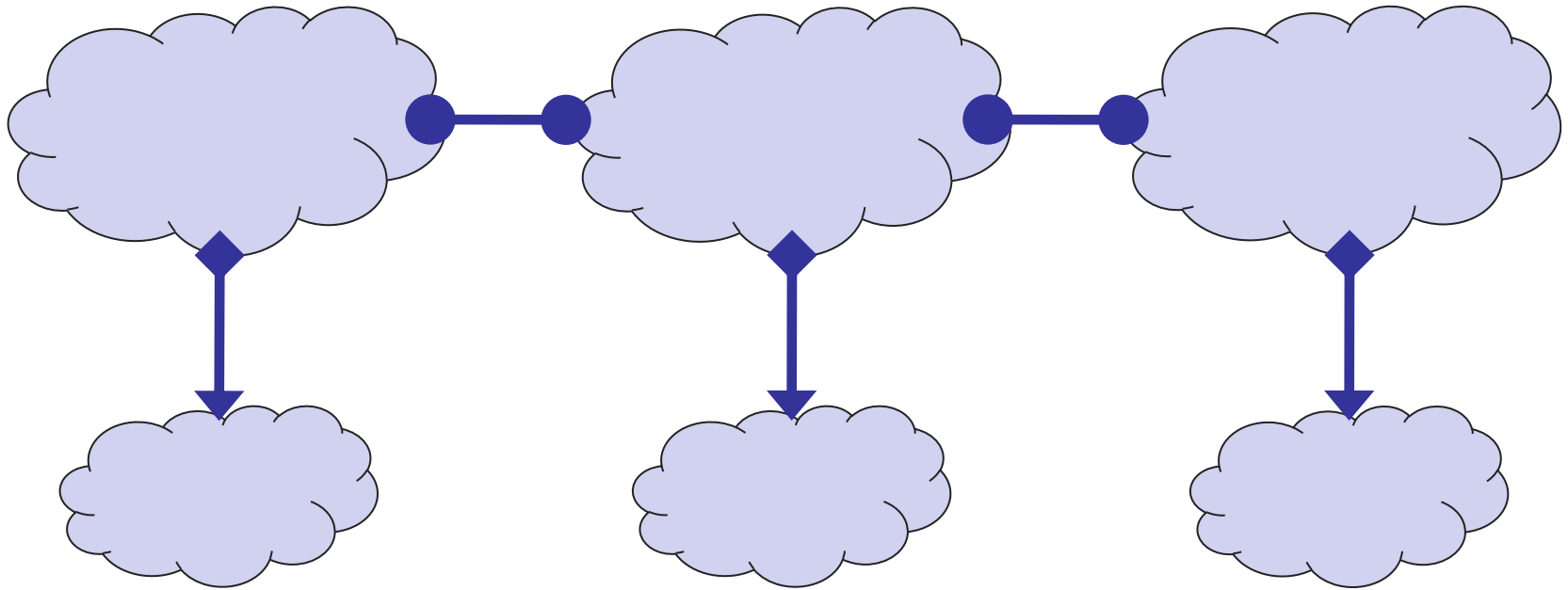
Administrative structure shapes Inter-domain routing

- ASes want freedom to pick routes based on policy
- ASes want autonomy
- ASes want privacy

Topology & policy shaped by inter-AS business relationship

- Three basic kinds of relationships between ASes
 - AS A can be AS B's customer
 - AS A can be AS B's provider
 - AS A can be AS B's peer
- Business implications
 - Customer pays provider
 - Peers don't pay each other
 - » Exchange roughly equal traffic

Business relationships



Relations between ASes

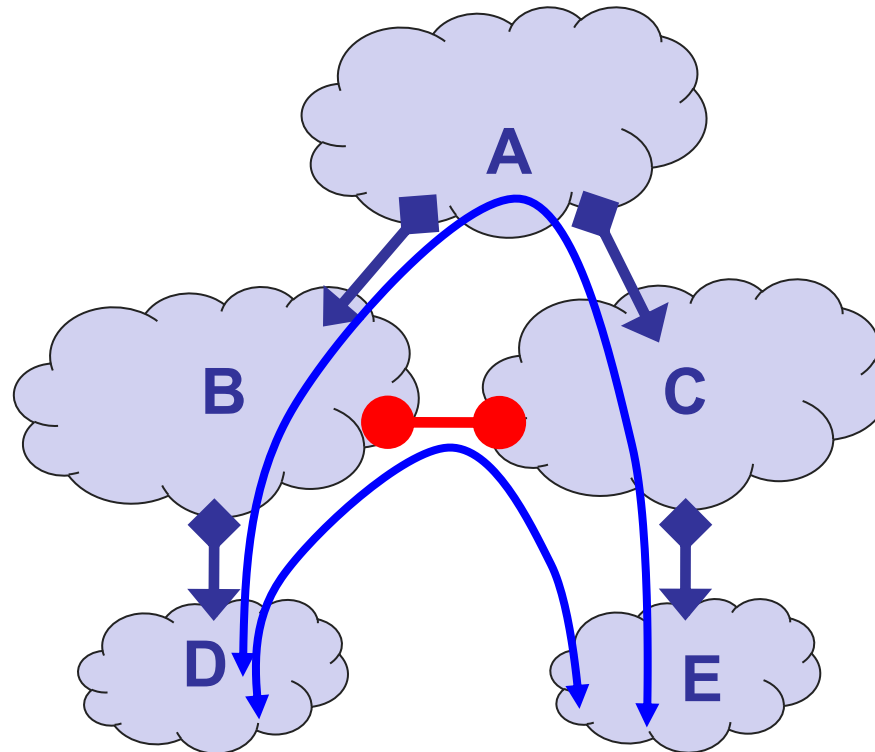
provider \longleftrightarrow customer

peer $\bullet\text{---}\bullet$ peer

Business implications

- Customers pay provider
- Peers don't pay each other

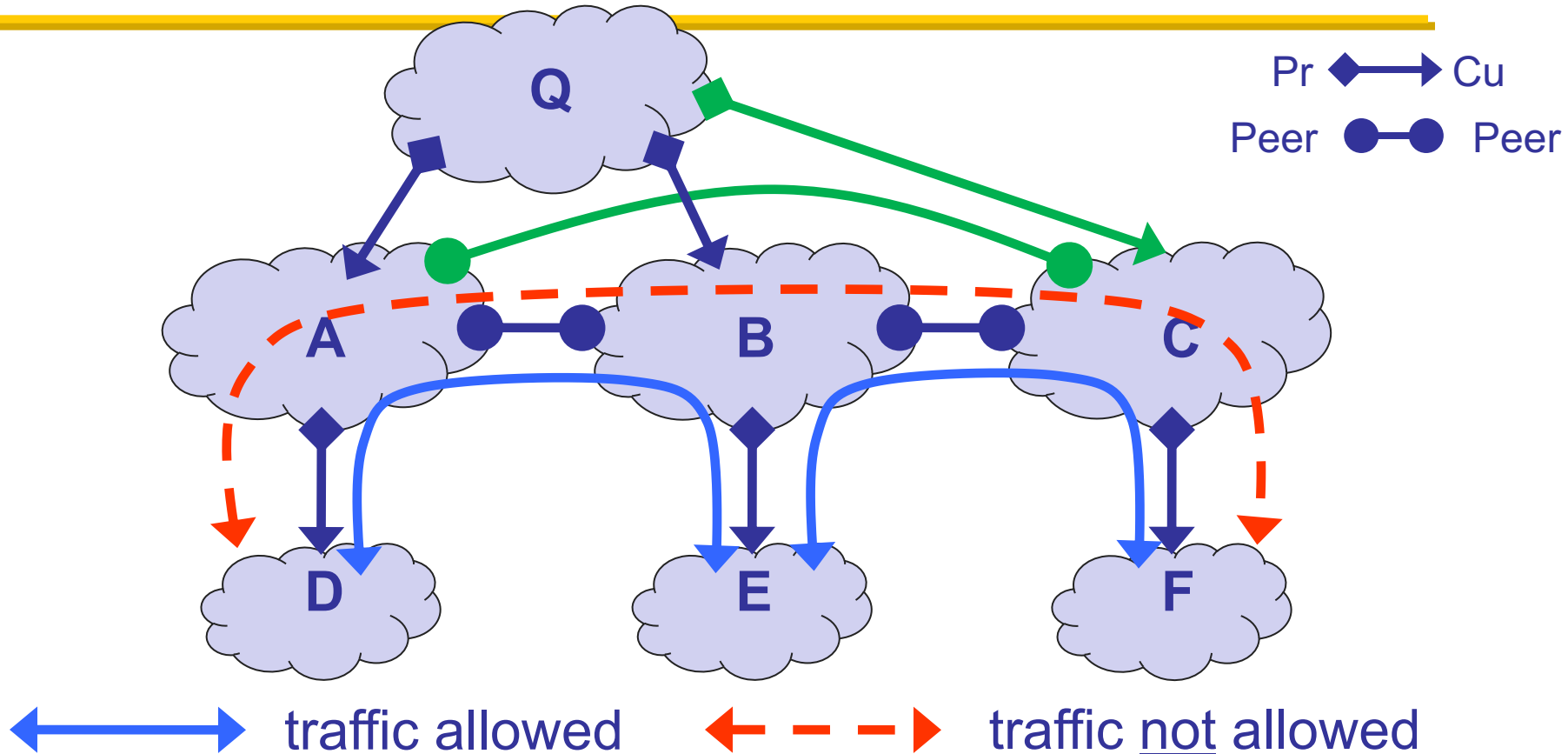
Why peer?



D and E
communicate a lot

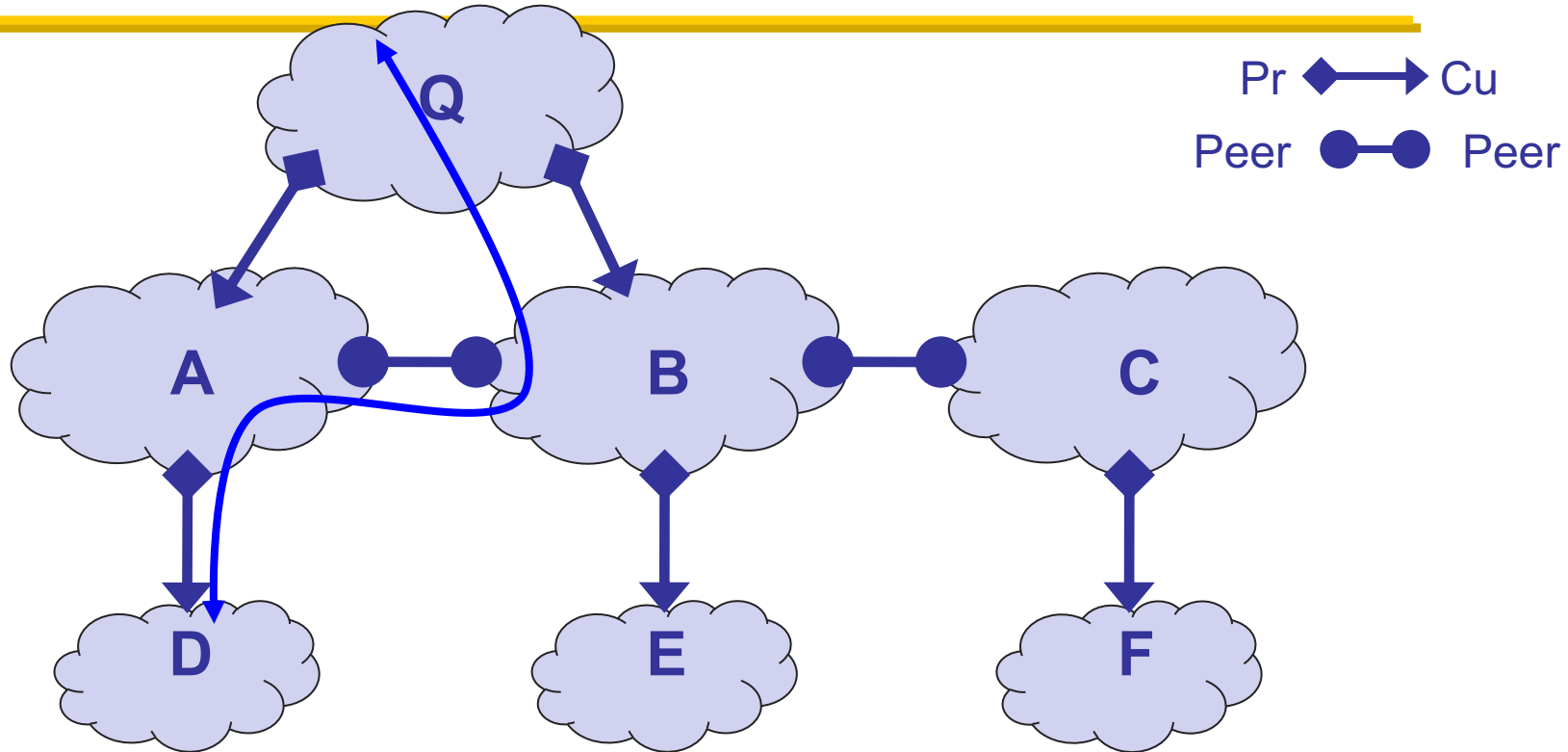
Peering saves
B and C money

Routing follows the money!



- ASes provide “transit” between their customers
- Peers do not provide transit between other peers

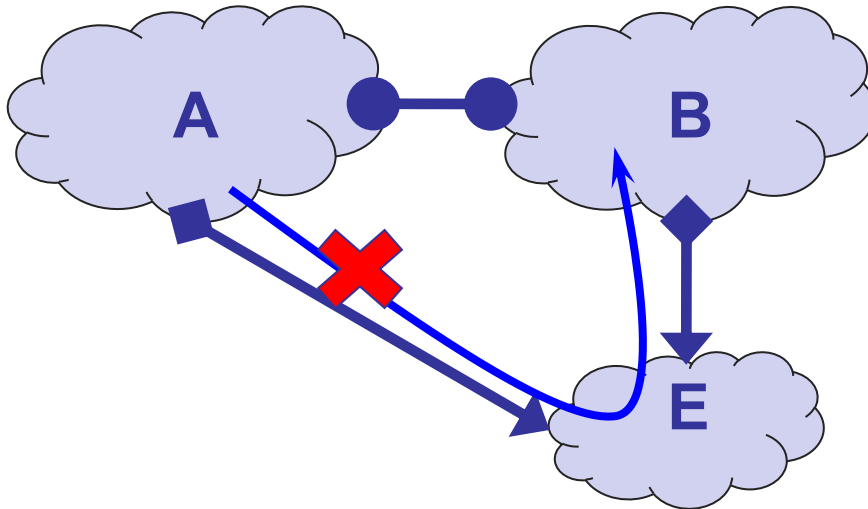
Routing follows the money!



- An AS only carries traffic to/from its own customers over a peering link

Routing follows the money!

Pr \longleftrightarrow Cu
Peer $\bullet\text{---}\bullet$ Peer



- Routes are “valley” free (more details later)

In short

- AS topology reflects business relationships between ASes
- Business relationships between ASes impact which routes are acceptable

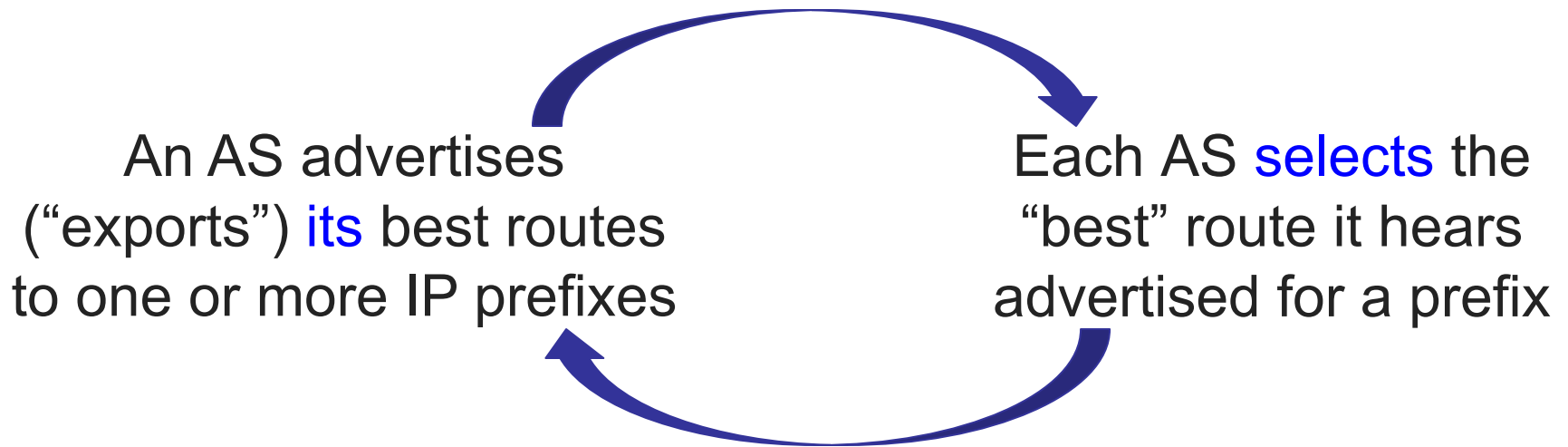
BGP (Today)

- The role of policy
 - What we mean by it
 - Why we need it
- Overall approach
 - Four non-trivial changes to DV

Inter-domain routing: Setup

- Destinations are IP prefixes (12.0.0.0/8)
- Nodes are Autonomous Systems (ASes)
 - Internals of each AS are hidden
- Links represent both physical links and business relationships
- BGP (Border Gateway Protocol) is the Inter-domain routing protocol
 - Implemented by AS border routers

BGP: Basic idea



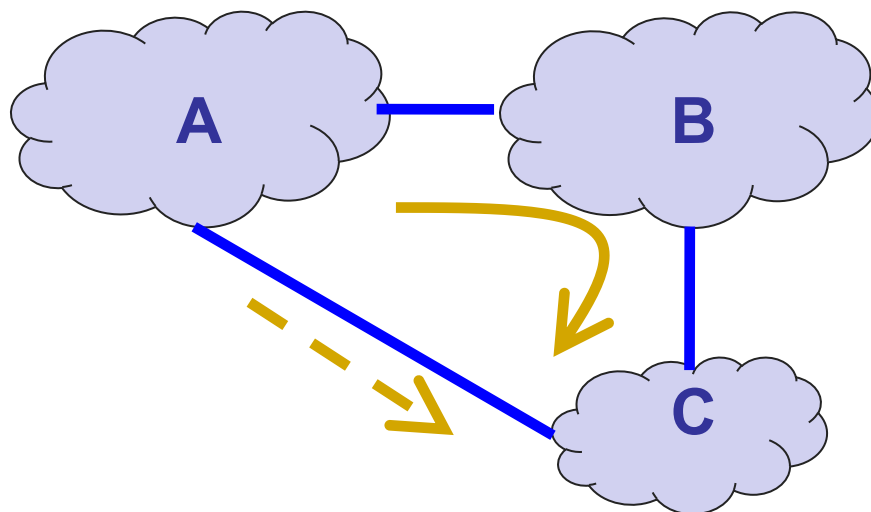
You've heard this story before!

BGP inspired by Distance-Vector

- Per-destination route advertisements
- No global sharing of network topology information
- Iterative and distributed convergence on paths
- With four crucial differences!

BGP & DV differences: (1) Not picking shortest-path routes

- BGP selects the best route based on policy, not shortest distance (i.e., least-cost)
- AS A may prefer “A,B,C” over “A,C”

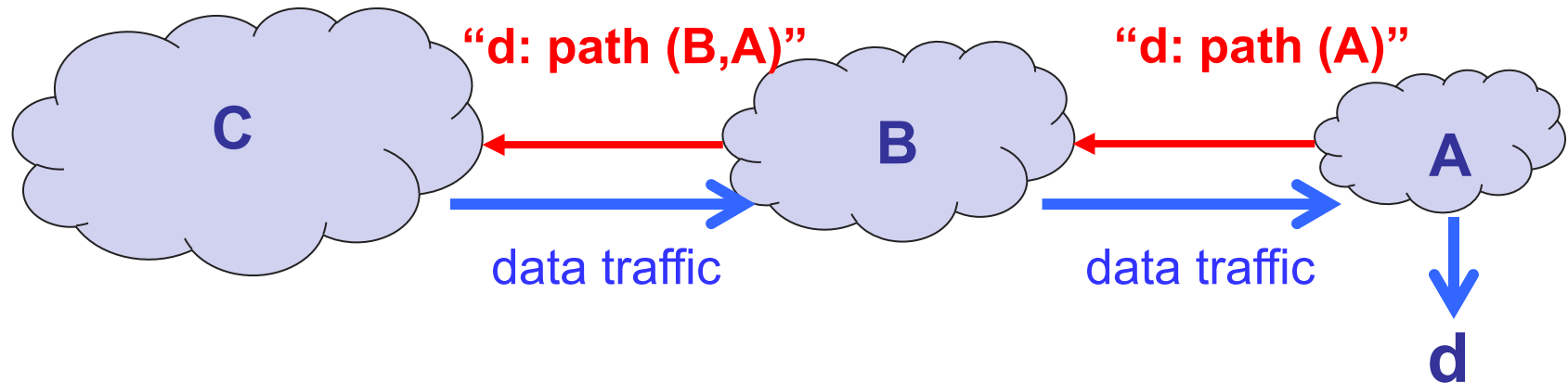


- How do we avoid loops?

BGP & DV differences:

(2) Path-Vector routing

- Key idea: advertise the entire path
 - Distance vector: send distance metric per dest d
 - Path vector: send the entire path for each dest d



BGP & DV differences:

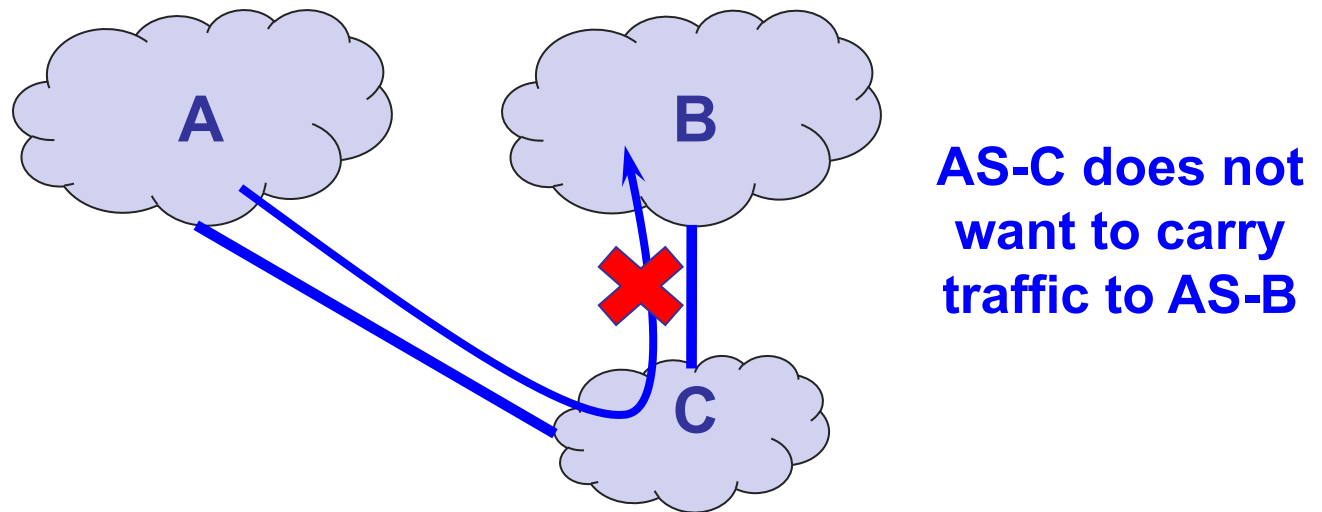
(2) Path-Vector routing

- Key idea: advertise the entire path
 - Distance vector: send distance metric per destination
 - Path vector: send the entire path for each destination
- Benefits
 - Loop avoidance is straightforward (simply discard paths with loops)
 - Flexible and expressive policies based on entire path

BGP & DV differences: (3)

Selective route advertisement

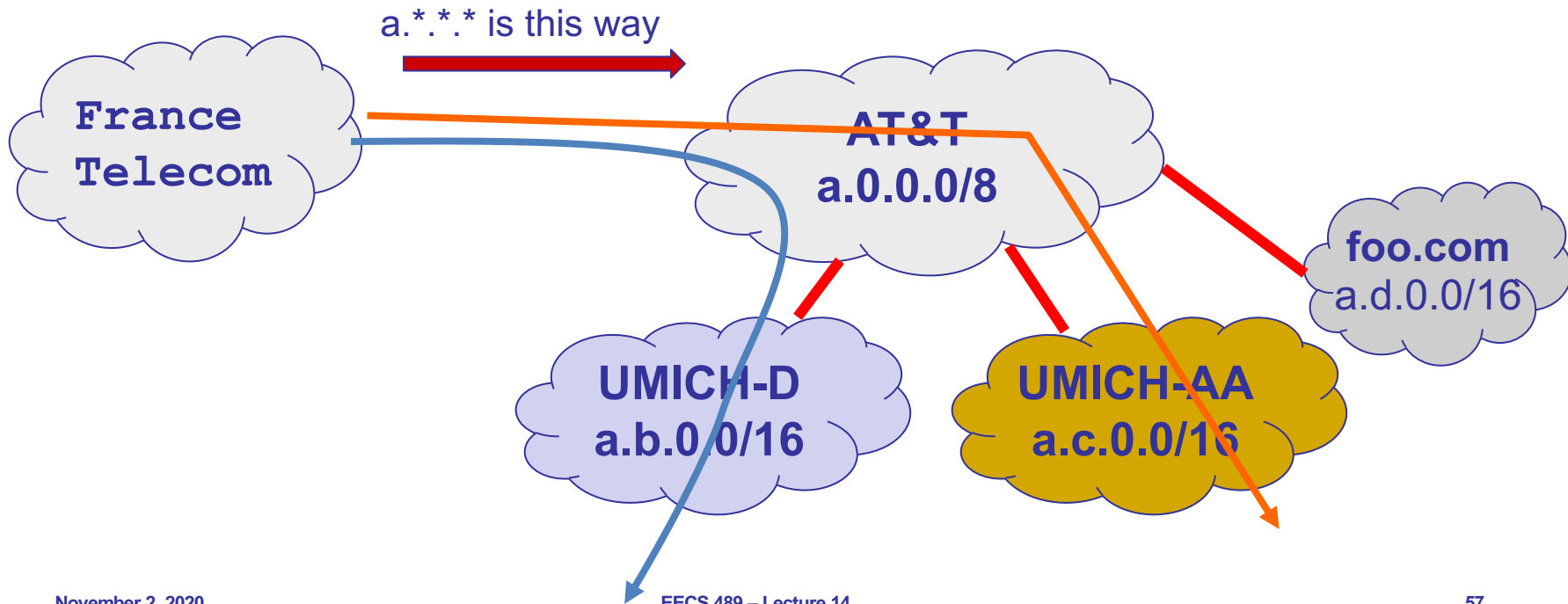
- For policy reasons, an AS may choose not to advertise a route to a destination
- Hence, **reachability is not guaranteed** even if graph is physically connected



BGP & DV differences:

(4) BGP may aggregate routes

- For scalability, BGP may aggregate routes for different prefixes



Summary

- Two key challenges in inter-domain routing
 - Scaling (Addressing)
 - Administrative structure (BGP)
 - » Issues of autonomy, policy, privacy
- **Next lecture:** BGP policies, protocol, and challenges