

CS161 HW5

1.

a. neither, three worlds satisfy the statement (satisfiable) but it is not always true in all worlds (not valid)

Smoke	Fire	Smoke \Rightarrow Fire	-Smoke	-Fire	-Smoke \Rightarrow -Fire	(Smoke \Rightarrow Fire) \Rightarrow (-Smoke \Rightarrow -Fire)
T	T	T	F	F	T	T
T	F	F	F	T	T	T
F	T	T	T	F	F	F
F	F	T	T	T	T	T

b. neither, statement is satisfiable but not valid since there is one world where it is not true

Smoke	Fire	Smoke \Rightarrow Fire	Heat	Smoke \vee Heat	(Smoke \vee Heat) \Rightarrow Fire	(Smoke \Rightarrow Fire) \Rightarrow ((Smoke \vee Heat) \Rightarrow Fire)
T	T	T	T	T	T	T
T	F	F	T	T	F	T
F	T	T	T	T	T	T
F	F	T	T	T	F	F
T	T	T	F	T	T	T
T	F	F	F	F	T	T
F	T	T	F	T	T	T
F	F	T	F	F	T	T

c. valid since it is true for all worlds in truth table

Smoke	Heat	Fire	Smoke \wedge Heat	((Smoke \wedge Heat) \Rightarrow Fire)	((Smoke \Rightarrow Fire) \vee (Heat \Rightarrow Fire))	((Smoke \wedge Heat) \Rightarrow Fire) \Leftrightarrow ((Smoke \Rightarrow Fire) \vee (Heat \Rightarrow Fire))
T	T	T	T	T	T	T
T	T	F	T	F	F	T
F	T	T	F	T	T	T
F	T	F	F	T	T	T
T	F	T	F	T	T	T
T	F	F	F	T	T	T
F	F	T	F	T	T	T
F	F	F	F	T	T	T

2.

a. Knowledge Base: mythical(unicorn) \Rightarrow immortal(unicorn),
 \neg mythical(unicorn) $\Rightarrow \neg$ immortal(unicorn) \wedge mammal(unicorn),
 immortal(unicorn) \vee mammal(unicorn) \Rightarrow horned(unicorn),
 horned(unicorn) \Rightarrow magical(unicorn),

b. Convert to CNF: (\neg mythical(unicorn) \vee immortal(unicorn)) \wedge

$$\begin{aligned}
& (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}) \wedge \text{mammal}(\text{unicorn}))) \wedge \\
& (\neg (\text{immortal}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn})) \\
& = \\
& (\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \wedge \\
& ((\neg \text{immortal}(\text{unicorn}) \wedge \neg \text{mammal}(\text{unicorn})) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn})) \\
& = \\
& (\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \wedge \\
& (\neg \text{immortal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{mammal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn}))
\end{aligned}$$

c. No, it is not possible to prove that the unicorn is mythical.

$$\begin{aligned}
& (\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \wedge \\
& (\neg \text{immortal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{mammal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn})) \wedge \\
& (\neg \text{mythical}(\text{unicorn})) \\
& \Rightarrow \\
& (\neg \text{immortal}(\text{unicorn})) \wedge \quad \text{because } (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \\
& \text{mammal}(\text{unicorn})) \quad \text{because } (\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \\
& \Rightarrow \\
& \neg \text{mythical}(\text{unicorn}) \wedge \quad \text{because } (\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \\
& \text{horned}(\text{unicorn})) \quad \text{because } (\neg \text{mammal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \\
& \Rightarrow \\
& \text{magical}(\text{unicorn}) \quad \text{done, cannot continue with resolution}
\end{aligned}$$

Yes, it is possible to prove magical by refutation.

$$\begin{aligned}
& (\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \wedge \\
& (\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \wedge \\
& (\neg \text{immortal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{mammal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge \\
& (\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn})) \wedge \\
& (\neg \text{magical}(\text{unicorn}))
\end{aligned}$$

\Rightarrow
 $\neg \text{horned}(\text{unicorn})$
 \Rightarrow
 $\neg \text{mammal}(\text{unicorn}) \wedge \neg \text{immortal}(\text{unicorn})$
 \Rightarrow
 $\text{mythical}(\text{unicorn}) \wedge$ because $(\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn}))$
 $\neg \text{mythical}(\text{unicorn})$ because $(\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn}))$
 \Rightarrow
 $\text{magical}(\text{unicorn})$ by contradiction

Yes, it is possible to prove horned.

$(\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn})) \wedge$
 $(\text{mythical}(\text{unicorn}) \vee (\neg \text{immortal}(\text{unicorn}))) \wedge$
 $(\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn})) \wedge$
 $(\neg \text{immortal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge$
 $(\neg \text{mammal}(\text{unicorn}) \vee \text{horned}(\text{unicorn})) \wedge$
 $(\neg \text{horned}(\text{unicorn}) \vee \text{magical}(\text{unicorn})) \wedge$
 $(\neg \text{horned}(\text{unicorn}))$
 $\Rightarrow \neg \text{mammal}(\text{unicorn}) \wedge \neg \text{immortal}(\text{unicorn})$
 $\text{mythical}(\text{unicorn}) \wedge$ because $(\text{mythical}(\text{unicorn}) \vee \text{mammal}(\text{unicorn}))$
 $\neg \text{mythical}(\text{unicorn})$ because $(\neg \text{mythical}(\text{unicorn}) \vee \text{immortal}(\text{unicorn}))$
 \Rightarrow
 $\text{horned}(\text{unicorn})$ by contradiction

3.

- a. $\{x/A, y/B, z/B\}$ unifies $P(A, B, B)$, $P(x, y, z)$
- b. unifier does not exist for $Q(y, G(A, B))$, $Q(G(x, x), y)$
- c. $\{x/\text{John}, y/\text{John}\}$ unifies $\text{Older}(\text{Father}(y), y)$, $\text{Older}(\text{Father}(x), \text{John})$ because $y = \text{John}$, $\text{Father}(x) = \text{Father}(y)$, $x = \text{John}$
- d. unifier does not exist for $\text{knows}(\text{Father}(y), y)$, $\text{Knows}(x, x)$ because $x = y$, $x = \text{Father}(y)$, $y = \text{Father}(y)$

4.

a.

$A x, \text{Food}(x) \Rightarrow \text{Likes}(\text{John}, x)$
 $\text{Food}(\text{Apple})$
 $\text{Food}(\text{Chicken})$
 $A x, \exists y, \text{Eats}(y, x) \wedge \neg \text{Killedby}(y, x) \Rightarrow \text{Food}(x)$
 $A y, x, \text{Killedby}(y, x) \Rightarrow \neg \text{Alive}(y)$
 $\text{Eats}(\text{Bill}, \text{Peanuts}) \wedge \text{Alive}(\text{Bill})$
 $A x, \text{Eats}(\text{Bill}, x) \Rightarrow \text{Eats}(\text{Sue}, x)$

b.

$(\neg \text{Food}(x) \mid \text{Likes}(\text{John}, x)) \wedge$

Food(Apple) &
Food(Chicken) &
-Eats(y, x) | Killedby(y, x) | Food(x) &
- Killedby(y, x) | -Alive(y) &
Eats(Bill, Peanuts) & Alive(Bill) &
-Eats(Bill, x) | Eats(Sue, x)

because $A x, E y, \text{Person}(y) \& \text{Eats}(y, x) \& \text{-Killedby}(y, x) \Rightarrow \text{Food}(x)$ can be changed to:
 $\text{-(Person}(y) \& \text{Eats}(y, x) \& \text{-Killedby}(y, x)) \mid \text{Food}(x)$
 $= \text{-Person}(y) \mid \text{-Eats}(y, x) \mid \text{Killedby}(y, x) \mid \text{Food}(x)$

c. Proof that John likes peanuts given knowledge base:

$\text{(- Food}(x) \mid \text{Likes}(\text{John}, x)) \&$
Food(Apple) &
Food(Chicken) &
-Eats(y, x) | Killedby(y, x) | Food(x) &
- Killedby(y, x) | -Alive(y) &
Eats(Bill, Peanuts) & Alive(Bill) &
-Eats(Bill, x) | Eats(Sue, x) &

Unification on $\text{(-Eats}(y, x) \mid \text{Killedby}(y, x) \mid \text{Food}(x))$ with Eats(Bill, Peanuts) by $\{x/\text{Peanuts}, y/\text{Bill}\}$:
 $\text{-Eats}(\text{Bill}, \text{Peanuts}) \mid \text{Killedby}(\text{Bill}, \text{Peanuts}) \mid \text{Food}(\text{Peanuts})$

Resolution with Eats(Bill, Peanuts):

$\text{Killedby}(\text{Bill}, \text{Peanuts}) \mid \text{Food}(\text{Peanuts})$

Unification with $\text{- Killedby}(y, x) \mid \text{-Alive}(y)$ by $\{x/\text{Peanuts}, y/\text{Bill}\}$:

$\text{-Killedby}(\text{Bill}, \text{Peanuts}) \mid \text{-Alive}(\text{Bill})$

Resolution with Alive(Bill):

$\text{-Killedby}(\text{Bill}, \text{Peanuts})$

Resolution with Killedby(Bill, Peanuts) | Food(Peanuts):

$\text{Food}(\text{Peanuts})$

Unification with $\text{- Food}(x) \mid \text{Likes}(\text{John}, x)$ by $\{x/\text{Peanuts}\}$:

$\text{- Food}(\text{Peanuts}) \mid \text{Likes}(\text{John}, \text{Peanuts})$

Resolution with Food(Peanuts):

$\text{Likes}(\text{John}, \text{Peanuts})$

d. What food does Sue eat:

Unification on Eats(Bill, Peanuts) with $\text{-Eats}(\text{Bill}, x) \mid \text{Eats}(\text{Sue}, x)$ by $\{x/\text{Peanuts}\}$

$\text{-Eats}(\text{Bill}, \text{Peanuts}) \mid \text{Eats}(\text{Sue}, \text{Peanuts})$

Resolution with Eats(Bill, Peanuts):

$\text{Eats}(\text{Sue}, \text{Peanuts})$

e. What food does Sue eat, with added knowledge: $A x, \text{-Eat}(x) \Rightarrow \text{Die}(x), \text{Die}(x) \Rightarrow \text{-Alive}(x), \text{Alive}(\text{Bill})$

New knowledge base:

(\neg Food(x) | Likes(John, x)) &

Food(Apple) &

Food(Chicken) &

\neg Eats(y, x) | Killedby(y, x) | Food(x) &

\neg Killedby(y, x) | \neg Alive(y) &

(Eats(x, y) | Die(x)) &

(\neg Die(x) | \neg Alive(x)) &

Alive(Bill) &

\neg Eats(Bill, x) | Eats(Sue, x)

Unification on Alive(Bill) and \neg Die(x) | \neg Alive(x) with {x/Bill}:

\neg Die(Bill) | \neg Alive(Bill)

Resolution with Alive(Bill):

\neg Die(Bill)

Unification on Eats(x, y) | Die(x) and \neg Die(Bill) with {x/Bill, y/k} where K is a skolem constant:

Eats(Bill, K) | Die(Bill)

Resolution with \neg Die(Bill):

Eats(Bill, K)

Unification on \neg Eats(Bill, x) | Eats(Sue, x) with {x/K}

\neg Eats(Bill, K) | Eats(Sue, K)

Resolution with Eats(Bill, K):

Eats(Sue, K) ie, Sue eats something