# Rebalancing Shards in Clickhouse

for self-hosted OSS clusters

By Shivji Kumar Jha & Pranav Mehta

### Safe Harbour Statement

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## ABOUT US



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- Interests: Databases, Streaming, Infra, App Backends
- Contributed code to MySQL, Pulsar, Clickhouse
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### Love Clickhouse

Simple (single executable) and speaks SQL

Fast, Very fast (see clickbench). Written in C++ like DBs should be ©

Huge adoption & a large community of users & contributors

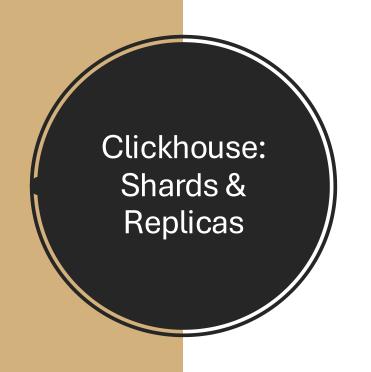
Very popular (Github: 41.6K \*, 7.4K forks)

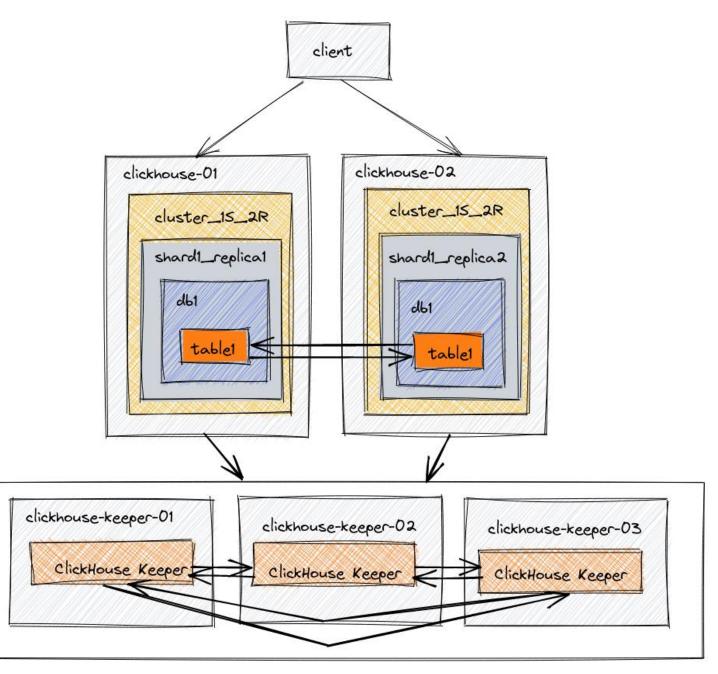
Flexible: Permissible, lots of configs, choices of deployments etc

Tons of integrations (Engines, views, I/O formats etc)

And.. Growing fast!

# Rebalancing Shards





https://clickhouse.com/docs/architecture/replication



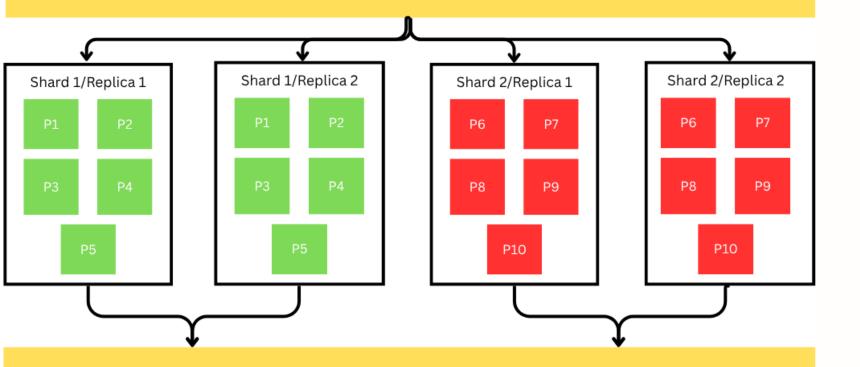
# DATA REBALANCING



- **Avoid hotspots** Prevent some nodes from becoming overloaded while others sit idle.
- Optimal use of resources evenly spread disk usage, faster querying
- **Maintain performance** Ensure consistent query speed and throughput.
- Enable scaling Distribute data evenly when adding or removing nodes

# Rebalancing Shards in Clickhouse

#### Distributed Table



Clickhouse Keeper

### Clickhouse Deployment

- 4 node cluster
  - 2 shard
  - 2 replica
- ReplicatedMergeTree
- Clickhouse Keeper

### Distributed Table Shard 1/Replica 2 Shard 2/Replica 2 Shard 1/Replica 1 Shard 2/Replica 1 P7 P7 P6 Р9 Р9 P8 P8 P10 P10 Clickhouse Keeper

# **Current Options**

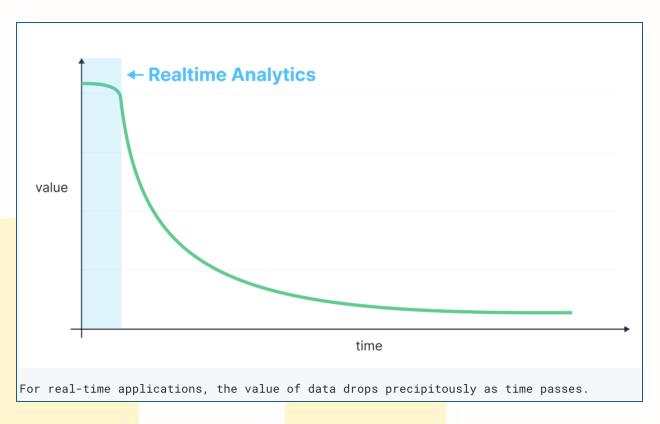
# Option 0: Use clickhouse cloud



- All data on S3 or object storage of choice
- Compute-compute segregation
- If you use the self-hosted version?
  - You can get quite far without needing any of this!
  - there are some other options

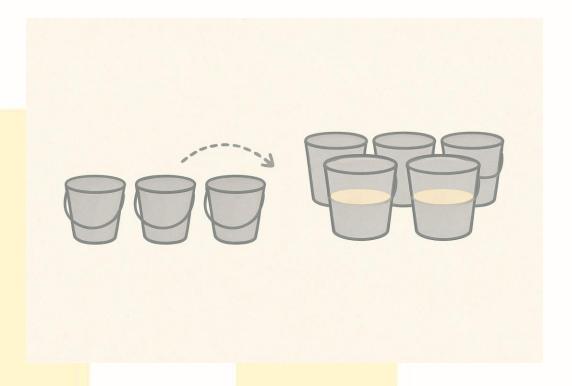
# Option 1: Let TTL do its magic over time

- Do Nothing!
- Realtime data looses value over time.
  - Eg: supply chain tracking, fraud detection, live call centre feedback etc
- If you have a time series:
  - Old data expires & rebalances over time.
  - Old data is rolled up into summary (MVs)



# Option 2: Move data to new fat cluster

- Export data using INSERT FROM SELECT into new cluster
  - Assuming you have extra resources for 2 clusters
  - Pros: Works even without partitioning.
  - Cons:
    - Poor performance on large datasets;
    - High IO and network overhead



# Option 3: Write to new shard temporarily

 Modify the cluster to write exclusively to the new shard until balance is restored

Pros: Easy

• Cons:

Causes potential load imbalance;

 Doesn't rebalance existing data. Queries may still be skewed



# Option 4: Skew ingest to new shards

- Adjust the distributed table's shard weights to bias new writes to a new shard.
  - Pros: Easy; no need to change write targets.
  - Cons:
    - Causes potential load imbalance
    - Doesn't rebalance existing data. Queries still skewed



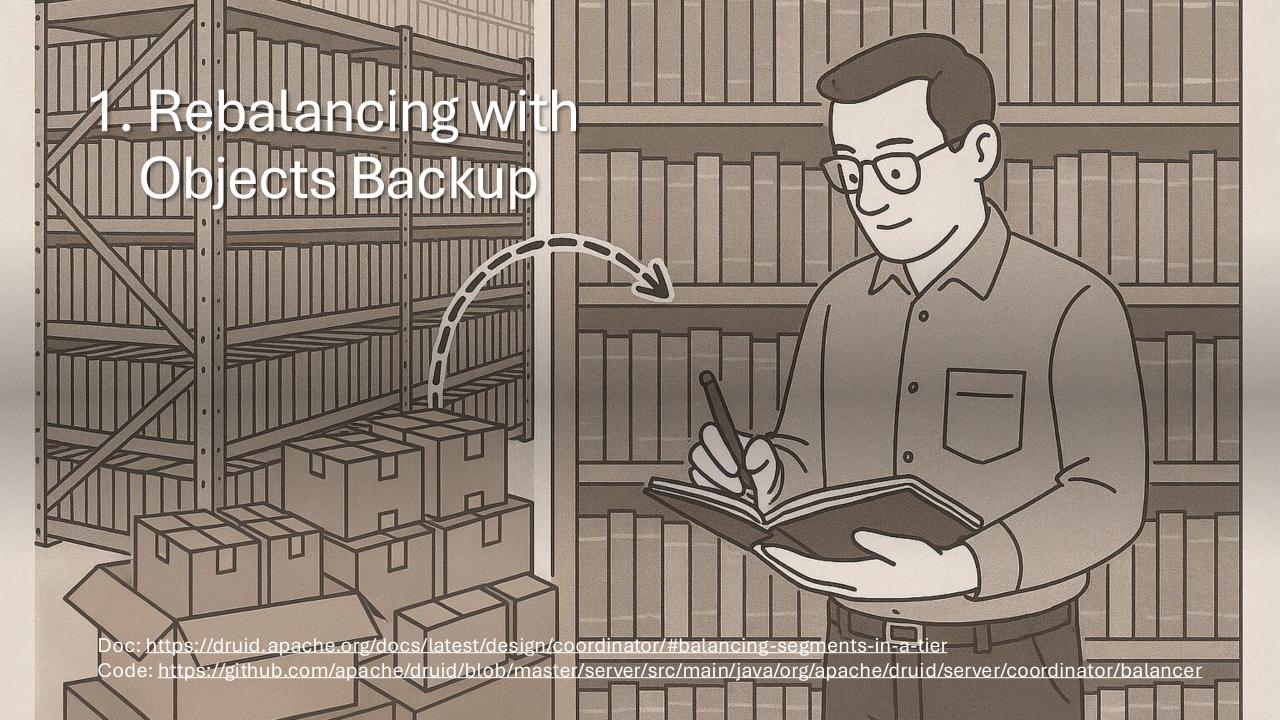
# Option 5: Move Partitions Manually

- For existing partitioned data,
  - Detach partition from loaded node
  - Move them to the new node, and
  - Reattach to the less loaded destination
- Pros:
  - More precise; avoids rewriting all data.
- Cons:
  - Manual and operationally intensive.
  - Inevitable maintenance window with partial responses / downtime.



Inspirations from the neighborhood





#### 1. Triggering the Balancing Process

- The Druid Coordinator ( DruidCoordinator. java ) periodically runs coordination cycles.
- · During each cycle, it checks if segment balancing is needed based on metrics or configuration.
- It invokes a BalancerStrategy to determine which segments to move and where.

#### **11** 2. Segment Selection

- Class Involved: SegmentToMoveCalculator.java
- Role: Applies rules to decide which segments should be eligible for moving (based on size, usage, etc.).
- May use ReservoirSegmentSampler for sampling if the dataset is large.

#### 3. Choosing a Strategy

- Factory Pattern Used: Strategy factories such as:
  - CachingCostBalancerStrategyFactory
  - CostBalancerStrategyFactory
  - RandomBalancerStrategyFactory
- These produce the actual strategy class implementing BalancerStrategy.

#### 4. Evaluating Segment Costs

- Key Class: CostBalancerStrategy.java (or CachingCostBalancerStrategy)
- Uses:
  - ServerCostCache: Caches how "expensive" it is to add segments to a server.
  - SegmentsCostCache: Maintains cost metrics for individual segments.
- Objective: Minimize movement cost while balancing load across Historical nodes.

#### **6** 5. Picking Target Server

- Method: findNewSegmentHomeBalancer() inside strategy class.
- Evaluates all available servers and chooses the best one for the segment based on:
  - Disk usage
  - Load rules
  - Balancing cost

#### **3** 6. Executing Moves

- Once a move decision is made, the Coordinator enqueues it as a segment move task.
- The Historical node will load the new segment from deep storage or peer, and drop it from the old node.

#### ▶ 7. Tier-Aware Balancing

- Class: TierSegmentBalancer.java
- If configured, segments are balanced within their designated tier (e.g., hot/warm/cold storage tiers).
- · Helps align rebalancing with resource priority.

#### **☑** 8. Post-Move Updates

- Once the segment is loaded and acknowledged by the target Historical node:
  - Coordinator updates the metadata.
  - Caches (ClusterCostCache) are updated for the next decision cycle.



#### # 1. SimpleLoadManagerImpl

- . What it is: The original, now-deprecated load balancer.
- Responsibilities:
  - Collects basic broker metrics (CPU, memory, I/O).
  - Implements straightforward bundle assignment/unloading logic.
- Limitations:
  - · Centralized and coarse-grained.
  - Lacks advanced splitting/unloading strategies pulsar.apache.org +9 .

#### 2. ModularLoadManagerImpl

- Role: The modern default implementation.
- Core features:
  - Modular strategy interface ( Modular Load Manager Strategy ): allows plug-in or configuration of balancing strategies pulsar.apache.org +3 .
  - Local broker data: stats on CPU, direct memory, bandwidth, bundle counts ( LocalBrokerData ).
  - Historical data: sliding-window averages of usage ( TimeAverageBrokerData , BundleData ).
  - Leadership election: only leader broker (written to ZooKeeper /loadbalance/leader) makes assignment/split/unload decisions.
  - Bundle strategies:
    - Assignment: e.g., least long-term message rate + resource weighting pulsar.apache.org +9 .
    - Splitting criteria: topic count, message rate, size thresholds.
    - · Shedding logic: offload high-load bundles to underutilized brokers.

#### 3. ModularLoadManagerStrategy and LeastLongTermMessageRate

- Interface: defines method selectBrokerForBundle(...).
- **Default implementation**: LeastLongTermMessageRateStrategy:
  - Normalizes brokers by message rate and max resource utilization.
  - Excludes overloaded brokers beyond configured threshold pulsar.apache.org pulsar.apache.org .
  - Ensures bundles are distributed to equalize long-term load.

#### 4. ExtensibleLoadManagerImpl (in extensions)

- Purpose: Pulsar 3.x+ evolution of the load balancer.
- Key differences pulsar.apache.org +9 :
  - Decentralized metadata store: uses system topics + table views instead of ZooKeeper.
  - Local broker/table views: each broker publishes top-k bundle loads.
  - Distributed assignment and splitting: not just leader-based; tasks delegated locally.
  - Persistent bundle ownership: replicated via system topics.
  - Improved upgradeability: modular code, better debug metrics, reduced ZK dependencies.

#### 5. Supporting classes for both implementations

- LoadData: Aggregates broker + bundle stats.
- BrokerData / LocalBrokerData : Real-time broker usage metrics.
- TimeAverageBrokerData / BundleData: Windowed history used for steady-state decisions.
- LoadManagerContextImpl: Glue between Broker and LoadManager: triggers updates, split/unload tasks.
- Strategy implementations: found in org.apache.pulsar.broker.loadbalance.impl.strategies.

#### 6. Client-visible components & configuration

- ZooKeeper nodes (/loadbalance/brokers/..., /loadbalance/leader):
  - Used by modular load balancer.
- Broker config options ( broker.conf ):
  - loadManagerClassName: choose between Simple, Modular, Extensible.
  - · Splitting/unload strategies: thresholds, intervals etc.
- Admin APIs/tools:
- Check "loadManagerClassName" dynamically.
- Manually trigger unloads/splits.

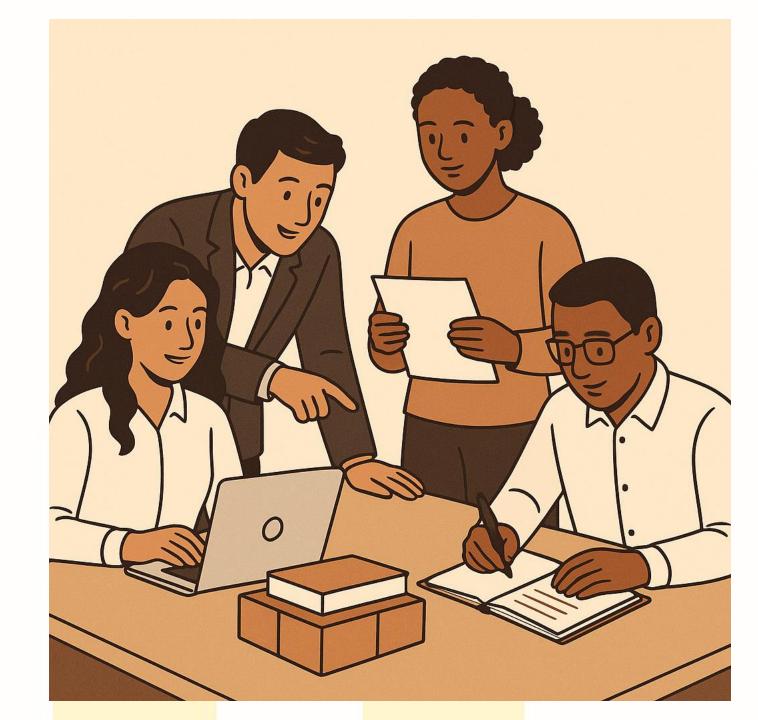
#### 3 7. Bundle lifecycle & orchestration

- . Bundle ownership: tracked via ZK (modular) or system topic (extensible)
- · Unloading & splitting:
  - Triggered when metrics exceed thresholds.
  - Splits create new bundles and redistribute.
  - Ownership change via topic close/open → transparent to producers/consumers pulsar.apache.org +2
     pulsar.apache.org +9 pulsar.apache.org +1 pulsar.apache.org +4 github.com +1 pulsar.apache.org .
- Bundle-state channel (extensible):
- Finite-state machine for split/unload ops.
- Broadcast over system topic for consistency streamnative.io

# The (ideal) rebalancing algo:

- Keeps data spread in a way that help parallelize queries
  - If two parts are more likely to be picked by a query, keep apart.
- Evaluates cost based on:
  - Server capacity (% disk, cpu, ram, network i/o etc in use)
  - Demand for data: read/write throughput etc
  - Recent movements so we don't keep moving same part over and over.
  - Historical averaged load data
- An assortment of algos to choose from!
  - And maybe an easy plug-and-play to add custom rebalance logic?

# Clickhouseongoing work



# Self-balancing architecture



#### Proposal:

https://github.com/ClickHouse/ClickHouse/issues/13574



#### **Introduces live part movement**

ALTER TABLE table\_name MOVE
PARTITION|PART partition\_expr TO
SHARD 'shard\_id'

# Self-balancing architecture

- Mark part to be moved with a unique identifier
- Copy data to the destination shard
- Remove data from source shard
- When reading data, use the unique identifier to discard duplicate part
- Open Issues #49574, #40814, #40189, #40188
- Disabled!
- nvartolomei.com/rebalancing-data-shared-nothing

# Self-balancing architecture – Challenges

- Race condition leading to duplicate data in SELECT #49574
  - proposal made #50777
- Tasks Cleanup- #40814
- Data loss along with zero copy replication #40189, #40188
- Assigning UUID to existing parts
- Disabled!

# Self-balancing architecture – Our efforts

- https://github.com/ClickHouse/ClickHouse/issues/65633
- https://github.com/ClickHouse/ClickHouse/pull/68020

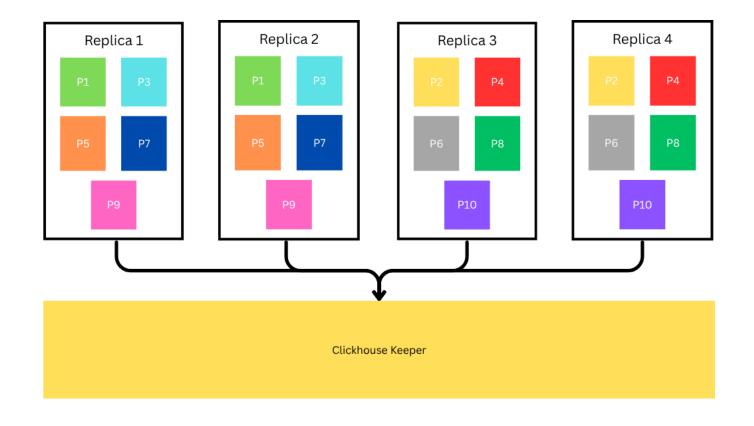
Shout out: <u>Tarun Annapareddy</u>

# **Cl**ickhouse Proposal

# Distributed Table Shard 1/Replica 1 Shard 1/Replica 2 Shard 2/Replica 2 Shard 2/Replica 1 Р9 P10 P10 Clickhouse Keeper

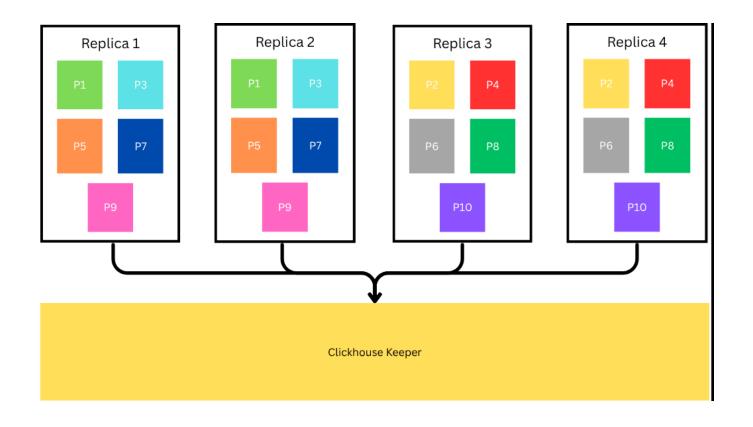
### ReplicatedMergeTree

- 2 shards
- 2 replicas



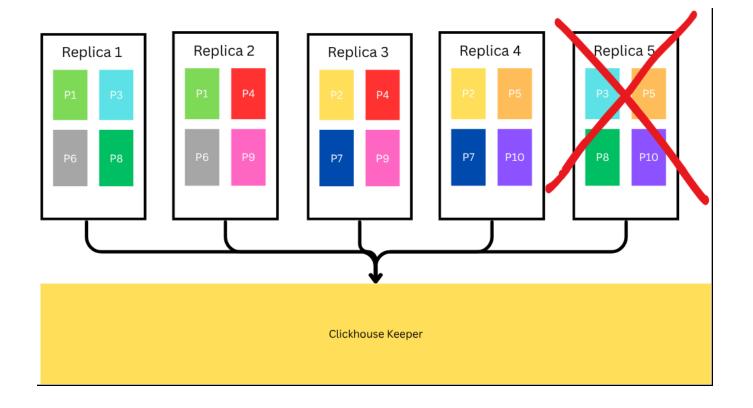
# ReplicatedMergeTree Extended

• Replication Factor: 2



# ReplicatedMergeTree Extended

- Replication Factor: 2
- Rebalancing after adding a Node



# ReplicatedMergeTree Extended

- Replication Factor: 2
- Rebalancing after removing a Node

# Open PR

- Basic SELECT/INSERT support
- Basic re-sharding support with a naïve algorithm
- Table Settings
  - Cluster
  - cluster\_replication\_factor
- Query Settings
  - cluster\_query\_shards
- Queries
  - SYSTEM SYNC REPLICA CLUSTER
  - SYSTEM DROP CLUSTER REPLICA
- System Table
  - system.cluster\_partitions

# Our Experiments

Tailored to our needs



### Our Usecase

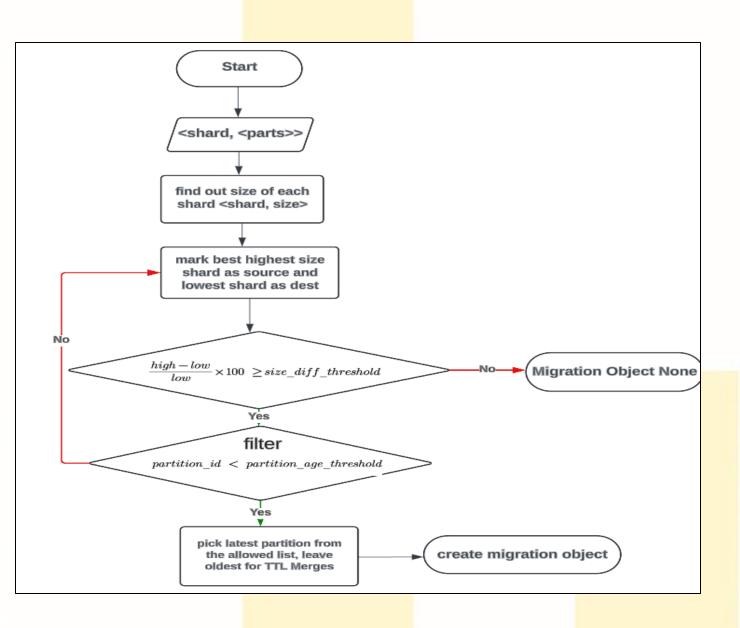
- Lots of on-prem deployments
- limited access
- Zero touch upgrades
- T-shirt size deployments
- Each size has fixed number of shards
- Bigger size means more clickhouse shards
- Re-balance during size upgrade



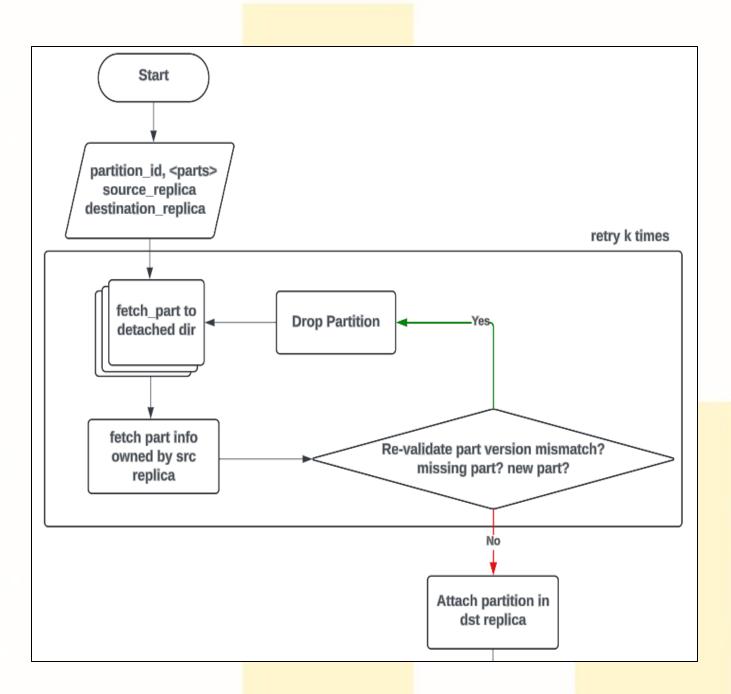
# An external utility

- Select source/destination shards for migration
- 2. Select partition least likely to be updated from source shard
- 3. Fetch parts from source shard
- 4. Attach partition to destination shard
- 5. Drop partition on source shard





## **Partition Selection**



# Optimistic Migration

# Challenges

- Expensive Migration Retries
  - https://github.com/ClickHouse/ClickHouse/issues/66408
- Duplicates in SELECT
  - Can be solved with MOVE PARTS TO SHARD feature
- Works with limited partition key expression

### Future Work

- For script-based rebalancing
  - Extend SYSTEM STOP MERGES to stop merges at PART|PARTITION level
  - Harden MOVE TO SHARD feature
- For fully automatic rebalancing
  - Implement Clickhouse rebalancing proposal
  - Extend K8s operator to rebalance data
  - Lets Collaborate!

# QUESTIONS

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#### Shivji Kumar Jha

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- tinyurl.com/shiv-slides

# Thank You!