

ANALYSIS OF ASSEMBLY CODE OF BASIC C PROGRAM (GCC 9.3.0)

Concepts of Programming Languages
Assignment -4

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10th April 2021

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*** C Program used for demonstration ***

- * Key features included in Program: For and While loop, Conditional Statement Switch Statement and assignment
- </> Below is the Code, it is just for demonstration purpose, switch case 1 calculates occurrences of 5 in array using linear search and case 2 finds sum of all array elements.

```
#include <stdio.h>
     int main() {
         int key=2;
         int i=0;
         int arr[10]={5,1,2,3,5,9,6,7,5,5};
         switch(key) {
             case 1:
                 int count=0;
                 for(i=0;i<10;i++){
11
                      if(arr[i]==5){
12
                          count++;
13
                      }
                 if(count>0)
15
                      printf("5 is found %d times.\n",count);
17
                 else
21
                      printf("5 not Found in array\n");
                 break;
             case 2:
                 int sum=0;
                 while(i<10) {
                      sum=sum+arr[i];
                      i++;
                 printf("Sum of all array elements : %d\n",sum);
                 break;
             default :
                 printf("Enter correct choice\n");
                 break;
```

*** Assembly Program ***

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Generated Assembly equivalent by compiler using the command gcc -S File_name.c

which generated a "File_name.s" file containing our required assembly code. The same is shown and explained briefly below in parts.

</> 1

Below shown snippet is the very beginning of the file, it contains all typed strings in printf statements, these will be used later when printf is called, also indicates main function below

```
.file
            "BT19CSE019_CPL_A4.c"
    .text
    .section
                .rodata
                                                             ; strings used in printf
.LC0:
    .string "5 is found %d times.\n"
    .string "5 not Found in array"
    .align 8
.LC2:
    .string "Sum of all array elements : %d\n"
    .string "Enter correct choice"
    .text
    .globl main
   .type main, @function
```

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```
main:
                                                           ; Main function
.LFB0:
   .cfi_startproc
   endbr64
   pushq %rbp
   .cfi_def_cfa_offset 16
   .cfi_offset 6, -16
   movq %rsp, %rbp
   .cfi_def_cfa_register 6
   subq $64, %rsp
           %fs:40, %rax
   movq
           %rax, -8(%rbp)
   movq
         %eax, %eax
   xorl
   movl
           $0, -64(%rbp)
           $5, -48(%rbp)
   movl
   movl
           $1, -44(%rbp)
           $2, -40(%rbp)
   movl
   mov1
           $3, -36(%rbp)
           $5, -32(%rbp)
   mov1
   movl
           $9, -28(%rbp)
           $6, -24(%rbp)
   mov1
           $7, -20(%rbp)
   movl
           $5, -16(%rbp)
   mov1
           $5, -12(%rbp)
$1, -52(%rbp)
   movl
   cmp1
   cmpl $2, -52(%rbp)
   jmp .L15
```

This part of code is the crux of the Main function's control flow, as mentioned in comments frame pointer (%rbp) is pushed onto program stack activation record and required offsets are also set, stack pointer (%rsp) is set to frame pointer.

After that required size according to data is allocated on stack which is 64 in this case, essential data like return address are stored.

Data initialising starts then, starting with integer key used for toggling switch statement which is set to value to 2 (just for demonstration) initially.

Key is at -52(%rbp) 52 here denotes offset and negative sign is because stack grows downwards.

Next integer variable initialised is i, which is used as loop counter is set to value 0, it is present at -64(%rbp), and finally our array of 10 integers is stored as continuos blocks starting from-48(%rbp). Then key is compared to 1, if it is equal to 1 then program control jumps to .L2, else if key is equal to 2, program control jumps to case 2 at .L3, else default is executed in switch and jumps to .L15 to print a statement.

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```
; Code under switch case 1
       $0, -60(%rbp)
mov1
movl $0, -64(%rbp)
jmp .L5
       -64(%rbp), %eax
mov1
                                           ; Converting %eax to sign extended quad word into %rax
clta
mov1
        -48(%rbp,%rax,4), %eax
cmpl
       $5, %eax
jne .L6
add1
       $1, -60(%rbp)
       $1, -64(%rbp)
```

When switch case 1 is used, this part of code works namely labels with L2,L7,L6,L5,L8 As mentioned in comments first it initialises a new count integer to 0 at -60(%rbp) and starts our loop. wherein iterator integer i is set to 0 again and control is transferred to L5 which is discussed below.

L7: it stores i in a temporary register and then sign extends that to a quad word word in register %rax, now frame pointer and offset it access array first element and then takes i from %rax and increments i*4 to address to obtain arr[i] value.

This is stored in register %eax again, now as per algorithm it compares arr[i] in %eax with \$5 and then based on result jump is conditionally called

i.e. if jne (jump if not equal) is called is arr[i]!=5, this transfers control to L6 which increments i. Else i.e. if jne is not called count is incremented.

Variables offset : name i -64(%rbp) i count count sum sum key -52(%rbp) -56(%rbp) arr (arr[0]) -48(%rbp)

```
$9, -64(%rbp)
cmp1
jle .L7
cmpl
        $0, -60(%rbp)
ile .L8
        -60(%rbp), %eax
mov1
mov1
        %eax, %esi
        .LC0(%rip), %rdi
                                           ; loading address from .LC0 (String is stored) to %rdi
lead
movl
        $0, %eax
       printf@PLT
call
jmp .L10
       .LC1(%rip), %rdi
       puts@PLT
call
jmp .L10
```

As discussed above here we see the code under label L5, where comparison for loop is taking place, at start it compares i with 9, if i<=9 then control is passed on to loop body at .L7.

After loop terminates, count is compared to 0, if count<=0 then control is shifted to .L8 where string from .LC1 is loaded into %rdi and puts (works as printf) is called and then unconditionally jump to .L10 is called indicating label where break statement is executed. If count is not <=0 then, count is stored in %eax and is sent as parameter to printf next along with string at LC0.

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```
.LC1(%rip), %rdi
                   puts@PLT
          jmp .L10
          movl
                   $0, -56(%rbp)
          jmp .L11
                    -64(%rbp), %eax
          mov1
          cltq
          mov1
                    -48(%rbp,%rax,4), %eax
          addl
                   %eax, -56(%rbp)
                  $1, -64(%rbp)
                   $9, -64(%rbp)
          cmpl
89
                                                             ; moving variable sum value to %eax, parameter to printf; moving %eax to %esi, passing as parameter to print
          mov1
                   %eax, %esi
                  .LC2(%rip), %rdi
          mov1
                   $0, %eax
                   printf@PLT
           jmp .L10
```

Here, L3 is shown which is first called when switch is under case 2, and it initialises integer sum to 0 at -56(%rbp) and unconditional transfer happens to L11 where while loop body is present.

In While loop (L11) first 'i' is compared with 9, if i<=9 loop is executed else terminated. Inside loop body L12 where our main algorithm works, in L12 'i' is temporarily stored in %eax and again sign extended to quad word in %rax, and now as we have seen in for loop similar access to array element is made using offset and frame pointer and arr[i] is stored into %eax, now arr[i] is added to sum and i is incremented.

After loop terminates sum is stored in %eax ad sent as parameter for printf along with string from .LC2 (obtained from address: leaq) and then printf is called and control transferred to L10 which is break statement.

```
.L15:
   leaq
            .LC3(%rip), %rdi
                                                ; No Parameter print hence calling puts standard function
   call
            puts@PLT
   nop
.L10:
   mov1
           $0, %eax
           -8(%rbp), %rdx
   movq
           %fs:40, %rdx
   xorq
                                                ; JUMP if equal to .L14 to return from Main function
   je .L14
   call
           __stack_chk_fail@PLT
                                                ; calls for function which checks Program Stack Fail
.L14:
   leave
   .cfi_def_cfa 7, 8
   ret
   .cfi_endproc
```

Now, L15 here is for default case in switch which prints a statement stored in LC3 and then puts function is called which works like printf and at the end nop indicates no operation. In L10, which is our break statement gets the return address from activation record and compares using XOR and then conditionally L14 is called which returns control to address pointed by return address and ends program process.

Else if jump L14 doesn't take place call to function is made which checks program stacks failure.

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```
main, .-main
.ident "GCC: (Ubuntu 9.3.0-17ubuntu1~20.04) 9.3.0"
.section .note.GNU-stack,"",@progbits
.section .note.gnu.property,"a"
.align 8
       1f - 0f
.long
        4f - 1f
.long
.long
.string "GNU"
.align 8
        0xc00000002
.long
        3f - 2f
.long
.long
.align 8
```

This data is regarding the compiler options and its details.