

# STACK

Q) Design a stack which supports push, pop, peek, isEmpty & minEle. And the time complexity of each of these operations should be  $O(1)$

Ex:

```
push(5)
push(1)
push(3)
peek() → 3
minEle → 1
pop()
minEle → 1
pop()
minEle → 5
```

We have a stack,  $A = [5\ 2\ 1\ 7]$  &  $B = []$

Taking a variable, minele = infinity, we will pop (7) from A and push (7) to B. We will check if the minimum value between infinity and 7. So the value of minele = 7 as  $7 < \text{infinity}$ .

Now, we will pop (1) from A and push (1) to B. We will check for minimum element between 1 & 7. The value of minele = 1.

This process will be done for remaining two elements. For this, the time complexity would be  $O(2N)$ , as we empty A to B and then return the value from B to A. So for A to B it is 'n' and from B to A it will be 'n'.

## How to get this done in a time complexity of $O(1)$ ?

We will have 2 stacks. For A we will push (5) and will do the same for B as well.

When 1 is pushed in A and B, minEle will become 1 as  $5 > 1$ .

We will push(3) in A. now the minEle = 1 in B, as  $3 > 1$ , so we will push (1) again in stack B.

For the peek value, we will check stack A and for minEle we will check for stack B.

To find minEle, all we have to do is check the peek() of B.

# LeetCode:155. Min Stack

## Input

["MinStack", "push", "push", "push", "getMin", "pop", "top", "getMin"]

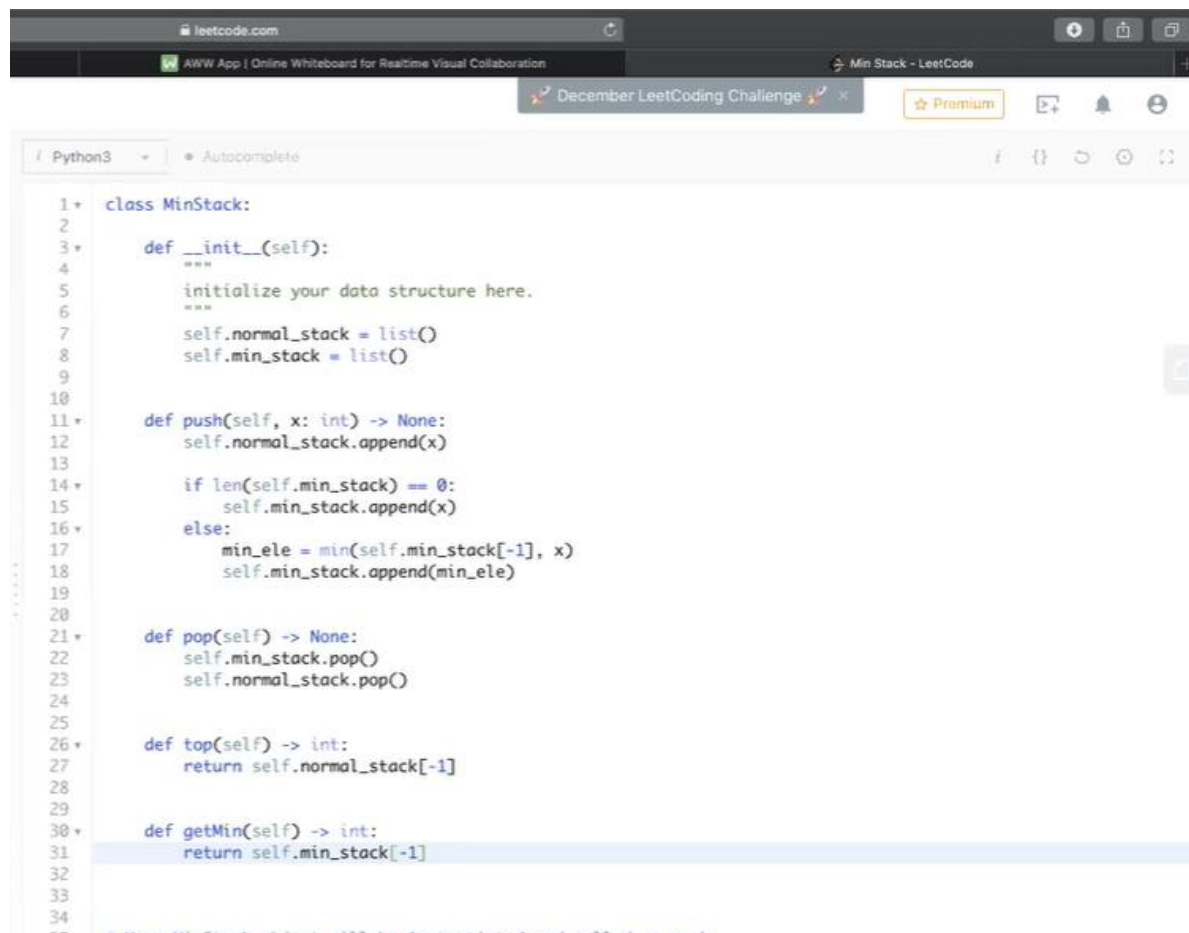
[[], [-2], [0], [-3], [], [], [], []]

## Output

[null, null, null, null, -3, null, 0, -2]

## Explanation

```
MinStack minStack = new MinStack();
minStack.push(-2);
minStack.push(0);
minStack.push(-3);
minStack.getMin(); // return -3
minStack.pop();
minStack.top();    // return 0
minStack.getMin(); // return -2
```



The screenshot shows a web browser window with the LeetCode website. The page title is "Min Stack - LeetCode". There is a "December LeetCode Challenge" banner and a "Premium" badge. The main content is a code editor for Python3. The code defines a class `MinStack` with the following methods:

```
1 class MinStack:
2
3     def __init__(self):
4         """
5         initialize your data structure here.
6         """
7         self.normal_stack = list()
8         self.min_stack = list()
9
10
11     def push(self, x: int) -> None:
12         self.normal_stack.append(x)
13
14         if len(self.min_stack) == 0:
15             self.min_stack.append(x)
16         else:
17             min_ele = min(self.min_stack[-1], x)
18             self.min_stack.append(min_ele)
19
20
21     def pop(self) -> None:
22         self.min_stack.pop()
23         self.normal_stack.pop()
24
25
26     def top(self) -> int:
27         return self.normal_stack[-1]
28
29
30     def getMin(self) -> int:
31         return self.min_stack[-1]
32
33
34
35 # Your MinStack object will be instantiated and called as such:
```

```
class MinStack:

    def __init__(self):
        """
        initialize your data structure here.
        """
        self.normal_stack = list()
        self.min_stack = list()

    def push(self, x):
        self.normal_stack.append(x)

        if len(self.min_stack) == 0:
            self.min_stack.append(x)
        else:
            min_ele = min(self.min_stack[-1], x)
            self.min_stack.append(min_ele)

    def pop(self):
        self.normal_stack.pop()
        self.min_stack.pop()

    def top(self):
        return self.normal_stack[-1]

    def getMin(self):
        return self.min_stack[-1]
```

We take two variables 'normal\_stack' and 'min\_stack'. In normal stack we will append all elements, but in min\_stack we will only append the minimum value.

## Find the next greater element in right for each element.

[2, 1, 7, 4, 6, 8, 1, 9]

For

2 → 7

1 → 7

7 → 8

4 → 6

8 → 9

1 → 9

9 → None

[7, 7, 8, 6, 8, 9, 9, None]

### Brute Force:

```
for i in range(n-1):
    for j in range(i+1,n):
        if A[j] > A[i]:
            print(A[j])
            break
```

### 2<sup>nd</sup> Approach:

We are 2 and stack is also empty, so we will add 2 to the stack. At 1, we have not yet found the greater element for 2 and 1 so we will push 1 as well into the stack.

At 7, we see that 7 is the next greater element of 1. So we will pop 1 and push 7. As 7 is also the next greater element of 2, we can pop 2 and push 7.

At 4, as  $7 > 4$ , we push 4 to stack as we need to find the next greater element for 4.

At 6.  $6 > 4$ , so we can pop 4 and push 6. Since  $6 < 7$  and we have to find the next greater element for 7 as well, we will push 6.

At 8,  $8 > 6$ , so we pop 6 and push 8.  $8 > 7$  so we will pop 7 and push 8.

The next element 1 is less than 8. So we need to find the next greater element for 8 and for 1.

At 9,  $9 > 1$ , so we pop 1 and push 9. As  $9 > 8$ , we pop 8 and push 9.

As there is no other element present, we will push None for 9.

```
def solve(A):
    n = len(A)
    stack = list()

    next_greater_element = [0] * len(A)
    # this give index and values as well.

    for idx, val in enumerate(A):
        if len(stack) == 0:
            stack.append(idx)
        else:
            cur = val

            while len(stack) != 0 and stack[-1] < cur:
                x = stack[-1]
                stack.pop()
                next_greater_element[x] = cur
            stack.append(idx)

    return next_greater_element

if __name__ == "__main__":
    A = [2, 1, 7, 4, 6, 8, 1, 9]
    ans = solve(A)
    print(ans)
```

