

GH RAISONI COLLEGE OF ENGINEERING

CAE - I

2020 - 2021 EVEN TERM

CAE - I EXAMINATION SUMMER - 2021 (ONLINE MODE)

DEPARTMENT: ARTIFICIAL INTELLIGENCE

SEMESTER/SECTION: 4<sup>th</sup> / A

DATE OF EXAMINATION: 13/02/2021

SUBJECT: GRAPH THEORY AND NUMBER SYSTEM

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CO1.

Q.1.

Simple Graph:

A simple graph  $(V, E)$  consists of a nonempty set representing vertices,  $V$ , and a set of unordered pairs of elements of  $V$  representing edges  $E$ .

A simple graph has:

→ no arrows

→ no loops

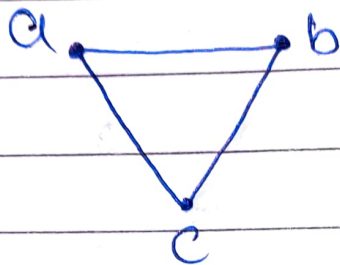
*Shivam*

Pg. (1)

→ cannot have multiple edges joining vertices.

Example:

$$V = \{a, b, c\}, E = \{\{a, b\}, \{b, c\}, \{a, c\}\}$$



Multigraph:

A multigraph is a set of vertices  $V$ , a set of edges  $E$  and a function:

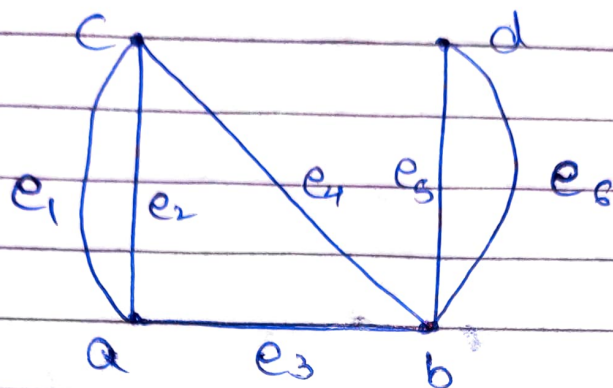
$$f: E \rightarrow \{\{u, v\} : u, v \in V \text{ and } u \neq v\}$$

Example:

$$V = \{a, b, c, d\}, E = \{e_1, e_2, \dots, e_6\},$$
$$f: E \rightarrow \{\{u, v\} : u, v \in V \text{ and } u \neq v\}$$

$e$	$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$
$f(e)$	$\{a, c\}$	$\{a, c\}$	$\{a, b\}$	$\{c, b\}$	$\{b, d\}$	$\{b, d\}$

~~Ans.~~ Pg. 2



Pseudograph :

A pseudograph is a set of vertices  $V$ , a set of edges  $E$  and a function  $f: E \rightarrow \{\{u, v\} : u, v \in V\}$ . If  $e \in E$  is such that  $f(e) = \{u, v\} = \{u\}$ , then we say  $e$  is a loop.

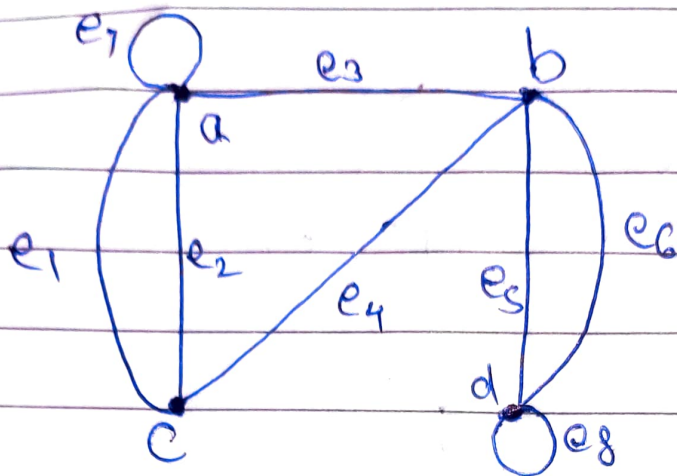
Example :

$V = \{a, b, c, d\}$ ,  $E = \{e_1, e_2, \dots, e_8\}$ ,  
 $f: E \rightarrow \{\{u, v\} : u, v \in V\}$

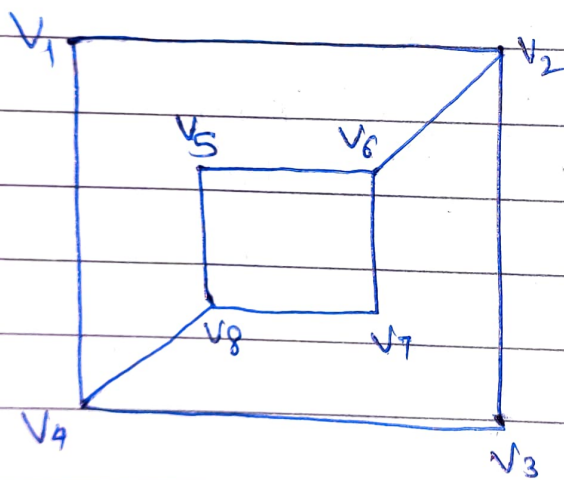
$e$	$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$
$f(e)$	$\{a, c\}$	$\{a, c\}$	$\{a, b\}$	$\{a, b\}$	$\{b, d\}$	$\{b, d\}$	$\{a, d\}$

$e$	$e_8$
$f(e)$	$\{a\}$

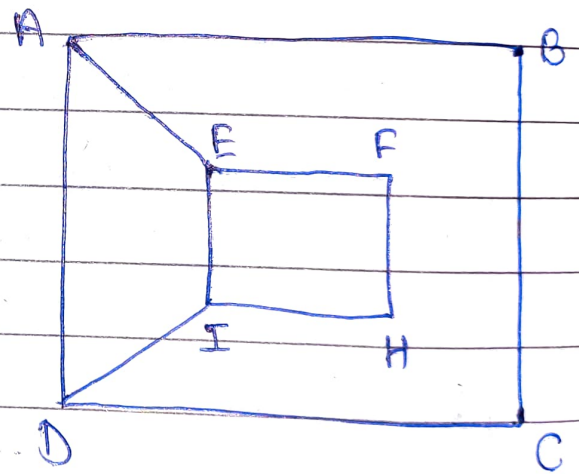




CO1.  
Q.2.



G



H

The graph G and H both have 8 vertices and 10 edges. They both have 4 vertices each of degree 3 and 4 vertices each of degree 2.

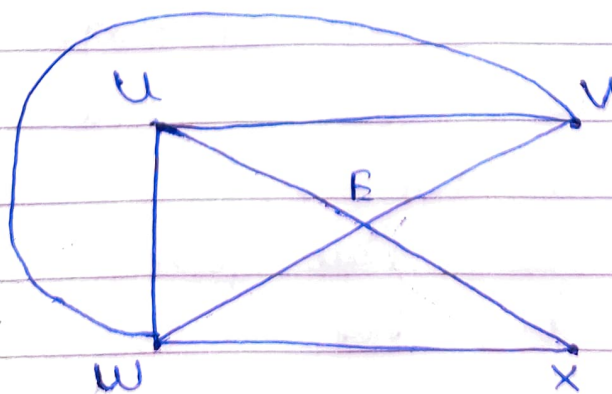
Now consider  $\text{degree}(v_1) = 2$  in  $G$ . Then  $v_1$  must correspond to either  $B, C, H, F$ . Since these are vertices of ~~deg~~ degree 2 in  $H$ .

However, each of these vertices in  $H$  is adjacent to another vertex of degree 2 in  $H$  viz  $B$  is adjacent to  $C$ ,  $F$  is adjacent to  $H$  but  $A$  is adjacent to  $v_2$  and  $v_4$  in  $G$  which are of degree 3. Thus the preservation of adjacency of the vertices is not maintained.

∴  $G$  and  $H$  are not isomorphic graphs.

Q.2.

Q.3. Given graph is :



~~Answer~~ Pg. no. 5



The vertex  $V = \{u, v, w, E, x\}$   
Edges  $E = \{(u, v), (u, w), (u, E), (E, x), (w, x), (v, w)\}$

Now, degree of each vertices are:

$$\deg(u) = 3$$

$$\deg(v) = 3$$

$$\deg(w) = 4$$

$$\deg(x) = 2$$

Now  $u$  and  $v$  have only two vertices of odd degree of 3.

According theorem, a connected graph contains a euler trail, but not an eulerian circuit, if and only if it has exactly two vertices of odd degree.

Thus, the path from by edge is.  
 $(u, v) - (v, w) - (w, u) - (u, x) - (x, w) - (w, v)$  is  
An eulerian trail.

Thus, the above ~~the~~ graph has  
eulerian trail but not eulerian circuit.