

Assignment no. 6

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Q.1.

1.

$S \rightarrow OBB$

$B \rightarrow OS/1S/0$

The PDA given as

$A = \{(q), (q, 1), (S, B, 0, 1), \delta, q, S, ?\}$

The production rule S can be

$R_1 - \delta(q, \epsilon, S) = \{(q, OBB)\}$

$R_2 - \delta(q, \epsilon, B) = \{(q, OS) \mid (q, 1S) \mid (q, 0)\}$

$R_3 = \delta(q, 0, 0) = \{(q, \epsilon)\}$

$R_4 = \delta(q, 1, 1) = \{(q, \epsilon)\}$

Testing 0100000 against PDA

$\delta(q, 0100000, S) = \delta(q, 0100000, OBB)$

$= \delta(q, 10000, BB)$ R_1

$= \delta(q, 10000, 1S B)$ R_3

$= \delta(q, 0000, SB)$ R_2

$= \delta(q, 0000, OBBB)$ R_1

$= \delta(q, 000, BBB)$ R_3

$= \delta(q, 00, OBB)$ R_2

$= \delta(q, 00, BB)$ R_3

$= \delta(q, 00, OB)$ R_2

$= \delta(q, 00, B)$ R_3

$$= \delta(q, 0, 0) \quad R_2$$

$$= \delta(q, \epsilon) \quad R_3$$

The 0100000 is accepted by PDA.

Q.2. Convert given CFG to CNF

$$G_1 = \{ S \rightarrow AB, S \rightarrow C, A \rightarrow a, B \rightarrow b \}$$

$$G_2 = \{ S \rightarrow AA, A \rightarrow a, B \rightarrow C \}$$

Step 1: Eliminate start symbol from RHS
 $S_1 \rightarrow S$

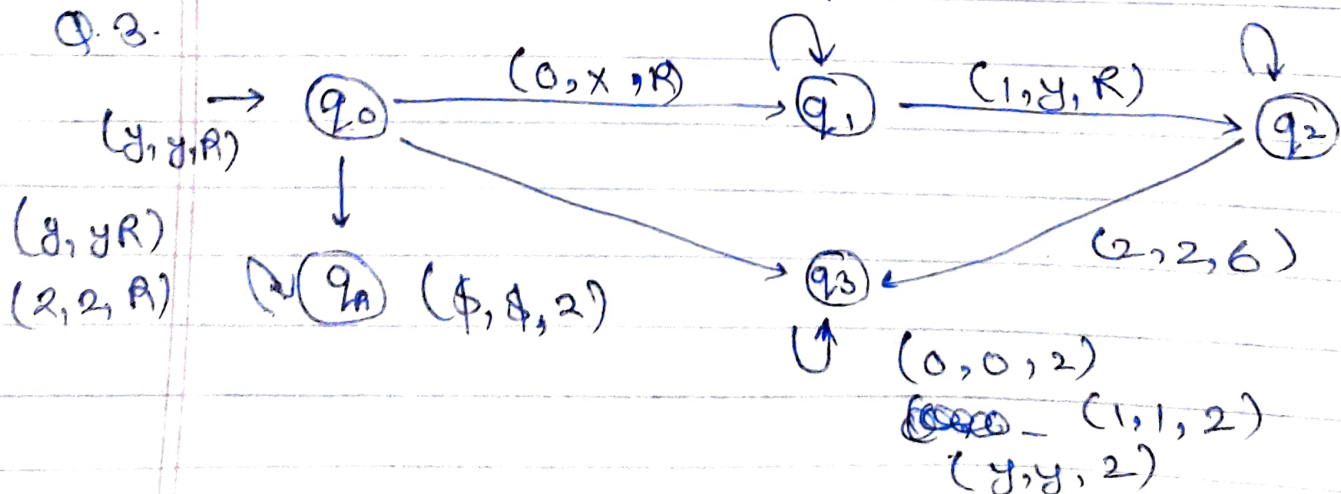
Step 2: In grammar, remove the null, unit and useless production.

Step 3: Eliminate terminals from the RHS of the production if they exist with other non-terminal or terminals

$$S \rightarrow RA \quad R \rightarrow a$$

Step 4: Eliminate RHS with more than two non-terminals.

$$S \rightarrow RS \quad R \rightarrow AS$$



To solve this approach firstly replace one 0 from front by 1, then keep moving right till we get one 1 & then replace this 1 by 0. Again keep moving right till we get 2, then replace it by 2 and move left.

Q4.

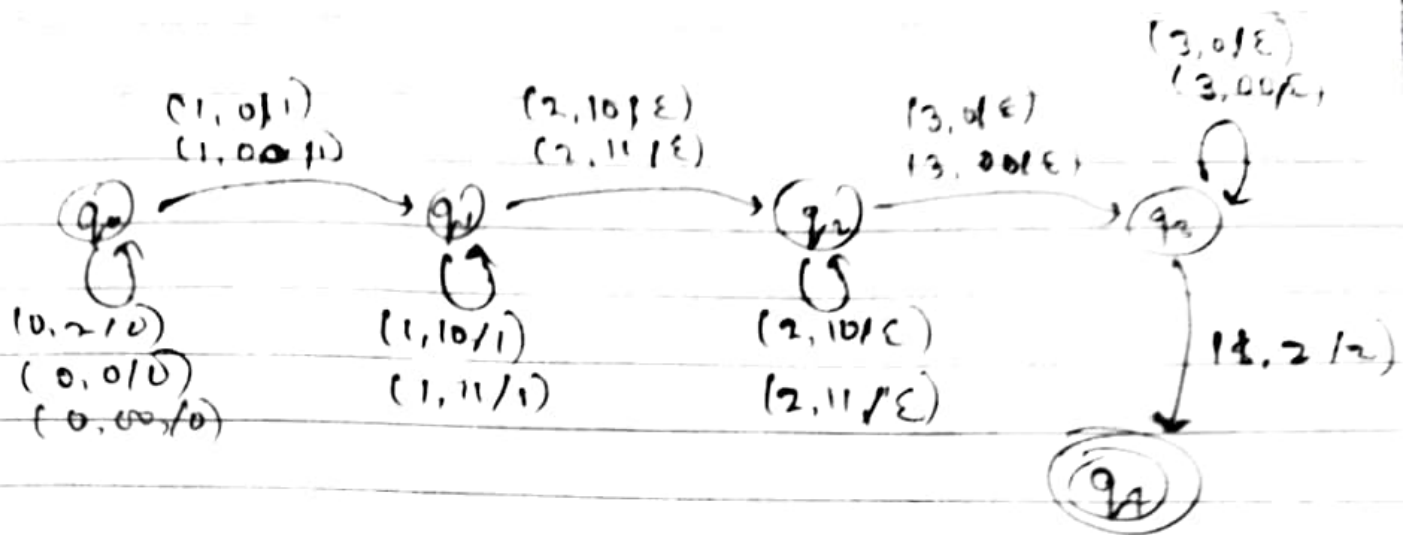
The Turing machine M can be constructed by following moves:

- Let q_1 be initial state
- Let if M is in q_1 on scanning '1', it enters the state q_2 and write B (blank)
- If M is in q_2 on scanning '1', it enters the state q_1 and write B (blank)
- From the above moves, we can observe that M enters the state q_1 if it scans an even number of x's and it enters q_2 if it scans odd number of x's. Hence q_1 is only accepting state.

Hence $M = \{ \{q_1, q_2\}, \{1\}, \{1, B\}, \delta, q_1, B, \{q_1\} \}$
 where δ is given by.

Tape Alphabet Symbol	Present state q_1	Present state q_2
1	$B R q_2$	$B R q_1$

Q5



Input : 001112233

Result : Accepted

Input : 000112233

Result : Not Accepted.