

Lecture-24 Recap

P①

Expectation of sum =
Sum of expectations.

→ coupon collecting problem

→ drunkard's random walk

Quick Sort Algorithm

Average case $n \log n$

$$\begin{array}{l}
 5, 9, 3, \underline{10}, 11, 14, 8, 4, 17, 6, 9 \text{ (2)} \\
 \text{3, 6, 1, 7, 8, 9, 5, 2, 10, 4} \\
 \{5, 9, 3, 8, 4, \underline{6}\} \quad 10 \{11, 14, 17\} \quad 5 \\
 \{5, 3, \underline{4}\} \quad 6 \{9, \underline{8}\} \quad 10 \{11, 14, 17\} \quad 2 \\
 \{3\} \quad 4 \{5\} \quad 6 \{9, \underline{8}\} \quad 10 \{11, 14, 17\} \quad 1 \\
 \{3\} \quad 4 \{5\} \quad 6 \{8\} \\
 \{3\} \quad 4 \{5\} \quad 6 \quad 8 \{9\} \quad 10 \{11, \underline{14}, 17\} \quad 2 \\
 \{3\} \quad 4 \{5\} \quad 6 \quad 8 \{9\} \quad 10 \{11\} \quad 14 \{17\} \\
 \{3\} \quad 4 \{5\} \quad 6 \quad 8 \{9\} \\
 X = \text{no. of comparisons} = 19
 \end{array}$$

For the sake of simplicity
 let us label the smallest
 element as 1, next smallest
 as 2 and so on. if i & j are
 ever compared
 otherwise
 $I(i, j) = \begin{cases} 1 & \text{if } i \text{ \& } j \text{ are} \\ & \text{ever compared} \\ 0 & \text{otherwise} \end{cases}$
 $I(1, 5) = 0, I(5, 7) = 1$ etc

$$X = \sum_{i \neq j} I(i, j)$$

③

$$= \sum_{i=1}^{n-1} \sum_{j=i+1}^n I(i, j)$$

$$E[X] = \sum \sum E[I(i, j)]$$

How do you compute

$$E[I(i, j)]?$$

$$E[I(i, j)] = P(I(i, j) = 1)$$

i or j must be selected as pivots.

As long as i & j are in the same bracket, there is a possibility that they get compared.

For any indicator random variable

$$E[X] =$$

$$P(X=1)$$

$$E[X] =$$

$$1 \cdot P(X=1) + 0 \cdot P(X=0)$$

possibility

If they are in different brackets, then they will never get compared. ①

Focus on these $j-i+1$ values

$i, i+1, i+2, \dots, j-1, j$

If a pivot is chosen from one of these values, then they will get separated into different brackets.

i & j will get compared to each other iff i or j is the pivot.

$$P(I(i,j)=1) = \frac{2}{j-i+1}$$

$$E[X] = \sum_{i=1}^{n-1} \sum_{\substack{j=i \\ j=i+1}}^n \frac{2}{j-i+1}$$

$$E[X] = \sum_{i=1}^{n-1} \left[\sum_{j=i+1}^n \frac{2}{j-i+1} \right] \quad (5)$$

$$\sum_{j=i+1}^n \frac{2}{j-i+1} \approx \int_{i+1}^n \frac{2}{x-i+1} dx$$

$$= 2 \log(x-i+1) \Big|_{i+1}^n$$

$$= 2 \left[\log(n-i+1) - \log(1) \right]$$

$$= 2 \log(n-i+1) - \frac{2 \log 2}{\text{ignore}}$$

$$\approx 2 \log(n-i+1)$$

$$E[X] \approx \sum_{i=1}^{n-1} 2 \log(n-i+1)$$

$\approx ?$

$$2 \int_1^{n-1} \log(n-x+1) dx \quad (6)$$

$$n-x+1=y$$

$$-dx = dy$$

$$= 2 \int_n^2 \log y (-dy)$$

$$= 2 \int_2^n \log y dy$$

$$= 2 \left(y \log y - y \right)_2^n$$

~~$$\sim 2n$$~~

$$\sim 2n \log n$$

$$= O(n \log n)$$

$$N \left(\frac{N}{1} + \frac{N}{2} + \frac{N}{3} + \dots + \frac{N}{N} \right) \quad (7)$$

$$6 \left(\frac{6}{1} + \frac{6}{2} + \dots + \frac{6}{6} \right) = 14.7$$

$$\left\{ \begin{array}{l} X_1 = \text{no. of times 1 appears} \\ X_2 = \text{no. of times 2 appears} \\ \vdots \\ X_6 = \text{no. of times 6 appears} \end{array} \right.$$

$$E[X_1] + E[X_2] + \dots + E[X_6] = 14.7$$

$$E[X_1] = \frac{14.7}{6}$$