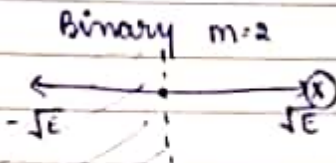


M-ary signalling

1. M-ary PSK

$$M = 2^k$$

k is the no of bits encoded for waveform



min noise power needed to cause error

⊗ Add a vector so that it goes to shaded region, min len. = \sqrt{E} to touch the decision boundary

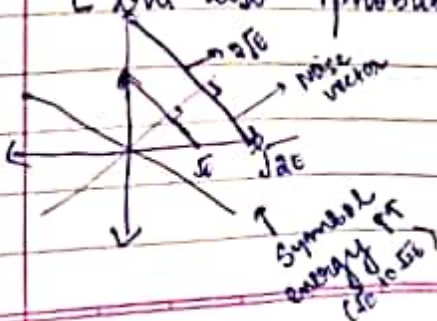
$$\min(\text{Power}) = (\sqrt{E})^2 = E$$

AS M ↑ PB ↑

Performance is degrading then why do we ↑ M?

[Bit Rate Goes up
Every waveform encodes more than 1 bit

[Same bandwidth as M ↑ but data rate goes up but also probability of bit error goes up.



Doubling the energy

$$(\sqrt{2E})^2 + (\sqrt{2E})^2 = 2E + 2E = 4E = E_b$$

$$E_b = 2E \quad \text{half of it} = \sqrt{E}$$

min noise power

when $M=4$ is $\frac{E}{2}$

project (Orthogonally)

\sqrt{E} on the perpendicular bisector.

$d(s_i, s_j) \downarrow$ as M ↑

noise power goes down

∴ more errors

but bandwidth

needed is less i.e. it is

04/11

If distance b/w signals \uparrow $P_B \downarrow$
 M increases in OFSK 2 orthogonal
 unit vectors distance is always $\sqrt{2}$ P_B same.

M-ary OFSK

P_B is majorly dependent on the distance
 b/w the signal vectors.

Claim: As $M \uparrow$ $P_B \downarrow$ Why?

As $M \uparrow$ the bit energy goes down.

$$\text{Bit energy} = \text{Symbol energy} / \text{bits}$$

If we are comparing 2 different M , we should compare
 same $\frac{E_b}{N_0}$, then P_B will go down.

Comparing Symbol energy $P_{\text{symbol error}}$ goes down
 as $M \uparrow$
 min noise power to cause error still remains same.
 Symbol error \uparrow as $M \uparrow$ but $P_B \downarrow$

$M=2$

1 \rightarrow 2
 2 \rightarrow 1

$M=4$

1 \rightarrow 2
 1 \rightarrow 3
 1 \rightarrow 4

04/11/19

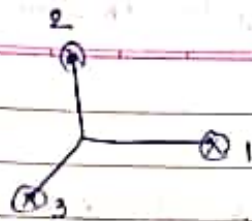
$M \uparrow$ $P_B \downarrow$ Same E_b/N_0

As M goes up symbol error probability goes up
 $M \uparrow$ $P_B \uparrow$

$$d(s_i, s_j) = k \quad i \neq j$$

the amount of noise to be added is same
 for misclassification.

As $M \uparrow$, there are more no of classes,
 the distance b/w 2 classes is same.
 The same noise vector can misclassify
 even if m is increased.



$$M=3$$

$$1-2$$

$$2-1$$

$$M=3$$

$$1 \rightarrow 2$$

$$1 \rightarrow 3$$

$$d(1,2) = d(1,3) = d(2,3)$$

The ways to introduce misclassification with same noise power increases as M increases when looking at same $\frac{E_b}{N_0}$ ratio that is $P_s \uparrow$ SNR

Each symbol is encoded by K bits
 M different symbols, $K = \log_2 M$ bits, every symbol (duration T),
 one with high value of M needs to have a higher value of K for keeping $\frac{E_b}{N_0}$ same, each symbol energy encoded for K bits,

$$P_s \rightarrow K \text{ bits}$$

$$E_{\text{bit}} \rightarrow \frac{P_s}{K}$$

As $K \uparrow$ bit energy \downarrow

$$K \uparrow \quad E_b \downarrow \quad E_s (\text{constant})$$

\therefore comparing @ same $\frac{E_b}{N_0}$, we get an increase in symbol energy, $d(s_i, s_j)$ between symbols dependent on symbol energy goes up, \therefore noise power needed for wrong classification \uparrow , and hence error \downarrow as $M \uparrow$ (bit)

Large value of M gives better performance, that is lower error but higher bandwidth.

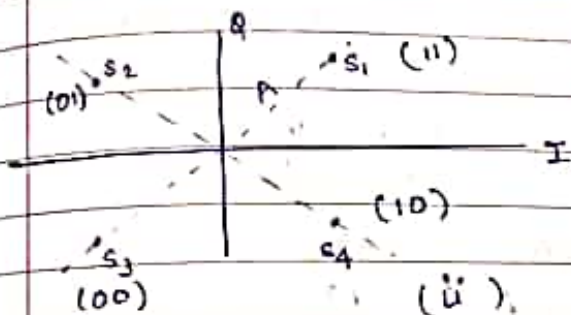
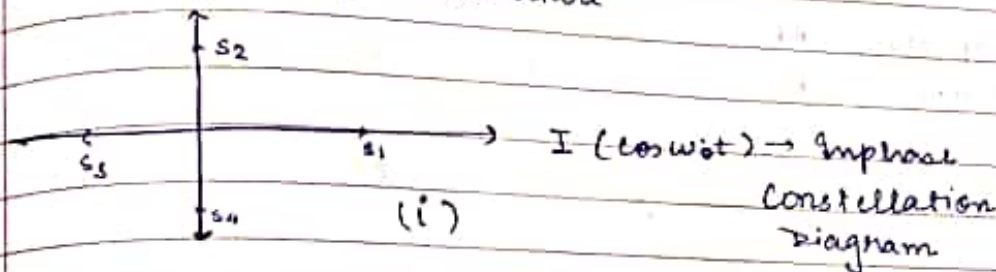
M waveforms

$$w_{i+1} - w_i = \pi/T$$

In QPSK

$$S_i(t) = \sqrt{\frac{2E}{T}} \cos\left(\omega t + \frac{\pi}{2}(i-1)\right) \quad 0 \leq t \leq T$$

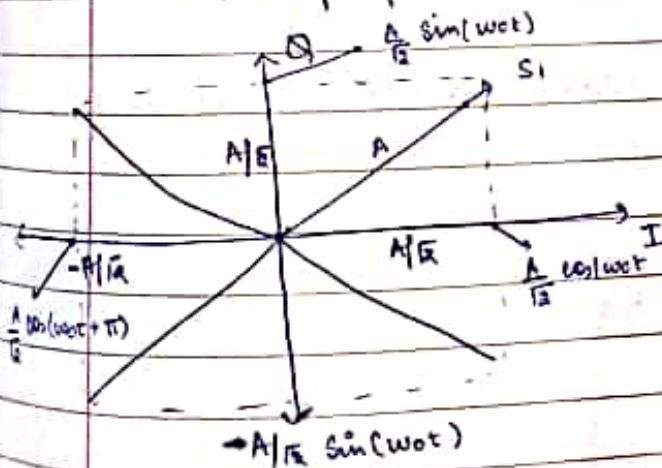
$Q(\sin \omega t) \rightarrow$ Quadrature $i = 1, 2, 3, 4$



they both are equivalent as everything depends on distance & bit energy

$$(ii) S_i(t) = \sqrt{\frac{2E}{T}} \cos\left(\omega t + \frac{\pi}{4} + \frac{\pi}{2}(i-1)\right)$$

$S_1(t)$ projection on I and Q axis will be $A/\sqrt{2}$



S_1	$\frac{A}{\sqrt{2}}$ (I axis)	$\frac{A}{\sqrt{2}}$ (Q axis)
S_2	$-\frac{A}{\sqrt{2}}$ (I axis)	$\frac{A}{\sqrt{2}}$ (Q axis)

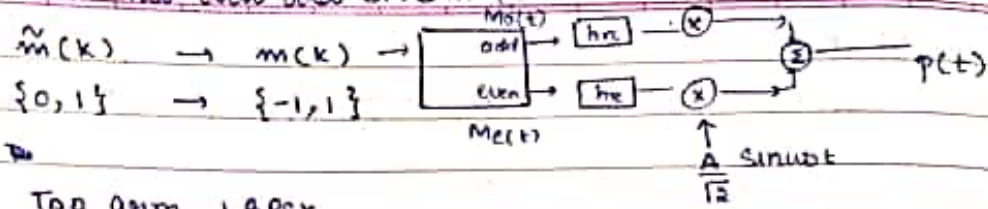
{ only looking @ 1 axis we say that it is BPSK
0 \rightarrow -ve
1 \rightarrow +ve

Basically QPSK can be looked as 2 BPSK modulations, there is no interference b/w the 2 BPSK schemes

Implement a QPSK as 2 BPSK.

Modulator

Divide into even odd stream



Top arm 1 BPSK

Bottom arm 1 BPSK

< implement demodulator > < HW >

E_{QPSK1} where you transmit s_1, s_2, s_3, s_4 E_{QPSK2} where QPSK implemented as 2 BPSK Schemes.

$$\frac{E_{QPSK2}}{N_0} = \frac{S_{QPSK2} \times T}{N/W} = \frac{S_{QPSK2}}{N} \left(\frac{W}{R} \right)$$

$$\text{In QPSK1 Signal Power } (A)^2 = A^2$$

$$\text{In QPSK2 Signal Power } \left(\frac{A}{\sqrt{2}} \right)^2 = \frac{A^2}{2}$$

$$S_{QPSK2} = \frac{S_{QPSK1}}{2}$$

$$\frac{E_{QPSK2}}{N_0} = \frac{S_{QPSK1}}{2N} \left(\frac{W}{R} \right)$$

even stream / odd stream arriving at half the rate of original stream.

$$\begin{aligned} \frac{E_{QPSK2}}{N_0} &= \frac{S_{QPSK1}}{2N} \left(\frac{2W}{R_{QPSK1}} \right) \\ &= \frac{E_{QPSK1}}{N_0} \end{aligned}$$

we are going to do coherent detection

$$P_B(QPSK2) = P_B(QPSK1) = P_B(BPSK) @$$

the corresponding E_b/N_0 ratio.

Just for $M=2$ & $M=4$, the fact that M or P or is not true because we reduced the rate by 2.

when

m-ary FSK

$$S_i(t) = \sqrt{\frac{2E}{T}} \cos(\omega_i t)$$

$$0 \leq t \leq T$$

$$i = 1, \dots, M$$

$$d(S_i, S_j) = K, \quad i \neq j$$

How many ways are there to misclassify?
 which of them is more likely?
 • distance b/w symbols are same i.e. all misclassifications are equally likely.
 there are $m-1$ ways to misclassify a certain symbol.

Let us say you misclassify a symbol (K bits),
 now how many bits were misclassified?
 = every symbol error, means all bits were wrong.

$$M = 2^K$$

one symbol can be misclassified $2^K - 1$ ways.

{ 0000 → 0101
 2 bit errors -

0000
 1111

□ bit misclassification → by 0.5 probability
 If I take any of these bits → half of the bits give correct detection & half of them give incorrect.

$$\text{Incorrect / misclassify} = \frac{2^K}{2}$$

$$\frac{P_b}{P_s} = \frac{2^{K-1}}{2^K - 1}$$

$$< \frac{1}{2} > \text{ min}$$

↑ Relating P_b with P_s is assumed if the misclassification has already occurred.
 Every symbol error gives a bit error.

$$K=3$$

00 01 11 10
 0/3

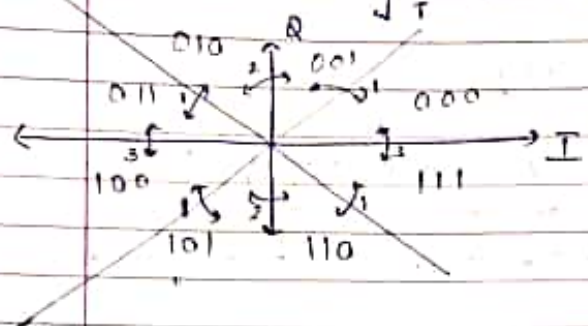
$$\frac{2^{3-1}}{2^3 - 1} = \frac{2}{3}$$

✓ \rightarrow \rightarrow \rightarrow

M-ary TSK

$$s_i(t) = \sqrt{\frac{2E}{T}} \cos(2\pi f_c t + \frac{2\pi}{M} (i-1)) \quad 0 \leq t \leq T$$

$i = 1 \dots M$



Higher probability of misclassification in close by classes (Adjacent)

Assuming P_s is small.

If misclassification occurs only over adjacent - Given that symbols are misclassified, then how much does symbol error contributes to bit error. Instead of above encoding scheme, as there are a lot of bit errors we can use Gray Code

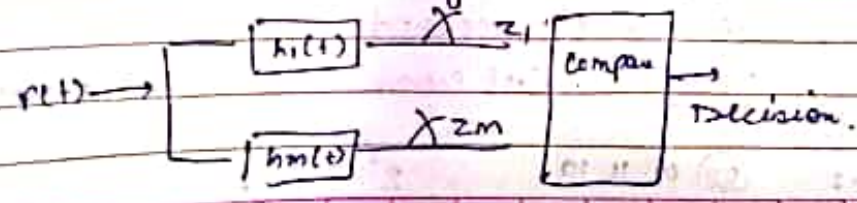
- | | |
|-----|------------------|
| 000 | } Anticlockwise: |
| 001 | |
| 011 | |
| 010 | |
| 110 | |
| 111 | |
| 101 | |
| 100 | |

$\frac{P_b(\text{Gray})}{P_s} = \frac{1}{k}$

Once the symbol error occurred, how many bits were misclassified.

Symbol error rate (P_s)

- FSK - coherent demodulator - using matched filter detection.

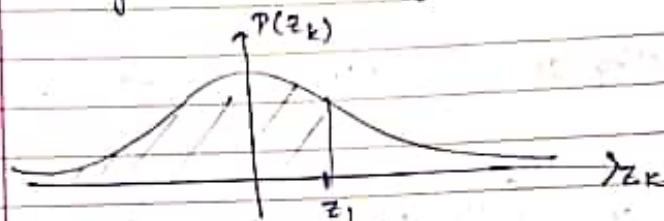


$Z_k \sim N(0, \sigma_0)$ $k \neq 1$ if S_1 was tx.

$$P(Z_1 > Z_k | Z_1) = \int_{-\infty}^{Z_1} N_{Z_k}(0, \sigma_0) dZ_k$$

= integrate out Z_k which have less value than Z_1

For a given Z_1 what is the probability $Z_k < Z_1$
integrate out the pdf of Z_k .



$$Z_1 \sim N(\mu, \sigma) \quad P(Z_1) =$$

Evaluate $P_S = 1 - P_C$

closed form solution does not exist.

Upper bounds for P_S

M waveforms $d(S_i, S_j) = K \quad i \neq j$

all misclassification equally likely

↓

$\exists k \in (2, \dots, m)$ s.t. $Z_k > Z_1$

E_k : event for $Z_k > Z_1$, $k = 2, \dots, m$

misclassification: any one of these $m-1$ events occur
 $1 - (\text{none } Z_k > Z_1)$

$P_S = P(\text{one of } E_k \text{ occurred})$ @ least one of $E_k =$

$$P\left(\bigcup_{k=2}^m E_k\right) \leq \sum_{k=2}^m P(E_k)$$

Z_2 & Z_3 both can be $>$ than Z_1

any event in union of $E_k \Rightarrow$ one of the event occurs

Assume one is T_{x_i} , it gets misclassified into $m-1$ classes

Assume S_1 was transmitted

$$P(Z_1) \sim N(\sqrt{E}, \sigma^2)$$

$$P(Z_k) \sim N(0, \sigma^2), k \neq 1$$

$k = 2 \dots M$

$$P(\text{correct detecting } S_1) = P(Z_1 > Z_2, Z_1 > Z_3 \dots Z_1 > Z_M | Z_1) \times P(Z_1)$$

$$P(\text{correct } S_1) = \int_{-\infty}^{\infty} P(Z_1 > Z_2 \dots Z_1 > Z_M | Z_1) \cdot P(Z_1) dZ_1$$

$$P(Z_1 > Z_2 \dots Z_1 > Z_M | Z_1) = \prod_{k=2}^M P(Z_1 > Z_k | Z_1)$$

02/11/19 Symbol error rates (P_s)

1. MFSK (coherent demodulation)

$$S_i(t) = \sqrt{\frac{2E}{T}} \cos(\omega t) \quad 0 \leq t \leq T \quad i = 1 \dots M$$

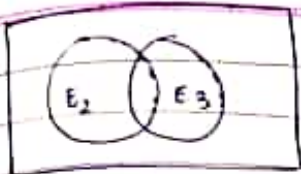
o/p of M correlators $\rightarrow Z_1, \dots, Z_M$

- Assume S_1 was transmitted

$$(a) - P_s = 1 - P_c$$

$$(b) - P_c = \int_{-\infty}^{\infty} P(Z_1 > Z_2, Z_1 > Z_3 \dots Z_1 > Z_M | Z_1) P(Z_1) dZ_1$$

$$P(Z_1 > Z_2 \dots Z_1 > Z_M | Z_1) = \prod_{k=2}^M P(Z_1 > Z_k | Z_1)$$



$$P(E_2 \cup E_3) = P(E_2) + P(E_3) - P(E_2 \cap E_3)$$

$P(E_k)$ = Each event is worried about $s_k \neq s_1$
 so what is the probability that E_k happens.
 it is exactly same as Binary FSK.

$$P(E_k) = Q\left(\sqrt{\frac{E}{N_0}}\right)$$

$$\sum_{k=2}^M P(E_k) = (M-1) Q\left(\sqrt{\frac{E}{N_0}}\right)$$

E - Symbol energy

$$E_s = K E_b$$

$$\therefore P_s \leq (M-1) Q\left(\sqrt{\frac{K E_b}{N_0}}\right)$$

$$K \rightarrow \infty \quad P_s \rightarrow 0$$

$$P_s \leq M Q\left(\sqrt{\frac{K E_b}{N_0}}\right)$$

$$P_s \leq M e^{-\frac{K E_b}{2 N_0}}$$

$$P_s \leq 2^k e^{-\frac{K E_b}{2 N_0}}$$

$$D(x) \leq e^{-x^2/2}$$

$$P_s \leq e^{K \left[\ln 2 - \frac{E_b}{2 N_0} \right]}$$

$$P_s < e^{\frac{K}{2} \left[2 \ln 2 - \frac{E_b}{N_0} \right]}$$

$$\text{we want : } \frac{E_b}{N_0} > 2 \ln 2$$

• as $K \uparrow$ error will go up if $2 \ln 2 > \frac{E_b}{N_0}$

But by $\uparrow K$ $P_s \downarrow$

$$K \rightarrow \infty \quad P_s \downarrow 0$$

$$\frac{E_b}{N_0} > 1.386$$

tighter upper bound

$$P_s \sim e^{-K \left(\sqrt{\frac{E_b}{N_0}} - \sqrt{\ln 2} \right)^2}$$

expression is valid only if $\sqrt{\frac{E_b}{N_0}} > \sqrt{\ln 2}$

$$\frac{E_b}{N_0} > \ln 2$$

Shannon's limit

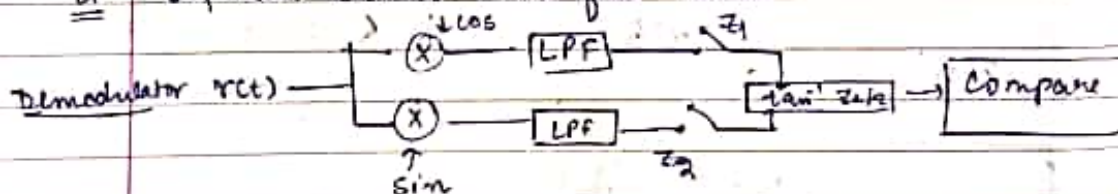
Greater noise power compared to bit energy $\left\{ \frac{E_b}{N_0} > 0.69 \right.$

? $\frac{KT}{P_s} \downarrow$ (upper bound)
(In theory)

$$\text{If } \frac{E_b}{N_0} > 0.69 \quad K \rightarrow \infty \quad P_s \downarrow 0$$

If you increase SNR increasing M is going to decrease P_s .

2. symbol error rate for MPSK



θ was transmitted

$$\sim N(\sqrt{E_b}, 0)$$

$$p(z_1, z_2) \rightarrow p(m, \beta)$$

$$m = \sqrt{z_1^2 + z_2^2}$$

$$\beta = \tan^{-1} \left(\frac{z_2}{z_1} \right)$$

for QPSK dec. boundary $-\pi/2$ to $\pi/2$
 $-\pi/4$ to $\pi/4$

$$P_s = 1 - \int_{-\pi/M}^{\pi/M} \int_0^{\infty} P(m, p) dm dp$$

Q. 44

$$L(\hat{s}) = \| \bar{r} - \bar{h} s \|^2$$

$$= \bar{r}^T \bar{r} - 2 \bar{r}^T \bar{h} s + \bar{h}^T \bar{h} s^2$$

let us somehow assume that s^2 remain const
 QPSK constellation $\text{dist}(s)$ same.

What I need?

$$\min(L(\hat{s})) \quad \text{or} \quad \max(\bar{r}^T \bar{h})$$

↑ max projection / or
 correlation / or
 matched filter receiver

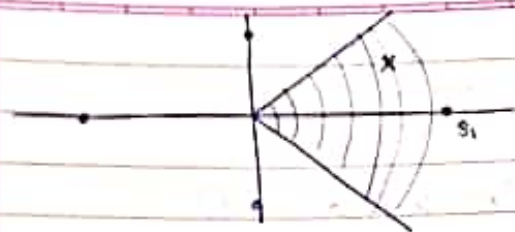
what if I $-L(\hat{s}) \cdot K$

$$P(\hat{s} | \bar{r}) = \underbrace{\propto e}_{\text{Gaussian T.D.F.}}$$

Conditional PDF/likelihood
 {Multivariate as s is a vector}

To minimize the cost; you maximize the likelihood.

↓
 Max. likelihood estimate (MLE) of parameter
 symbol vector.



$\left\{ \begin{array}{l} d(x, s_1) \text{ is minimum} \\ \text{decision} \end{array} \right.$

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MFSK

$T \rightarrow$ Symbol duration

$$R_s = \frac{1}{T}$$

$$W = R_s = \frac{R_b}{\log_2 M}$$

R_b - bits/sec

$$\text{Bandwidth efficiency} : \frac{R_b}{W} = \text{bits/sec/Hz} \\ = \log_2 M$$

For every Hz of channel bandwidth you are able to transmit R_b bits

The idea to transfer more & more bits

As you $\uparrow M$, channel bandwidth efficiency \uparrow

M-FSK

coherent

$$W = \frac{(M-1)}{2T} \approx \frac{M}{2T}$$

$\frac{1}{2T}$ comes from separation

$$\approx \frac{M \cdot R_s}{2}$$

$$R_s = \frac{1}{T}$$

The aim is to compute $\frac{R_b}{W}$

$$W \approx \frac{M \cdot R_b}{2 \log_2 M}$$

$$\frac{R_b}{W} = \frac{2 \log_2 M}{M}$$

As $M \uparrow$ the bandwidth efficiency decreases.

Suppose we want to maintain same P_b in both case)
what effects occur in MPSK & MSK

• MPSK: - As $M \uparrow$; \uparrow Bandwidth efficiency

If $M \uparrow$ $P_b \uparrow$ maintain same E_b/N_0 ratio
but in order to maintain constant P_b we will
increase E_b (Power \uparrow).
No

MSK: - As $M \uparrow$; \downarrow Bandwidth efficiency
 $\downarrow E_b$ ratio
No

If $M \uparrow$ $P_b \downarrow$ maintain same E_b/N_0 ratio

If Bandwidth is a constraint you use MPSK; i.e.
if B/W is large / i.e. efficiency is high)

you need to choose one of the 2
trade off b/w Bandwidth efficiency & E_b (Power
Ratio).

Shannon - Hartley 'Capacity' Theorem.

AWGN channel

Max. data rate that channel can support i.e. demodulation
errors are ignorable (or are very very small), theoretically?

It gives an upper bound to the idea of max
data rate.

upper bound on data rate = channel capacity = C

$$C = W \cdot \log_2 \left(1 + \frac{S}{N} \right)$$

upper bound

$\frac{S}{N}$: Signal to noise ratio (SNR)

Energy

: Avg SNR received over the W Hz

W : Bandwidth of the channel

C : Max. Bit Rate

Shannon believed there exists some scheme of demodulation with $\left(W \propto \frac{S}{N} \right)$ as parameter.

In terms of $\frac{E_b}{N_0}$ ratio

$$C = W \cdot \log_2 \left(1 + \frac{S}{N} \right)$$

$$\frac{E_b}{N_0} = \frac{S}{NR}$$

Let us assume that channel rate & Bit Rate are same. ($C = R$)

$$\frac{E_b}{N_0} \frac{R}{W} = \frac{S}{N}$$

$$\frac{R}{W} = \log_2 \left(1 + \frac{E_b}{N_0} \left(\frac{R}{W} \right) \right)$$

$$\frac{E_b}{N_0} = \frac{2^{R/W} - 1}{R/W}$$

upper bound.

what if $\frac{R}{W} \rightarrow 0$ Bandwidth \uparrow (FSK)

$$\frac{E_b}{N_0} = \ln 2$$

Shannon's limit

Increasing Bandwidth \rightarrow Increasing M in FSK

- Quadrature Mod / Demod (non coherent)
- Error performance P_b & P_s
:- what happens as $M \uparrow$
- Bandwidth efficiency