

# Termination

There are 4 aspects of the termination

- i) Alternative optimum
- ii) Unboundedness
- iii) Infeasibility
- iv) Cycling

## Alternative optimum

consider the following LPP.

$$\text{Max } Z = 2x_1 + 3x_2$$

$$\text{s.t. } 5x_1 + 3x_2 \leq 15 \quad \text{--- ①}$$

$$4x_1 + 6x_2 \leq 12 \quad \text{--- ②}$$

$$x_1, x_2 \geq 0$$

We introduce slack variables  $x_3$  and  $x_4$  to ① and ② then we have,

$$\text{Max } Z = 2x_1 + 3x_2 + 0x_3 + 0x_4$$

$$\text{s.t. } 5x_1 + 3x_2 + x_3 = 15$$

$$4x_1 + 6x_2 + x_4 = 12$$

$$x_1, x_2, x_3, x_4 \geq 0.$$

$C_B$	B	$x_B$	b	$a_1$	$a_2$	$a_3$	$a_4$	min ratio	operation
0	$a_3$	$x_3$	15	5	3	1	0	$15/3=5$	
0	$a_4$	$x_4$	12	4	6	0	1	$12/6=2 \rightarrow$	
$Z_j - C_j$			0	-2	-3 $\uparrow$	0	0		
0	$a_3$	$x_3$	9	3	0	1	$-1/2$	$9/3=3 \rightarrow$	$R_1' = R_1 - 3R_2'$
3	$a_2$	$x_2$	2	$2/3$	1	0	$1/6$	$2/(2/3)=3$	$R_2' = R_2/6$
$Z_j - C_j$			6	0 $\uparrow$	0	0	$1/2$		
2	$a_1$	$x_1$	3	1	0	$1/3$	$-1/6$		$R_1' = R_1/3$
3	$a_2$	$x_2$	0	0	1	$-2/9$	$5/18$		$R_2' = R_2 - \frac{2}{3}R_1'$
$Z_j - C_j$			6	0	0	0	$1/2$		

Observation:

- In iteration 2 the non-basic variable  $a_1$  has  $Z_j - C_j$  is zero
- If we enter this variable we get another optimum solution with same objective function value.
- If we enter  $a_3$  again we get the same values in iteration 2.

$$\max Z = 2x_1 + 3x_2$$

$$\text{s.t. } 5x_1 + 3x_2 \leq 15 \quad \text{--- (1)}$$

$$4x_1 + 6x_2 \leq 12 \quad \text{--- (2)}$$

$$x_1, x_2 \geq 0$$

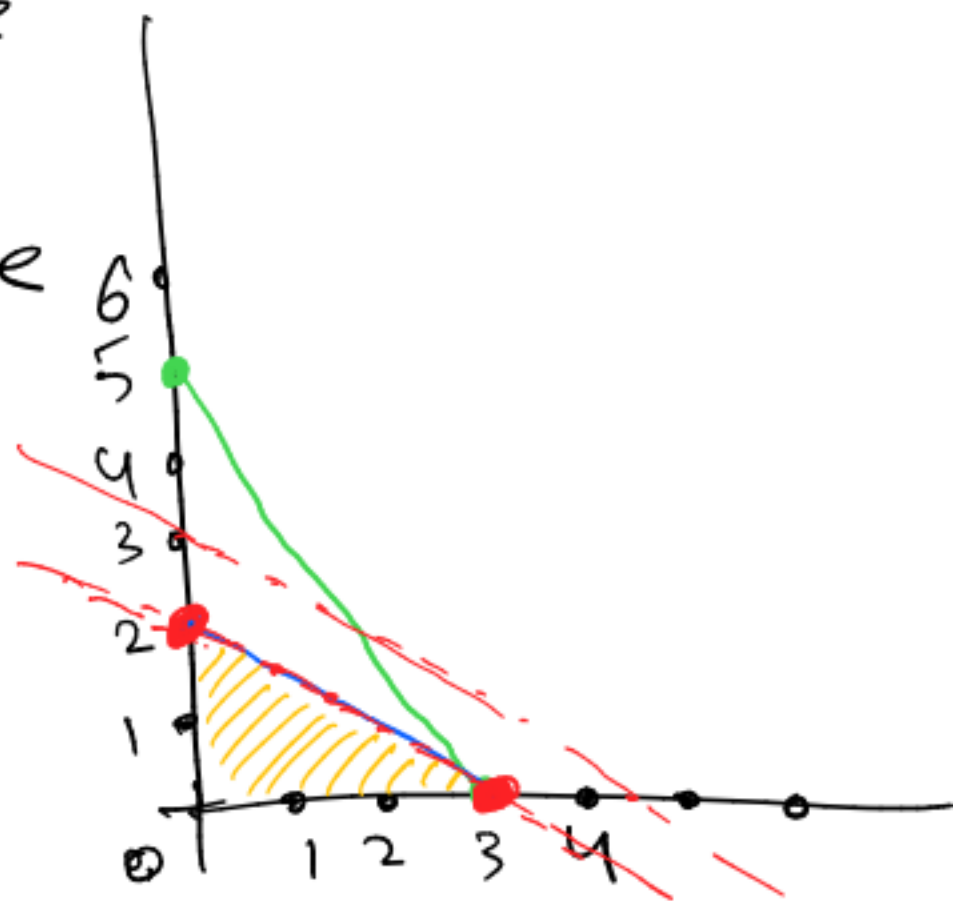
Two points on (1) are  
 $(0, 5)$  and  $(3, 0)$

Two points on (2) are  
 $(0, 2)$  and  $(3, 0)$

Two points on  
 objective function line  
 are,

$$2x_1 + 3x_2 = 6$$

$(0, 2)$  and  $(3, 0)$



$(0, 2)$  and  $(3, 0)$

## Unboundedness

Consider the following LPP.

$$\text{Max } Z = 3x_1 + 4x_2$$

$$\text{s.t. } -x_1 + x_2 \leq 0 \quad \text{--- (1)}$$

$$-x_1 + 3x_2 \leq 3 \quad \text{--- (2)}$$

$$x_1, x_2 \geq 0.$$

After introducing slack variables we have,

$$\text{max } Z = 3x_1 + 4x_2 + 0x_3 + 0x_4$$

$$\text{s.t. } -x_1 + x_2 + x_3 = 0$$

$$-x_1 + 3x_2 + x_4 = 3$$

$$x_1, x_2, x_3, x_4 \geq 0.$$



			$C_j$	$3$	$4$	$0$	$0$		
$C_B$	$B$	$x_B$	$b$	$a_1$	$a_2$	$a_3$	$a_4$	min ratio	operations
$0$	$a_3$	$x_3$	$0$	$-1$	<span style="border: 1px solid red;">1</span>	$1$	$0$	$0/1 = 0$ <span style="color: red;">→</span>	
$0$	$a_4$	$x_4$	$3$	$-1$	$3$	$0$	$1$	$3/3 = 1$	
$Z_j - C_j$			$0$	$-3$	$-4$ <span style="color: red;">↑</span>	$0$	$0$		
$4$	$a_2$	$x_2$	$0$	$-1$	$1$	$1$	$0$	-----	$R_1' = R_1 / 1$
$0$	$a_4$	$x_4$	$3$	<span style="border: 1px solid red;">2</span>	$0$	$-3$	$1$	$3/2 = 1.5$ <span style="color: red;">→</span>	$R_2' = R_2 - 3 \cdot R_1$
$Z_j - C_j$			$0$	$-7$ <span style="color: red;">↑</span>	$0$	$4$	$0$		
$4$	$a_2$	$x_2$	$3/2$	$0$	$1$	$-1/2$	$1/2$	---	$R_1' = R_1 + R_2'$
$3$	$a_1$	$x_1$	$3/2$	$1$	$0$	$-3/2$	$1/2$	-----	$R_2' = R_2 / 2$
$Z_j - C_j$				$0$	$0$	$-13/2$ <span style="color: red;">↑</span>	$7/2$		

### Observation

- the algorithm terminates but  $x_3$  tries to enter the basis, and there is no leaving variable
- This is called unboundedness.

# Infeasibility

$$\max Z = 5x_1 + 3x_2$$

$$\text{s.t. } 3x_1 + x_2 \leq 1 \quad \text{--- (1)}$$

$$3x_1 + 4x_2 \geq 12 \quad \text{--- (2)}$$

$$x_1, x_2 \geq 0$$

Introduce slack variable  $x_3$  to (1) and surplus and artificial variables  $x_4$  and  $x_5$  to (2), then we have.

$$\max Z = 5x_1 + 3x_2 + 0x_3 + 0x_4 - Mx_5$$

$$\text{s.t. } 3x_1 + x_2 + x_3 = 1$$

$$3x_1 + 4x_2 - x_4 + x_5 = 12$$

$$x_1, x_2, x_3, x_4, x_5 \geq 0$$

$C_j$	B	$X_B$	b	$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	min ratio	operations
0	$a_3$	$x_3$	1	3	1	1	0	0	$1/1=1$	
-1	$a_5$	$x_5$	12	3	4	0	-1	1	$12/4=3$	
$Z_j - C_j$			-12	-3	-4	0	1	0		
0	$a_2$	$x_2$	1	3	1	1	0	0		$R'_1 = R_1/1$
-1	$a_5$	$x_5$	8	-4	0	-4	-1	1		$R'_2 = R_2 - 4R'_1$
$Z_j - C_j$			-8	9	0	4	1	0		

### observations

- all  $Z_j - C_j$ 's are  $\geq 0$ . Thus optimal condition is reached.
- The artificial variable still present in the basis.
- maximum objective value is  $-8 < 0$
- The solution is not feasible.



## Cycling

Consider the LPP.

$$\text{minimise } Z = -\frac{3}{4}x_1 + 20x_2 - \frac{1}{2}x_3 + 6x_4$$

$$\text{s.t. } \frac{1}{4}x_1 - 8x_2 - x_3 + 9x_4 \leq 0 \quad \text{--- (1)}$$

$$\frac{1}{2}x_1 - 12x_2 - \frac{1}{2}x_3 + 3x_4 \leq 0 \quad \text{--- (2)}$$

$$x_3 \leq 1 \quad \text{--- (3)}$$

$$x_1, x_2, x_3, x_4 \geq 0$$

## Reference

E. m. L. Beale, 1995.

Naval Research Logistic Quarterly, 2,  
pp. 269-275

consider the standard form.

$$\text{max } Z = \frac{3}{4}x_1 - 20x_2 - \frac{1}{2}x_3 - 6x_4 + 0x_5 + 0x_6 + 0x_7$$

$$\text{s.t. } \frac{1}{4}x_1 - 8x_2 - x_3 + 9x_4 + x_5 = 0$$

$$\frac{1}{2}x_1 - 12x_2 - \frac{1}{2}x_3 + 3x_4 + x_6 = 0$$

$$x_3 + x_7 = 1$$

$$x_1, x_2, \dots, x_7 \geq 0.$$

$C_j$	B	$x_B$	$b$	$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$	min ratio	operations
0	$a_5$	$x_5$	0	$\frac{1}{4}$	-8	-1	9	1	0	0	$0/\frac{1}{4} = 0$	} $0/\frac{1}{2} = 0$ .....
0	$a_6$	$x_6$	0	$\frac{1}{2}$	-12	$-\frac{1}{2}$	3	0	1	0		
0	$a_7$	$x_7$	1	0	0	1	0	0	0	1		
$Z_j - C_j$				$-\frac{3}{4}$	20	$-\frac{1}{2}$	6	0	0	0		
$\frac{3}{4}$	$a_1$	$x_1$	0	1	-32	-4	36	4	0	0	.....	$R'_1 = R_1/\frac{1}{4}$
0	$a_6$	$x_6$	0	0	$\frac{1}{4}$	$\frac{3}{2}$	-15	-2	1	0	$0 \rightarrow$	$R'_2 = R_2 - \frac{1}{2}R_1$
0	$a_7$	$x_7$	1	0	0	1	0	0	0	1	.....	$R'_3 = R_3$
$Z_j - C_j$				0	-4	$-\frac{7}{2}$	33	3	0	0	$0 \rightarrow$	$R'_1 = R_1 + 32R'_2$
$\frac{3}{4}$	$a_1$	$x_1$	0	1	0	$\frac{1}{8}$	-84	-12	8	0	0	$R'_2 = R'_2/\frac{1}{4}$
-20	$a_2$	$x_2$	0	0	1	$\frac{3}{8}$	$-\frac{15}{4}$	$-\frac{1}{2}$	$\frac{1}{4}$	0	0	$R'_3 = R'_3$
0	$a_7$	$x_7$	1	0	0	1	0	0	0	1	1	
$Z_j - C_j$				0	0	-2	18	1	1	0		

If we proceed in this way we see that the 7th iteration is same as iteration 1.