

IT486 v3.0: Blockchains and Cryptocurrencies

Bitcoin blockchain

Proof of Work

- Puzzle: find a nonce such that

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- The hash function is “puzzle hard”:
 - given d and (small enough) T , it is hard to find n such that $\text{hash}(n \parallel d) < T$.

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- Bitcoin's approach
 - Set the puzzle difficulty so that the expected time to find a solution is 10 mins (a constant)
 - Periodically change the difficulty so as to try to maintain this expected time

The coinbase transaction

- the first transaction in the block
- does not consume (spend) previous unspent outputs contained in the blockchain
- Creates bitcoins from nothing
 - a single dummy input not linked to any output
 - the sum of the outputs is equal to the block reward + the transaction fees
 - output addresses: one or more of the miner's own bitcoin address

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 - 2012 → 25 BTC
 - 2016 → 12.5 BTC
 - 2020 → 6.25 BTC
- By 2056, the reward will be about 0.0125 BTC

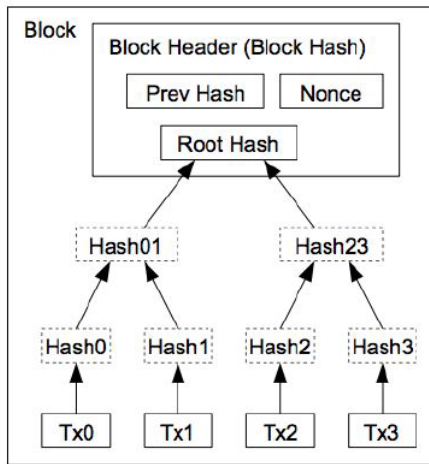
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- By 2056, the reward will be about 0.0125 BTC
- At that point, miners will mainly profit from collecting transaction fees

Structure of a block

- a header (80 bytes) consisting of three sets of metadata
 - a hash of the block header of the previous (last added) block of the blockchain
 - information related to PoW
 - root of the Merkle tree summarizing all the transactions in the block
- a list of transactions that make up the bulk of its size
 - almost 1000 times larger than the block header

Structure of a block



Transactions Hashed in a Merkle Tree

Identifying blocks

- How to identify a block uniquely within the blockchain?
- By the block hash
 - computed by each node as the block is received from the network
 - not actually stored inside the block
 - not computed when the block is transmitted on the network

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- Bitcoin's genesis block
 - mined at 2009-01-03 18:15:05
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 - created the first 50 BTC of mining reward
- Satoshi included a message in the coinbase txn of the genesis block

How to mine

- miners construct a candidate block filling it with transactions
- transactions are inserted in the block as they are received from the network
- solve PoW:
 - calculate the hash of this block header, including the root of the Merkle tree, the hash of the previous block and the nonce and check if it is smaller than the current target
 - Since the nonce is a 32-bit integer, a miner has to try approx. 4 billion values; trying all values may not work

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 - the miner can also reorder the transactions to create a different merkle tree
 - since there can be several thousand transactions in a block, this creates an inconceivable number of nonces that can be tried

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- with a Merkle tree, rehashing is more efficient when a new transaction is received
- instead of rehashing all the transactions, compute the new root of the merkle tree, and execute the PoW by considering this new root
- only a few operations to recompute all the hashes along the branches on the path from the new txn to the root

Tamper resistance

- imagine an attacker tries to change the content of a block, for example a txn
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- an attacker would have to recompute the PoW for the entire chain