Math 55, Handout 10.

RECURSIVE DEFINITIONS.

1.1. Recursively defined functions. A function f on \mathbb{Z}_+ can be recursively defined using the following two steps:

Basis step: Specify the value of the function at zero
Recursive step: Give a rule for tinding its value at an integer from
His values at smaller integers

Q1. Is this a valid recursive definition of a function $f: \mathbb{Z}_+ \to \mathbb{Z}$?

 $\begin{array}{lll} \textbf{n}=\textbf{0} & \textbf{f(n)=2} & f(0)=2, & f(n)=\left\{ \begin{array}{ll} f(n-1) & \text{if n is odd and $n\geq 1$} \\ 2f(n-2) & \text{if $n\geq 2$}. \end{array} \right. \\ \textbf{n=1} & \textbf{f(n)=2} & \textbf{form solution:} \\ \textbf{n=2} & \textbf{f(n)=4} & \textbf{L2n)?} & \textbf{no} \\ \textbf{n=3} & \textbf{f(n)=4} & \textbf{2n}? & \textbf{No} \\ \textbf{2n}? & \textbf{No} & \textbf{defintion as the function} \\ \textbf{16} & \textbf{well-defined} \end{array}$

- 1.2. Properties of recursively defined functions can be typically proved by
- Q2. Prove that the Fibonacci sequence $\{f_n\}$ satisfies $\sum_{j=1}^n f_j^2 = f_n f_{n+1}$. Balls: n=1, $\sum_{j=1}^n f_j^2 = f_j^2 = f_j^2 = 1$ and $f_n f_2 = |\cdot| = 1$ = $f_{n+1}(f_n + f_{n+1}) \rightarrow f_{n+1}$ by inductive my pothering = $f_{n+1}(f_n + f_{n+1}) \rightarrow f_{n+1}$ by the righter fiberacci = $f_{n+1} + f_{n+2} \vee f_{n+1}$

this shows the equivalence for n+1 whenever it holds for n. this proves the inductive stop hence proving the equivalence for any any n=1H

1.3. Recursively defined sets. A set S can be recursively defined using the following two steps: Basis step: $A \in \Sigma^*$ Cwhere A is the empty string containing no symbols) Recursive step: If $W \in \Sigma^*$ and $X \in \Sigma$, then $W * E^*$

Q3. Show that the set S defined by $1 \in S$ and $s+t \in S$ whenever $s, t \in S$ is the set $I\!N$. IES S+t $E\!S$ whenever s $E\!S$ and t $E\!S$

proof: nes for every positive integer

Basis step n=1 1 ES thus P(1) istrue

inductive step p(k) is true KES. Prove that p(k+D) is also true since IES and kes by the def

thus p(k+1) is true. P(n) is true when n is a positive integer n

Q4. Give a recursive definition of the set of positive integer powers of 5.

basis inductive step $s \in S$ and $s \in S \longrightarrow Ss \in S$ basis

rooted trees 1

Recursively defined, such as

1.4. Besides functions and sets, other structures can be recursively defined, such as



Strings & rooted trees

STRUCTURAL AND GENERALIZED INDUCTION.

3.1. Structural induction principle. Results about recursively defined sets (and other structures) can be proved using the following two steps:

Basis step: Show that the result holds for all elements specified in the basis step of the elements in the recursive step of the definition, the result holds for these new olements.

- 3.2. Generalized induction is using (weak or strong) induction to prove results about sets that have the well-ordering property. The exicultable **ordering** is often used for that purpose.
- 3.3. One can extend generalized induction even further to partially ordered sets.
- 3.4. Transfinite induction can be used for ordinals larger than IN. This is beyond the scope of Math 55.
- Q5. Use generalized induction to show that if $a_{m,n}$ is defined recursively by $a_{0,0} = 0$ and

$$a_{m,n} = \begin{cases} a_{m-1,n} + 1 & \text{if } n = 0 \text{ and } m > 0 \\ a_{m,n-1} + 1 & \text{if } n > 0, \end{cases}$$

then $a_{m,n} = m + n$ for all $(m,n) \in \mathbb{Z}_+ \times \mathbb{Z}_+$.

(0,3) • conjecture: am, n = m + n (1,2)四·

induction on M+n

BASS: $m+n=0 \longrightarrow m=n=0$ and $a_{m,n}=m+n=0+0=0$

inductive step: Assume am, n=m+n for m+n=k

Thow that mtn=k+1

case 1: n=0, m=k+1. then am, n=am-1, n+1 (m-1)+n=k+1-1=k. The inductive hypothesis applies to am-1,n and giver am-1,n=m-1+n=k. am,n=k+1V

<u>case 2</u>: n>0, m+n=k+1. Then $am,n=a_{m,n-1}-1$. And $m+(n+1)=k+1\cdot k=k$. To the inductive hypothesis applies to am,n-1 and gives am,n-1=m+(n-1)=k. To am,n=am,n-1+1=k+1

both cases the inductive step is proven.