

## Data Structure

A mathematical and logical model of data is known as *Data Structure*.

**Primitive data structure:** The data structure, which is available in the compiler, is known as a primitive data structure. Example: Array

**Non-primitive data structure:** The data structure, which is not available in the compiler, is known as non-primitive data structure. Examples: Stack, Queue, Linked List

**Linear Data Structure:** The data structure in which each element can access maximum one predecessor element and one successor element is known as *linear data structure*. Example: Stack, Queue etc.

**Non-linear Data Structure:** The data structure in which each element can access any number of predecessor elements and any number of successor elements is known as *Non-linear data structure*. Example: Tree, Graphs, etc.

**Static Data Structure:** The data structure in which the number of elements is fixed, is known as *Static Data Structure*. Example: Arrays

**Dynamic Data Structure:** The data structure in which the number of elements is not fixed, is known as *Dynamic Data Structure*. Example: Linked List.

## Array

It is a static primitive data structure. It is a homogeneous collection of data. The elements in the array are stored on consecutive memory locations. Array is also known as a subscripted variable, e.g.,  $A[i]$  is  $i^{\text{th}}$  element of the array  $A$  (i.e.,  $i$  is the subscript with variable  $A$ ). As we know that the elements of the array are stored on consecutive memory locations, it becomes convenient to find out the address of memory location of  $i^{\text{th}}$  element, for given base address (address of first element) and  $W$  (i.e. the number of memory locations required by one element).

### One dimensional array

Number of elements  $N = UB - LB + 1$  Where  $LB$ -Lower Bound  $UB$ -Upper Bound

Memory Location of  $A[i]$ ;  $\text{Loc}(A[i]) = \text{Base}(A) + W * (i - LB)$

$$\text{Loc}(A[i]) = \text{Base}(A) + W * i$$

Where  $N$  is given as in C++,  $LB$  is assumed as 0

### Two dimensional array

Number of elements =  $ROWS \times COLS = (UBI - LBI + 1) \times (UBJ - LBJ + 1)$

**Row Major:**  $\text{Loc}(A[I][J]) = \text{Base}(A) + W * (\text{COLS} * (I - LBI) + (J - LBJ))$

where  $\text{Base}(A)$  is address of first element's memory location

$\text{COLS}$  is number of columns =  $UBJ - LBJ + 1$

$W$  is number of memory locations required by one element

$LBI$  is Lower Bound of row

$UBJ$  is upper bound of column

$LBJ$  is Lower bound of column

$\text{Loc}(A[I][J]) = \text{Base}(A) + W * (\text{COLS} * I + J)$

where  $\text{COLS}$  is the number of columns,  $LBI=0$  and  $LBJ=0$

**Column Major:**  $\text{Loc}(A[I][J]) = \text{Base}(A) + W * (\text{ROWS} * (J - LBJ) + (I - LBI))$

where  $\text{Base}(A)$  is address of first element's memory location

$\text{ROWS}$  is the number of rows =  $UBI - LBI + 1$

$W$  is number of memory locations required by one element

$LBI$  is Lower Bound of Column

$UBI$  is upper bound of Row

$LBJ$  is Lower bound of Row

$\text{Loc}(A[I][J]) = \text{Base}(A) + W * (\text{ROWS} * J + I)$

where  $R$  is number of Rows,  $LBI=0$  and  $LBJ=0$

**Exercise 1** A one-dimensional array P[100] is stored in memory with a base address as 5000. Find out addresses of P[15] and P[40], if each element of this array requires 4 bytes.

Given,

$$\begin{aligned}
 \text{Base}(P) &= 5000 \\
 W &= 4 \\
 \text{Loc}(P[I]) &= \text{Base}(P) + W*I \\
 \text{Loc}(P[15]) &= 5000 + 4*15 \\
 &= 5000 + 60 \\
 &= 5060 \\
 \text{Loc}(P[40]) &= 5000 + 4*40 \\
 &= 5000 + 160 \\
 &= \underline{\underline{5160}}
 \end{aligned}$$

**Exercise 2** A one-dimensional array A[-5..25] is stored in memory with each element requiring 2 bytes. If the base address is 8000, find out the following:

- a) Address of A[5] and A[-3]
- b) Total no. of elements present in the array

Given,

$$\begin{aligned}
 \text{Base}(A) &= 8000 \\
 W &= 2 \\
 LB &= -5 \\
 \text{Loc}(A[I]) &= \text{Base}(A) + W*(I-LB) \\
 \text{Loc}(A[5]) &= 8000 + 2*(5-(-5)) \\
 &= 8000 + 20 \\
 &= 8020 \\
 \text{Loc}(A[-3]) &= 8000 + 2*(-3-(-5)) \\
 &= 8000 + 4 \\
 &= \underline{\underline{8004}}
 \end{aligned}$$

$$\text{Total} \quad \text{No. of Elements} = UB-LB+1 = 25 - (-5) + 1 = 31$$

**Exercise 3** A two-dimensional array Q[5][15] is stored in memory along the row with each element requiring 2 bytes. If the base address is 6500, find out the following:

- a) Addresses of Q[5][10] and Q[3][5]
- b) Total no. of elements present in the array

Given,

$$\begin{aligned}
 \text{Base}(Q) &= 6500 \\
 W &= 2 \\
 COLS &= 5 \\
 \text{Row Major, Loc}(Q[I][J]) &= \text{Base}(Q) + W*(COLS*I+J) \\
 \text{Loc}(Q[5][10]) &= 6500 + 2*(15*5+10) \\
 &= 6500 + 170 \\
 &= 6670 \\
 \text{Loc}(Q[3][5]) &= 6500 + 2*(15*3+5) \\
 &= 6500 + 100 \\
 &= \underline{\underline{6600}}
 \end{aligned}$$

$$\text{Total} \quad \text{No. of Elements} = ROWS*COLS = 5*15 = 75$$

**Exercise 4** R[-4..4,7..17] is a two-dimensional array, stored in the memory along the column with each element requiring 4 bytes. If the base address is 5000, find out the following:

- a) Addresses of R[2][10] and R[3][15]
- b) Total no. of elements present in the array

Given,

$$\begin{aligned}
 \text{Base}(R) &= 5000 \\
 W &= 4 \\
 ROWS &= UBI-LBI+1 = 4 - (-4) + 1 = 9 \\
 COLS &= UBJ-LBJ+1 = 17 - 7 + 1 = 11
 \end{aligned}$$

$$\begin{aligned}
 \text{Column Major, Loc}(R[I][J]) &= \text{Base}(R) + W*(ROWS*(J-LBJ)+(I-LBI)) \\
 \text{Loc}(R[2][10]) &= 5000 + 4*(9*(10-7)+(2-(-4))) \\
 &= 5000 + 4*27 \\
 &= 5132 \\
 \text{Loc}(R[3][15]) &= 5000 + 4*(9*(15-7)+(3-(-4))) \\
 &= 5000 + 4*72 \\
 &= \underline{\underline{5316}}
 \end{aligned}$$

$$\begin{aligned}
 \text{Total} \quad \text{No. of Elements} &= ROWS*COLS = 9*11 \\
 &= 99
 \end{aligned}$$