CSCE 625

Assignment -1

Fall 2017

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Objective:

To implement A* search and develop a heuristic that will solve the Blocksworld problem more efficiently.

Implementation:

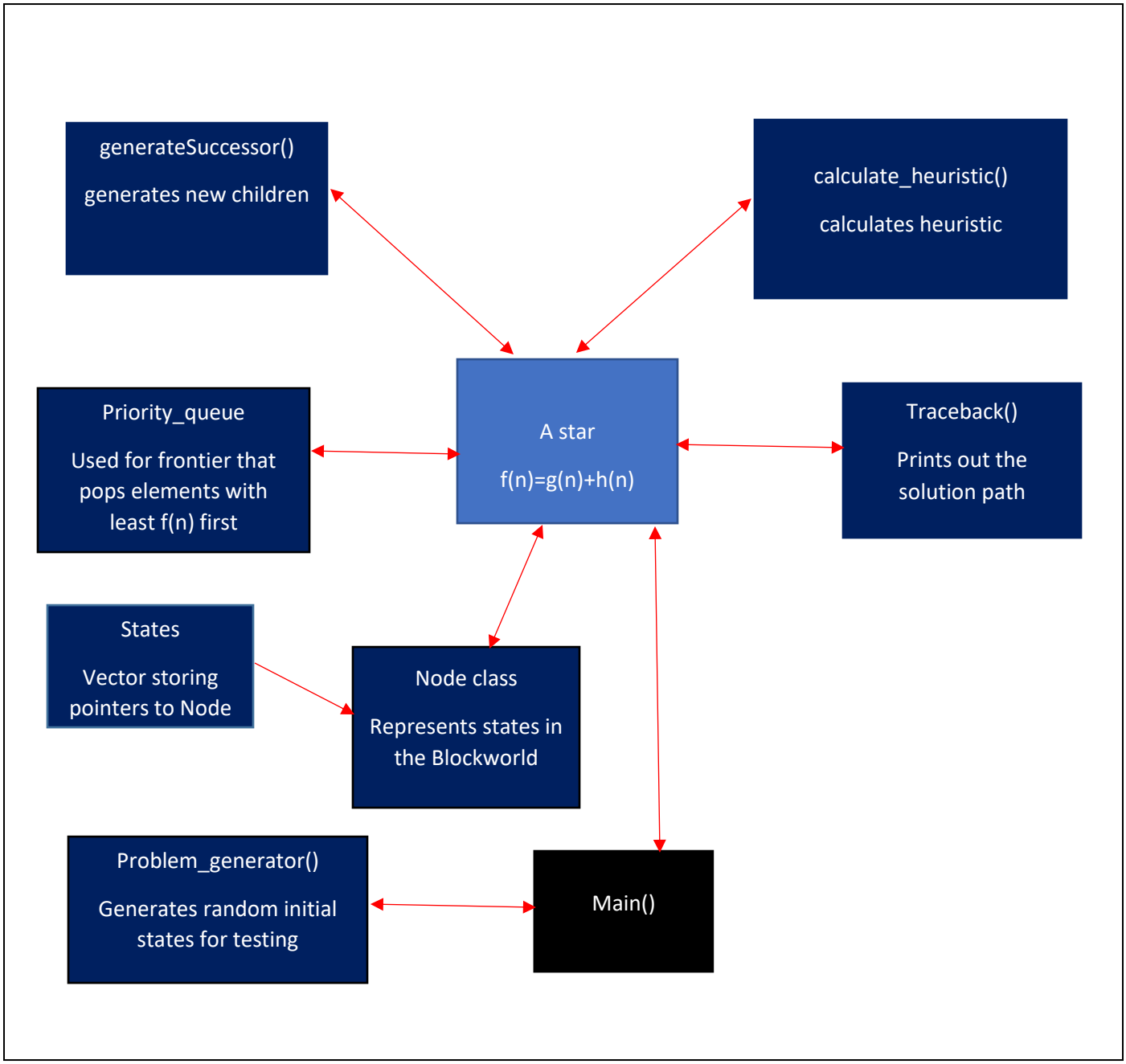

Table of Performance Metrics:

Blocks	Stacks	Total Goal Tests/ Number of Iterations	Max Queue Size	Depth	Time (Seconds)	
5	3	9	18	9	0.0014	Success
5	10	7	67	7	0.0634	Success
8	4	13	78	13	0.015	Success
10	5	15	157	15	0.059	Success
10	10	14	481	14	0.64	Success
15	10	24	966	24	3.60	Success
15	15	21	1531	21	9.24	Success
18	15	25	1684	25	11.44	Success
20	15	29	2616	29	20.64	Success
23	17	39	3687	39	52.36	Success
25	18	38	4126	38	82.24	Success
26	20	41	5930	41	174.18	Success
26	26	41	7208	41	378.02	Success

Heuristic Description:

The below heuristic is used to solve the problem efficiently.

h = 22.5 * **p_0** - 12.5 * **p_2** - 1.5***p_1** + 5.5***sum_1** + 10.5***sum_2**

p_0 -> Total Number of blocks out of place on all the stacks.

p_1 -> Number of blocks out of place on the first stack compared to Goal state.

p_2 -> Total Number of blocks on all the stacks except the 1st stack.

sum_1 -> Each block 'b1' on the stacks other than 1st stack is compared with its below blocks on the same stack and if ASCII (b1) > below block, then 1 point is incremented for b1. This check for b1 will be done with all of its below blocks one by one and then moving to next stack and repeating this process for all the elements on that stack.

sum_2 - > This is only for the blocks on 1st stack and its opposite to that of the above sum_1 calculation procedure. If for each block on stack 1 let's say 'b1' if ASCII(b1) < below block then 1 point is incremented. This check for b1 will be done with all the blocks which are below it.

- Coefficients are added by analyzing the weightage ratio that needs to be given for each and their values are assigned by trial and error after checking with multiple test cases.

A* Algorithm used:

(f,g are initially set to infinite, h=0 for every node.)

function AStar returns a solution or failure

States <- explored set

frontier <- priority queue ordered by PATH-COST (f(n)), with only one initial element.

loop do

IF *frontier* empty **then return** failure

node <-POP from frontier /* Chooses the lowest cost node in the frontier based on f(n)*/

IF GOAL-TEST(*node*.STATE) **then return** Solution(node)

IF iterations > 1000 **return** failure

child_node <- Generate children

For each generated children:

g'= (node->g) + 1;

IF *child_node* is not in explored or frontier **then**

Child_node - >g =g'

Calculate heuristic and f=g+h

Add child_node.state to explored (States)

Frontier <- Insert(child_node,frontier)

ELSE IF *child_node* is in explored or frontier with higher Path Cost **then**

Replace that with new *child_node*

Discussion:

- For most of my test cases, I observed that **Total Goal tests = Depth**. That means I'm directly moving towards the goal without visiting unnecessary nodes. Hence I strongly think that my Heuristic is **admissible** as its not overestimating and also in many cases its prediction is same as the true path length to the goal.

- I was able to solve 26 blocks and 25 stacks with an average depth of 41, which seems pretty good for me. Although there are few cases where Total Goal tests is little more than the depth but still final path is optimal.

- During implementation I felt that estimating sum_1 and sum_2 are crucial for the heuristic to work effectively. Initially when I was using a simple heuristic with number of blocks out of place it was doing >1000 Goal tests for 7 blocks 5 stacks. So adding sum_1 and sum_2 to heuristic have improved the search by a great factor.

- Using the C++ STL library's Priority Queue, I was unable to iterate through its elements so I have extended its functionality in order to iterate through its elements.

- Number of Iterations/ Goal Tests is NOT exceeding 41 for any number of blocks and stacks input.

Ideas to Improve:

- Queue Size is >2000, if more than 20 blocks is given as input. Although Heuristic is working perfectly because of the excess queue size program run time is increasing as it needs to iterate through all the elements of Priority queue to check if the node already exists. This checking will be done for every node. Which would be of the order O(n^2).

- Additionally, Priority queue is built using heap and delete operation in heap causes a lot of trouble as heap doesn't support search by value. This would be O(n*log(n)). Hence I would suggest using a Hash Map/Unordered set/ 2-3 tree which does search, delete, insert operations in O(1), O(log n) time respectively. This way I think we can reduce the Running time.

Thank You so much, Patrick