

Assignment-Unit-3

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PART-A

Q.1 - What do you mean by partial order set?

Ans: A partially order set (Normally Poset) is a set A , together with a relation, \leq , that obeys below properties

(i) For all $a, b, c \in A$ (Reflexive), $(a, a) \in A$ is also

(ii) For all $a, b \in A$ if $a R b$ and $b R a$ then $a = b$ (Anti-Symmetric Property)

(iii) for $a, b, c \in A$ if $a R b$ and $b R c$ then $a R c$ (Transitive property)

Q.2 Define Hasse Diagram.

Ans: A Hasse diagram is a graphical representation of a finite partially order set (POSet) that shows the ordering of elements without displaying all the relations explicitly.

Q.3 What are Lattices?

Ans A Lattice is a special kind of poset, in which every pair of elements has both;

(i) A least upper bound (LUB) - Called the join (denoted $a \vee b$)

(ii) A greatest lower bound (GLB) - Called the meet (denoted $a \wedge b$)

Q-4 - Define Karnaugh maps.

Ans - A Karnaugh map (k-map) is a graphical tool used in digital logic design to simplify boolean algebra expressions.

It helps to reduce logic functions visually.

Q-5 - What is complemented Lattices?

Ans - A complemented lattices is a type of bounded

Lattices in which every element has a complement.

Q-6 - Define Modular and Complete Lattices.

Ans - Modular Lattices - A Lattices (A, \vee, \wedge) is called

a modular lattices if it satisfies the modular law:

For all $a, b, c \in A$ if, $a \leq c$ then $a \vee (b \wedge c) = (a \vee b) \wedge c$

Complete Lattices - A Lattices (A, \vee, \wedge) is called

a complete lattices if every subset of A has:

(1) A least upper bound (LUB) : $\forall S \exists u \in A$ such that

(2) A greatest lower bound (GLB) : $\forall S \exists l \in A$ such that

(join) $\forall s \in S \ s \leq u$
(meet) $\forall s \in S \ s \geq l$

Q-7 - What do you mean by the Morphisms of Lattices?

Ans - Morphisms are structure-preserving functions between lattices. They help to compare or map one lattice to another while maintaining lattice operations.

(Ans) $f: (A, \vee, \wedge) \rightarrow (B, \vee, \wedge)$ is a morphism if $f(a \vee b) = f(a) \vee f(b)$ and $f(a \wedge b) = f(a) \wedge f(b)$

Part-B: Short Questions and Answers

Q-8 - Define posets, lattices and Boolean algebra.

Ans - Poset - A poset is a set 'A' equipped with a binary relation ' \leq ' that satisfies:

(i) Reflexive - $a \leq a$

(ii) Anti-Symmetric - if $a \leq b$ and $b \leq a$ then $a = b$

(iii) Transitive - if $a \leq b$ and $b \leq c$ then $a \leq c$

Lattice - A lattice is a poset in which every pair of elements has:

(i) A least upper bound (join), denoted $a \vee b$.

(ii) A greatest lower bound (meet), denoted $a \wedge b$.

Boolean Algebra - A boolean algebra is a complemented distributive lattice with:

(i) A least element (0) and greatest element (1)

(ii) complements for every element a : a' such that:

$$(a) a \vee a' = 1$$

$$(b) a \wedge a' = 0$$

Q-9 - If the set $P = \{1, 2, 3, 4, 6, 12\}$ with the relation \leq

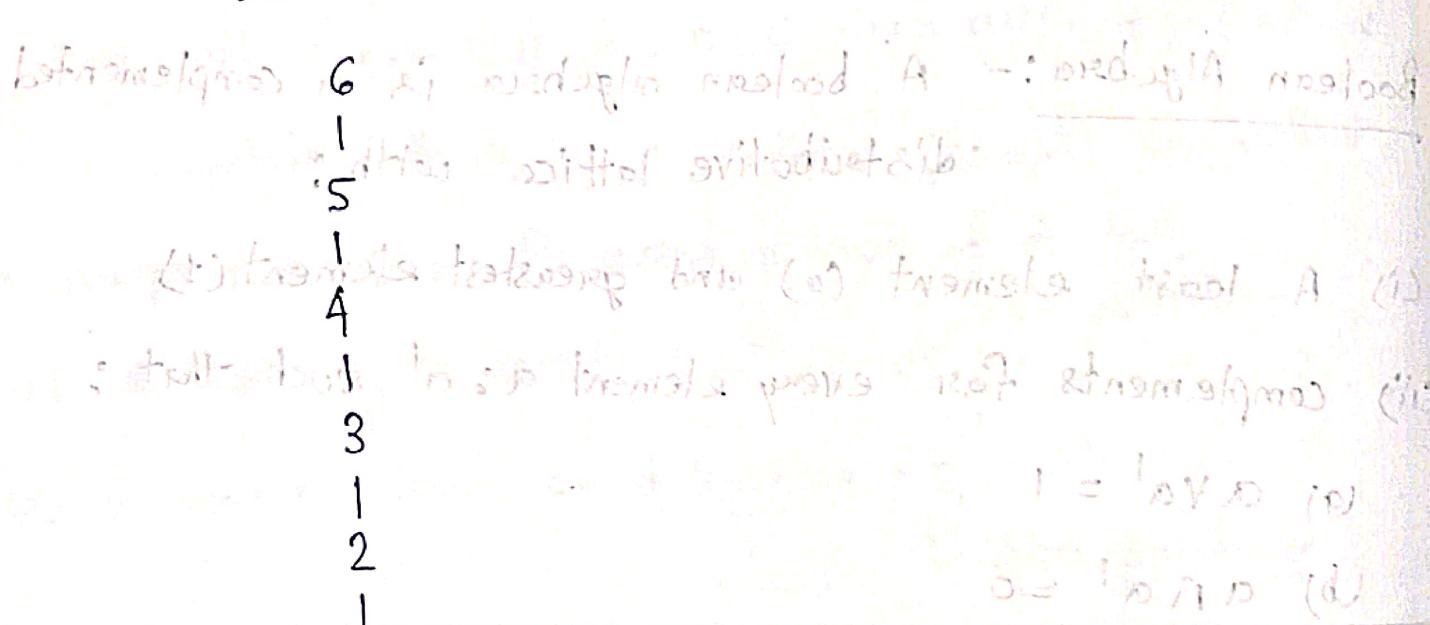
defined as: $a \leq b \Leftrightarrow a \text{ divides } b$. Is P poset under the given operation?

Ans - Yes, $P = \{1, 2, 3, 4, 6, 12\}$ with relation \leq defined as:

$$a \leq b \Leftrightarrow a \text{ divides } b$$

- (i) Reflexivity - Every element divides itself.
- $\forall a \in P, a|a \Rightarrow a \leq a$
- (ii) Anti-Symmetry - If a/b and b/a , then $a=b$.
This is true because the only way two positive integers divide each other is if they are equal.
- (iii) Transitivity - If a/b and b/c then a/c .
This is known as property of divisibility.

- Q-10 - Let $A = \{1, 2, 3, 4, 5, 6\}$ with relation \leq defined as "a is less than or equal to b". draw Hesse diagram of (A, \leq) .
- Solve - draw Hesse diagram



\leq relation is reflexive, antisymmetric, transitive, showing the ordering relation between elements.

$$d \text{ divides } a \Leftrightarrow a \geq d$$

Q-11. If $a, b \in D_{18}$ and $a+b = \text{lcm}(a, b)$, $a \cdot b = \text{gcd}(a, b)$

and $a' = \frac{18}{a}$. Show that $\{D_{18}, +, \cdot, ', 1, 18\}$ is

not a Boolean algebra.

$\therefore D_{18}$ represents the divisors of 18

Solve - $\therefore D_{18} = \{1, 2, 3, 6, 9, 18\}$

Need to show that $\{D_{18}, +, \cdot, ', 1, 18\}$ is not a Boolean algebra.

Given, $a+b = \text{lcm}(a, b)$

$a \cdot b = \text{gcd}(a, b)$

$$a' = \frac{18}{a}$$

\therefore we know that A Boolean algebra must satisfy certain properties, including distributivity, existence of complement, and existence of identity elements for both operations.

Let $a = 6 \in D_{18}$. Then

$$a' = \frac{18}{6} = 3 \in D_{18} \quad (\text{satisfied})$$

$$a+a' = \text{lcm}(6, 3) = 6 \neq 18 \Rightarrow a+a' \neq 1 \quad (\text{Not satisfied})$$

$$a \cdot a' = \text{gcd}(6, 3) = 3 \neq 1 \Rightarrow a \cdot a' \neq 0 \quad (\text{Not satisfied})$$

The complement operation does not satisfy the required properties:

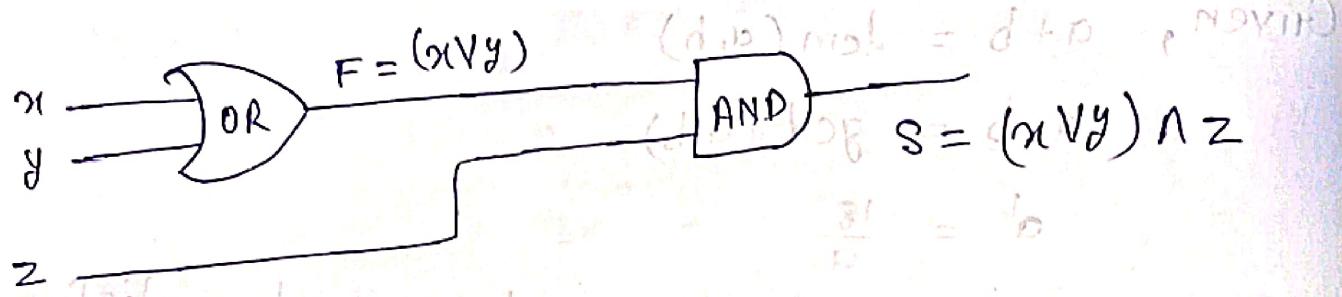
$$a+a' \neq 1 \quad \text{and} \quad a \cdot a' \neq 0$$

So $\{D_{18}, +, \cdot, ', 1, 18\}$ is not a Boolean Algebra.

Q-12 - Draw the circuit represented by the Boolean function $f(x, y, z) = (x \vee y) \wedge z$.

Ans - The circuit for the boolean function $f(x, y, z) = (x \vee y) \wedge z$ can be drawn using two logic gates.

- (1) OR gate for $(x \vee y)$ - x and y are input.
- (2) AND gate for $(x \vee y) \wedge z = (x \vee y) \text{ and } z$ are input.



Q-13 - Draw the logic networks corresponding to the following Boolean expressions $xy + x\bar{y}$.

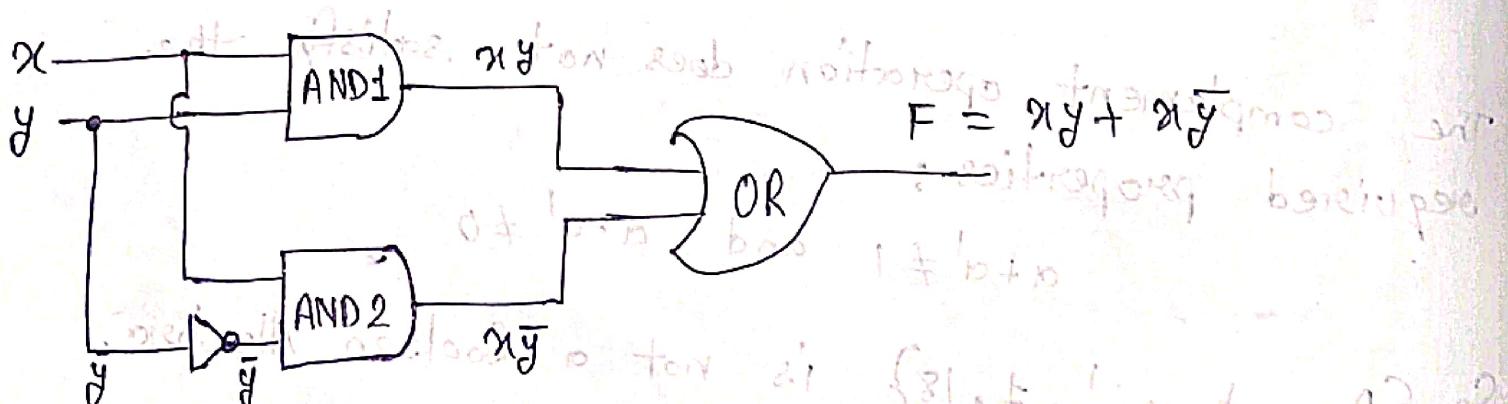
Ans - $F = xy + x\bar{y} = x(y + \bar{y}) = x \cdot 1 = x$ as per complement

(1) AND gate 1 for (xy) - x, y are input

(2) NOT gate for (\bar{y}) is input

(3) AND gate 2 for $(x\bar{y})$ - x, \bar{y} are input

(4) OR gate for $(xy + x\bar{y})$ - $xy, x\bar{y}$ are input



Q-14 - Draw the logic networks corresponding to the following Boolean expressions $xy'z' + x'y'z + xy'$

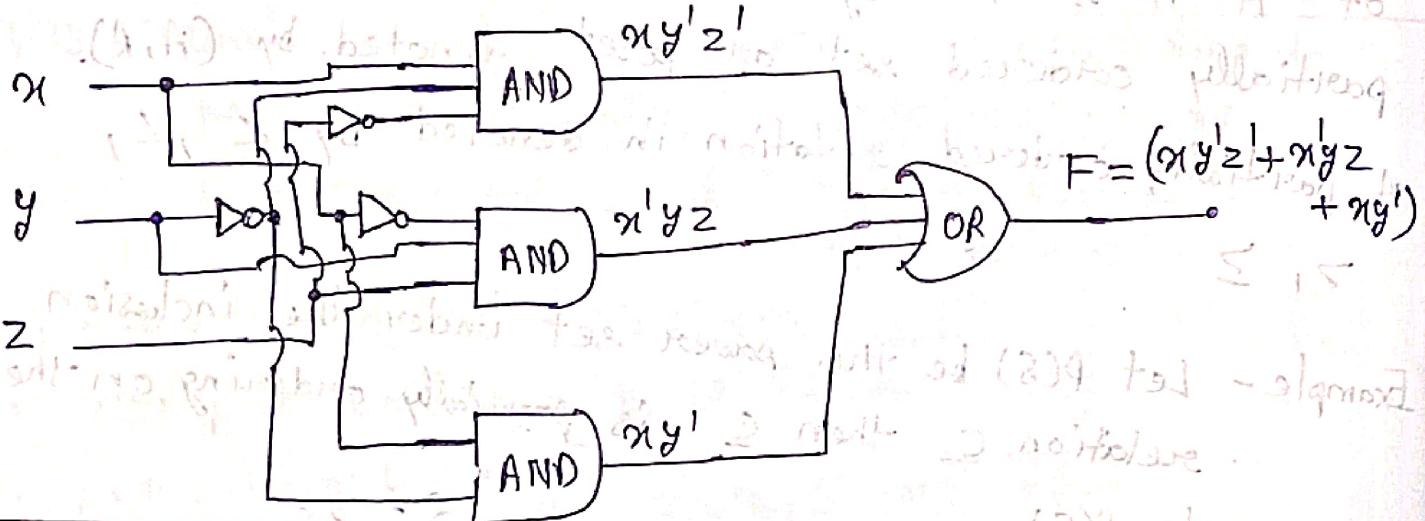
Ans - The circuit or logic networks for $xy'z' + x'y'z + xy'$ can be drawn using 3 (p, q), A \Rightarrow p AND q

(i) 3 NOT gates : for $\bar{x}, \bar{y}, \bar{z}$

(ii) 3 AND gates : one for each product term ($x'y'z'$, $x'y'z$, xy')

(iii) 1 OR gate : to combine all three terms

$F = (x'y'z' + x'y'z + xy')$



Part C - Long Questions and Answers

Q-15 - Which of the following relations on $A = \{0, 1, 2, 3\}$ are partial ordering? Determine the properties of a partial ordering that the others lack?

- $R_1 = \{(0,0), (1,1), (2,2), (3,3)\}$
- $R_2 = \{(0,0), (1,1), (2,0), (2,3), (2,2), (3,2), (3,3)\}$
- $R_3 = \{(0,0), (1,1), (1,2), (2,2), (3,3)\}$

Ans - To determine if a relation on $A = \{0, 1, 2, 3\}$ is a partial ordering, it must satisfy three properties:

1. Reflexivity : For all a in A , (a, a) is in R .
2. Antisymmetry : For all a and b in A , if $a R b$ and $b R a$, then $a = b$.
3. Transitivity : For all a, b , and c in A , if $a R b$ and $b R c$, then $a R c$.

(a) $R_1 = \{(0,0), (1,1), (2,2), (3,3)\}$

1. Reflexivity - Yes, (a, a) is in R_1 for all a in A .
2. Antisymmetry - Yes, there are no pairs (a, b) and (b, a) in R_1 where $a \neq b$.
3. Transitivity - Yes, there are no pairs (a, b) and (b, c) in R_1 where $a \neq b$ and $b \neq c$, so transitivity holds vacuously.

- Therefore, R_1 is a partial ordering.

(b) $R_2 = \{(0,0), (1,1), (2,0), (2,3), (2,2), (3,2), (3,3)\}$

1. Reflexivity :- Yes, (a, a) is in R_2 for all a in A .
2. Antisymmetry :- No, $(2, 3)$ and $(3, 2)$ are in R_2 , but $2 \neq 3$.
3. Transitivity :- Not applicable due to failure of antisymmetry.

- Therefore, R_2 is not a partial ordering due to lack of antisymmetry.

(c) $R_3 = \{(0,0), (1,1), (1,2), (2,2), (3,3)\}$

1. Reflexivity :- Yes, (a, a) is in R_3 for all a in A .

2. Antisymmetric :- Yes, if there are no pairs (a, b) and (b, a) in R_3 where $a \neq b$

3. Transitive :- Yes, for $(1, 2)$ and no other pair starting with 2, transitivity holds and hence R_3 is transitive.

Therefore, R_3 is a partial ordering.

Q-16 - Let D_m denote the positive divisor of m ordered by divisibility. Draw the Hasse diagram of

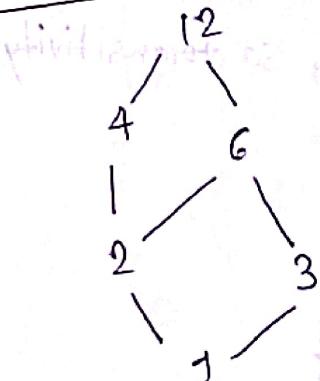
$$(a) D_{12} \quad (b) D_{15}$$

$$(c) D_{16} \quad (d) D_{17}$$

Ans -

~~(a) $D_{12} = \{1, 2, 3, 4, 6, 12\}$~~

Hasse Diagram -



$$(b) D_{15}$$

$$D_{15} = \{1, 3, 5, 15\}$$

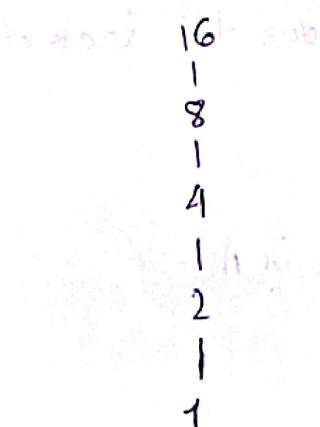
Hasse Diagram -



$$(c) D_{16}$$

$$D_{16} = \{1, 2, 4, 8, 16\}$$

Hasse Diagram -



$$(d) D_{17}$$

$$D_{17} = \{1, 17\}$$

Hasse Diagram -



Q-17 - Let $L_1 = \{a, b, c\}$ and $L_2 = \{1, 2, 3, 5, 6, 10, 15, 30\}$
 be the sets of two alphabets and multiplies of
 30 respectively. Prove that the set $\{P(L_1), \cap, \cup\}$
 and $\{L_2, \gcd, \text{lcm}\}$ are isomorphic.

Solve - Given, $L_1 = \{a, b, c\}$, $L_2 = \{1, 2, 3, 5, 6, 10, 15, 30\}$

Need to prove,

$$\{P(L_1), \cap, \cup\} \cong \{L_2, \gcd, \text{lcm}\}$$

$$P(L_1) = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$$

Each subset of L_1 can be mapped to a divisor of 30, L_2

$$\{\emptyset\} \rightarrow 1 \quad \{c\} \rightarrow 5 \quad \{b, c\} \rightarrow 15$$

$$\{a\} \rightarrow 2 \quad \{a, b\} \rightarrow 6 \quad \{a, b, c\} \rightarrow 30$$

$$\{b\} \rightarrow 3$$

$$\{a, c\} \rightarrow 10$$

This mapping is bijective and preserves:

$$\text{Intersection } (\cap) \longleftrightarrow \gcd$$

$$\text{Union } (\cup) \longleftrightarrow \text{lcm}$$

Example -

$$\{a\} \cap \{b\} = \emptyset \rightarrow \text{maps to } \gcd(2, 3) = 1$$

$$\{a\} \cup \{b\} = \{a, b\} \rightarrow \text{maps to lcm}(2, 3) = 6$$

Since the structure is preserved and mapping is one to one

$$\{P(L_1), \cap, \cup\} \cong \{L_2, \gcd, \text{lcm}\}$$

Hence, they are isomorphic. Proved

Q-18 - If $a, b \in D_{20}$ and $a+b = \text{lcm}(a, b)$, $a \cdot b = \text{gcd}(a, b)$
 and $a' = \frac{20}{a}$ shows that $\{D_{20}, +, \cdot, ', 1, 20\}$
 is not a Boolean algebra.

Solve - $\because D_{20}$ represents the divisors of 20

$$\therefore D_{20} = \{1, 2, 4, 5, 10, 20\}$$

Need to show that $\{D_{20}, +, \cdot, ', 1, 20\}$ is not a
 Boolean algebra.

Given, $a+b = \text{lcm}(a, b)$ and $a \cdot b = \text{gcd}(a, b)$

$$a \cdot b = \text{gcd}(a, b)$$

$$a' = \frac{20}{a}$$

\because we know that A Boolean algebra must satisfy certain properties
 including distributivity, existence of complement, and
 existence of identity elements for both operations.

Let, $a = 2 \in D_{20}$; Then satisfied in given set

$$a' = \frac{20}{2} = 10 \in D_{20} \quad (\text{satisfied})$$

$$a+a' = \text{lcm}(2, 10) = 10 \neq 20 \quad (\text{not satisfied})$$

$$a \cdot a' = \text{gcd}(2, 10) = 2 \neq 1 \quad (\text{not satisfied})$$

The complement operation does not satisfy boolean algebra required properties.

So, $\{D_{20}, +, \cdot, ', 1, 20\}$ is not a Boolean Algebra

Q-19 - Consider the Boolean algebra $\{D_{210}, +, \cdot, ', 1, 210\}$
 where $a+b = \text{lcm}(a,b)$, $a \cdot b = \text{gcd}(a,b)$, $a' = \frac{210}{a}$

Then find atoms of D_{210} ?

Is $X = \{1, 2, 105, 210\}$ a subalgebra?

Is $Y = \{1, 2, 3, 6\}$ a subalgebra?

Find two sub algebras of 8 elements.

Solved $\therefore D_{210}$ represents the divisors of 210

$$\therefore D_{210} = \{1, 2, 3, 5, 6, 7, 10, 14, 15, 21, 30, 35, 42, 70, 105, 210\}$$

(i) Finding Atoms of D_{210} -

\because we know that Atoms in a Boolean algebra are the minimal elements that cover the least element (1 in this case). In the context of divisors and the operations defined:

(1) Prime Divisors : The atoms of D_{210} are the prime divisors of 210

(2) Prime factorization of 210 $\Rightarrow 210 = 2 \times 3 \times 5 \times 7$

\therefore Atoms of D_{210}

The atoms are the prime factors: $\{2, 3, 5, 7\}$.

(ii) check if $X = \{1, 2, 105, 210\}$ is a subalgebra.

For X to be a subalgebra, it must be closed under $+, \cdot$, and $'$ operations.

$$2+105 = \text{lcm}(2, 105) = 210 \quad (\text{in } X)$$

$$2 \cdot 105 = \text{gcd}(2, 105) = 1 \quad (\text{in } X)$$

$$2' = \frac{210}{2} = 105 \text{ (in } X)$$

$$105' = \frac{210}{105} = 2 \text{ (in } X) \Rightarrow (5, 8, 10) \in S \text{ (2)}$$

Therefore, X is a subalgebra. $\Rightarrow (5, 8, 10) \in S$ (1)

(iii) check if $Y = \{1, 2, 3, 6\}$ is a subalgebra. \exists (2)

For Y to be subalgebra, it must be closed under the operations:

$$2+3 = \text{lcm}(2, 3) = 6 \text{ (in } Y)$$

$$2 \cdot 3 = \text{gcd}(2, 3) = 1 \text{ (in } Y)$$

$$\frac{2+3}{2} = \frac{210}{2} = 105 \text{ (not in } Y)$$

Therefore, Y is a subalgebra.

(iv) Find two subalgebra of 8 elements

To find subalgebra of 8 elements, we look for subsets of D_{210} that are closed under the operations and contain 8 elements, including 1 and 20, and also that are closed under lcm gcd, and the complement operation defined by $d = \frac{210}{a}$

One potential subalgebra is $\{1, 2, 3, 6, 35, 70, 105, 210\}$.

Another potential subalgebra is $\{1, 2, 5, 10, 21, 42, 105, 210\}$.

Q-20 - Express $E(x,y,z)$ in its complete sum-of-products form.

$$(a) E(x,y,z) = (x'(y+z))' + x'y = \frac{012}{231} = 1201$$

$$(b) E(x,y,z) = y(x'+yz)' \text{ is a different}$$

$$(c) E(x,y,z) = x(xy+y'+x'y) = xy+x'y$$

Solve - To express each Boolean function $E(x,y,z)$ in its complete sum of products (SOP) form,

we follow these steps:

1. Simplify the expression using Boolean algebra.
2. Find the truth table.
3. Write the minterms (product term where output=1).
4. Express the function as a sum of those minterms.

$$(a) E(x,y,z) = (x'+y) + x'y \quad \text{---(i)}$$

apply De Morgan's Theorem.

$$(x'+y)' = x'' \cdot y' = x \cdot y' \quad \text{Put in Eq-(i)}$$

$$E(x,y,z) = x \cdot y' + x' \cdot y$$

Truth Table

x	y	z	$E(x,y,z)$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	0

Minterms (rows with 1)

$$(0,1,0) \rightarrow x'y'z'$$

$$(0,1,1) \rightarrow x'y'z$$

$$(1,0,0) \rightarrow xy'z'$$

$$(1,0,1) \rightarrow xy'z$$

Complete SOP form -

$$E(x,y,z) = x'y'z' + x'y'z + xy'z' + xy'z$$

(b) $E(x,y,z) = y(x+yz)'$

apply De Morgan's Theorem

$$(x+yz)' = x' \cdot (yz)' = x' \cdot (y'+z') \quad \text{Put in Eqn (ii)}$$

$$E(x,y,z) = y \cdot x' \cdot (y'+z') = yx'y' + yx'z' = 0 + yx'z'$$

$$E(x,y,z) = x'y'z' \quad \text{To prepare, 1 is the minterm}$$

Truth Table

x	y	z	$E(x,y,z)$
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

Only 1 in row $(0,1,0) \rightarrow x'y'z'$

Complete SOP form -

$$E(x,y,z) = x'y'z'$$

$$(c) E(x,y,z) = x(xy + y' + x'y)$$

Simplify the inner part

$$\begin{aligned} (xy + y' + x'y) &= y' + xy + x'y = y' + y(x+x') \\ &= y' + y(1) = y' + y = 1 \end{aligned}$$

Now,

$$E(x,y,z) = x \cdot 1 = x$$

Truth Table

x	y	z	E(x,y,z)
0	x	x	0
1	x	x	1

whenever $x=1$, regardless of y, z the output is 1,

so there are 4 rows with output 1:

$$(1,0,0) \rightarrow \text{minterm: } xy'z'$$

$$(1,0,1) \rightarrow \text{minterm: } x'y'z$$

$$(1,1,0) \rightarrow \text{minterm: } xy'z'$$

$$(1,1,1) \rightarrow \text{minterm: } xyz$$

x	y	z	E(x,y,z)
0	0	0	0
0	1	0	0
1	0	1	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

Even though the logic only depends on x , the complete SOP must include y and z for each case where the output is 1.

Complete SOP form -

$$E(x,y,z) = xy'z' + x'y'z + xy'z' + xyz$$

Q-21 - Construct circuit that produce the following output $F(x,y,z) = (x+y+z)(\bar{x}\bar{y}\bar{z})$

Ans - The circuit or logic networks for

$$F(x,y,z) = (x+y+z)(\bar{x}\bar{y}\bar{z})$$

can be draw using:

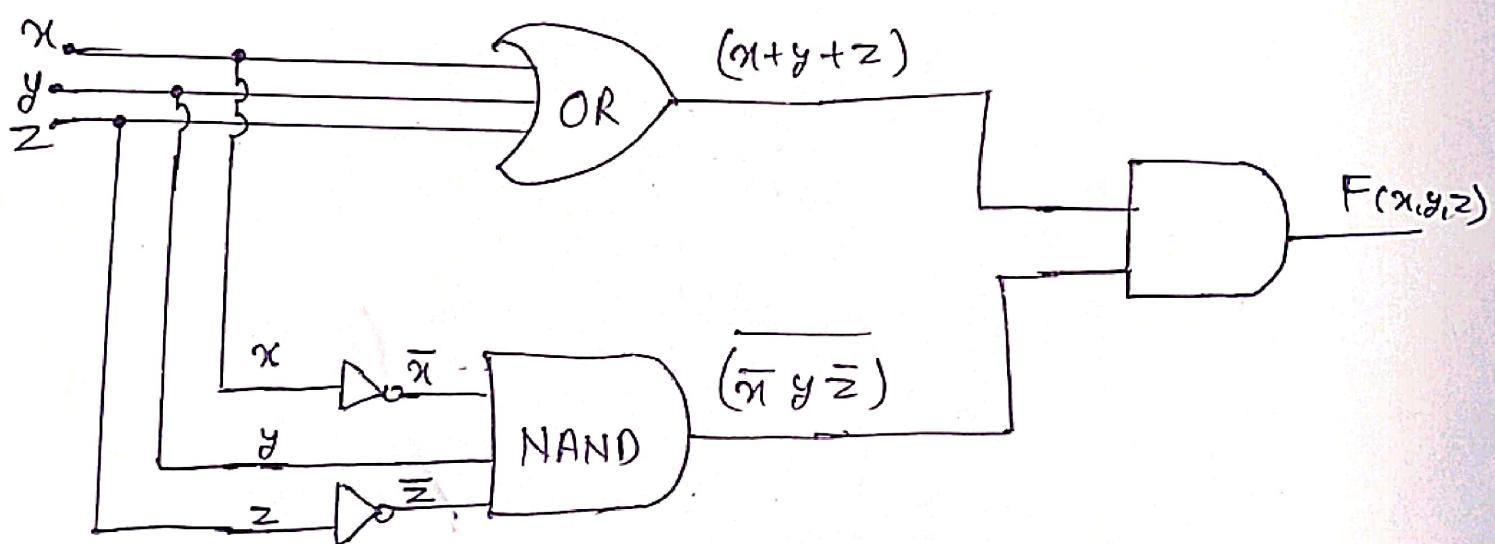
(i) 1 OR gate : for $x+y+z$

(ii) 2 NOT gate : for \bar{x}, \bar{y}

(iii) 1 NAND gate : for $(\bar{x}\bar{y}\bar{z})$

(iv) 1 AND gate : for $(x+y+z)(\bar{x}\bar{y}\bar{z})$

Circuit :-



Final Output $F(x,y,z) = (x+y+z)(\bar{x}\bar{y}\bar{z})$