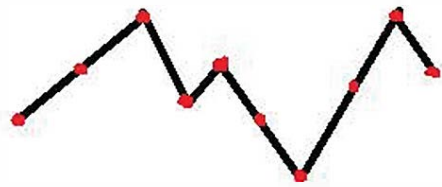


Stack

Approaches

1. Next greater/smaller element to right/left
2. Creating a graph of elements by their value and analyzing based on the top of the mountain value



```
//We are removing top of every mountain created by the trend
while(!s.empty() && s.top() > num[i] && k > 0){
    s.pop();
    k--;
}

while(!s.empty() && s.top() < nums[i]){
    thirdElement = s.top();
    s.pop();
}
```

Stack using arr

top = -1

empty -> if top = -1

full -> if top = size-1

Queue using arr

f = r = -1

empty -> if f == r

full -> if (rear+1) % size == f
(In circular queue)

-> if rear == size-1
(In normal queue)

3. Problems which need both NGETR and NGETL (Like Trapping) -> Where NGETL of s.top is 2nd element in stack and NGETR of s.top is current element
and NSETL and NSETR (Like Histogram area) -> Where NSETL of s.top is 2nd element in stack and NSETR of s.top is current element

```
int ans = 0;
stack<int> s;
for(int i = 0; i < height.size(); i++){
    while(!s.empty() && height[s.top()] < height[i]){
        int currHt = height[s.top()];
        s.pop();

        if(s.empty()) break;
        int lmax = s.top();
        int rmax = i;

        int width = rmax-lmax-1;
        int ht = min(height[lmax], height[rmax])-currHt;
        ans += ht*width;
    }
    s.push(i);
}
return ans;
```

```
int ans = 0;
for(int i = 0; i <= n; i++){
    int val = (i == n ? 0 : heights[i]);
    while(!s.empty() && heights[s.top()] >= val){
        int height = heights[s.top()];
        s.pop();
        int rb = i; //NSETR of s.top is current idx
        int lb = (!s.empty() ? s.top() : -1); //NSETL of s.top is next element in the stack

        ans = max(ans, height * (rb-lb-1));
    }
    s.push(i);
}
return ans;
```



```
int* NGETR(int *arr, int n){
    stack<int> s;
    int *ans = new int[n];
    ans[n-1] = -1;
    s.push(arr[n-1]);
    for(int i=n-2; i>=0; i--){
        while(!s.empty() && arr[i] >= s.top()){ /* Pop all smaller elements to right(which is stored in stack)
            s.pop();
        }
        if(s.empty()) /* IF stack is empty, it means there is no element greater than current element
            ans[i] = -1;
        else /* If not empty, it means there is an element greater than current element
            ans[i] = s.top();
            s.push(arr[i]);
    }
    return ans;
}
```

NGETR

```
//Find NGETR of all elements
int ngetr[n];
stack<int> s;
s.push(n-1); /* In NGETR problem, we used default value as -1, but here we use ARR LENGTH as default value
ngetr[n-1] = n;
for(int i=n-2; i>=0; i--){
    while(!s.empty() && arr[s.top()] <= arr[i]){
        s.pop();
    }
    if(s.empty())
        ngetr[i] = n;
    else
        ngetr[i] = s.top();
    s.push(i);
}
/*Since we are storing the index of NGETR in ngetr[] array, we check if the index of NGETR lies in the window(k) then jump
to NGETR and check the same for it, else the current index in the ngetr[] is the max of current window
for(int i=0; i<=n-k; i++){
    int j = i;
    while(ngetr[j] < i+k){
        j = ngetr[j];
    }
    cout<<arr[j]<<endl;
}
}
```

Sliding Window

Largest area of rectangle in histogram

```
int largestRectangleArea(vector<int>& heights) {
    vector<int> rb = NSETR(heights);
    vector<int> lb = NSETL(heights);

    int ans = 0;
    for(int i = 0; i < heights.size(); i++){
        int width = rb[i] - lb[i] - 1;
        int area = width * heights[i];
        ans = max(ans, area);
    }
    return ans;
}
```

Calculating NSETL and NSETR separately

```
/* The approach is same but here we are calculating NGETR while calculating NGETL
void NSETL(vector<int>& heights, vector<int>& rb, vector<int>& lb){
    int n = heights.size();
    stack<int> s;

    s.push(0);
    for(int i = 1; i < n; i++){
        while(!s.empty() && heights[s.top()] >= heights[i]){
            rb[s.top()] = i; /* NSETR of s.top is current element
            s.pop();
        }
        if(s.empty()) lb[i] = -1;
        else lb[i] = s.top();

        s.push(i);
    }
}

int largestRectangleArea(vector<int>& heights) {
    int n = heights.size();
    vector<int> rb(n, n), lb(n, -1);
    NSETL(heights, rb, lb);

    int ans = 0;
    for(int i = 0; i < heights.size(); i++){
        int width = rb[i] - lb[i] - 1;
        int area = width * heights[i];
        ans = max(ans, area);
    }
    return ans;
}
```

```
/*In this approach we are finding area of rectangle formed by top element of the stack (Not each element of array as before)
int largestRectangleArea(vector<int>& heights) {
    int n = heights.size();
    stack<int> s;

    int ans = 0;
    for(int i = 0; i <= n; i++){
        int val = (i == n ? 0 : heights[i]);
        while(!s.empty() && heights[s.top()] >= val){
            int height = heights[s.top()];
            s.pop();
            int rb = i; /* NSETR of s.top is current idx
            int lb = (s.empty() ? s.top() : -1); /* NSETL of s.top is next element in the stack

            ans = max(ans, height * (rb-lb-1));
        }
        s.push(i);
    }
    return ans;
}
```

The point we need to note here is

At any point of time the NSETR of s.top is current element and NSETL of s.top is 2nd element in stack


```

int minEle;
void pop(stack<int> &s){
    if(s.top() < minEle)
        minEle = 2*minEle - s.top();

    s.pop();
}

void push(stack<int> &s, int x){
    if(s.empty()){
        s.push(x);
        minEle = x;
    }
    else if(x < minEle){
        s.push(2*x - minEle);
        minEle = x;
    }
    else{
        s.push(x);
    }
}

int getMin(stack<int> &s){
    return minEle;
}

int getMin(stack<int> &s){
    return minEle;
}

int main(){
    stack<int> s;
    int choice;
    do{
        int x;
        cin>>x;

        push(s,x);
        cin>>choice;
    }while(choice);

    cout<<getMin(s)<<endl;
    pop(s);
    pop(s);
    cout<<getMin(s)<<endl;
    return 0;
}

```

getMin

Push(x) : Inserts x at the top of stack.

If stack is empty, insert x into the stack and make minEle equal to x.

If stack is not empty, compare x with minEle. Two cases arise:

1. If x is greater than or equal to minEle, simply insert x.
2. If x is less than minEle, insert ($2*x - \text{minEle}$) into the stack and make minEle equal to x.

For example, let previous minEle was 3. Now we want to insert 2. We update minEle as 2 and insert $2*2 - 3 = 1$ into the stack.

Pop() : Removes an element from top of stack.

Remove element from top. Let the removed element be y. Two cases arise:

1. If y is greater than or equal to minEle, the minimum element in the stack is still minEle.
2. If y is less than minEle, the minimum element now becomes ($2*\text{minEle} - y$), so update ($\text{minEle} = 2*\text{minEle} - y$).

This is where we retrieve previous minimum from current minimum and its value in stack.

For example, let the element to be removed be 1 and minEle be 2. We remove 1 and update minEle as $2*2 - 1 = 3$.

How it works

In push : IN here we check if the new value is lesser then minEle, if so, we modify the value and push to stack

$$x < \text{minEle}$$

$$x - \text{minEle} < 0$$

$$2x - \text{minEle} < 0 + x$$

$$2x - \text{minEle} < x$$

ie modified val < newMinValue --> Here the modified value is lesser than the min value.

Therefore modified value is the only value that is less than minVal (which is the identity of modified value)

In pop: In here we check if the top is a modified value, if so we get the prev min by ($\text{top} < \text{minEle}$)

we assign $\text{minEle} = 2*\text{minEle} - \text{top}$

$$2*\text{minEle} - \text{modified}$$

$$2*\text{minEle} - 2x - \text{prevMinEle}$$

$$2*\text{minEle} - 2*\text{minEle} - \text{prevMinEle}$$

$$\text{prevMinEle}$$

Therefore we get the prev min value