C++ Programming

Who?

Shiv Shankar Dayal

When?

March 10, 2018





You hopefully know this after the previous three talks...

General formalization of haplotyping.

Inputs

- A genotype matrix G.
- The rows of the matrix are taxa / individuals.
- The columns of the matrix are SNP sites / characters.

Outputs

- A haplotype matrix H.
- Pairs of rows in H explain the rows of G.
- The haplotypes in H are biologically plausible.

Our formalization of haplotyping.

Inputs

- A genotype matrix G.
- The rows of the matrix are individuals / taxa.
- The columns of the matrix are SNP sites / characters.
- The problem is directed: one haplotype is known.
- The input is biallelic: there are only two homozygous states (0 and 1) and one heterozygous state (2).

Outputs

- A haplotype matrix *H*.
- Pairs of rows in H explain the rows of G.
- The haplotypes in *H* form a perfect phylogeny.

We can do perfect phylogeny haplotyping efficiently, but ...

- 1 Data may be missing.
 - This makes the problem NP-complete . . .
 - ... even for very restricted cases.

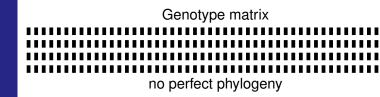
Solutions:

- Additional assumption like the rich data hypothesis.
- No perfect phylogeny is possible.
 - This can be caused by chromosomal crossing-over effects.
 - This can be caused by incorrect data.
 - This can be caused by multiple mutations at the same sites.

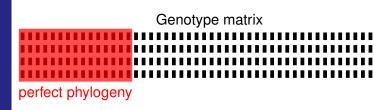
Solutions:

- Look for phylogenetic networks.
- Correct data.
 - Find blocks where a perfect phylogeny is possible.

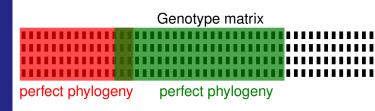
- Partition the site set into overlapping contiguous blocks.
- Compute a perfect phylogeny for each block and combine them.
- 3 Use dynamic programming for finding the partition.



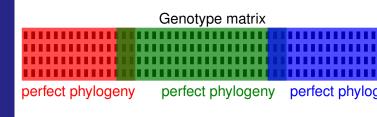
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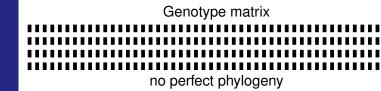


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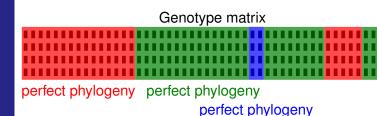
Objective of the integrated approach.

- Partition the site set into noncontiguous blocks.
- Compute a perfect phylogeny for each block and combine them.
- Compute partition while computing perfect phylogenies.



Objective of the integrated approach.

- 1 Partition the site set into noncontiguous blocks.
- Compute a perfect phylogeny for each block and combine them.
- Compute partition while computing perfect phylogenies.



The formal computational problem.

We are interested in the computational complexity of the function χ_{PP} :

- It gets genotype matrices as input.
- It maps them to a number k.
- This number is minimal such that the sites can be covered by k sets, each admitting a perfect phylogeny.
 (We call this a pp-partition.)

Finding pp-partitions of haplotype matrices.

We start with a special case:

- The inputs M are already haplotype matrices.
- The inputs *M* do not allow a perfect phylogeny.
- What is $\chi_{PP}(M)$?

Example

```
M:
0 0 0 1
0 1 0 0
1 0 0 0
1 0 0 0
0 1 0 0
1 0 0 0
0 1 0 1
1 1 0 0
0 0 1 0
```

No perfect phylogeny is possible.

Finding pp-partitions of haplotype matrices.

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- \blacksquare The inputs M are already haplotype matrices.
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- What is $\chi_{PP}(M)$?

Example



Perfect phylogeny Perfect phylogeny $\chi_{PP}(M) = 2$.

Bad news about pp-partitions of haplotype matrices.

Theorem

Finding optimal pp-partition of haplotype matrices is equivalent to finding optimal graph colorings.

Proof sketch for firm direction

Let G be a graph.

Build a matrix with a column for each vertex of G.

3

For each edge of G add four rows inducing the submatrix $\begin{pmatrix} 0 & 0 \\ 0 & 1 \\ 1 & 0 \end{pmatrix}$.

4

The submatrix enforces that the columns lie in different perfect phylogenies.

Implications for pp-partitions of haplotype matrices.

Corollary

If $\chi_{PP}(M) = 2$ for a haplotype matrix M, we can find an optimal pp-partition in polynomial time.

Corollary

Computing χ_{PP} for haplotype matrices is

- NP*-hard*,
- not fixed-parameter tractable, unless P = NP,
- very hard to approximate.

Finding pp-partitions of genotype matrices.

Now comes the general case:

- \blacksquare The inputs M are genotype matrices.
- The inputs *M* do not allow a perfect phylogeny.
- What is $\chi_{PP}(M)$?

Example

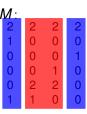
```
M:
2 2 2 2 2
1 0 0 0 0 No perfect phylogeny is possi-
0 0 0 1 ble.
0 0 2 2 0
```

Finding pp-partitions of genotype matrices.

Now comes the general case:

- The inputs *M* are genotype matrices.
- The inputs *M* do not allow a perfect phylogeny.
- What is $\chi_{PP}(M)$?

Example



Perfect phylogeny Perfect phylogeny $\chi_{PP}(M) = 2$.

Bad news about pp-partitions of haplotype matrices.

Theorem

Finding optimal pp-partition of genotype matrices is at least as hard as finding optimal colorings of 3-uniform hypergraphs.

Proof sketch.

- 1 Let G be a 3-uniform hypergraph.
- Build a matrix with a column for each vertex of G.
- For each hyperedge of G add four rows inducing the submatrix $\begin{pmatrix} 2 & 2 & 2 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$.
- The submatrix enforces that the three columns do not all lie in the same perfect phylogeny.

Implications for pp-partitions of genotype matrices.

Corollary

Even if we know $\chi_{PP}(M) = 2$ for a genotype matrix M, finding a pp-partition of any fixed size is still

- NP-hard,
- not fixed-parameter tractable, unless P = NP,
- very hard to approximate.

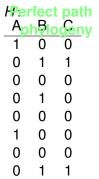
Automatic optimal pp-partitioning is hopeless, but...

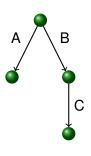
- The hardness results are worst-case results for highly artificial inputs.
- Real biological data might have special properties that make the problem tractable.
- One such property is that perfect phylogenies are often perfect path phylogenies: In HapMap data, in 70% of the blocks where a perfect phylogeny is possible a perfect path phylogeny is also possible.

Example of a perfect path phylogeny.

Genotype matrix

G:Haplotype
A B matrix
2 2 2 2
0 2 0
2 0 0
0 2 2





The modified formal computational problem.

We are interested in the computational complexity of the function χ_{PPP} :

- It gets genotype matrices as input.
- It maps them to a number k.
- This number is minimal such that the sites can be covered by k sets, each admitting a perfect path phylogeny.
 (We call this a ppp-partition.)

Good news about ppp-partitions of genotype matrices.

Theorem

Optimal ppp-partitions of genotype matrices can be computed in polynomial time.

Algorithm

- Build the following partial order:
- Can one column be above the other in a phylogeny?
 Can the columns be the two children of the root of a perfect path phylogeny?
- Cover the partial order with as few compatible chain pairs as possible.
 - For this, a maximal matching in a special graph needs to be computed.

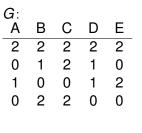
▶ The algorithm in action

Summary

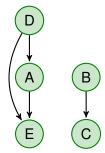
- Finding optimal pp-partitions is intractable.
- It is even intractable to find a pp-partition when just two noncontiguous blocks are known to suffice.
- For perfect path phylogenies, optimal partitions can be computed in polynomial time.

Computation of the partial order.

Genotype matrix



Partial order

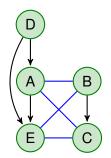


Partial order: →

Computation of the partial order.

Genotype matrix

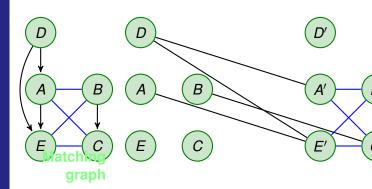
Partial order



Partial order: →

Compatible as children of root:

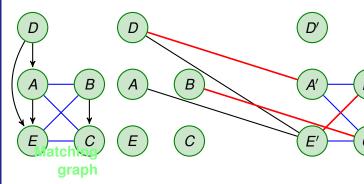
The matching in the special graph.



Partial order



The matching in the special graph.

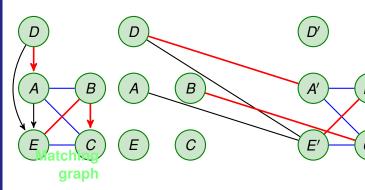


Partial order

A maximal matching in the matching graph



The matching in the special graph.



Partial order

A maximal matching in the matching graph induces perfect path phylogenies.

Return
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