



HEM: A Hardware-aware Event Matching Algorithm for Content-based Pub/Sub Systems

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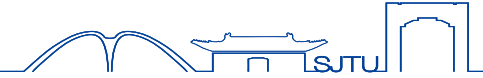
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Outline

- 1 Background
- 2 Problem
- 3 Motivation
- 4 Design
- 5 Evaluation
- 6 Future Work



1 Background



- **Content-based publish/subscribe system**
 - **composed of publishers, subscribers, and brokers.**
 - **on-demand data distribution.**
 - **realizes the decoupling of time, space and synchronization of communication parties.**
- **Event Matching Algorithm**
 - **find all matching subscriptions for each event as soon as possible.**

2 Problem--Matching Semantics



- **Event:** attribute-value pairs.
- **Subscription:** predicates/constraints.
- **Predicate/Constraint:** closed interval defined in an attribute.

$$E_1 = \{(a_1, 3), (a_3, 7)\}$$

$$S_1 = \{(a_1, [1, 5]), (a_2, [5, 8])\}$$

$$S_2 = \{(a_1, [3, 3]), (a_3, [1, 10])\}$$

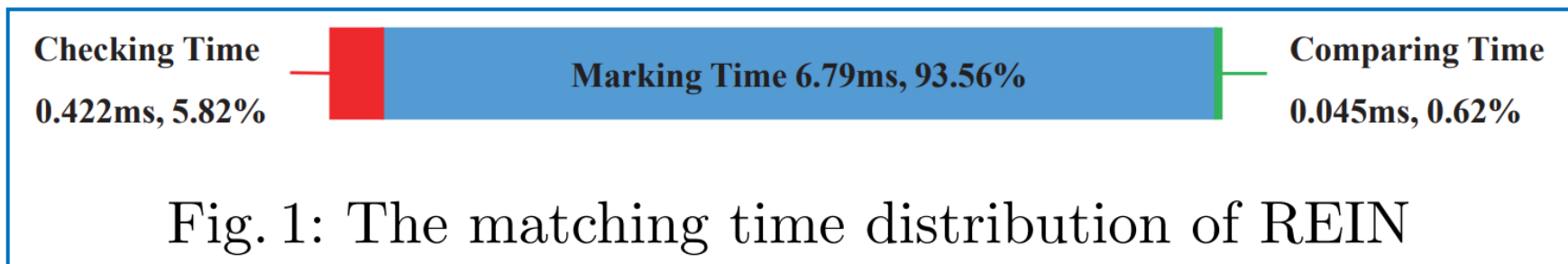
E_1 matches S_2

Size of a subscription \leq Size of an event \leq Dimension

3 Motivation



- The **essence** of **backward marking** based event matching algorithm
 1. searches and marks all the unmatching subscriptions in each attribute
 2. unmarked subscriptions are the matches of the event
- The **inherent defects** of backward marking
 1. tends to mark the same unmatching subscription for multiple times → redundant operations
 2. heavily dependent on dimension



3 Motivation

- Discovery
 - efficient OR operation on bitset with millions of bits
 - large memory capacity → trade space for time
- Breakthrough
 1. inserting: pre-mark subscriptions in bitsets according to their predicates' value distribution
 2. matching: select bitsets to execute several OR operations instead of millions of assignment operations

SubID: 1 2 3 4 5
(00110)

OR op

(10011)



(10111)



S₂ is the unique matches.

4 Design--Insertion

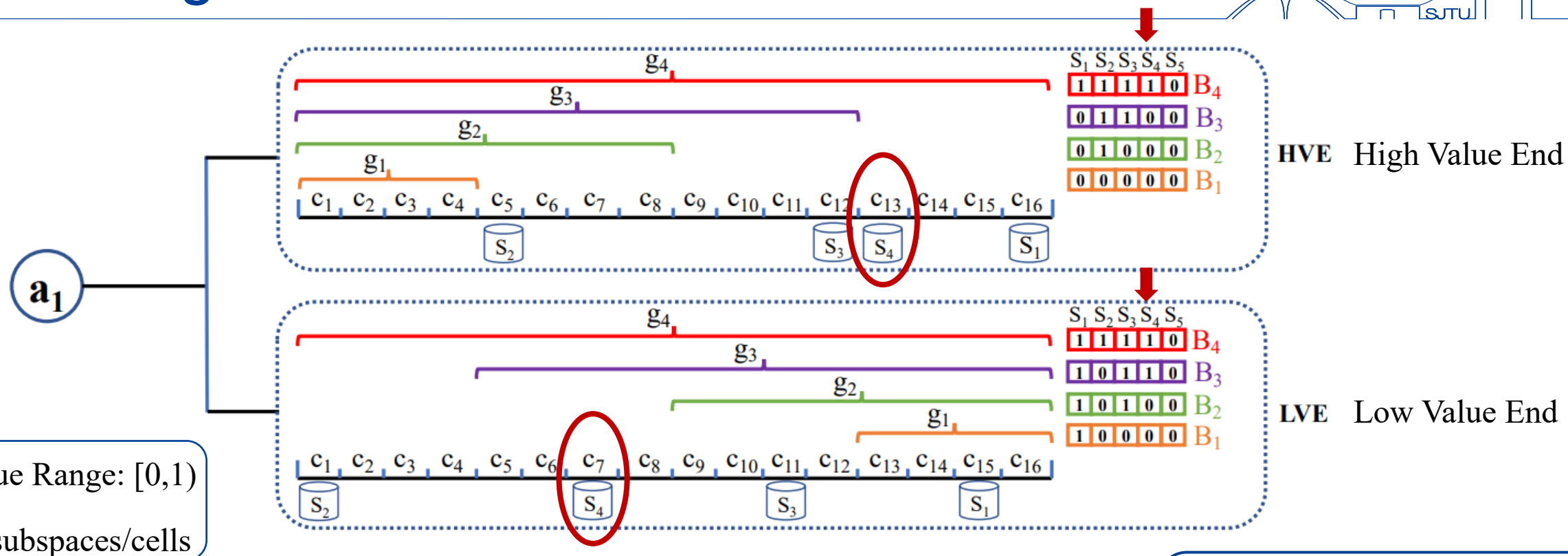


Table 1: Sample subscriptions

| ID | a_1 | a_2 | ID | a_1 | a_2 | ID | a_1 | a_2 |
|-------|----------------|--------------|-------|--------------|---------------|-------|----------------|--------------|
| S_1 | $[0.9, 0.95]$ | $[0.8, 0.9]$ | S_2 | $[0.0, 0.3]$ | $[0.5, 0.7]$ | S_3 | $[0.63, 0.69]$ | $[0.1, 0.2]$ |
| S_4 | $[0.38, 0.76]$ | - | S_5 | - | $[0.4, 0.57]$ | | - | |

$S_4 : (a_1, [0.38, 0.76])$

LVE: $\lfloor 0.38 * 16 + 1 \rfloor = 7$

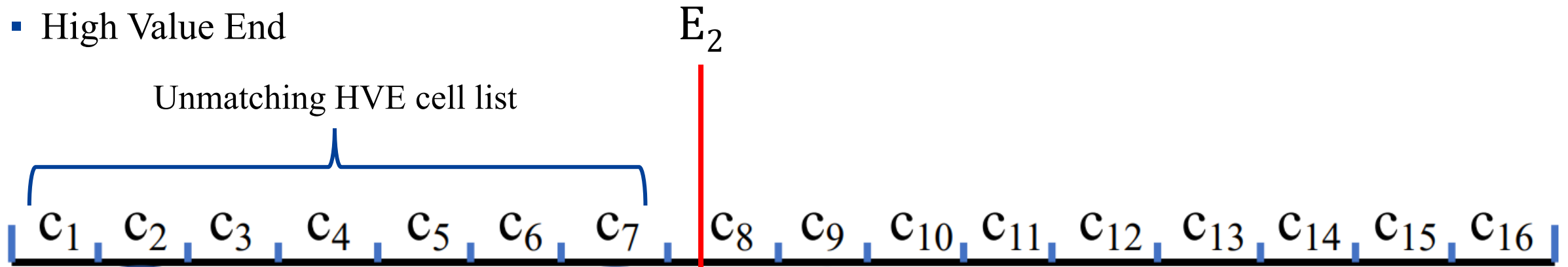
HVE: $\lfloor 0.76 * 16 + 1 \rfloor = 13$



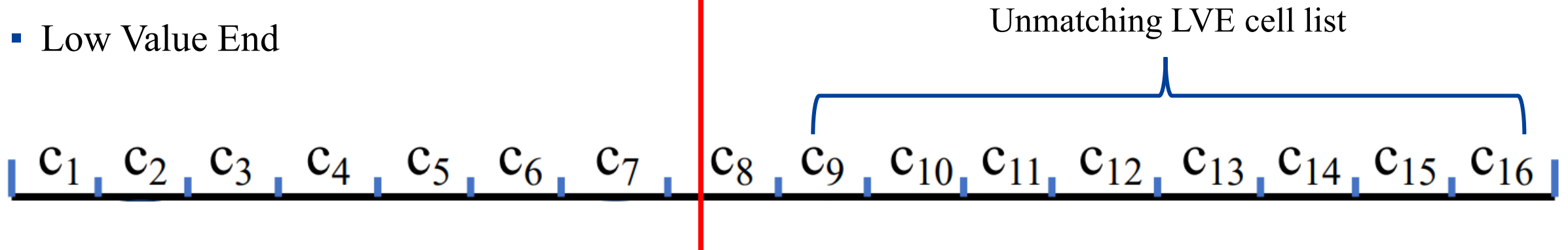
4 Design—Matching without bitsets



- Event $E_2 = \{(a1, 0.48)\} \rightarrow c_8$
- High Value End

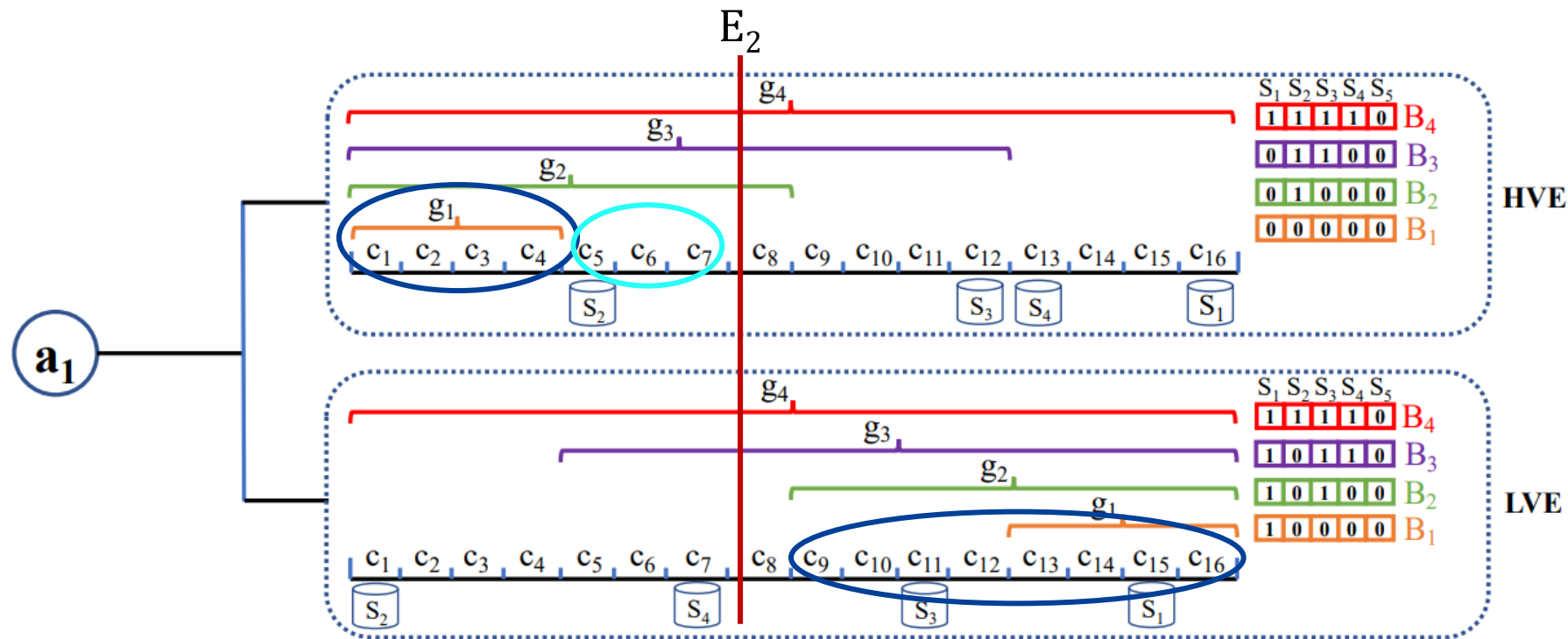


- Low Value End



4 Design—Matching with bitsets

- Event $E_2 = \{(a1, 0.48)\} \rightarrow c_8$



The number of marking cells is 3. (Cell Number/Group Number - 1)

SubID: 1 2 3 4 5
HVE $B_1(00000)$

OR op

LVE $B_2(10100)$



(10100)



Mark op

(11100)

S_4 and S_5 are the matches.

4 Design—Analysis



Theorem 1. *Given n, ψ_E, ψ_S, d and g , when $\psi_E = d, g > 1$ and the predicate values are uniformly distributed in the value domain $[0, 1]$, the improvement ratio of the marking time of HEM is $1 - \frac{1}{g}$ with the SPC method.*

Group Number $g = 32 \rightarrow 96.875\%$

5 Evaluation--Setting



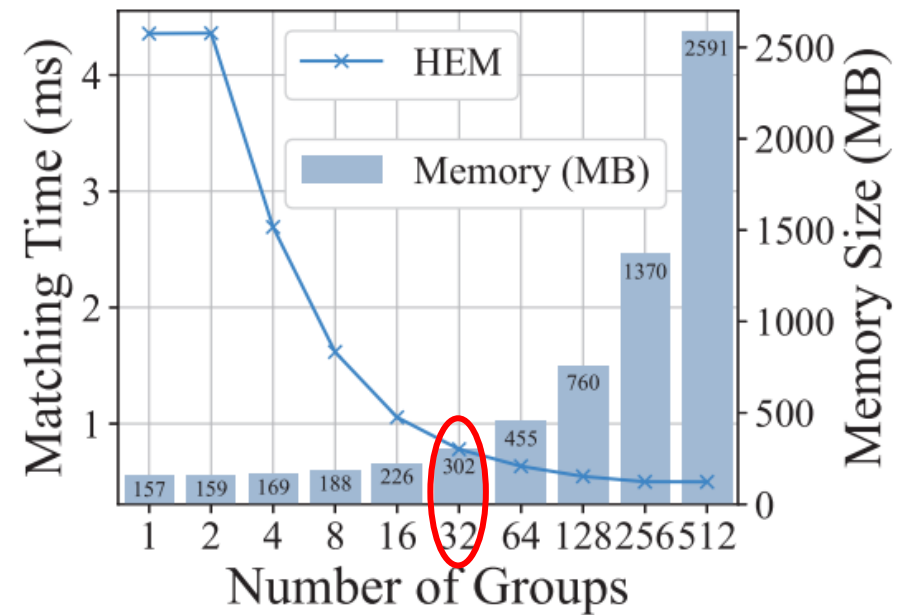
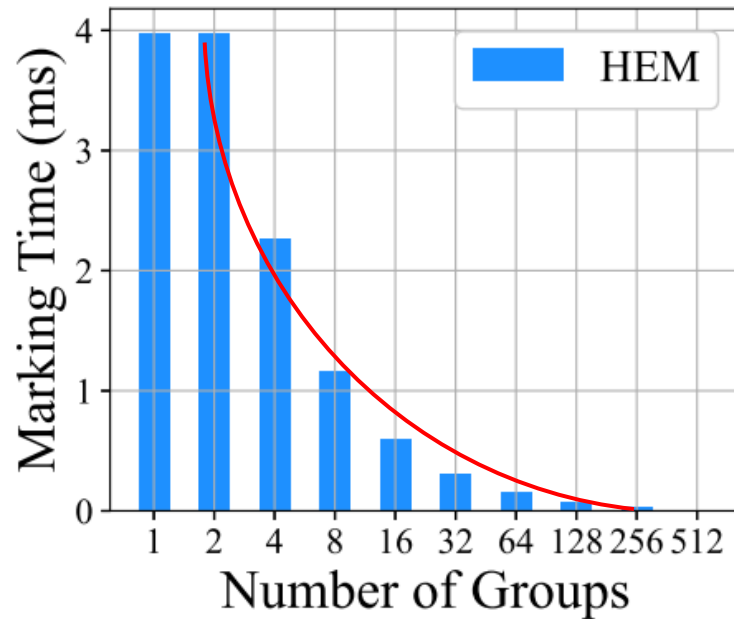
- Language: C++ / g++9.3.0 -O3
- Baseline: REIN, TAMA, Ada-REIN, OpIndex
- Testbed:
 - Ubuntu 20.04
 - AMD 3.7 GHz
 - 64GB RAM
- Codes: <https://github.com/shiwanghua/HEM>

Table 2: Parameter settings in the experiments

| Name | Description | Experimental Values |
|---------------|--------------------------------|-------------------------------------|
| \mathcal{R} | The value domain of attribute. | $[1, 1M]$ |
| α | Parameter of Zipf. | 0 , $1 \sim 5$ |
| d | Number of attributes. | 20 , 30, 100, $300 \sim 900$ |
| n | Number of subscriptions. | 0.3M, 1M , $3M \sim 9M$ |
| ψ_E | Event Size. | 20 , $30 \sim 80$ |
| ψ_S | Subscription size. | 5, 10 , $15 \sim 30$ |
| w | Predicate width. | 0.1, 0.2, 0.3 ~ 0.9 |

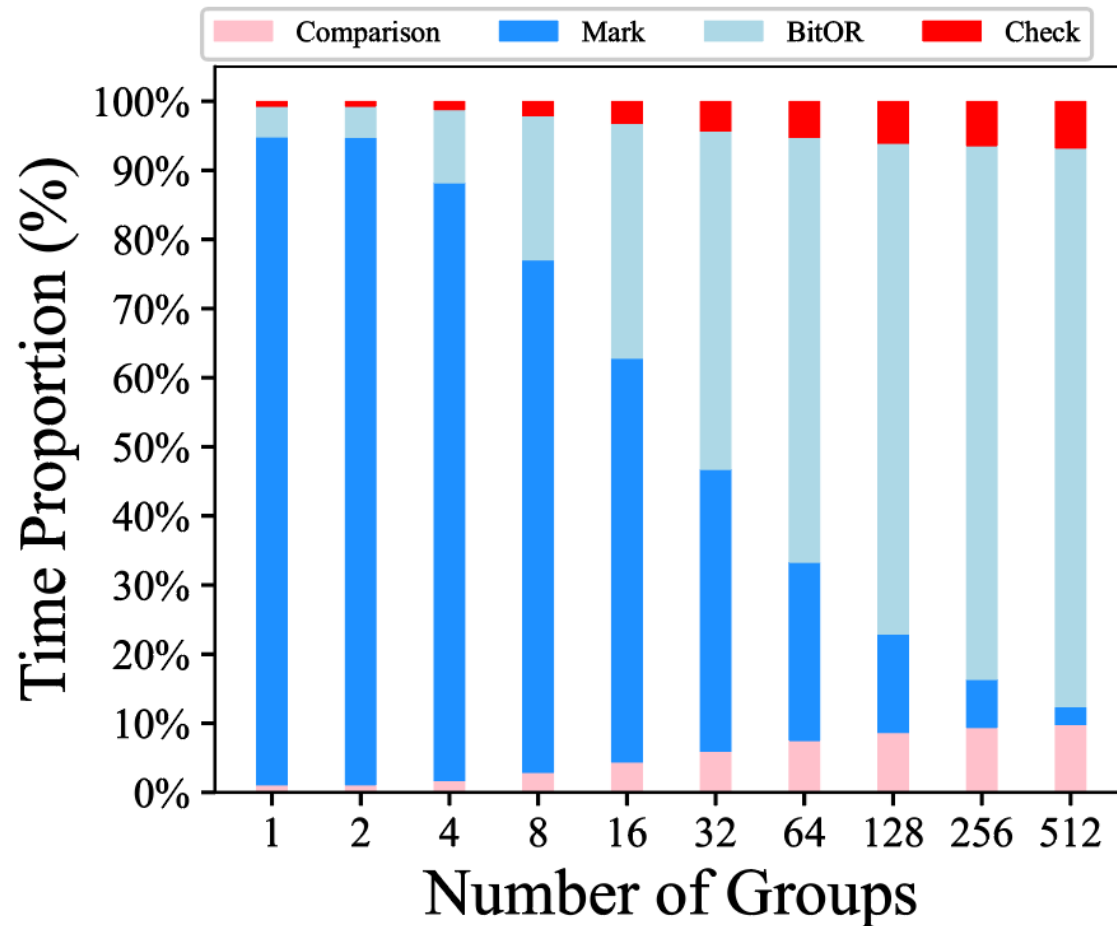


5 Evaluation-- Verification



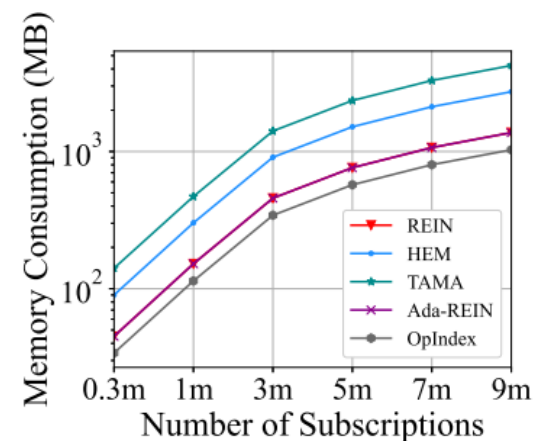
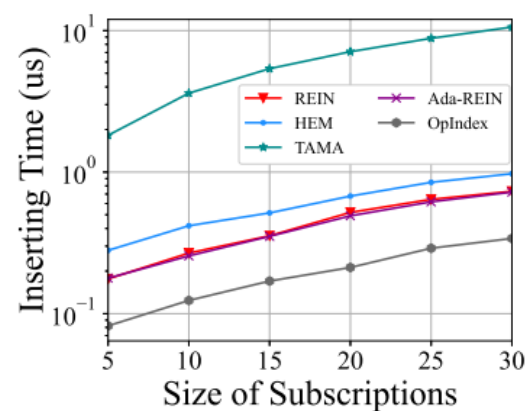
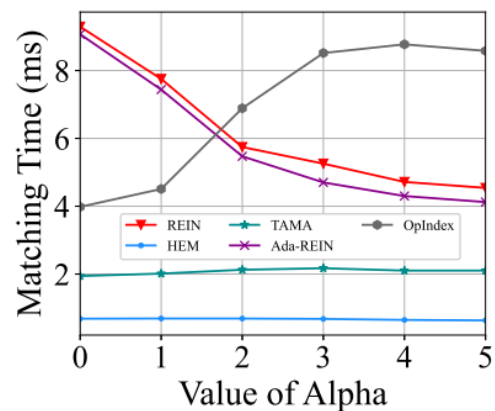
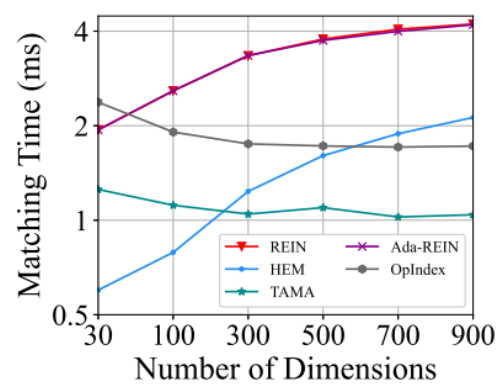
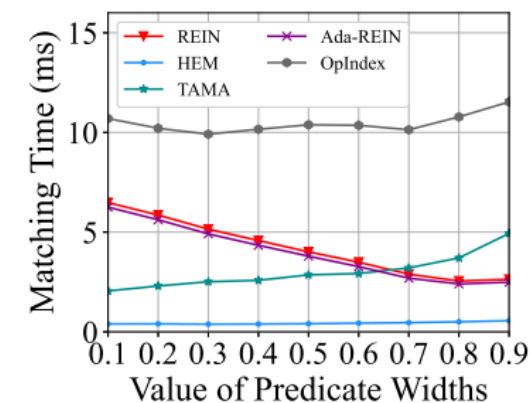
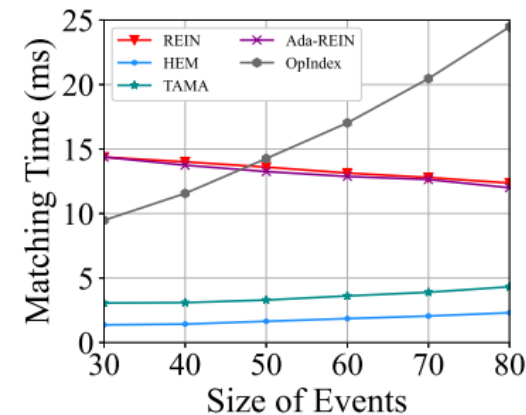
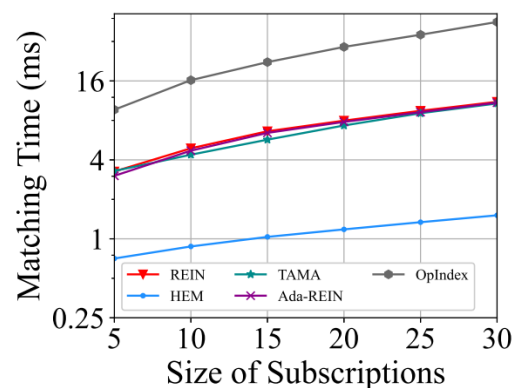
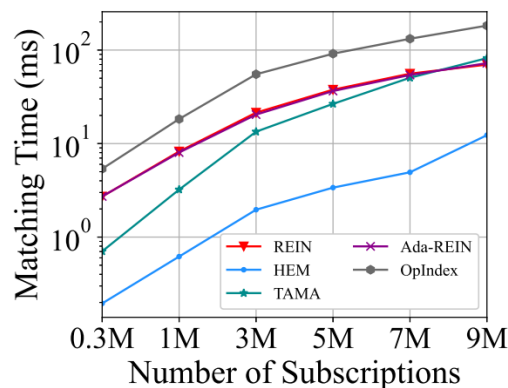
$$\text{Group Number } g = 32 \rightarrow 1 - \frac{0.31ms}{6.79ms} = 95.4\% \approx 96.875\%$$

5 Evaluation-- Verification



- Bottleneck is alleviated.
- Choose $g = 32$.

5 Evaluation--Metric



6 Future Work—series of optimizations



1. Double Reverse Optimization (Virtual Memory): how to select bitsets ?
 2. Dynamic Group Optimization (Load Balancing): how to construct groups ?
 3. Virtual Attribute Group Optimization
 4. Real Attribute Group Optimization
- } high dimensional space
5. OR Operation Optimization (avx2, avx512): 512-bits OR op
- Other directions:
approximate matching / response time / thread pool parallelization

Thanks for listening!

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