Lab1 Report

This report describes our work on Lab 1. Section 1 presents a list of modified files. Section 2 explained our implementation and modification to the source files. Next, section 3 presents how the syscall is processed. Finally, section 4 gives a summary of the contribution of each member. Here is the link to our video https://drive.google.com/file/d/1_8Wb-Zpq0F4DW5TCNLYu4_Z0OlQ0l7Ks/view?usp=sharing. Shixun Wu and Enzuo Zhu contribute to this project equally.

List of modified files:

Implementation Details:

CLI output of lab1_test:

```
hart 1 starting
hart 2 starting
init: starting sh
$ ass1 65536 2
[sysinfo] active proc: 3, syscalls: 44, free pages: 32564
[procinfo 4] ppid: 3, syscalls: 10, page usage: 21
[procinfo 5] ppid: 3, syscalls: 10, page usage: 21
[sysinfo] active proc: 5, syscalls: 236, free pages: 32510
$ ass1 65000 1
exec ass1 failed
$ ass1 65000 1
[sysinfo] active proc: 3, syscalls: 367, free pages: 32564
[procinfo 8] ppid: 7, syscalls: 10, page usage: 20
[sysinfo] active proc: 4, syscalls: 495, free pages: 32538
```

Detailed explanation of our implementations

Add syscall number in kernel/syscall.h

```
#define SYS_sysinfo 22
#define SYS_procinfo 23
```

Add syscall mapping array in kernel/syscall.c

```
static uint64 (*syscalls[])(void) = {
  \\ other syscalls ...
[SYS_sysinfo] sys_sysinfo,
[SYS_procinfo] sys_procinfo
};
```

Declare the following variable defined in other files in kernel/syscall.c

```
extern struct proc proc[NPROC];
char *states[] = {
[UNUSED]
           "unused",
[USED]
           "used",
[SLEEPING] "sleep"
[RUNNABLE] "runble",
[RUNNING] "run ",
[ZOMBIE]
           "zombie"
};
extern struct {
struct spinlock lock;
struct run *freelist;
} kmem;
struct run {
struct run *next;
};
```

- Define sys sysinfo in kernel/syscall.c
 - arg = 0, number of processes: Traverse along the linked list of PCB, proc. When a process is not UNUSED, increase the number of processes by one.
 - arg = 1, number of syscalls: We define an intger variable to record the number of syscalls so far has been made. The variable is increased by one after a system call is finished in the function syscall() in kernel/syscall.c.
 - arg = 2, number of free pages: We go over the linked list of kmem.freelist. kmem.lock is held during the traverse. The number of free pages equal to the length of the kmem.freelist.

o arg = otherwise: return -1.

```
uint64
sys_sysinfo(void)
int arg = 0;
int ret_val = ∅;
argint(∅, &arg);
if(arg == 0){
    struct proc *p;
    int num_proc = 0;
    for(p = proc; p < &proc[NPROC]; p++){</pre>
    if(p->state == UNUSED)
        continue;
    if(p->state >= 0 && p->state < NELEM(states) && states[p->state])
        num_proc += 1;
    }
    ret_val = num_proc;
}
else if(arg == 1){
    ret_val = num_syscall;
}
else if(arg == 2){
    struct run *r = kmem.freelist;
    int num_free_page = 0;
    acquire(&kmem.lock);
    while(r){
    r = r - next;
    num_free_page += 1;
    }
    release(&kmem.lock);
    ret_val = num_free_page;
}
else{
    ret_val = -1;
}
return ret_val;
```

- Define sys_procinfo in kernel/syscall.c
 - ppid: ppid is obatained with myproc()->parent->pid. wait lock is held.
 - syscall_count: We add an int variable num_syscall in PCB. The num_syscall is initiliazed with zero in allocproc() in kernel/proc.c. Whenever a system call is finished, we update the num_syscall of a process with myproc()->num_syscall += 1 in the function syscall().
 - page_usage: We first obtain the memory usage from myproc()->sz, then compute the page usage with (myproc()->sz+4095) / 4096. The page usage is rounded up.

pinfo is copied back to user space with the function copyout. The first parameter, pagetable, is obtained from myproc(). The second parameter dest is obtained from argument. The third parameter is data source. The last parameter is data length.

```
uint64
sys_procinfo(void)
uint64 dest;
argaddr(0, &dest);
struct {
    int ppid;
    int syscall_count;
    int page_usage;
}pinfo;
struct proc *p = myproc();
acquire(&wait_lock);
pinfo.ppid = p->parent->pid;
release(&wait_lock);
pinfo.syscall_count = p->num_syscall;
pinfo.page_usage = (int)((p->sz+4095) / 4096);
if(copyout(p->pagetable, dest, (char*)&pinfo, sizeof(pinfo))<0)</pre>
return -1;
return 0;
}
```

Declare sysinfo and procinfo in user/user.h

```
int sysinfo(int);
int procinfo(void*);
```

Redirect sysinfo and procinfo to syscall sys_sysinfo and sys_procinfo in user/usys.pl

```
entry("sysinfo");
entry("procinfo");
```

The process of syscall

• User space program makes a system call sysinfo. The sysinfo is then redirected to the following assemly in user/usys.S:

```
.global sysinfo
sysinfo:
li a7, SYS_sysinfo
```

```
ecall
ret
```

The pseudoinstruction li load the macro SYS_sysinfo into register a7. The assembly ecall then traps the user mode into the kernel mode.

• The function usertrap() handles the interruption. usertrap() is called from trampoline.S with the following lines:

```
# load the address of usertrap(), from p->trapframe->kernel_trap
ld t0, 16(a0)
```

usertrap() then calls the syscall().

• In syscal1(), the syscall number is fetched from register a7:

```
int num;
struct proc *p = myproc();
num = p->trapframe->a7;
if(num > 0 && num < NELEM(syscalls) && syscalls[num]) {
  // Use num to lookup the system call function for num, call it,
  // and store its return value in p->trapframe->a0
  p->trapframe->a0 = syscalls[num]();
}
```

• After the syscall finished, function usertrapret() is called in usertrap(). It returns to user mode from kernel mode.