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# ACM Template

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## Contents

## 0 Header

```
1 #include <bits/stdc++.h> // #include <algorithm>
2 using namespace std;
3 #define clr(a, x) memset(a, x, sizeof(a))
4 #define pb(x) push_back(x)
5 #define X first
6 #define Y second
7 typedef pair<int, int> PII;
8 typedef vector<int> VI;
9 typedef long long ll;
10 const int INF = 0x3f3f3f3f;
11 const int minINF = 0xc0c0c0c0;
12 const int mod = 1e9 + 7;
13 const double eps = 1e-6; //若题目要求保留到小数点后6位, 则为1e-7
14 const double PI = acos(-1.0);
15 const double E = exp(1);
16
17 //关闭同步/解除绑定
18 ios_base::sync_with_stdio(false);
19 cin.tie(0);
20
21 // 字符串转数字要考虑负号
```

# 1 Math

## 1.1 Prime

### 1.1.1 Euler Sieve

原理：对于任意合数，必定可以有最小质因子乘以最大因子的分解方式  $O(n)$  因此，对于每个合数，只要用最大因子筛一遍，枚举时只要枚举最小质因子即可。因为一般有  $prime[i] * prime[i] \leq n$ ，所以筛选范围为  $\sqrt{n}$  1e7-665000 1e6-80000 1e5-1e4

```

1  int prime[cnt]; //prime[0]储存素数的个数，素数下标从1开始
2  bool heshu[maxn];
3  void getPrime(int n)
4  {
5      for (int i = 2; i <= n; ++i)
6      {
7          if (!heshu[i])
8              prime[++prime[0]] = i;
9          for (int j = 1; j <= prime[0] && i * prime[j] <= n; ++j)
10             {
11                 heshu[i * prime[j]] = true; //找到的素数的倍数不访问
12                 if (i % prime[j] == 0)
13                     break;
14             }
15     }
16 }
```

### 1.1.2 有关素数的基础算法

素数判定

```

1  bool isPrime(int n)
2  {
3      if (n == 1)
4          return false;
5      if (n == 2 || n == 3)
6          return true;
7      if (n % 6 != 1 && n % 6 != 5)
8          return false;
9      for (int i = 5; i * i <= n; i += 6)
10         if (n % i == 0 || n % (i + 2) == 0)
11             return false;
12     return true;
13 }
```

约数枚举

```

1  vector<int> res;
2  void divisor(int n)
3  {
4      for (int i = 1; i * i <= n; ++i)
5         if (n % i == 0)
6             {
7                 res.push_back(i);
8                 if (i != n / i)
9                     res.push_back(n / i);
10             }
11 }
```

## 1.1.3 Miller Rabin

```

1  typedef long long ll;
2  ll Mul(ll a, ll b, ll mod)
3  {
4      if (mod <= 1000000000)
5          return a * b % mod;
6      else if (mod <= 1000000000000LL)
7          return ((a * (b >> 20) % mod) << 20) + (a * (b & ((1 << 20) - 1))) % mod;
8      else
9      {
10         ll d = (ll)floor(a * (long double)b / mod + 0.5);
11         ll res = (a * b - d * mod) % mod;
12         if (res < 0)
13             res += mod;
14         return res;
15     }
16 }
17
18 bool Miller_Rabin(ll n, int s) //s >= 8
19 {
20     if (n == 2)
21         return true;
22     if (n < 2 || !(n & 1))
23         return false;
24     int t = 0;
25     ll u = n - 1;
26     while ((u & 1) == 0)
27         ++t, u >>= 1;
28     for (int i = 0; i < s; ++i)
29     {
30         ll a = rand() % (n - 1) + 1;
31         ll x = mod_pow(a, u, n);
32         for (int j = 0; j < t; ++j)
33         {
34             ll y = Mul(x, x, n);
35             if (y == 1 && x != 1 && x != n - 1)
36                 return false;
37             x = y;
38         }
39         if (x != 1)
40             return false;
41     }
42     return true;
43 }

```

## 1.1.4 区间素数筛

```

1  bool is_prime[1000]; //is_prime[i-a]=true <=> i是素数
2  bool is_prime_small[100];
3  ll prime[maxn]; //只求素数个数时可省略
4  int segment_sieve(ll a, ll b) //对区间[a,b]内的整数执行筛法
5  {
6      for (ll i = 0; i * i < b; ++i) //直接 i*i 会溢出
7          is_prime_small[i] = true;
8      for (ll i = 0; i < b - a; ++i)
9          is_prime[i] = true;
10     for (ll i = 2; i * i < b; ++i)

```

```

11     if (is_prime_small[i])
12     {
13         for (ll j = 2 * i; j * j < b; j += i) //筛[2,sqrt(b))
14             is_prime_small[i] = false;
15         for (ll j = max(2LL, (a + i - 1) / i) * i; j < b; j += i) //筛[a,b)
16             is_prime[j - a] = false;
17     }
18     int cnt = 0;
19     for (ll i = 0; i < b - a; ++i)
20         if (is_prime[i])
21             prime[cnt++] = i + a; // ++cnt;
22     return cnt;
23 }

1 ans = segment_sieve(a, b + 1) - (a == 1)); //[a,b]

```

## 1.2 Euler

### 1.2.1 欧拉函数

欧拉函数的值等于不超过  $m$  并且和  $m$  互素的数的个数

```

1 ll euler(ll n) //求欧拉函数值
2 {
3     ll res = n;
4     for (ll i = 2; i * i <= n; ++i)
5         if (n % i == 0)
6         {
7             res = res / i * (i - 1);
8             while (n % i == 0)
9                 n /= i;
10        }
11    if (n != 1)
12        res = res / n * (n - 1);
13    return res;
14 }
15 int phi[maxn];
16 void euler_phi(int n) //筛出欧拉函数值的表
17 {
18     for (int i = 0; i <= n; ++i)
19         phi[i] = i;
20     for (int i = 2; i <= n; ++i)
21         if (phi[i] == i)
22             for (int j = i; j <= n; j += i)
23                 phi[j] = phi[j] / i * (i - 1);
24 }

```

### 1.2.2 欧拉降幂

Euler Theorem 费马小定理 ( $p$  为素数)  $a^b = a^{b \bmod (p-1)} \bmod p$  扩展欧拉定理 (欧拉降幂) ( $a$  和  $p$  不互质)  $a^b \bmod p = a^{b \bmod P (b < \phi(p))} a^{b \bmod p} = a^{b \bmod \phi(p) + \phi(p)} \bmod p (b \geq \phi(p))$  指数循环节: 从  $a^0$  到  $a^{\phi(p)-1}$  不是重复的, 从  $a^{\phi(p)}$  开始出现循环节, 长度为  $\phi(p)$   $a^b \bmod p (1 \leq a \leq 1e9, 1 \leq b \leq 10^2e7, 1 \leq p \leq 1e6)$

```

1 ll solve(ll a, string &str, ll p)
2 {
3     int phi = euler_phi(p);
4     bool flag = false;
5     ll b = 0;
6     for (int i = 0; i < str.length(); ++i)
7     {

```