

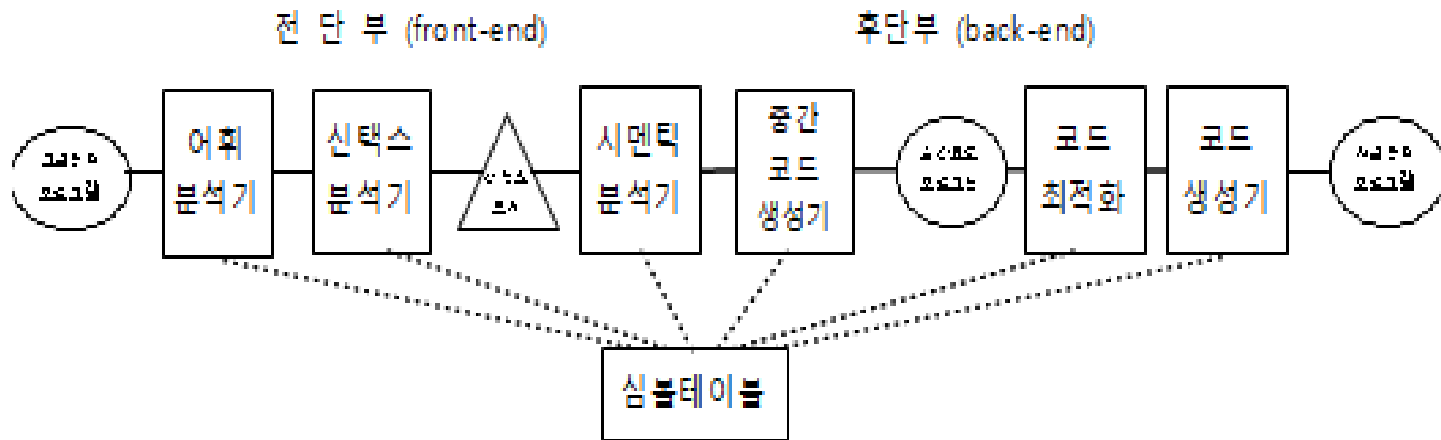
C 언어와 컴파일러

소프트웨어 학부

2023 봄학기

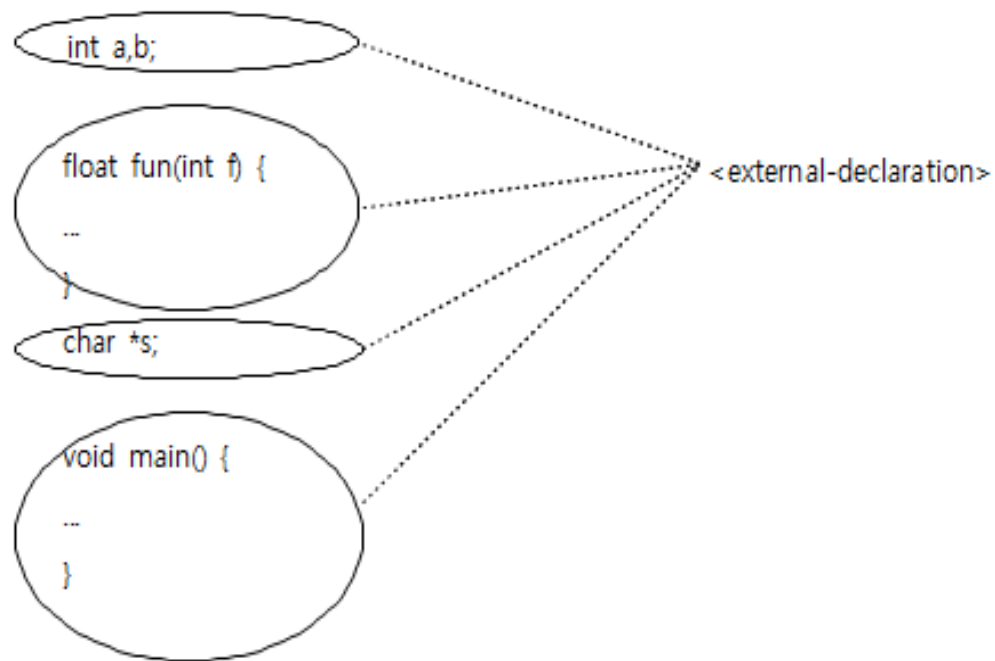
유재우 교수

컴파일러의 구조



문법과 언어

- Grammar
- Syntax, Semantics
- Language

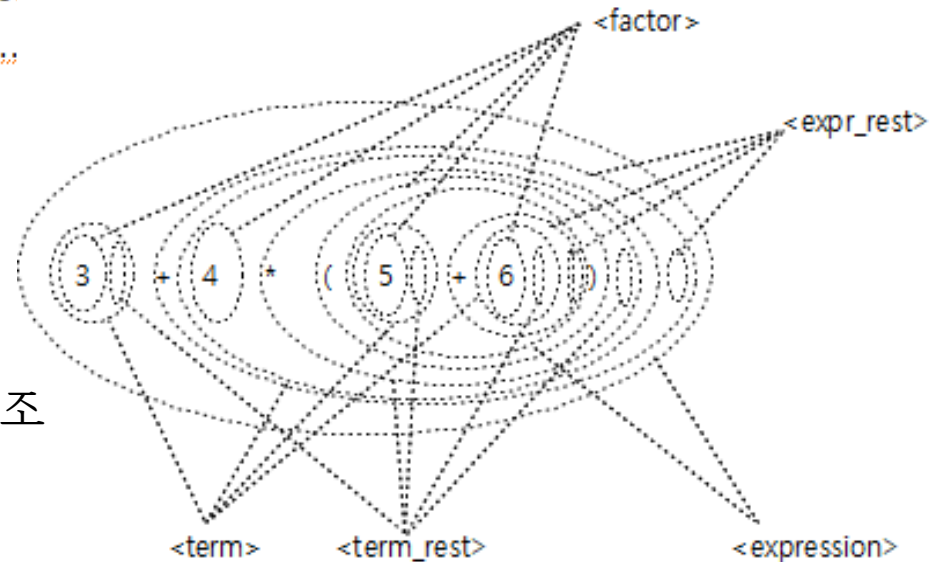


간단한 수식을 위한 문법

G_0 :

- $\langle \text{expression} \rangle ::= \langle \text{term} \rangle \langle \text{expr_rest} \rangle$
- $\langle \text{expr_rest} \rangle ::= + \langle \text{term} \rangle \langle \text{expr_rest} \rangle$
 $\quad \quad \quad | \quad \lambda$
- $\langle \text{term} \rangle ::= \langle \text{factor} \rangle \langle \text{term_rest} \rangle$
- $\langle \text{term_rest} \rangle ::= * \langle \text{factor} \rangle \langle \text{term_rest} \rangle$
 $\quad \quad \quad | \quad \lambda$
- $\langle \text{factor} \rangle ::= \langle \text{number} \rangle$
 $\quad \quad \quad | \quad (\langle \text{expression} \rangle)$
- $\langle \text{number} \rangle ::= 0 \mid 1 \mid \dots$

수식 “3+4*(5+6)”의 문법적구조



Context-Free Grammar

- $G = (T, N, P, S)$

- T: terminal symbols
- N: nonterminal symbols
- P: production rules ' $A ::= \alpha$ ', $A \in N$, $\alpha \in (N \cup T)^*$
- S: start symbol, $S \in N$

$G_1 = (\{+, *, (,), 0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}, \{E, R, T, Q, F, N\}, P_1, E)$

P_1 :

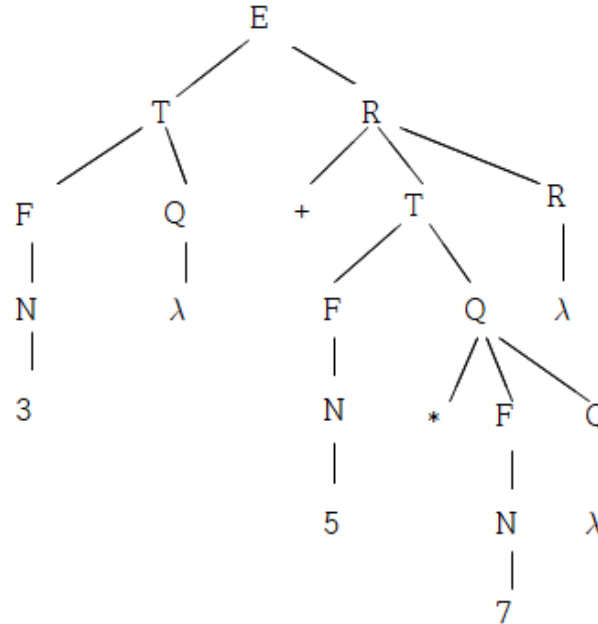
E	::=	T	R
R	::=	+	T R
			λ
T	::=	F	Q
Q	::=	*	F Q
			λ
F	::=	N	
			(E)
N	::=	0	1 ... 9

Sentence & Language of a Grammar

derivations

$E \Rightarrow T R \Rightarrow F Q R \Rightarrow N Q R \Rightarrow 3 Q R \Rightarrow 3 R \Rightarrow 3 + T R$
 $\Rightarrow 3 + F Q R \Rightarrow 3 + N Q R \Rightarrow 3 + 5 Q R \Rightarrow 3 + 5 * F Q R$
 $\Rightarrow 3 + 5 * N Q R \Rightarrow 3 + 5 * 7 Q R \Rightarrow 3 + 5 * 7 R \Rightarrow 3 + 5 * 7$

syntax tree

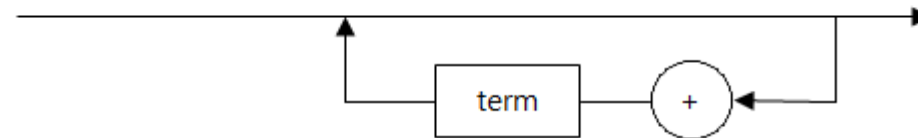


Syntax Graph

expression



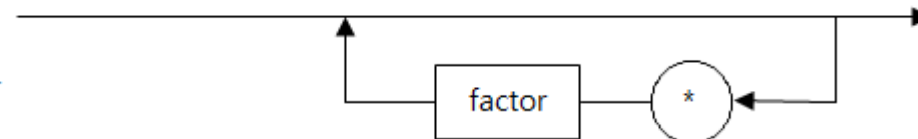
expr_rest



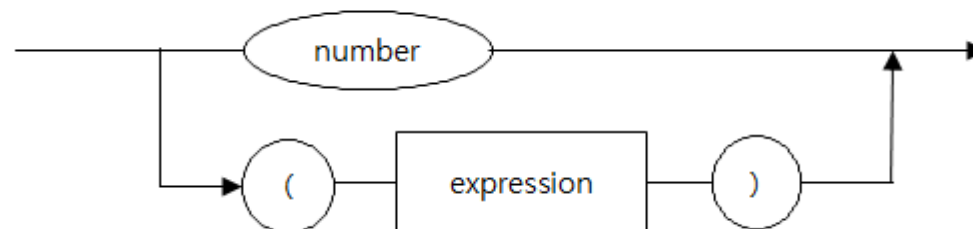
term



term_rest

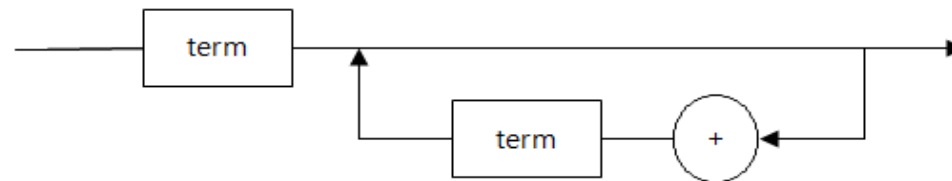


factor

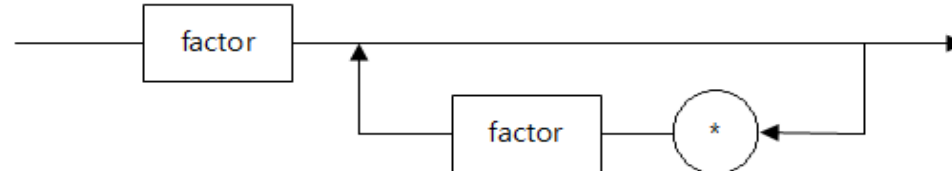


Syntax Graph for Simple Expressions

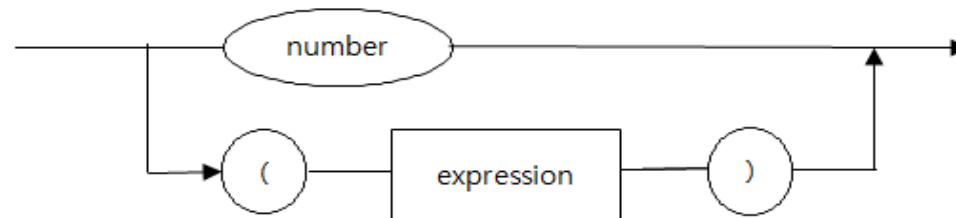
expression



term



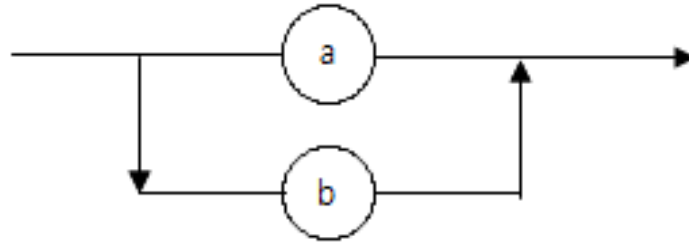
factor



Syntax Graph and Recursive-Descent Parser

- { “a”, “b” }

expression_1

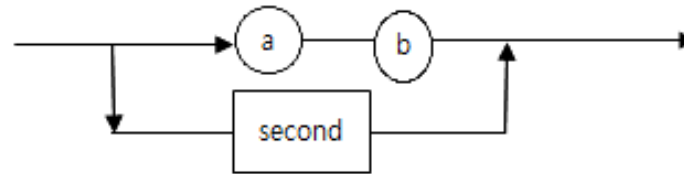


```
char ch;
void expression_1( ) {
    if (ch=='a')
        ch=getchar( );
    else if (ch=='b')
        ch=getchar( );
    else
        error();
}
```

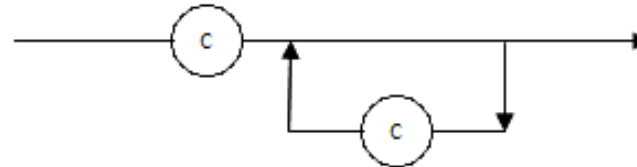
```
void main() {
    ch=getchar( );
    expression_1( );
    if (ch!=EOF)
        error();
}
void error() {
    .....
}
```

- “ab” 나 “c...c” 인식

expression_2



second



```

char ch=' ';
char get_token() {
    while (ch==' ')
        ch=getchar();
    return (ch); }
void expression_2() {
    if (ch=='a')
        ch=get_token();
    if (ch=='b')
        ch=get_token();
    else
        error();
    else
        second();}
  
```

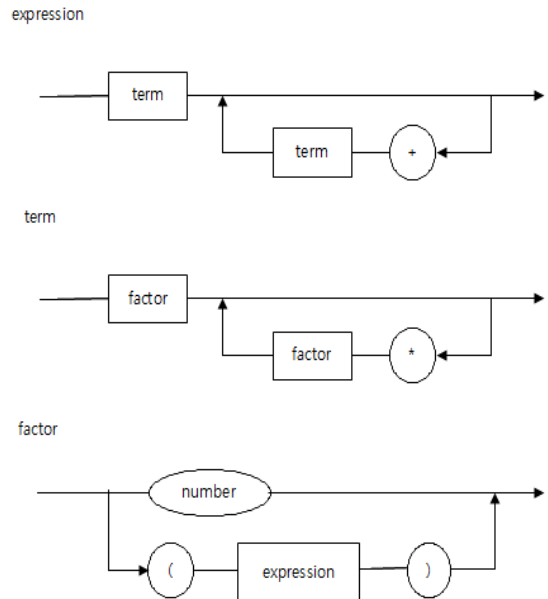
```

void second () {
    if (ch=='c') {
        ch=get_token();
        while (ch=='c')
            ch=get_token(); }
    else
        error();
}
void main () {
    ch=get_token();
    expression_2();
    if (ch!=EOF)
        error();
}
  
```

간단한 수식을 위한 파싱 프로그램

```
enum {NULL, PLUS, STAR, NUMBER,
      LPAREN, RPAREN, END} token;
void get_token ( ) {
    // token = next token of input
}
void expression ( ) {
    term();
    while (token==PLUS) {
        get_token();
        term( ); }
}
void term ( ) {
    factor();
    while (token==STAR) {
        get_token();
        factor( ); }
}
void factor ( ) {
    if (token==NUMBER)
```

```
        get_token();
    else if (token==LPAREN) {
        get_token();
        expression();
        if (token==RPAREN)
            get_token();
        else
            error( ); }
    else error();
}
main ( ) {
    get_token();
    expression();
    if (token!=END)
        error( );
}
error ( ) {
    // error handling
}
```



get_token() 함수 프로그램

입력문자	토큰종류 (token)
'0'~'9'	NUMBER
'+'	PLUS
'*'	STAR
'('	LPAREN
)'	RPAREL
EOF	END
기타	NULL

integer number : [0-9] [0-9]*

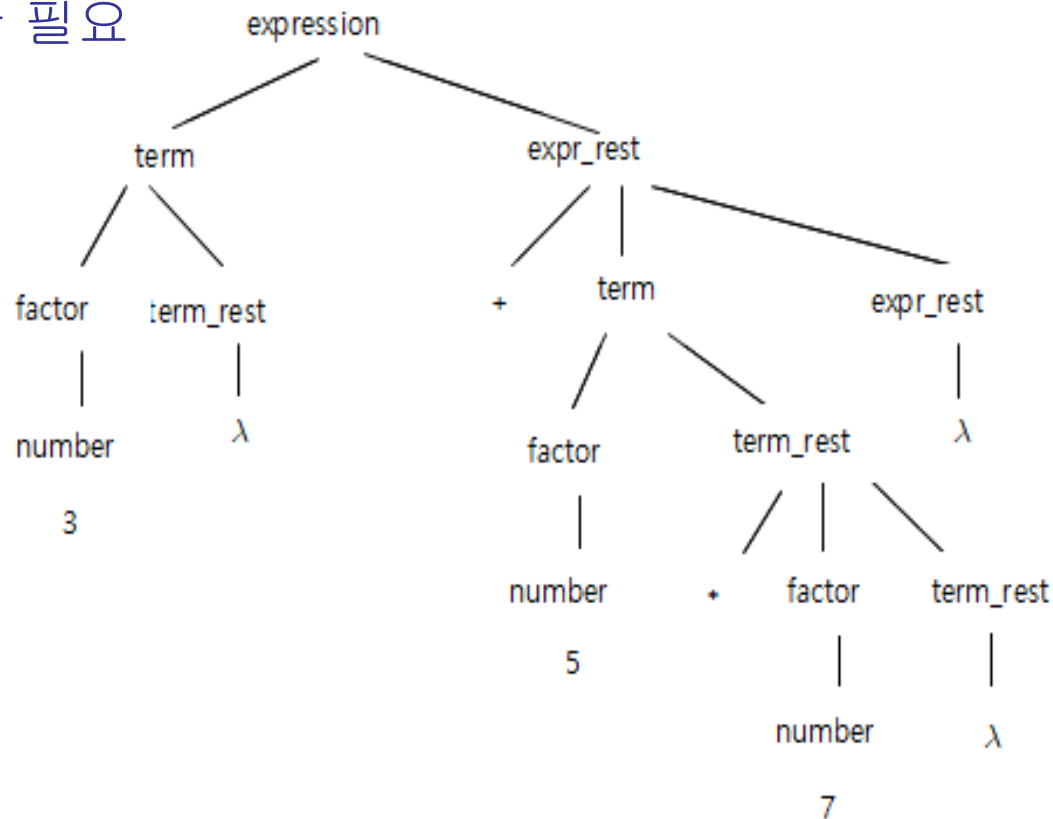
identifier: letter (letter | digit)*

Recursive-descent Parser 의 작동

- 시작기호로 부터 하향식으로 호출

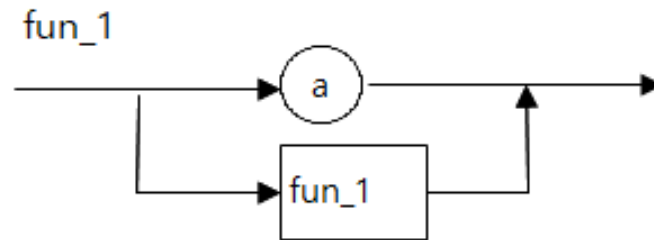
main() --> expression() --> term() --> factor() --> expression() ...

- 좌측 먼저 처리
- 한개의 토큰씩만 필요



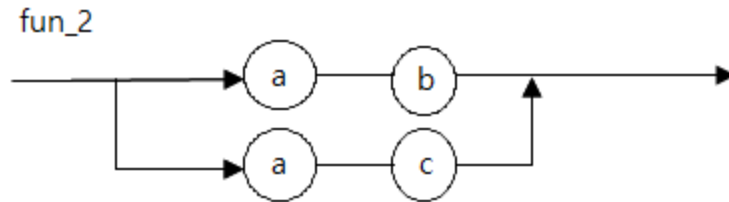
신택스 그래프의 제한사항

<경우-1>

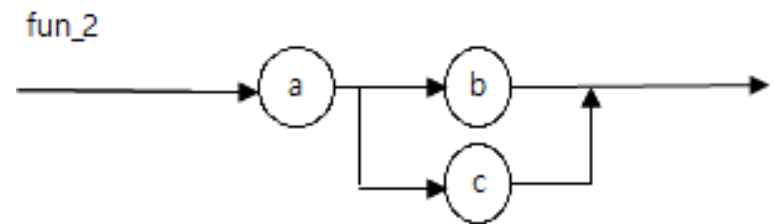


```
void fun_1 ( ) {  
    if (token=='a')  
        get_token( );  
    else  
        fun_1( );  
}
```

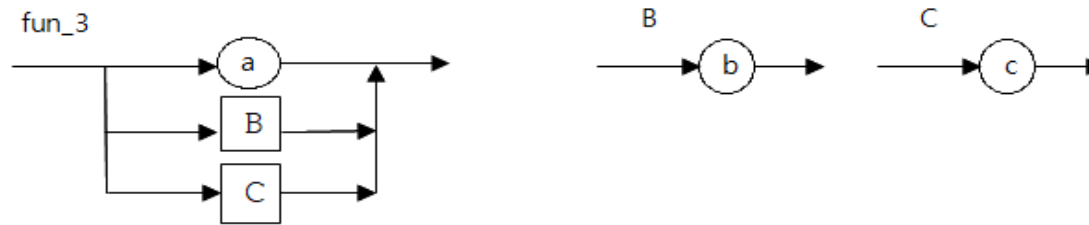
<경우-2> “ab” 나 “ac” 인식



```
void fun_2 ( ) {  
    if (token=='a') {  
        get_token( );  
        if (token=='b')  
            get_token( );  
        else    error( ); }  
    else if (token=='a') {  
        get_token( );  
        if (token=='c')  
            get_token( );  
        else    error( ); }  
    else    error( );  
}
```

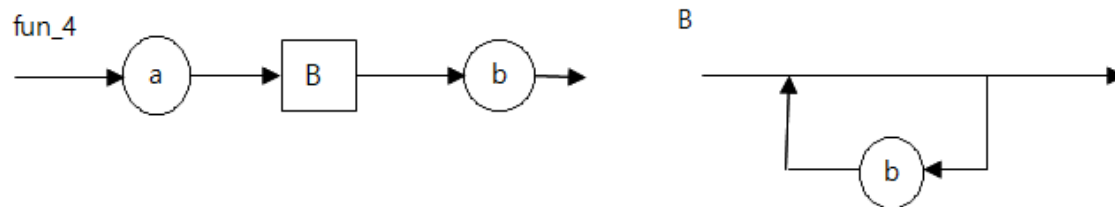


<경우-3> “a” 나 “b” 나 “c” 인식



```
void fun_3 ( ) {  
    if (token=='a')  
        get_token( );  
    else  
        B( );  
    else                // 작성 오류 ?  
        C( );  
    .....  
}  
void B( ) {  
    if (token=='b')  
        get_token( );  
    else  
        error( );  
}  
void C( ) {  
    // 생략  
}
```


<경우-4> “ab....b”



```
void fun_4( ) {  
    if (token=='a')  
        get_token( );  
        B( );  
        if (token=='b')  
            get_token( );  
        else  
            error( );  
    else  
        error( );  
}  
void B( ) {  
    while (token=='b') {  
        get_token( );  
    }  
}
```

수식의 값 계산

```
int num;
enum NULL,NUMBER,PLUS,
      STAR,LP,RP,END} token;
```

```
void main () {
    int result;
    get_token();
    result=expression();
    if (token!=END)
        error(3);
    else
        printf("%d \n",result);
}
```

```
int expression () {
    int result;
    result=term();
    while (token==PLUS) {
        get_token();
        result=result+term( ); }
    return (result);
}
```

```
int term () {
    int result;
    result=factor();
    while (token==STAR) {
        get_token();
        result=result*factor( ); }
    return (result);
}
```

```
int factor () {
    int result;
    if (token==NUMBER) {
        result=num
        get_token(); }
    else if (token==LP) {
        get_token();
        result=expression();
        if (token==RP)
            get_token();
        else
            error(2); }
    else
        error(1);
    return (result);
}
```

```
void get_token () {
    // next token --> token
    // number value --> num
}
```

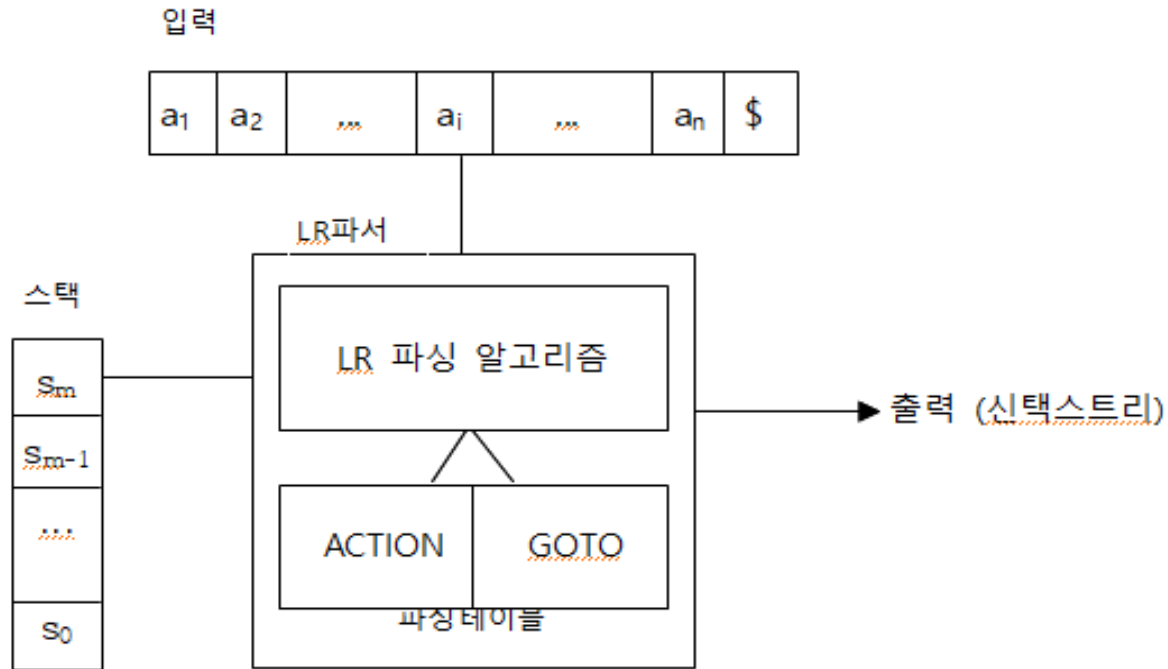
```
void error (int i) {
    switch (i) {
        case 1: ... break;
        case 2: ... break;
        case 3: ... break;
    }
    exit(1);
}
```

수식의 확장

- 연산자: 곱하기(*), 나누기(/)
- 피연산자: 정수 및 실수
- 결과값 인쇄:
 - 정수형 수식은 정수값, 실수형 수식은 실수값
- Warning messages
 - 혼합연산
- Error messages
- 함수나 token의 타입

```
typedef enum {INT,FLT} kind;  
struct { kind t;  
        union { int i; float f;} val; }
```

LR Parser



상태번호

스택 : $s_0 \dots s_m$

end-marker '\$'

파싱테이블: ACTION[], GOTO[]

LR Parsing

파싱 동작:

- “shift i”, “reduce j”, “accept”, “error”

파서상황 (configuration) $Q : (S_0 X_1 S_1 \dots X_m S_m, a_i \dots a_n \$)$

초기 파서상황 $Q_0 : (0, a_0 \dots a_n \$)$

스택 꼭대기: S_m , 입력 토큰이 a_i 인 경우, $ACTION[S_m][a_i]$ 을 참조

- “shift j”:

$(S_0 X_1 S_1 \dots X_m S_m, a_i a_{i+1} \dots a_n \$) \rightarrow (S_0 X_1 S_1 \dots X_m S_m a_i j, a_{i+1} \dots a_n \$)$

- “reduce j”:

문법의 j 번째 규칙: $A ::= \alpha$, $r=|\alpha|$,

$(S_0 X_1 S_1 \dots X_m S_m, a_i \dots a_n \$) \rightarrow$

$(S_0 X_1 S_1 \dots X_{m-r} S_{m-r}, a_i \dots a_n \$) \rightarrow$

$(S_0 X_1 S_1 \dots X_{m-r} S_{m-r} A_k, a_i \dots a_n \$)$, 여기서 $GOTO[s_{m-r}][A]=k$

- “accept” : 성공적 종료
- “error” : 에러 처리

LR Parsing Algorithm

```
push(0)
token=get_token()
while (1) { s =stack top의 상태번호
    switch (ACTION[s][token]) {
        case "shift i" :
            push(token)
            push(i)
            token=get_token()
            break
        case "reduce j" :
            j번째 생성규칙을 "A ::=  $\alpha$ " 라고 가정
            스택에서  $2|\alpha|$  만큼의 기호를 삭제
            t=삭제후 stack top 의 상태번호
            push(A)
            push(GOTO[t][A])
            break
        case "accept" :
            파서 종료처리
        default :
            error() 함수호출
    }
}
```

LR문법과 LR Parsing Tables

- ① $E \rightarrow E + T$
 ② $\quad \quad | T$
 ③ $T \rightarrow T * F$
 ④ $\quad \quad | F$
 ⑤ $F \rightarrow (E)$
 ⑥ $\quad \quad | n$

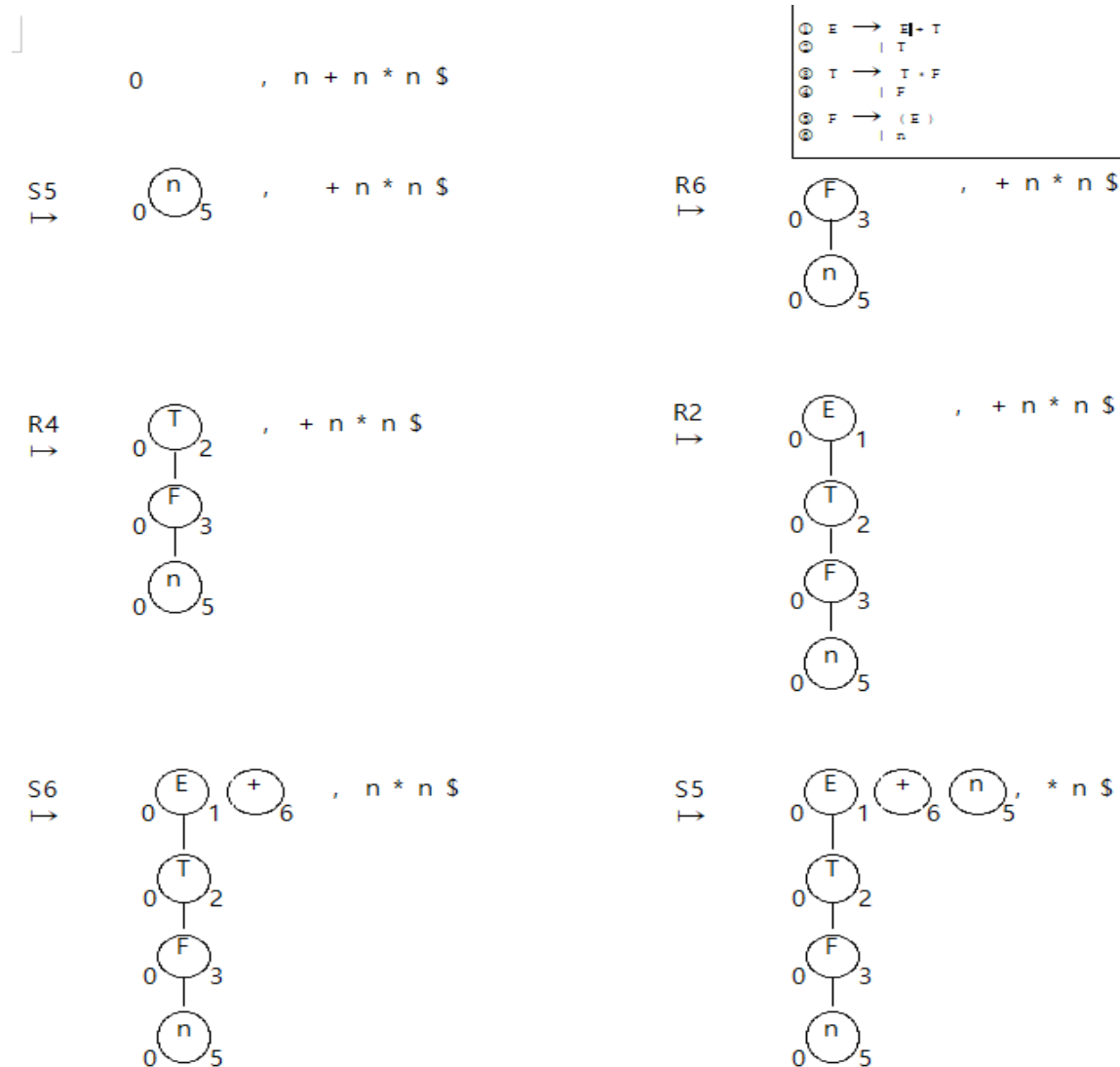
테이블 상태번호	ACTION 테이블						GOTO 테이블		
	n	+	*	()	\$	E	T	F
0	S5			S4			1	2	3
1		S6				acc			
2		R2	S7		R2	R2			
3		R4	R4		R4	R4			
4	S5			S4			8	2	3
5		R6	R6		R6	R6			
6	S5			S4				9	3
7	S5			S4					10
8		S6		s11					
9		R1	S7		R1	R1			
10		R3	R3		R3	R3			
11		R5	R5		R5	R5			

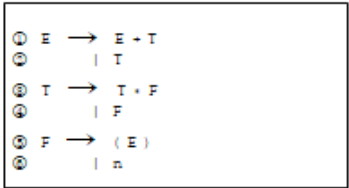
$n + n * n$ 의 파싱과정

	<u>스택</u>	<u>입력</u>
	(0,	$n + n * n \$$)
S5 \mapsto	(0 n 5,	$+ n * n \$$)
R6 \mapsto	(0 F 3,	$+ n * n \$$)
R4 \mapsto	(0 T 2,	$+ n * n \$$)
R2 \mapsto	(0 E 1,	$+ n * n \$$)
S6 \mapsto	(0 E 1 + 6,	$n * n \$$)
S5 \mapsto	(0 E 1 + 6 n 5,	$* n \$$)
R6 \mapsto	(0 E 1 + 6 F 3,	$* n \$$)
R4 \mapsto	(0 E 1 + 6 T 9,	$* n \$$)
S7 \mapsto	(0 E 1 + 6 T 9 * 7,	$n \$$)
S5 \mapsto	(0 E 1 + 6 T 9 * 7 n 5,	$\$$)
R6 \mapsto	(0 E 1 + 6 T 9 * 7 F 10,	$\$$)
R3 \mapsto	(0 E 1 + 6 T 9,	$\$$)
R1 \mapsto	(0 E 1,	$\$$)

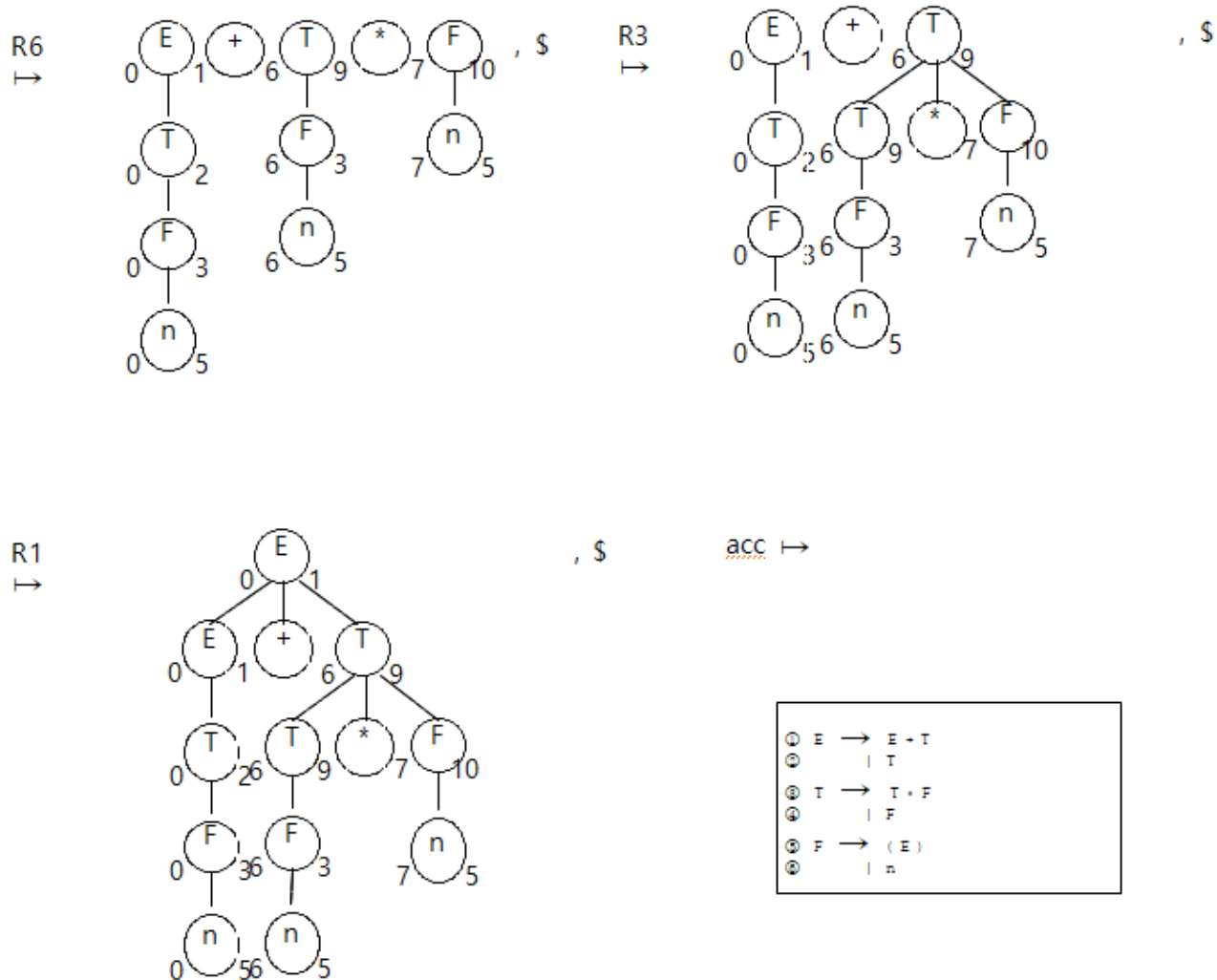
acc

Parse Tree 생성 개념(1)





Parse Tree 생성 개념 (3)



LR Parser 특징

- SLR, LALR, 및 LR 파서 : 파싱 알고리즘은 동일
- LR 파싱 테이블
문법으로부터 제작
파싱 테이블 크기: ‘12줄*심볼갯수’
C 언어: 500*500 정도 이상

LR Parser Program (1)

```
#define NUMBER 256      #define PLUS 257
#define STAR 258        #define LPAREN 259
#define RPAREN 260     #define END 261
#define EXPRESSION 0   #define TERM 1
#define FACTOR 2       #define ACC 999

int action[12][6]={
    {5, 0, 0, 4, 0, 0}, {0, 6, 0, 0, 0, ACC}, {0,-2, 7, 0,-2,-2},
    {0,-4,-4, 0,-4,-4}, {5, 0, 0, 4, 0, 0},    {0,-6,-6, 0,-6,-6},
    {5, 0, 0, 4, 0, 0}, {5, 0, 0, 4, 0, 0},    {0, 6, 0, 0,11, 0},
    {0,-1, 7, 0,-1,-1}, {0,-3, -3, 0,-3,-3},    {0,-5,-5, 0,-5,-5} };

int go_to[12][3]={
    {1,2,3},{0,0,0}, {0,0,0},{0,0,0},{8,2,3},{0,0,0},
    {0,9,3},{0,0,10},{0,0,0},{0,0,0},{0,0,0},{0,0,0} };

int prod_left[7]={0,EXPRESSION,EXPRESSION,TERM,TERM,FACTOR,FACTOR};
int prod_length[7]={0,3,1,3,1,3,1};

int stack[1000]; int top=-1; int sym;

void main () {
    yyparse();
}
```

LR Parser Program (2)

```
int yyparse() {  
    int i;  
    stack[++top]=0;                // initial state  
    sym=yylex();  
    do { i=action[stack[top]][sym-256];    // get relation  
        if (i==ACC)  
            printf("success !\n");  
        else if (i>0) shift(i);          // shift  
        else if (i<0) reduce(-i);        // reduce  
        else yyerror(); }  
    while (i!=ACC);  
}  
void push(int l) {  
    stack[++top]=i; }  
void shift(int l) {  
    push(i); sym=yylex(); }  
void reduce(int l) {  
    int old_top;  
    top-=prod_length[i];  
    old_top=top;  
    push(go_to[stack[old_top]][prod_left[i]]);}  
void yyerror() {  
    printf("syntax error\n");  
    exit(1); }
```

LR Parser Program (3)

```
int yylex() {  
    static char ch=' '  
    int i=0;  
    while (ch==' '||ch=='\t'||ch=='\n') ch=getchar();  
    if (isdigit(ch)) {  
        do  
            ch=getchar();  
        while (isdigit(ch));  
        return(NUMBER); }  
    else if (ch=='+'){ ch=getchar(); return(PLUS);}  
    else if (ch=='*'){ ch=getchar(); return(STAR);}  
    else if (ch=='('){ ch=getchar(); return(LPAREN);}  
    else if (ch==')'){ ch=getchar(); return(RPAREN);}  
    else if (ch==EOF) return(END);  
    else lex_error();  
}
```

LR Parser 에서 수식의 값계산 방법(1)

	<u>stack</u>	<u>input</u>								
	<div>트랙-1</div> <table><tr><td>0</td></tr><tr><td></td></tr></table> <div>트랙-2</div>	0		, (1 + 2) * 3 \$						
0										
S4 \mapsto	<table><tr><td>0</td><td>(4</td></tr><tr><td></td><td></td></tr></table>	0	(4			, 1 + 2) * 3 \$				
0	(4									
S5 \mapsto	<table><tr><td>0</td><td>(4</td><td>n 5</td></tr><tr><td></td><td></td><td>1</td></tr></table>	0	(4	n 5			1	, + 2) * 3 \$		
0	(4	n 5								
		1								
R6 \mapsto	<table><tr><td>0</td><td>(4</td><td>F 3</td></tr><tr><td></td><td></td><td>1</td></tr></table>	0	(4	F 3			1	, + 2) * 3 \$		
0	(4	F 3								
		1								
R4 \mapsto	<table><tr><td>0</td><td>(4</td><td>T 2</td></tr><tr><td></td><td></td><td>1</td></tr></table>	0	(4	T 2			1	, + 2) * 3 \$		
0	(4	T 2								
		1								
R2 \mapsto	<table><tr><td>0</td><td>(4</td><td>E 8</td></tr><tr><td></td><td></td><td>1</td></tr></table>	0	(4	E 8			1	, + 2) * 3 \$		
0	(4	E 8								
		1								
S6 \mapsto	<table><tr><td>0</td><td>(4</td><td>E 8</td><td>+ 6</td></tr><tr><td></td><td></td><td>1</td><td></td></tr></table>	0	(4	E 8	+ 6			1		, 2) * 3 \$
0	(4	E 8	+ 6							
		1								

LR Parser 에서 수식의 값계산 방법(2)

└

S5 \mapsto

0	(4	E 8	+ 6	n 5
		1		2

,) * 3 \$

R6 \mapsto

0	(4	E 8	+ 6	F 3
		1		2

,) * 3 \$

R4 \mapsto

0	(4	E 8	+ 6	T 9
		1		2

,) * 3 \$

R1 \mapsto

0	(4	E 8
		3

,) * 3 \$

S11 \mapsto

0	(4	E 8) 11
		3	

, * 3 \$

R5 \mapsto

0	F 3
	3

, * 3 \$

R4 \mapsto

0	T 2
	3

, * 3 \$

LR Parser 에서 수식의 값계산 방법(3)

」

S7 \mapsto

0	T 2	. 7
	3	

, 3 \$

S5 \mapsto

0	T 2	. 7	n 5
	3		3

, \$

R6 \mapsto

0	T 2	. 7	F10
	3		3

, \$

R3 \mapsto

0	T 2
	9

, \$

R2 \mapsto

0	E 1
	9

, \$

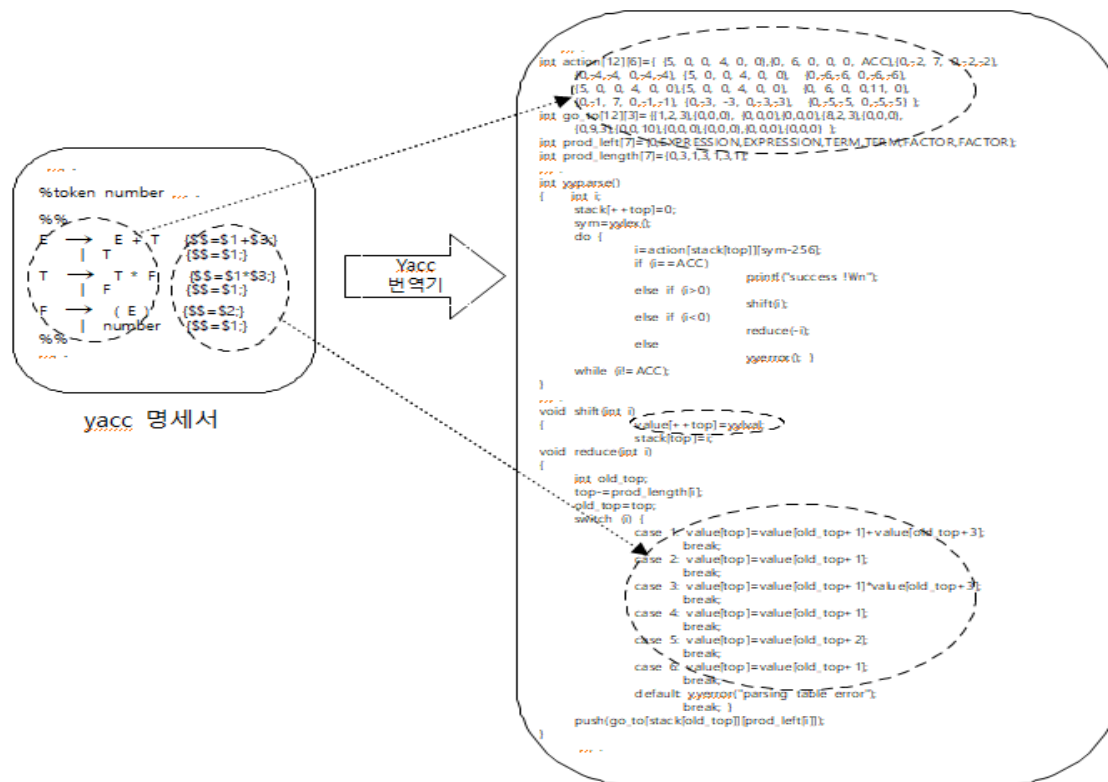
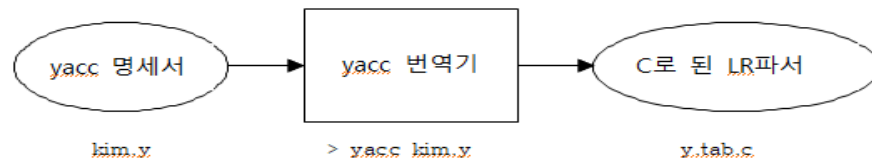
acc

C Program(1)

```
char yytext[32];
int yylval;
// 생략 ...
void shift(int l) {
    push(i);
    value[top]=yylval;
    sym=yylex(); }
void reduce(int l) {
    // 생략
    switch (i) { // 규칙번호에 따른 수식 값 계산
        case 1: value[top]=value[old_top+1]+value[old_top+3];
                break;
        case 2: value[top]=value[old_top+1];
                break;
        case 3: value[top]=value[old_top+1]*value[old_top+3];
                break;
        case 4: value[top]=value[old_top+1];
                break;
        case 5: value[top]=value[old_top+2];
                break;
        case 6: value[top]=value[old_top+1];
                break;
        default: yyerror("parsing table error");
                break;
    } }
}
```

C Program(2)

```
int yylex() {  
    // 생략 ...  
    int i=0;  
    if (isdigit(ch)) {  
        do {  
            yytext[i]=ch;  
            ch=getchar();  
        } while (isdigit(ch));  
        yytext[i]=0;  
        yylval=atoi(yytext);  
        return(NUMBER); }  
    // 생략 ...  
}
```



Yacc 프로그램

- Yacc 명세서

- 선언부 (declarations)

- %%

- 문법과 번역규칙부 (rules)

- %%

- 보조 프로그램부 (supporting programs)

Yacc 작성 (1)

```
%start E
%token number
%%
    E    : E '+' T
          | T
          ;
    T    : T '*' F
          | F
          ;
    F    : '(' E ')'
          | '0'
          | '9'
          ;
%%
yylex( ) {
    char ch;
    ch=getchar();
    return ch;
}
void main ( ) {
    yyparse();
}
```

문법 G_2

$$\begin{array}{lcl} E & \rightarrow & E + T \\ & | & T \\ T & \rightarrow & T * F \\ & | & F \\ F & \rightarrow & (E) \\ & | & 0 \\ & | & 9 \end{array}$$

Yacc 작성 (2)

```
%{
#include 'y.tab.h'
}%
%start E
%token  number
```

```
#define PLUS 1
#define STAR 2
#define LP 3
#define RP 4
#define number 5
```

```
%%
S      : E      { printf("%d", $1); }
      ;
E      : E PLUS T { $$ = $1 + $3; }
      | T      { $$ = $1; }
      ;
T      : T STAR F { $$ = $1 * $3; }
      | F      { $$ = $1; }
      ;
F      : LP E RP  { $$ = $2; }
      | number { $$ = $1; }
      ;
%%
```

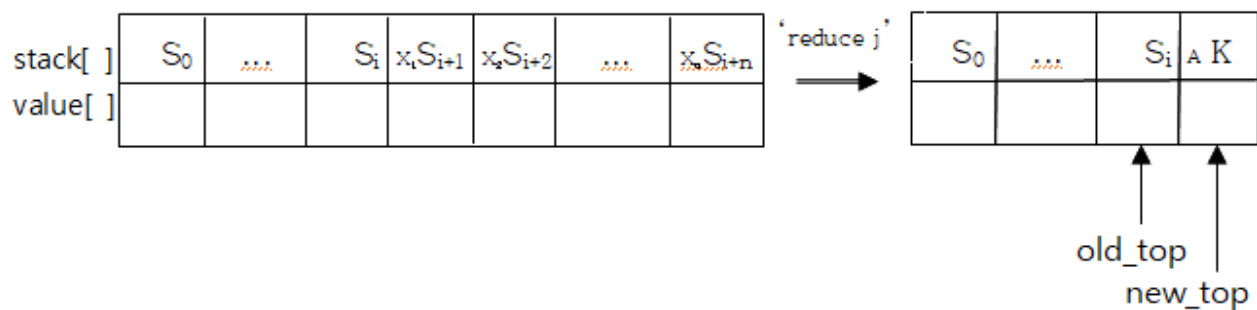
```
yylex() {
    // '+' : return (PLUS);
    // '*' : return (STAR);
    // '(' : return (LP);
    // ')' : return (RP);
    // 정수 : yylval=정수값; return (number);
}
void main () {
    yyparse();
}
```


Yacc 사용방법

번역규칙: $A ::= \alpha_1 \mid \alpha_2 \mid \dots \mid \alpha_n$

A : α_1 {번역규칙₁}
 $\mid \alpha_2$ {번역규칙₂}
 $\mid \dots$ \dots
 $\mid \alpha_n$ {번역규칙_n}

‘reduce’ 동작



$\$, \$1, \$2, \$3 \dots$ 등의 의미