

THEORY:

Linear search is one of the simplest searching algorithm. It is sequentially checked whether each item in the list.

It is worst searching algorithm with worst time complexity. It is a linear approach on the given complexity. It is a linear search, instead of searching linear increase of an ordered list, instead of searching the list in sequence. A binary search is used when the list is sorted by examining the middle term linear search is a technique to compare each item and every element after it.

PRACTICAL NO: 1Ans: Linear searchST 1: Lasted arrayALGORITHM

ST 1: Define a function with two parameter like for conditional statement with range in length of array to find index

ST 2: Now use if conditional statement to check whether user input equal to the elements in array

ST 3: If the conditional in step 2 satisfies, return the index no of the given array if the condition doesn't satisfies, then get out of loop

ST 4: Now initialize a variable to enter elements in the array from user. Now use split() method to split the values.

STEPS: Now initialize a variable as empty array

ST 5: Now we for conditional statement to append the elements given as input by user in the empty array

Step 7: Now again initialize another variable to arr to find the element in array.

Step 8: Again initialize a variable to call the defined function.

Step 9: Use if condition statement to check if the variable in step 8 matches with the Element you want to find.

34

```
# sorted array
## CODE
def linear(arr, n):
    for i in range(len(arr)):
        if arr[i] == n:
            return i
    print("Element not found")
array = input("Enter elements in array").split()
array = [int(i) for i in array]
array.sort()
n = int(input("Enter element to be searched"))
res = linear(array, n)
if res == -1:
    print("Element not found")
else:
    print("Element found at position", res)
```

>> Enter Element in array: 1 2 3 4 5
 >> Elements in array are: [1, 2, 3, 4, 5]
 >> Enter Element to be searched: 2
 >> Enter Element to be searched at position
 The Element is found at position 2

>> Enter Element in array: 3 2 5 1 4
 >> Enter Element in array are: [1, 2, 3, 4, 5]
 >> Elements in array are: [1, 2, 3, 4, 5]
 >> Enter Element to be searched: 6
 >> Enter Element to be searched
 Element not found

Inserted array (arr, n):
 def linear (arr, n):

for i in range(0, n):

if arr[i] == x:

return i

inp = input ("Enter Elements in array") . split ()

array = []

for word in inp:

array.append (int (word))

print ("Elements in array are: ", array)

x = int (input ("Enter the Element to be searched"))

st = linear (array, n)

Q2] Unsorted array:

Algorithm

Step 1: Define a function with two parameters we for conditional statement with range is length of array to find index.

Step 2: Now we off conditional statement to check whether the given statement is equal to element in array.

Step 3: At the condition in step 2 satisfied return the index no. of the given array. If the condition doesn't satisfied we get out of loop.

Step 4: Now initialize a variable to enter elements in the array from user now we split new 2) to split the values.

Ex: We declare another variable to use when we find the element in array.

Op: Again initialize a variable to call the defined function.

if arr == None:

 print("Element found at position : ", pos)

else:

 print("Element not found")

>>> Element elements in array : [2 4 5]

>>> Element to be searched : 6

Element not found at location 2

>>> Element in array one : [2, 4, 5, 3, 7]

>>> Element to be searched : 6

Element not found

HE

binary search!

create
def binary (arr, key):

start = 0

end = len (arr)

while start < end :

mid = (start + end) // 2

if arr[mid] > key:

end = mid

else arr[mid] < key:

start = mid + 1

else: return mid

return -1

array = input ("Enter the sorted list of numbers: ")

arr = []

for ind in array:

arr.append (int (ind))

key = int (input ("Enter element to search: "))

index = binary (arr, key)

if index ==

print ("Element not found")

else:
print ("Element not found at index ", index)

Practical no: 2

Aim: Binary Search

Algorithm:

Step 1: Define a function with two parameters. Now initialise variable with 0. Use while conditional statements to find mid value.

Step 2: Use if condition statement to determine at which position the mid value should point.

Step 3: If the condition doesn't satisfies then return -1

Step 4: Now initialize a variable to enter the elements in the array

Step 5: Use for conditional statement to append the elements in empty array.

Step 6: Now initialize a variable to find the elements in array.

Step 7: Now initialise a variable to call the defined func

Step 8: Now use if conditional statements to determine the index value and print the index value

```
#include <iostream>
#include <iostream>
using namespace std;

int binarySearch(int arr[], int n, int key) {
    int start = 0;
    int end = arr.length();
    int mid = (start + end) / 2;
    while (arr[mid] != key) {
        if (arr[mid] > key) {
            end = mid;
        } else {
            start = mid + 1;
        }
        mid = (start + end) / 2;
    }
    return mid;
}

int main() {
    cout << "Enter the sorted list of numbers: ";
    vector<int> arr;
    arr.append(1, 2, 3, 4, 5, 6, 7, 8, 9);
    cout << "Enter element to search: ";
    int key = int(input("Enter element to search: "));
    int index = binarySearch(arr, key);
    if (index < 0) {
        cout << "Element not found";
    } else {
        cout << "Element found at index " << index;
    }
}
```

>>> Enter the element in array :

2 5 10 12 15 20

>>> Element to be searched = 12

>>> The element was found at index - 3

>>> Enter the element in array :

30 15 6 7 8

>>> Element to be searched = 2

>>> Element was not found

Theory :

Binary search is also known as half interval search. Logarithmic search or binary is a search algorithm that finds the position of target value within a sorted array. If you are looking for numbers which is at the end of list then you need to search entire list in linear search which is time consuming. It can be avoided by using binary search.

AIM: Implementation of bubble sort program in given list.

THEORY: Bubble sort is based on the idea of repeatedly comparing pairs of adjacent elements and then swapping their positions till simplest form is sorting available. In this, we sort the given elements in ascending or descending order by comparing two adjacent elements at a time.

ALGORITHM:

STEP 1: bubble sort algorithm starts by comparing first two elements of an array and swapping if necessary.

STEP 2: If we want to sort the elements of array in ascending order then first element is greater than second then we need to swap the element.

STEP 3: If the given element is smaller than second element then do not swap the element.

```
# code
inp = input("Enter elements") .split()
arr = [ ]
for i in inp:
    arr.append(int(i))
print ("Elements of array before sorting : ", arr)
n = len(arr)
for i in range (0,n):
    for j in range (0,n-i):
        if arr[j] > arr[j+1]:
            temp = arr[j]
            arr[j] = arr[j+1]
            arr[j+1] = temp
print ("Elements of array after bubble sort: ", arr)
```

>>> Enter Element: 2 3 6 1 5
>>> Elements of array before sorting: [2, 3, 6, 1, 5]

>>> Elements of array after sorting:
[1, 2, 3, 5, 6]

STEP 4: Again second & third elements were compared & swapped if it is necessary & this process go on until last 2 record last element is compared & swapped.

STEPS: If there are n elements to be sorted then first unsorted n-1 above should be mentioned to get required output

STEP 5: Stick the output & input of above algorithm at bubble sort stepwise

Practical No: 4

code:
stack

42

Aim: Implementation of stack using python list.

Theory: A stack is linear data structure that can be represented in a real-world form by physical stack or pile. The element in the stack are the top most position. Thus, it follows LIFO principle (Last in First Out). It has 3 basic operation namely: push, pop, peek.

Algorithm:

Step 1: Create a class stack with instance variable item.
Step 2: Define the init method with self argument & initialize the initial value & then initialize to an empty list.

Step 3: Define methods push & pop under the class stack.

Step 4: Use if statement to give the condition that if length of given list is greater than the range of list the print stack is full.

Steps: Use the 'else' statement to print a statement as input the elements into the stack & initialize the value.

If Output :

```
>>> x.push(10)
>>> x.push(37)
>>> x.push(18)
>>> x.push(79)
>>> x.push(72)
>>> x.push(69)
```

The stack is full.

```
>>> x.pop()
```

72

```
>>> x.l
```

[10, 7, 8, 79, 72, 69]

STEP 6: Push method used to insert the element but pop method used to delete the element from the stack.

STEP 7: If in pop method, value is less than 1 the element from stack at topmost position

STEP 8: Assign the element value in push method & print the given value.

STEP 9: Match the input & output of above algo.

STEP 10: First condition checks whether the no. of elements are zero while second case when top i.e. assigned any value. If top is not assigned any value then print that the stack is empty.

HARDWARE PRACTICAL NOTES

```
Print ("Quick Sort")
def partition (arr, low, high):
    i = low - 1
    pivot = arr[high]
    for j in range (low, high):
        if arr[j] < pivot:
            i += 1
            arr[i], arr[j] = arr[j], arr[i]
    arr[i+1], arr[high] = arr[high], arr[i+1]
    return i+1
```

Aim: Implement quick sort to sort the given list.

Theory: The quick sort is a recursive algorithm based on the divide & conquer technique.

ALGORITHM:

Step 1: Choose sort first value, which is called pivot value first. Value element same as our first pivot value and we know that first will eventually end up as last in that list.

Step 2: The position places will happen next. It will find a split point & at the same move other items to the appropriate side of the list either less than or greater than pivot value.

Step 3: Partition begins by locating two position markers, lets call them leftmark & rightmark. To right mark at the beginning & end of remaining item in the list. The goal to wrong side will repeat to first value while also converging on the split point.

Step 4: We begin by increasing leftmark we locate a value but is greater than the pivot value, we then decrement right mark until we find value that is less than the pivot value.

```
I1 = input ("Enter Element in the list : ") - split()
alist1 = [I1]
for bin I1 =
    alist1.append (int (bin))
print ("Element input list are : ", alist1)
n = len(alist1)
quickSort (alist1, 0, n-1)
print ("Element after quicksort are : ", alist1)
```

11

Output:

Quick sort

Enter elements in the list : 21, 22, 20, 30, 24, 56

Elements in list are: [21, 22, 20, 30, 24, 56]

Elements after adjustment are [20, 21, 22, 24, 30, 56]

At the point where quick sort function is called. Adjustment has been made. Leftmost part of the list is now the sorted part.

Step 6: The pivot value can be exchanged with the content of split point and pivot value is placed in place.

Step 7: In addition all the items to left of split point are less than pivot value & all the items left to the right of split point are greater than it. The list can now be divided at split point and quarters turn. But the list can now be divided at split point & quick sort can be applied on either side of the pivot value.

Step 8: The quick sort function invokes a recursive function quick sort help

Step 9: Quick sort keeps beginning with some base case as the number of steps: If length of the list is less than or equal to 1 it is already sorted.

Step 10: If at a point when list can be partitioned & recursive function

Step 11: The recursive function implements the process described earlier.

P.B
PRACTICAL NO:6.

CODE:

46

- * Aim: Implementing a queue using python list.

Theory: Queue is a linear data structure which has two references front & rear.

Implementation of queue using python list is the simplest as the python list provides built-in function to perform the specified operation of queue. It is based on principle that a new element is inserted in rear & element of queue is deleted which is at front. It is called F/FIFO principle.

Queue C: Create a new empty queue

Enqueue () : Insert an element at the rear of the queue and similar to that of insertion of linked using tail.

Dequeue () : Returns Element which was at the front the front is moved to the successive Element In dequeue operation cannot remove Element if the queue is empty.

```

class queue:
    global r
    global f
    def __init__(self):
        self.r=0
        self.f=0
        self.l=[0,0,0,0,0]
    def add(self,data):
        n = len(self.l)
        if self.r < n:
            self.l[self.f] = data
            self.r += self.r + 1
        else:
            print("Queue is full")
        del remove(self)
        n = len(self.l)
        if self.r < n - 1:
            print(self.l[self.f])
            self.f = self.f + 1
        else:
            print("Queue is Empty")
    Q = Queue()
    Q.add(10)
    Q.add(20)
    Q.add(30)
    Q.add(40)

```

Q.add (10)
Q.add (20)

Q.remove ()
Q.remove(0)
Q.remove(1)
Q.remove()
Q.remove(1)
Q.remove()

ALGORITHM:

STEP1: Define a class queue & assign global variable
 two define init() method with self argument
 in init() function or initialize the init value
 with help of self argument

STEP2: Define empty list and define enqueue() method
 with 2 arguments , assign length of empty list

STEP3: Use if statements that length is equal to size
 then queue is full or else insert element
 in empty list or display that queue element
 added successfully & increment.

STEP4: Define dequeue() with self argument and this
 check if front element next front is equal
~~to length of list then display queue is~~
~~empty or else given front is at 0 &~~
~~writing front, delete element from front size~~
~~2 increment by 1.~~

STEPS: Now, all Queue() funct. & give element that has
 to be added in Empty list by using Enqueue
 & end print the list after adding & some few
 deletions & displaying new list after deleting the
 element in list.

Practical No: 7

48

Evaluation of positive expression

Aim: Program on Evaluation of given strings by using stack in python environment i.e. Python.

THEORY: One common expression in form of any numbers function we took one of the numbers of the expression in program. Reading the expression always from left to right in positive.

ALGORITHM:

STEP1: Define evaluate() function then create an empty stack in python

STEP2: Convert the string to a list by using list string method split.

STEP3: Calculate range of string & print it.

STEP4: Use for loop to assign the range of string then give condition using if statement

```

def evaluate(s):
    k = 0
    n = len(k)
    stack = []
    for i in range(n):
        if k[i] == '+':
            stack.append(int(k[i]))
        else:
            a = stack.pop()
            b = stack.pop()
            stack.append(str(int(b) - int(a)))
    if k[1] == '*':
        a = stack.pop()
        b = stack.pop()
        stack.append(str(int(b) * int(a)))
    else:
        a = stack.pop()
        b = stack.pop()
        stack.append(str(int(b) / int(a)))
    return stack.pop()
s = "869 * +"
v = evaluate(s)
print ("The evaluated value is: ", v)
print ("marked answer")

```

Output:

>>> She evaluated value is: 62
She evaluated value is: 62
which quantum

STEPS: Scan the token list from left to right till taken
in an equation, convert it from a string to an
integer & push the value onto the 'P'.

Step 1: If the token is an operator *, /, +, -, ^, it will
need 2 operands. On the 'P' twice the first
pop is 2nd operand and 1st pop is the first operand

Step 2: Perform the arithmetic operation push the result
back on the 'P'

Step 3: When the input conversion has been completely
finished the result is on the stack pop the 'P'
and return the value

Step 4: Print the result of string after evaluation of postfix

Step 5: Attach tokens and input of above algorithm

CODE:

Q1: Implementation of single linked list by adding
from last position

```
class Node
{
    public int data;
    public Node next;
}
```

5.

SUPER: A linked list is a series of data structures
where each element is a node
having both value members containing data
and individual address of the address list is also
node. Node consists of 2 parts ① Data
Data gives information of Element
value
and value is the link made by pointing to
it. so one of design list if we add/insert
new element from the list, all elements of
list has to adjust itself. Every time one add/
one element then the linked list is referred to
having new type of link.

IMPLEMENTATION:

```
class Node
{
    public int data;
    public Node next;
}

Node head = null;

void head (Node head, int data)
{
    Node newnode = new Node();
    newnode.data = data;
    newnode.next = head;
    head = newnode;
}
```

5.

Q2: Inserting of linked list can visit all the nodes

in linked list in order to program from a
function to main.

Q3: The function linked list returns can be modified with
help of this function make the list in form of
separated by blank spaces & new list.

```
def display(self):
```

```
    head = self.s
```

```
    while head != None:
```

```
        print(head.data)
```

```
    head = head.next
```

```
print("Linked list")
```

```
head = head.next
```

```
print(head.data)
```

```
start = linkedlist()
```

```
start.addL(50)
```

```
start.addL(60)
```

```
start.addL(70)
```

```
start.addB(40)
```

```
start.addB(30)
```

```
start.addB(20)
```

```
start.display()
```

```
Print("Linked Queue")
```

OUTPUT

```
>>>
```

```
20
```

```
30
```

```
40
```

```
50
```

```
60
```

```
70
```

>>>

Mallesh Gururani,

STEP 3: Thus, Entire list can be traversed using node which is referenced by head pointer of the linked list.

STEP 4: Now that we know that we can traverse the entire linked list using head pointer, we should only use it to refer the first node of list only.

STEP 5: We should not use head pointer to traverse entire linked list because the head pointer is our only reference to 1st node in the linked list, modifying reference to the head pointer can lead to changes which we cannot reverse back.

STEP 6: We may lose the reference to the 1st node in our linked list & hence must do our linked list so as to avoid making some unwanted changes to the 1st node, we will use a temporary node to traverse the entire linked list.

STEP 7: We will use this temporary node as a copy of the node we are currently traversing. Since we are walking constantly making a copy of temporary node the delta type of the temporary node should alone be node

STEP 8: Now that we're trying to the first node, if we want to access 2nd node of list follow it on next node of the 1st node

then the file is made in reflected by copies to
other file but in its mode as being
one because we can have one of notes
and file, taking some modified by others
and then now we do find permanent notes
as stable info

Ques) The file note as stored but is referred to
by last note has last note it stored in
one line only next note has value of the
file at the last note in store.

Ques) the case when the last note if file is
written by him he would say how to store him
last file or how to identify situation in
which his last note is last or not
Answer) File making from input & output option

25

```

def sort (arr, l, m, r):
    n1 = m - l + 1
    n2 = r - m
    L = [0] * n1
    R = [0] * n2
    for i in range (0, n1):
        L[i] = arr [l+i]
    for j in range (0, n2):
        R[j] = arr [m+1+j]
    i=0
    j=0
    k=l
    while i < n1 & j < n2:
        if L[i] <= R[j]:
            arr [k] = L[i]
            i+=1
        else:
            arr [k] = R[j]
            j+=1
        k+=1

```

PRACTICAL NO: 9

53

MERGE SORT: ~~discusses time complexity of merge sort~~

Aim: Implementation of merge sort by using python.

THEORY: Merge sort is a divide & conquer algorithm. It divides input array in two halves, calls itself for the two halves and then merges the two sorted halves. The merge (m, r, v) arr [$m+1, n$] are sorted and merges the two sorted sub-arrays.

ALGORITHM:

STEP1: The list is divided into left & right in each recursive call until two adjacent elements are obtained.

STEP2: Now begins the sorting process. The i and j iterator traverse the two halves in each call. The k iterator traverses the entire list & makes changes along the way.

STEP3: If the value at j , $L[i] \leq R[j]$ average to the arr [$i:j$] sort R is incremented if not the $R[i:j]$ is closed.

Step 4: this way, the values being assigned through arr_j are all sorted

Step 5: At the end of this loop, one of the halves may or have been traversed completely, remaining short in the list.

Step 6: Thus the merge sort has been implemented

```

while rk < n2:
    curr [rk] = R [kj]
    jk += 1
    k = k + 1

def mergesort curr, l, r):
    if l < r:
        m = int ((l + (r - 1)) / 2)
        mergesort curr, l, m)
        mergesort curr, m + 1, r)
        sort (curr, l, m, r)

curr = [12, 23, 34, 56, 38, 36, 98, 42]
print (curr)
m = len (curr)
mergesort (curr, 0, m - 1)
print (curr)

```

Output:

~~12, 23, 34, 56, 78, 45, 38, 98, 42~~
~~{12, 23, 34, 56, 42, 45, 78, 86, 98}~~

```
+-----+
set1 = set()
set2 = set()

for i in range(1,15):
    set1.add(i)
    set2.add(i)

print("set1: ", set1)
print("set2: ", set2)
print("\n")

set3 = set1 / set2

print("Union of set1 & set2: set3")
set4 = set1 & set2

print("Intersection of set1 and set2: set4", set4)
print("\n")

if set3 > set4:
    print("Set 3 is superset of set 4")
else:
    print("Set 3 is subset of set 4")
    if set4 < set3:
        print("Set 4 is subset of set 3")
    print("Set 4 is subset of set 4")
    print("\n")
```

Aim: Implementation of sets using python

ALGORITHM:

STEP 1: Define two empty set as set1 and set2 now we for statement provide the range of above 2 sets

STEP 2: Now add() method is used for addition the element according to given range then print the sets for addition

Print("Union of set1 & set2: set3")
set3 = set1 / set2
print("Intersection of set1 and set2: set4", set4)
print("\n")

STEP 3: Find the union & intersection of above 2 sets by using &, | method print the sets of union & intersection of sets.

STEP 4: Use if statement to find out the subset & superset of set3 and set4. Display the above set

STEP 5: Display the element in set3 is not in set4 using mathematical operation

Steps: Use if disjoint() to check that anything is common or elements is present or not or not then display. Thus it is mutually exclusive and non disjoint.

STEP 6: Use clear() to remove or delete few elements from the set after clearing the elements present in the set.

56

```
if set 4 ⊂ set 3:  
    print ("set 4 is subset of set 3")  
    print ("\n")  
set 5 = set 3 - set 4  
print ("Elements in set 3 and not in set 4: set 5", set 5)  
print ("\n")  
if set 4 is disjoint (set 5):  
    print ("set 4 and set 5 are mutually exclusive\n")  
set 4.clear()  
print ("set After applying clear, set 5 is empty set:")  
print ("set 5 =", set 5)
```

OUTPUTS

Set1 : {8, 9, 10, 11, 12, 13, 14}
Set2 : {1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11}

Union of set 1 & set 2: set 3

{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14}
Intersection of set 1 and set 2 : {8, 9, 10, 11}

Set 3 is superset of set 4

Elements in set 3 and not in set 4: set 5

{1, 2, 3, 4, 5, 6, 7, 8, 13, 14}

Set 4 & set 5 are mutually exclusive
After applying clear, set 5 is empty set
set 5 = set 1

#CODE

```

class node:
    def __init__(self, value):
        self.left = None
        self.val = value
        self.right = None

class BST:
    def __init__(self):
        self.root = None

    def add(self, value):
        p = node(value)

        if self.root == None:
            self.root = p
        else:
            self._add(self.root, p)

    def _add(self, root, p):
        if p.val < root.val:
            if root.left == None:
                root.left = p
            else:
                self._add(root.left, p)
        else:
            if root.right == None:
                root.right = p
            else:
                self._add(root.right, p)

    def print(self):
        print("Root is added successfully", self.root)

```

THEORY: Binary tree is a tree which supports maximum of 2 children for any node within the tree.
 Thus any node partition node can have either 0 or 1 or 2 children. There is another identify of binary tree that it is ordered such that one child is identified as left child and other as right child.

INORDER: Traverse the left subtree. The left subtree inform might have left and right subtree.

PREORDER: Visit the root node traversal the left subtree and right subtree because the right subtree.

POSTORDER: Traverse the left subtree then left subtree inform might have left & right subtree.

ALGORITHM

Step 1: Define class node & define init() with 2 args
Initialize the value in this method

Step 2: Again, define a class bst that is binary search tree with func with self argument & only the root is none.

Step 3: Define add() after adding the node, Define a variable p that p=node(value)

Step 4: Use if statement for checking if the conditional statement for it node is less than the main node then put on arrange them in leftside

Step 5: Use while loop for detecting the node is less than or greater than the main node & break the loop if it is not sufficient.

Step 6: Use if statements written that else statement for deciding last node is greater than main root then put it into rightside

if right == 0
Print ("val: "none)
Print ("p.val: "none)
Print ("succesfully")

added to rightside successfully
break
else:
right
if right == 0
Print ("val: "none)
Print ("p.val: "none)
Print ("succesfully")
else:
if root == None:
return
else:
if root == None:
return
else:
Print ("val: "none)
Print ("p.val: "none)
Print ("succesfully")
else:
if root == None:
return
else:
Print ("val: "none)
Print ("p.val: "none)
Print ("succesfully")
else:
if root == None:
return
else:
Print ("val: "none)
Print ("p.val: "none)
Print ("succesfully")

use .post-order (root, left)
 post-order (root, right)
 print (root, val)

bst0

#output:

>>> add()

root is added successfully

>>> add(2)

2 node is added to rightside successfully

>>> add(4)

4 node is added to rightside successfully

>>> add(5)

5 node is added to rightside successfully

>>> add(2)

2 node is added to rightside successfully

>>> print ("No. nodes:", nodes (root))

nodes:

1
2
3
4
5

nodes: None

STEP 7: After this, left-side tree & right subtree repeat this method to arrange the node accordingly to binary search tree.

Steps: Define inorder(), preorder() & postorder() with root as argument & use if statement based root is none & return that all.

Step 8: In inorder, else statement used for giving this condition at first left root & then right node.

Step 9: For preorder we have to give condition in else that first root, left & then right node.

Step 10: For postorder in else part assign left & right of root

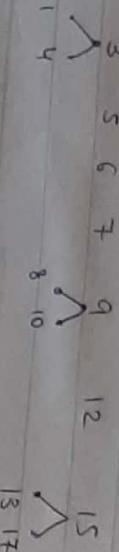
Step 12: Display the output & input

INORDER : (LVR)

STEP 1 :



STEP 2:



STEP 3: 1 3 4 5 6 7 8 9 10 12 13 15 17

• PREORDER : (VLR)

```
>>> Print ("\\n Preorder:", Preorder(root))  
Preorder:  
1  
3  
5  
4  
2  
7  
8  
10  
15  
13  
17
```

Preorder: None

```
>>> Print ("\\n Postorder:", Postorder(+root))  
Postorder:  
3  
5  
4  
2  
1  
7  
8  
10  
15  
13  
17
```

Postorder: None

* Binary tree:

↳ search

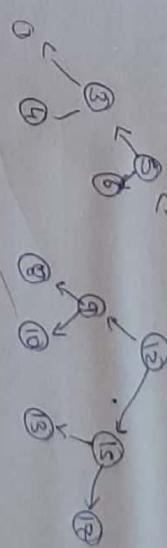
STEP 1: 7

A binary tree with root node 7. Node 7 has left child 5 and right child 12. Node 5 has left child 1 and right child 4. Node 12 has left child 9 and right child 15. Node 9 has left child 3 and right child 10. Node 15 has left child 13 and right child 17.

STEP 2:

7 5 3 6 12 9 15

STEP 3: 7 5 3 1 4 6 12 9 8 10 13 15 17

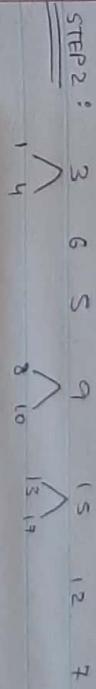


POSTORDER : (LRU)

STEP 1 :



STEP 2 :



STEP 3 : 1 4 3 6 5 8 10 9 12 13 17 15