

DATA STRUCTURES & ALGORITHMS

Sorting - Merge Sort

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Divide-and-Conquer

- **Divide** the problem into a number of sub-problems
 - Similar sub-problems of smaller size
- **Conquer** the sub-problems
 - Solve the sub-problems **recursively**
 - Sub-problem size small enough \Rightarrow solve the problems in straightforward manner
- **Combine** the solutions of the sub-problems
 - Obtain the solution for the original problem

Merge Sort Approach

- To sort an array $A[\text{beg} \dots \text{end}]$:
- **Divide**
 - Divide the n -element sequence to be sorted into two subsequences of $n/2$ elements each
- **Conquer**
 - Sort the subsequences recursively using merge sort
 - When the size of the sequences is 1 there is nothing more to do
- **Combine**
 - Merge the two sorted subsequences

Merge Sort

MERGE_SORT(arr, beg, end)

if beg < end

set mid = (beg + end)/2

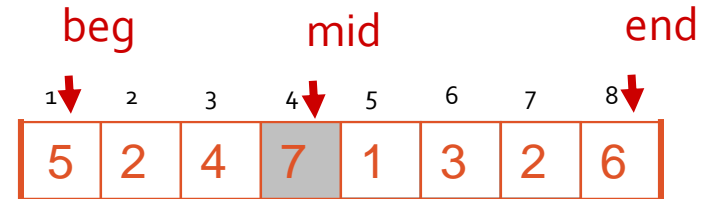
MERGE_SORT(arr, beg, mid)

MERGE_SORT(arr, mid + 1, end)

MERGE (arr, beg, mid, end)

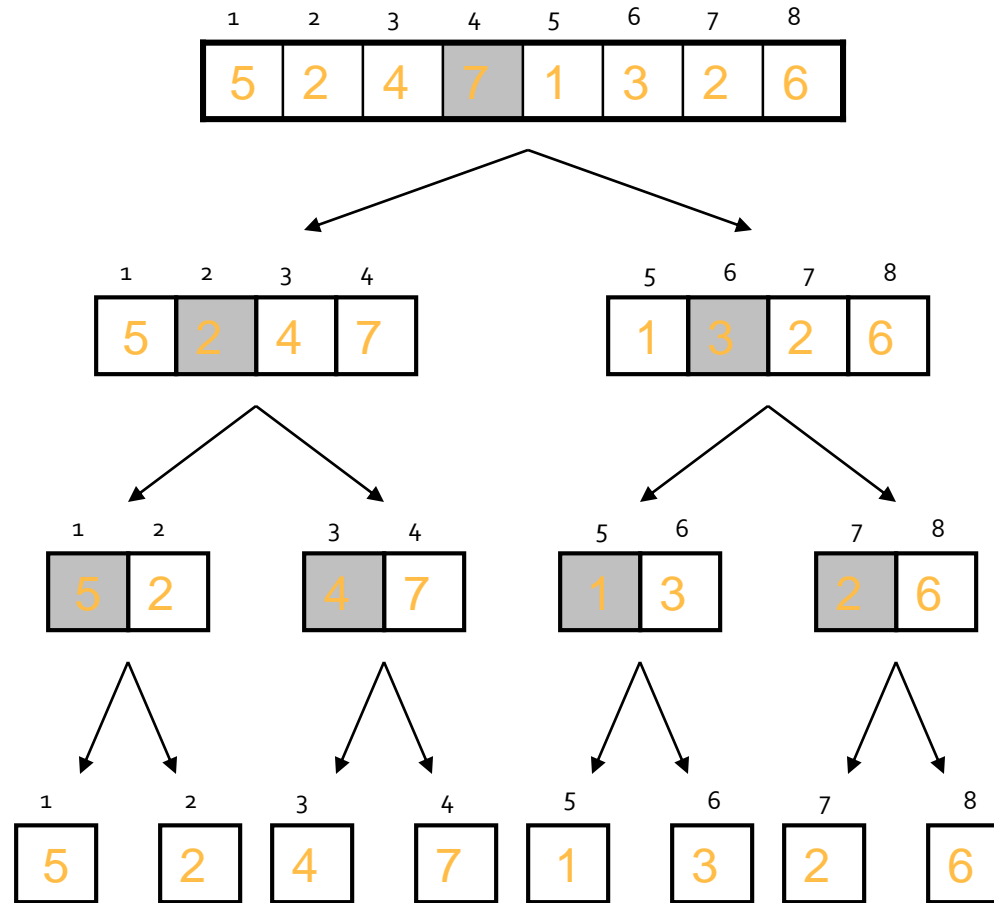
end of if

END MERGE_SORT



Example – n Power of 2

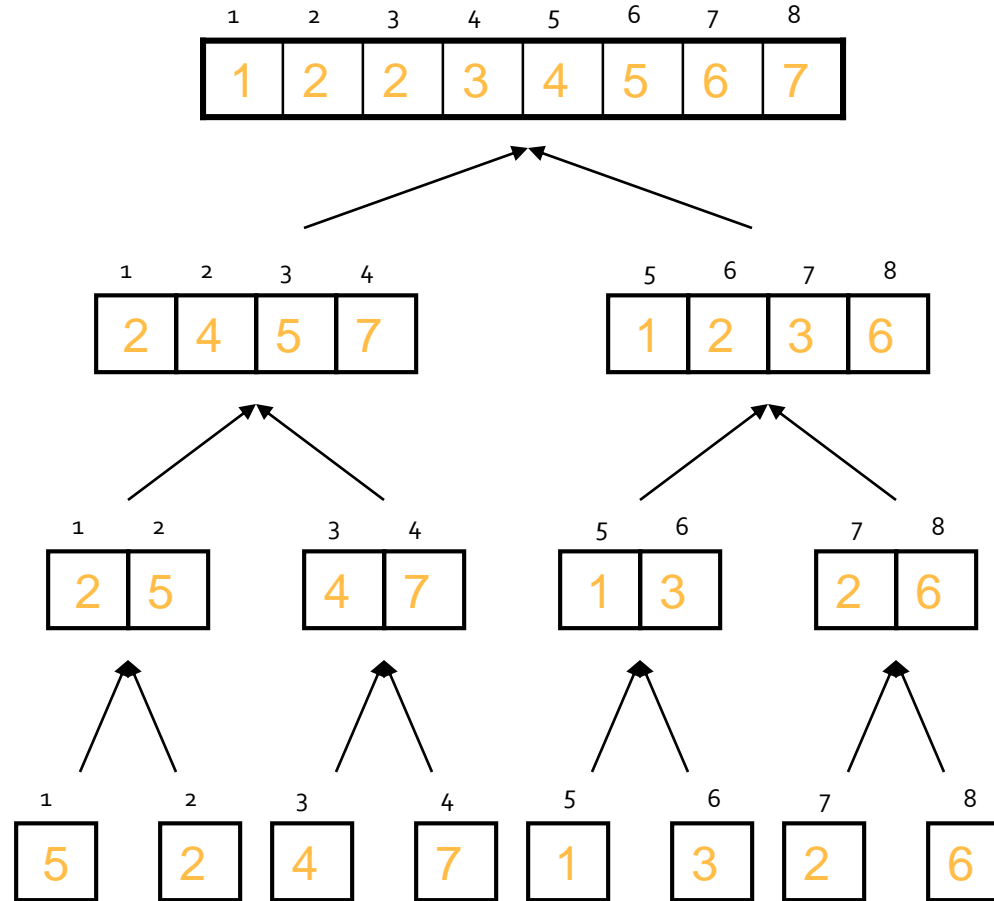
Divide



mid = 4

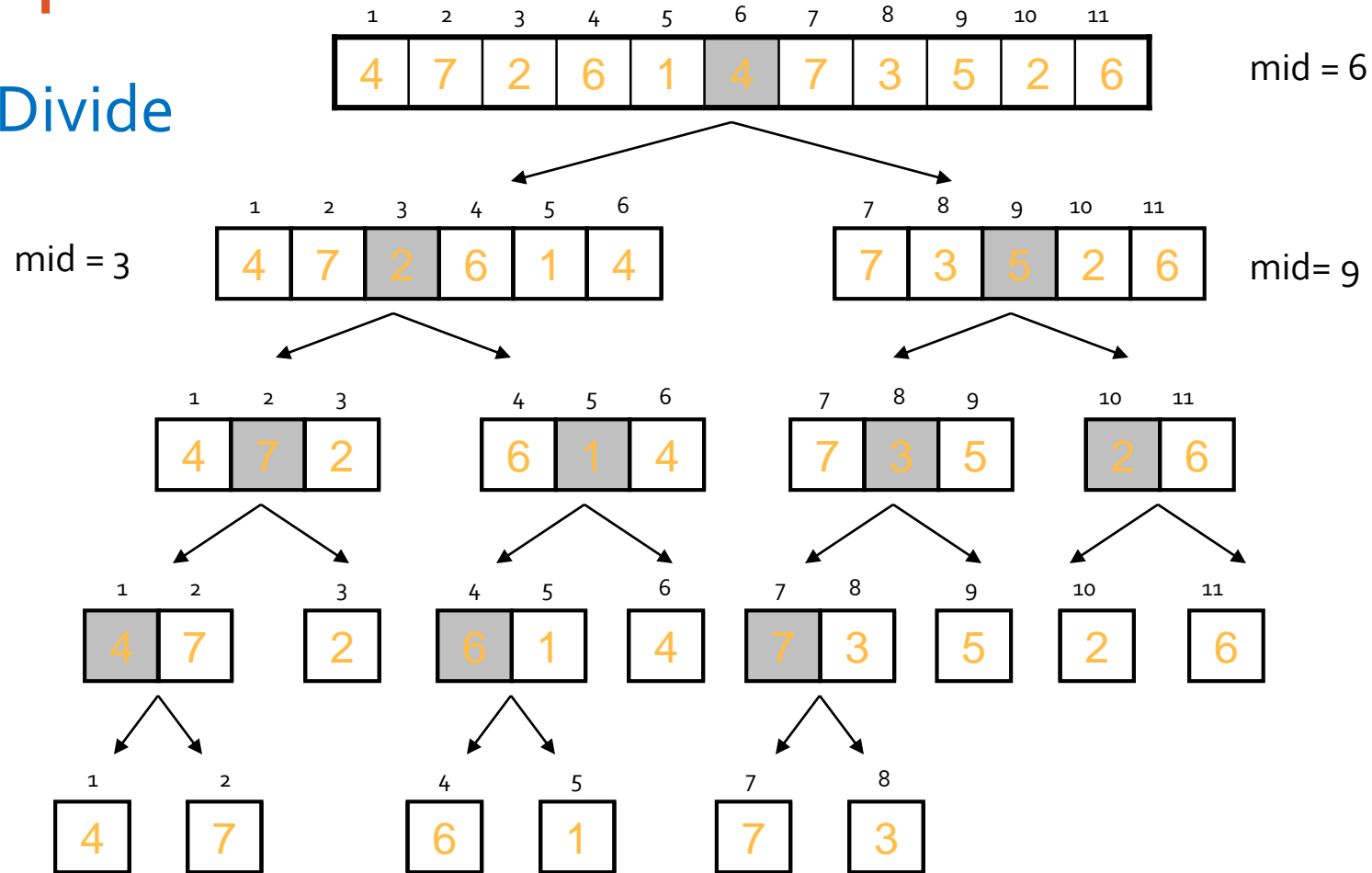
Example – n Power of 2

Conquer
and
Merge



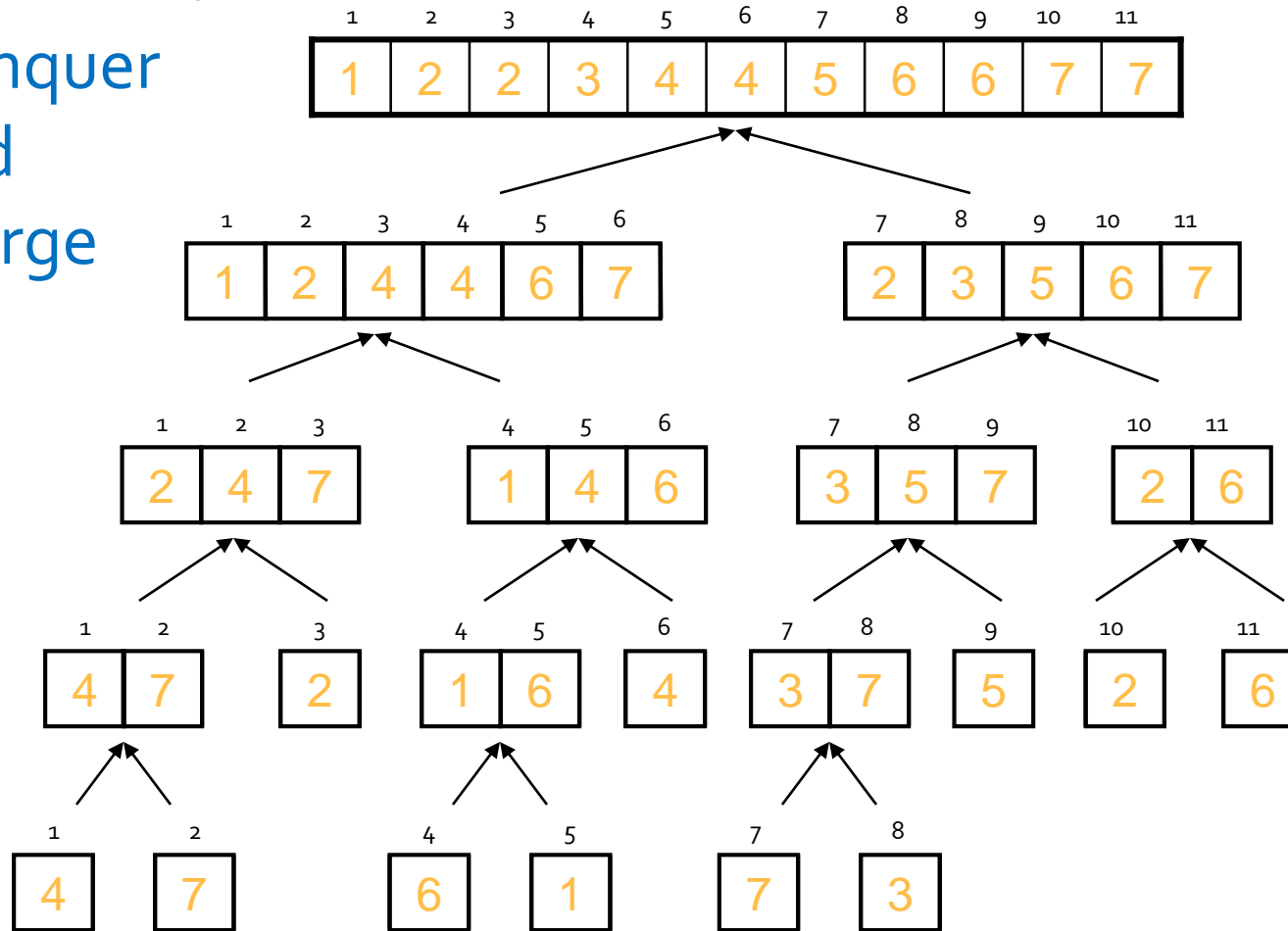
Example – n Not a Power of 2

Divide

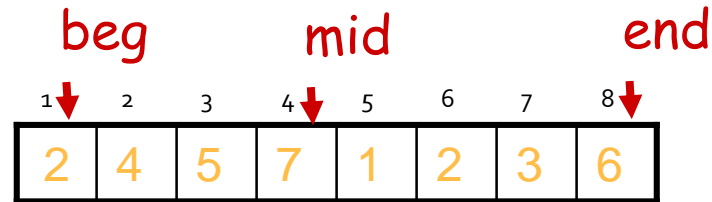


Example – n Not a Power of 2

Conquer
and
Merge



Merging



The important part of the merge sort is the **MERGE** function. This function performs the merging of two sorted sub-arrays that are **A[beg...mid]** and **A[mid+1...end]**, to build one sorted array **A[beg...end]**. So, the inputs of the **MERGE** function are **A[], beg, mid, and end**.

Merging

The algorithm for merge maintains three pointers, one for each of the two arrays and one for maintaining the current index of the final sorted array.

Have we reached the end of any of the arrays?

No:

Compare current elements of both arrays

Copy smaller element into sorted array

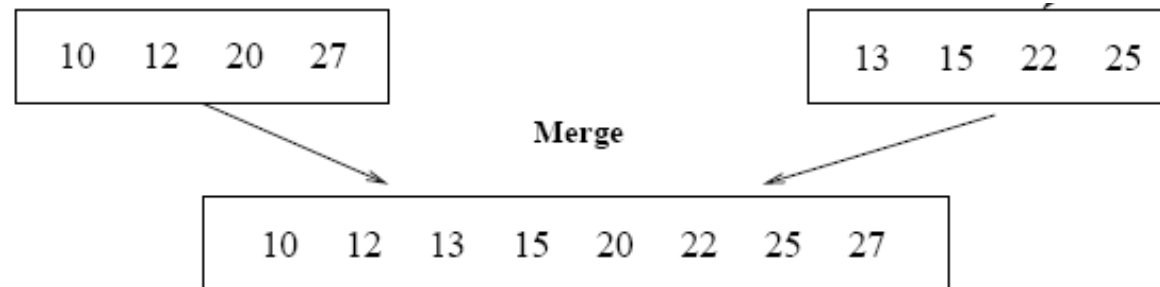
Move pointer of element containing smaller element

Yes:

Copy all remaining elements of non-empty array

Running Time of Merge (assume last **for** loop)

- Initialization (copying into temporary arrays):
 - $\Theta(n_1 + n_2) = \Theta(n)$
- Adding the elements to the final array:
 - n iterations, each taking constant time $\Rightarrow \Theta(n)$
- Total time for Merge:
 - $\Theta(n)$



Analyzing Divide-and Conquer Algorithms

- The recurrence is based on the three steps of the paradigm:
 - $T(n)$ – running time on a problem of size n
 - **Divide** the problem into a subproblems, each of size n/b : takes $D(n)$
 - **Conquer** (solve) the subproblems $aT(n/b)$
 - **Combine** the solutions $C(n)$

$$T(n) = \begin{cases} \Theta(1) & \text{if } n \leq c \\ aT(n/b) + D(n) + C(n) & \text{otherwise} \end{cases}$$

MERGE-SORT Running Time

- **Divide:**
 - compute `mid` as the average of `beg` and `mid`: $D(n) = \Theta(1)$
- **Conquer:**
 - recursively solve 2 subproblems, each of size $n/2 \Rightarrow 2T(n/2)$
- **Combine:**
 - MERGE on an n -element subarray takes $\Theta(n)$ time $\Rightarrow C(n) = \Theta(n)$

$$T(n) = \begin{cases} \Theta(1) & \text{if } n = 1 \\ 2T(n/2) + \Theta(n) & \text{if } n > 1 \end{cases}$$

Merge Sort - Discussion

- Running time insensitive of the input
- Advantage
 - Guaranteed to run in $O(n \log n)$
- Disadvantage
 - Requires extra space $\approx N$

Sorting Challenge 1

Problem: Sort a file of huge records with tiny keys

Example application: Reorganize your MP-3 files

Which method to use?

- A. merge sort
- B. selection sort
- C. bubble sort
- D. a custom algorithm for huge records/tiny keys
- E. insertion sort

Sorting Files with Huge Records and Small Keys

- Insertion sort or bubble sort?
 - NO, too many exchanges
- Selection sort?
 - YES, it takes **linear** time for exchanges
- Merge sort or custom method?
 - Probably not: selection sort simpler, does less swaps

Sorting Challenge 2

Problem: Sort a huge randomly-ordered file of small records

Application: Process transaction record for a phone company

Which sorting method to use?

- A. Bubble sort
- B. Selection sort
- C. Mergesort
- D. Insertion sort

Sorting Huge, Randomly - Ordered Files

- Selection sort?
 - NO, always takes quadratic time
- Bubble sort?
 - NO, quadratic time for randomly-ordered keys
- Insertion sort?
 - NO, quadratic time for randomly-ordered keys
- Mergesort?
 - YES, it is designed for this problem

Sorting Challenge 3

Problem: sort a file that is already almost in order

Applications:

- Re-sort a huge database after a few changes
- Double check that someone else sorted a file

Which sorting method to use?

- A. Mergesort
- B. Selection sort
- C. Bubble sort
- D. A custom algorithm for almost in-order files
- E. Insertion sort

Sorting Files That are Almost in Order

- Selection sort?
 - NO, always takes quadratic time
- Bubble sort?
 - NO, bad for some definitions of “almost in order”
 - Ex: B C D E F G H I J K L M N O P Q R S T U V W X Y Z A
- Insertion sort?
 - YES, takes linear time for most definitions of “almost in order”
- Mergesort or custom method?
 - Probably not: insertion sort simpler and faster