

# DATA STRUCTURES & ALGORITHMS

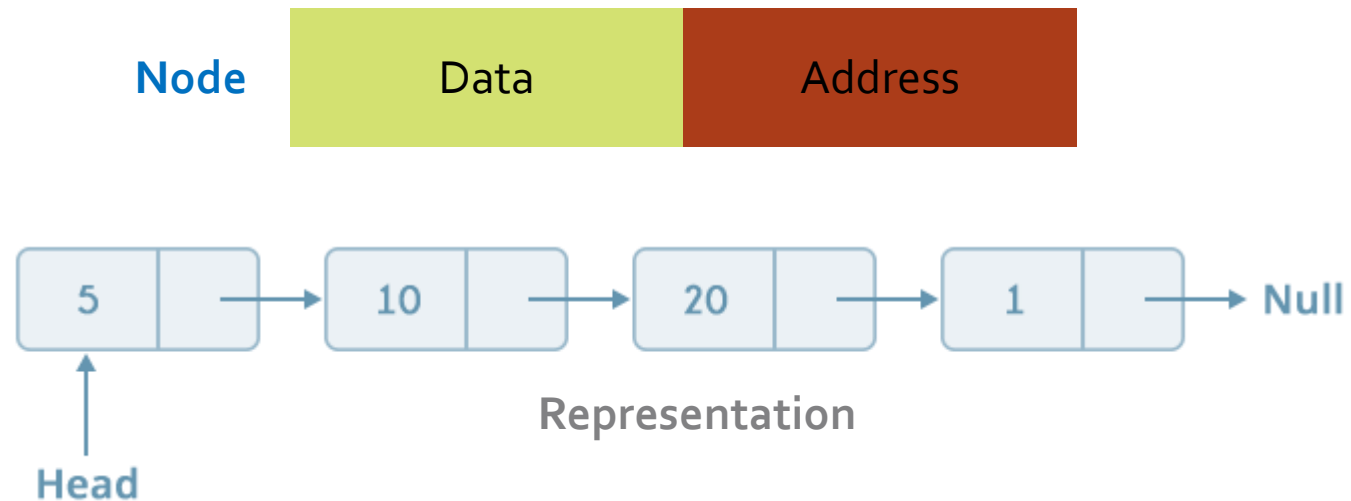
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Linked Lists

Instructor: Engr. Laraib Siddiqui

# Linked list

- It's a linear data structure, consist of nodes.
- Nodes includes:
  - ✓ Data - elements
  - ✓ Address/Link - address of another node.

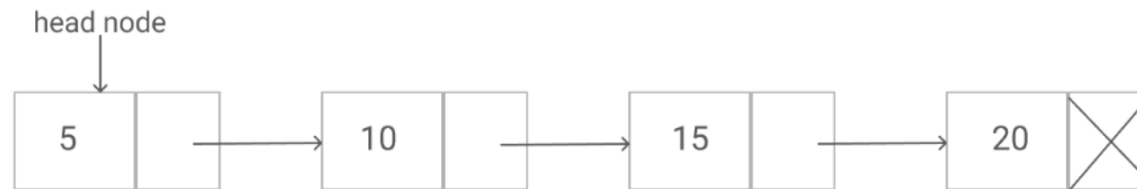


# Linked list vs Arrays

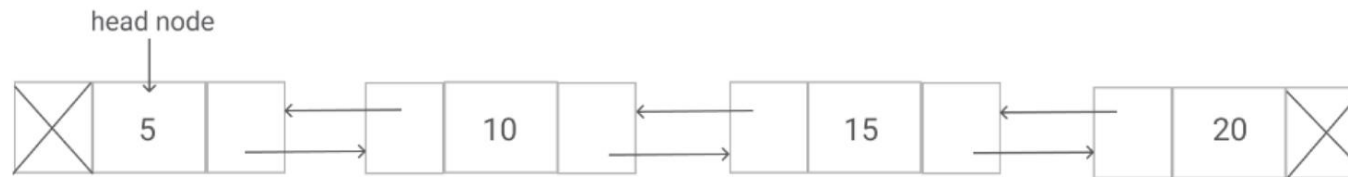
| Array   | Linked List  |
|---|--|
| Size is fixed; resizing is expensive              | Size is dynamic  |
| Data can be access randomly                       | No random access   |
| Insertion and Deletion operation takes more time. | Insertion and Deletion operations are fast in linked list. |
| Data is stored in consecutive location.           | Data is randomly stored.                                   |

# Types

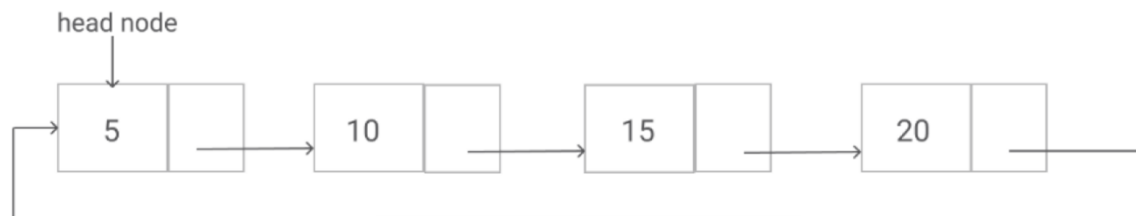
- Singly linked list – Item navigation is **forward** only.



- Doubly linked list – Items can be navigated **forward** and **backward**.



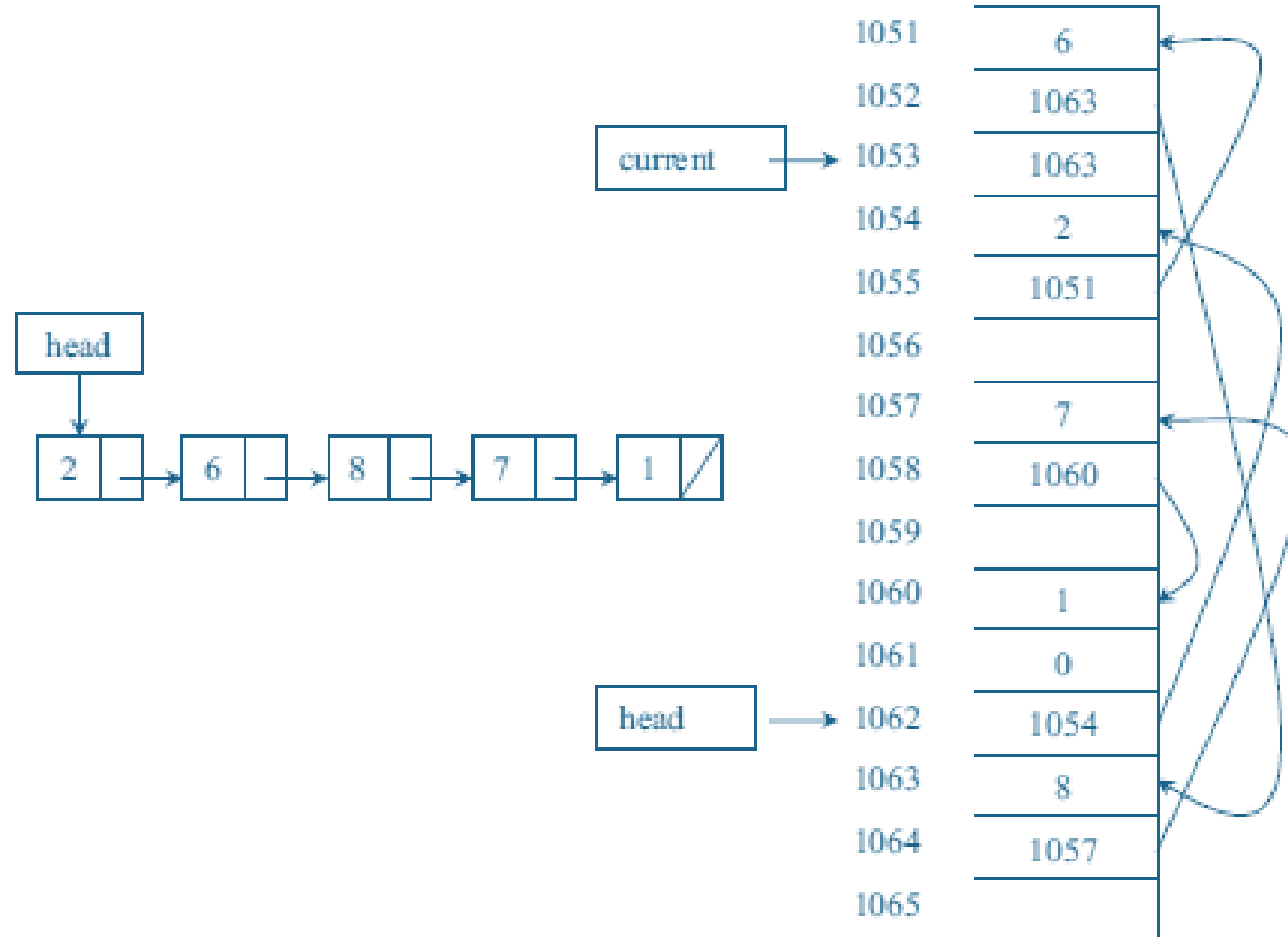
- Circular Linked List – Last item contains link of the **first element** as next and the first element has a link to the **last element** as previous.



# Comparison

|          | Singly                                   | Doubly                                    |
|----------|--|---|
| Strength | Consume less memory<br>Easy to implement | Traversing can be done in both directions |
| Weakness | Access previous elements not easy        | Consume more memory                       |

# Memory Layout



# Creation of a node

Using Structure

```
struct Node {  
    int data;  
    struct Node* next;  
};
```

Using Class

```
class Node {  
public:  
    int data;  
    Node* next;  
};
```

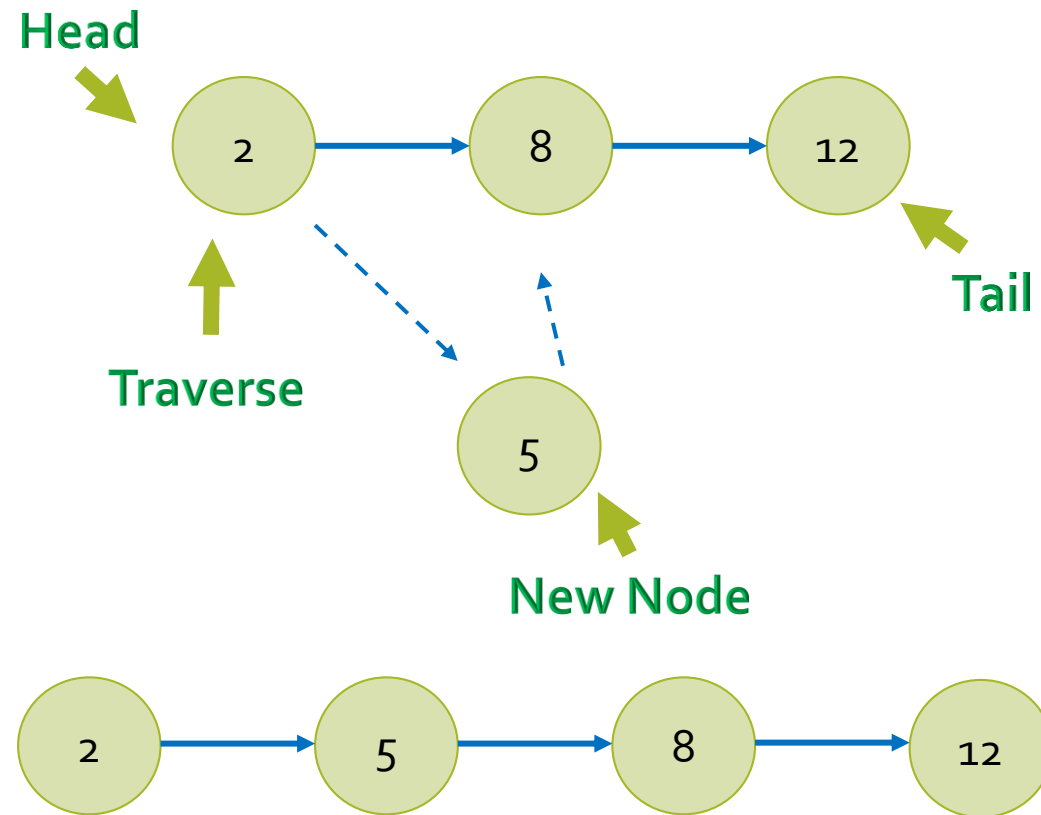
# Traversing

Let LIST be a linked list in memory. Following algorithm traverses LIST, applying an operation PROCESS to each element of LIST. The variable PTR points to the node currently being processed.

1. Set PTR: = START. [Initializes pointer PTR]
2. Repeat Steps 3 and 4 while PTR  $\neq$  NULL.
3.     Apply PROCESS to INFO [PTR]
4.     Set PTR:= LINK[PTR]. [PTR now points to the next node.]
- [End of Step 2 loop.]
5. Exit



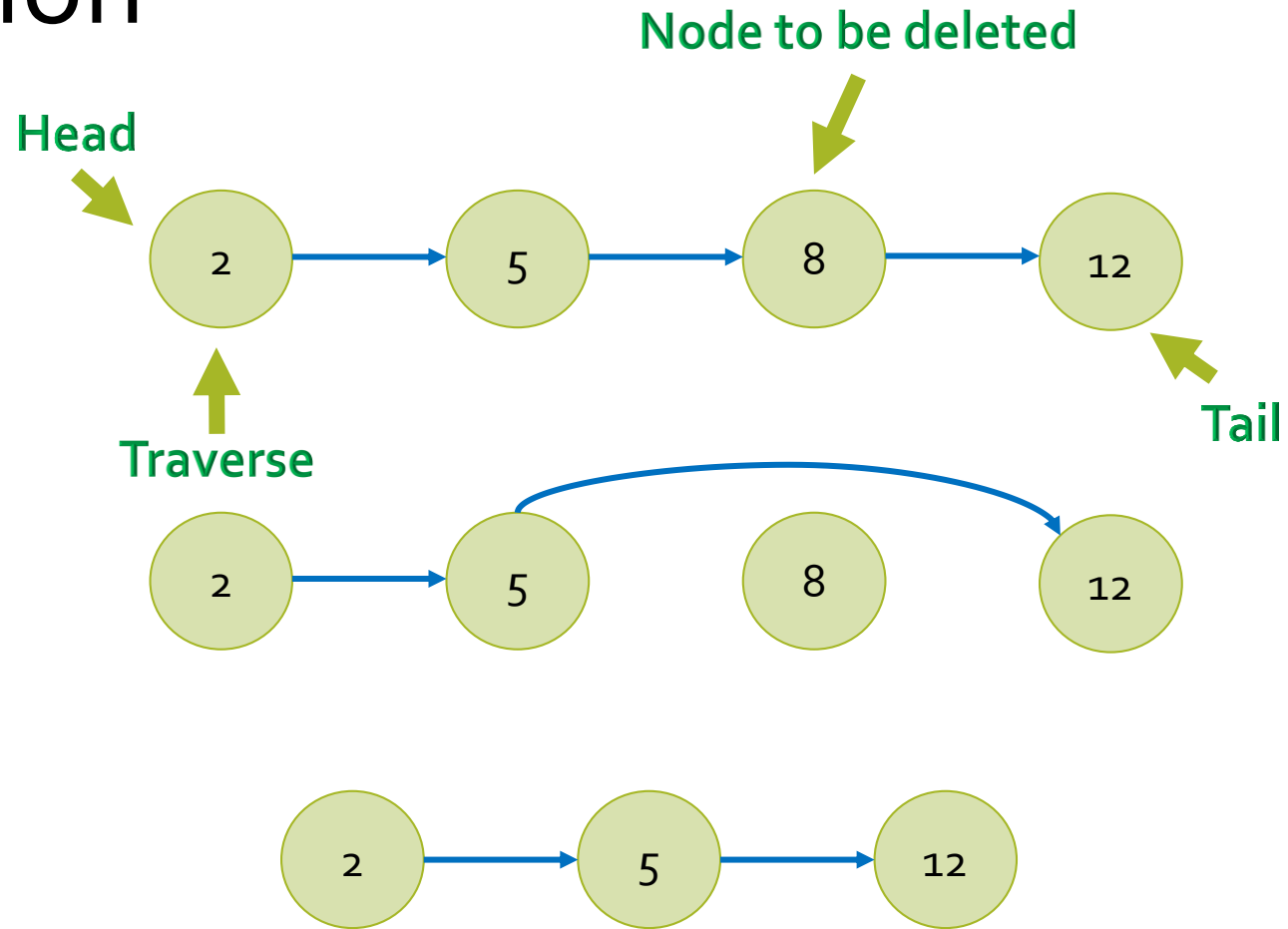
# Insertion



# Insertion at the beginning

1. Create a new node and assign the address to any node say ptr.
2. OVERFLOW,IF(PTR = NULL)  
write : OVERFLOW and EXIT.
3. ASSIGN INFO[PTR] = ITEM
4. IF(START = NULL)  
ASSIGN NEXT[PTR] = NULL  
ELSE  
ASSIGN NEXT[PTR] = START
5. ASSIGN START = PTR
6. EXIT

# Deletion



# Deletion

1. Check whether list is Empty ( $\text{head} == \text{NULL}$ )
2. If it is Empty then, display 'List is Empty!!! Deletion is not possible' and terminate the function.
3. If it is Not Empty then, define a Node pointer 'temp' and initialize with head.
4. Check whether list is having only one node ( $\text{temp} \rightarrow \text{next} == \text{NULL}$ )
5. If it is TRUE then set  $\text{head} = \text{NULL}$  and delete temp (Setting Empty list conditions)
6. If it is FALSE then set  $\text{head} = \text{temp} \rightarrow \text{next}$ , and delete temp.

# Complexity

|        | Singly | Doubly |
|--------|--------|--------|
| Access | $O(n)$ | $O(n)$ |
| Search | $O(n)$ | $O(n)$ |
| Delete | $O(1)$ | $O(1)$ |
| Insert | $O(1)$ | $O(1)$ |

# Applications

- Queue & stack implementation.
- Model real world objects.
- Model repeating event cycles.
- Separate chaining.
- Adjacency list for graphs.