DATA STRUCTURES & ALGORITHMS

Graph

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Introduction

A graph **G** is consist of two things:

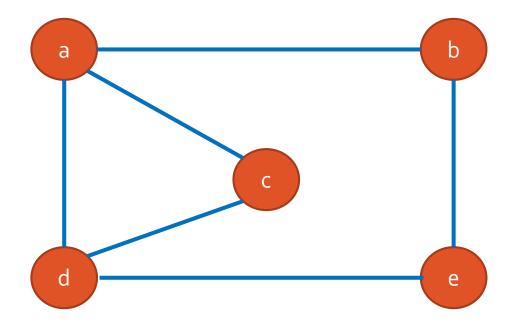
- A set V of elements called nodes/vertices.
- A set of E of edges such that each edge in E is identified with a unique pair {u, v} of nodes in V.

Graph can be indicated as G = (V, E)

Example

 $V = \{a, b, c, d, e\}$

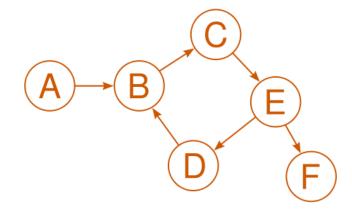
 $E = \{(a,b), (a,c), (a,d), (b,e), (c,d), (c,e), (d,e)\}$

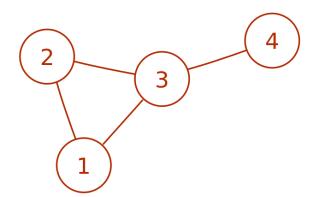


Types

- Directed
 - When edges have direction

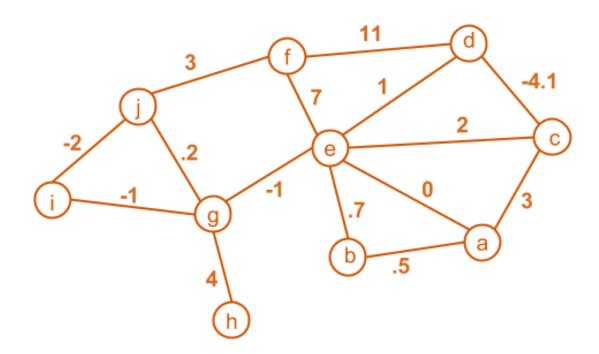
- Undirected
 - When edges have no direction





Weighted Graph

A graph in which each edge contains a value



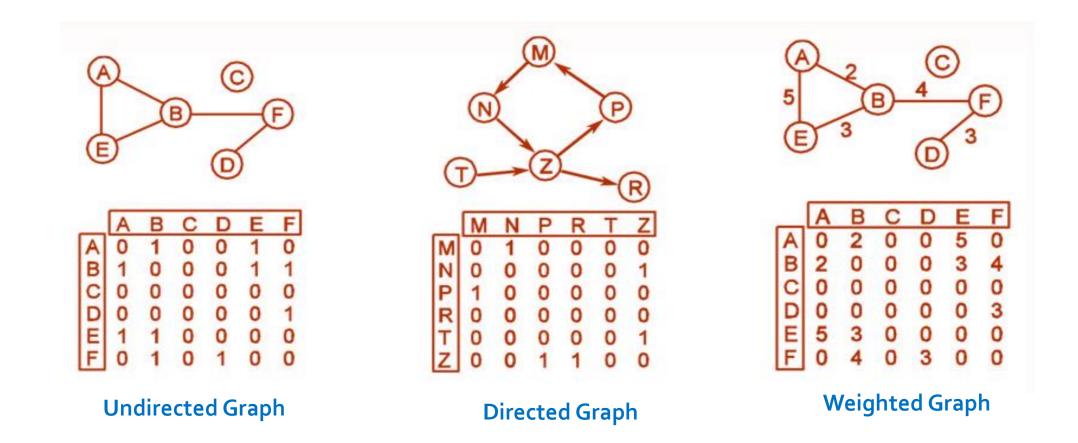
Key Terms

- Adjacent nodes: two nodes are adjacent if they are connected by an edge
- Path: sequence of vertices that connect two nodes in a graph
- Complete graph: in which every vertex is directly connected to every other vertex
- Multi Graph: in which numerous edges between a pair of vertices in a graph G= (V, E) occurs. There are no self-loops.
- Cycle: is a simple path with the same start and end vertex
- Degree of a vertex: number of edges that touch it
- In degree: no of incoming edges
- Out degree: no of outgoing edges

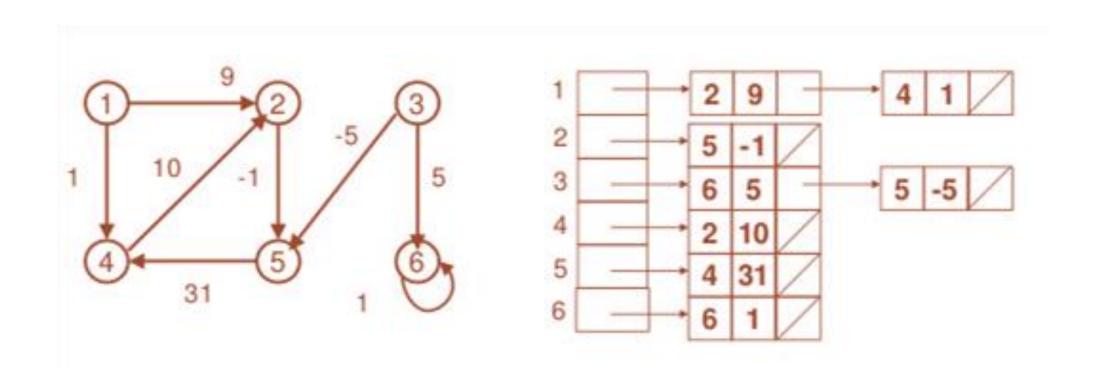
Graph Representation

- Adjacency matrix representation
 - Square grid of Boolean values
 - If the graph contains N vertices then the grid contains N rows N columns
 - For two vertices numbered I and J the element at row I and column J is true if there is an edge from I to J, otherwise false
- Adjacency lists representation
 - A graph of n nodes is represented by a one dimensional array L of linked lists, where
 - L[i] is the linked list containing all the nodes adjacent from node i.
 - The node in the list L [i] are in no particular order

Adjacency Matrix Representation



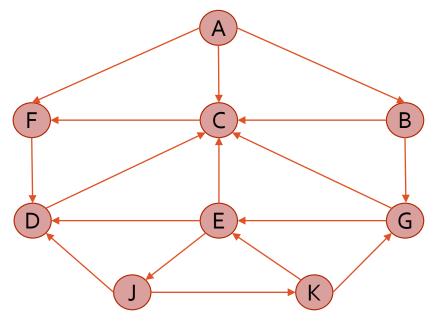
Adjacency Lists Representation



Breadth First Search (BFS) algorithm

Traversing a graph in a breadth ward motion and uses a queue to remember to get the next vertex to start a search, when a dead end occurs in any iteration.

Consider a given graph G. Suppose G represents the daily flights between cities of some airlines and suppose we want to fly from city A to city J with minimum no. of stops. In other words, we want the minimum path p from A to J.



| Adjacency List | | | | |
|----------------|---------|--|--|--|
| Α | F, C, B | | | |
| В | G, C | | | |
| C | F | | | |
| D | C | | | |
| Е | D, C, J | | | |
| F | D | | | |
| G | C, E | | | |
| J | D, K | | | |
| K | E, G | | | |
| | | | | |

Breadth First Search (BFS) algorithm

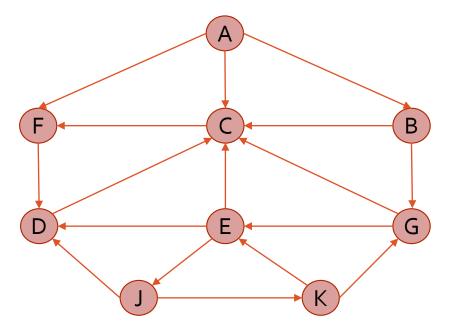
| Adjacency List | | | |
|----------------|---------|--|--|
| Α | F, C, B | | |
| В | G, C | | |
| C | F | | |
| D | С | | |
| Е | D, C, J | | |
| F | D | | |
| G | C, E | | |
| J | D, K | | |
| K | E, G | | |

| Step 1 | | Step 5 | | | | | |
|----------------|----------------------|----------------|--------------------------------|--|--|--|--|
| Fı | Queue A | F 5 | Queue A, F, C, B, D, G | | | | |
| R 1 | Origin Φ | R 6 | Origin Φ, A, A, A, F, B | | | | |
| Step 2 | | Step 6 | | | | | |
| F 2 | Queue A, F, C, B | F 6 | Queue A, F, C, B, D, G | | | | |
| R 4 | Origin Φ, A, A, A | R 6 | Origin Φ, A, A, A, F, B | | | | |
| Step 3 | | Step 7 | | | | | |
| F ₃ | Queue A, F, C, B, D | F ₇ | Queue A, F, C, B, D, G, E | | | | |
| R 5 | Origin Φ, A, A, A, F | R 7 | Origin Φ, A, A, A, F, B, G | | | | |
| Step 4 | | Step 8 | | | | | |
| F ₄ | Queue A, F, C, B, D | F 8 | Queue A, F, C, B, D, G, E, J | | | | |
| R 5 | Origin Φ, A, A, A, F | R 8 | Origin Φ , A, A, A, F, B, G, E | | | | |
| J < E < B < A | | | | | | | |

Depth First Search (DFS) algorithm

Traversing a graph in a depth ward motion and uses a stack to remember to get the next vertex to start a search, when a dead end occurs in any iteration.

Consider a given graph G. Suppose we want to find and print all nodes reachable from node J (including J itself).



Depth First Search (DFS) algorithm

| Adjacency List | | | |
|----------------|---------|--|--|
| Α | F, C, B | | |
| В | G, C | | |
| C | F | | |
| D | С | | |
| Е | D, C, J | | |
| F | D | | |
| G | C, E | | |
| J | D, K | | |
| K | E, G | | |

| Step 1 | | Step 5 | | | | |
|---------------------|---------------|--------|---------------|--|--|--|
| | Stack J | | Print C | | | |
| Step 2 | | | Stack D, E, F | | | |
| | Print J | Step 6 | | | | |
| | Stack D, K | | Print F | | | |
| Step 3 | | | Stack D, E | | | |
| | Print K | Step 7 | | | | |
| | Stack D, E, G | | Print E | | | |
| Step 4 | | | Stack D | | | |
| | Print G | Step 8 | | | | |
| | Stack D, E, C | | Print D | | | |
| | | | Stack | | | |
| J. K. G. C. F. E. D | | | | | | |