



Bahria University
Discovering Knowledge

Computer Architecture and Logic Design (CALD)

Lecture 15

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State Reduction and Assignment

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State Reduction and Assignment

- State Reduction
- Reductions on the number of flip-flops and the number of gates.
 - A reduction in the number of states may result in a reduction in the number of flip-flops.
 - An example state diagram showing in Fig. 5.25.

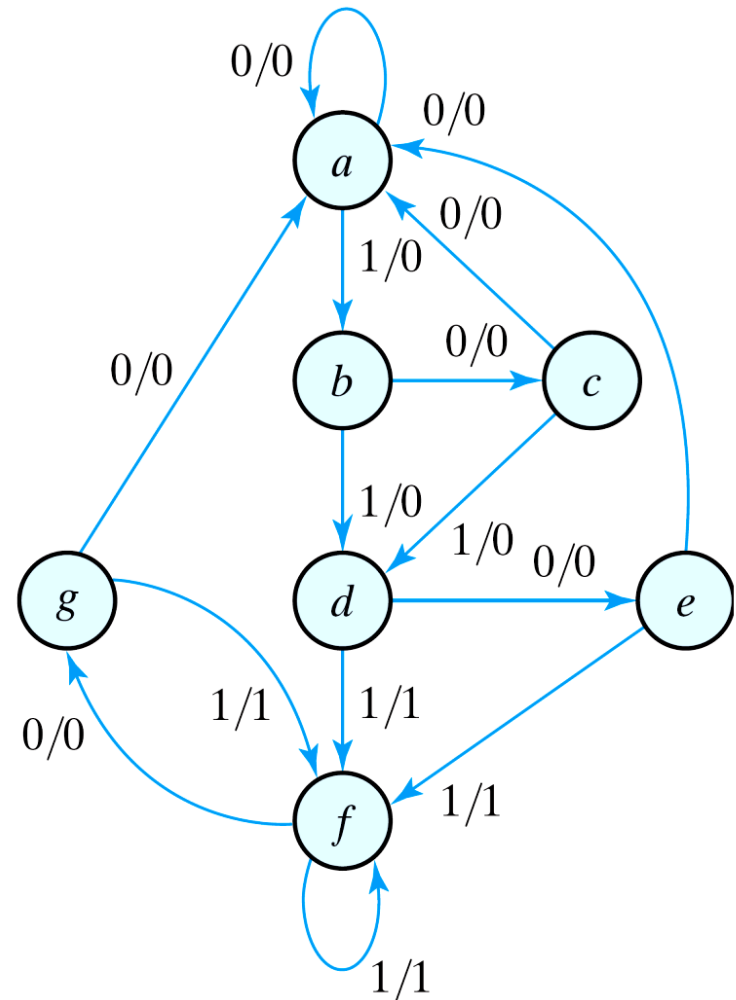


Fig. 5.25 State diagram

State Reduction

State: a a b c d e f f g f g a
Input: 0 1 0 1 0 1 1 0 1 0 0
Output: 0 0 0 0 0 1 1 0 1 0 0

- Only the input-output sequences are important.
- Two circuits are **equivalent**
 - Have identical outputs for all input sequences;
 - The number of states is not important.

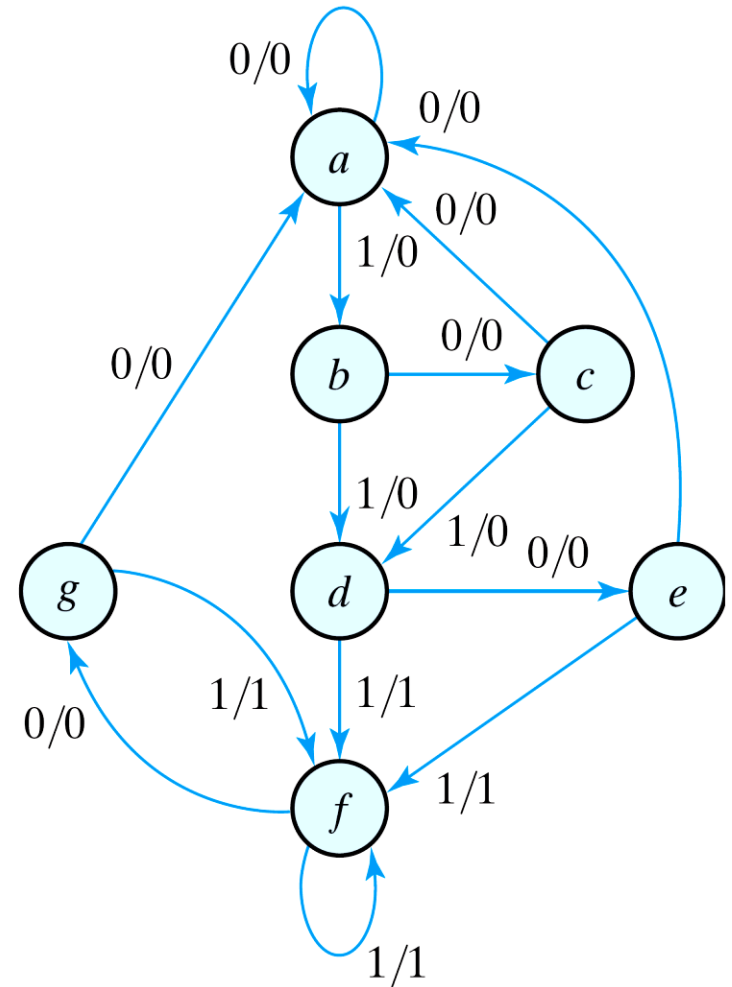


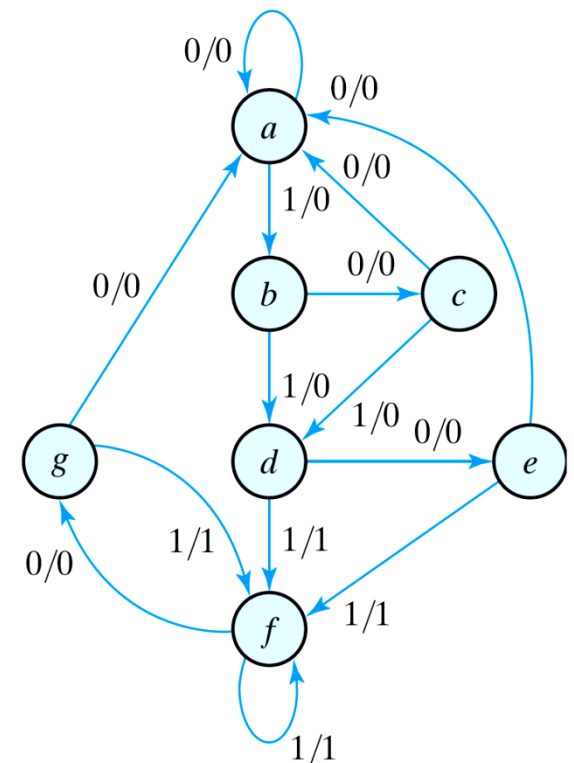
Fig. 5.25 State diagram

■ Equivalent states

- Two states are said to be equivalent
 - For each member of the set of inputs, they give exactly the same output and send the circuit to the same state or to an equivalent state.
 - One of them can be removed.

Table 5.6
State Table

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
<i>a</i>	<i>a</i>	<i>b</i>	0	0
<i>b</i>	<i>c</i>	<i>d</i>	0	0
<i>c</i>	<i>a</i>	<i>d</i>	0	0
<i>d</i>	<i>e</i>	<i>f</i>	0	1
<i>e</i>	<i>a</i>	<i>f</i>	0	1
<i>f</i>	<i>g</i>	<i>f</i>	0	1
<i>g</i>	<i>a</i>	<i>f</i>	0	1



■ Reducing the state table

- $e = g$ (remove g);
- $d = f$ (remove f);

Table 5.7
Reducing the State Table

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
a	a	b	0	0
b	c	d	0	0
c	a	d	0	0
d	e	f	0	1
e	a	f	0	1
f	e	f	0	1

- The reduced finite state machine

Table 5.8
Reduced State Table

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
<i>a</i>	<i>a</i>	<i>b</i>	0	0
<i>b</i>	<i>c</i>	<i>d</i>	0	0
<i>c</i>	<i>a</i>	<i>d</i>	0	0
<i>d</i>	<i>e</i>	<i>d</i>	0	1
<i>e</i>	<i>a</i>	<i>d</i>	0	1

State: a a b c d e d d e d e a

Input: 0 1 0 1 0 1 1 0 1 0 0

Output: 0 0 0 0 0 1 1 0 1 0 0

- The checking of each pair of states for possible equivalence can be done systematically using **Implication Table**.
- The unused states are treated as don't-care condition \Rightarrow fewer combinational gates.

Table 5.8
Reduced State Table

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
a	a	b	0	0
b	c	d	0	0
c	a	d	0	0
d	e	d	0	1
e	a	d	0	1

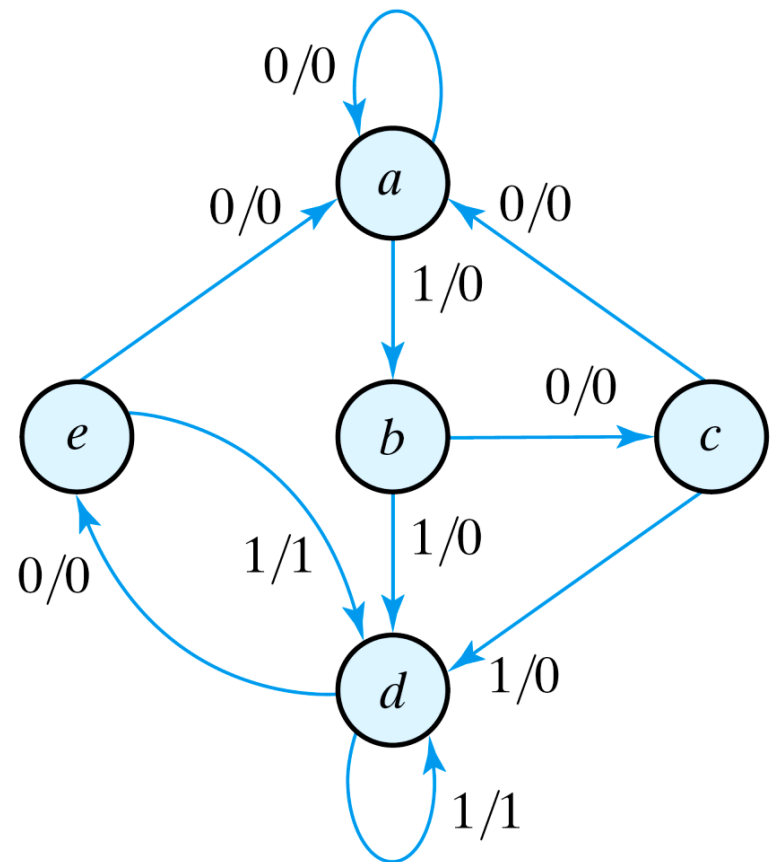


Fig. 5.26 Reduced State diagram

State Assignment

- State Assignment
- To minimize the cost of the combinational circuits.
 - Three possible binary state assignments. (m states need n -bits, where $2^n > m$)

Table 5.9
Three Possible Binary State Assignments

State	Assignment 1, Binary	Assignment 2, Gray Code	Assignment 3, One-Hot
<i>a</i>	000	000	00001
<i>b</i>	001	001	00010
<i>c</i>	010	011	00100
<i>d</i>	011	010	01000
<i>e</i>	100	110	10000

- Any binary number assignment is satisfactory as long as each state is assigned a unique number.
- Use binary assignment 1.

Table 5.10

Reduced State Table with Binary Assignment 1

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
000	000	001	0	0
001	010	011	0	0
010	000	011	0	0
011	100	011	0	1
100	000	011	0	1