Department *of*
Software Engineering
BAHRIA UNIVERSITY
Discovering Knowledge

Lecture 2-Algorithm Analysis

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Algorithm

- An ***algorithm*** is a set of instructions to be followed to solve a problem.
  - There can be more than one solution (more than one algorithm) to solve a given problem.
  - An algorithm can be implemented using different programming languages on different platforms.
- An algorithm must be correct. It should correctly solve the problem.
  - e.g. For sorting, this means even if (1) the input is already sorted, or (2) it contains repeated elements.
- Once we have a correct algorithm for a problem, we have to determine the efficiency of that algorithm.

Algorithmic Performance

There are *two aspects* of algorithmic performance:

- Time
    - Instructions take time.
    - How fast does the algorithm perform?
    - What affects its runtime?
- Space
    - Data structures take space
    - What kind of data structures can be used?
    - How does choice of data structure affect the runtime?

➢ We will focus on time:

- How to estimate the time required for an algorithm
- How to reduce the time required

Analysis of Algorithms

- ***Analysis of Algorithms*** is the area of computer science that provides tools to analyze the efficiency of different methods of solutions.
- How do we compare the time efficiency of two algorithms that solve the same problem?

  Naïve Approach: implement these algorithms in a programming language (C++), and run them to compare their time requirements. Comparing the programs (instead of algorithms) has difficulties.
  - *How are the algorithms coded?*
    - Comparing running times means comparing the implementations.
    - We should not compare implementations, because they are sensitive to programming style that may cloud the issue of which algorithm is inherently more efficient.
  - *What computer should we use?*
    - We should compare the efficiency of the algorithms independently of a particular computer.
  - *What data should the program use?*
    - Any analysis must be independent of specific data.

Analysis of Algorithms

- When we analyze algorithms, we should employ mathematical techniques that analyze algorithms independently of *specific implementations, computers, or data.*

- To analyze algorithms:
  - First, we start to count the number of significant operations in a particular solution to assess its efficiency.
  - Then, we will express the efficiency of algorithms using growth functions.

The Execution Time of Algorithms

- Each operation in an algorithm (or a program) has a cost.

  ➔ Each operation takes a certain of time.

`count = count + 1;` ➔ take a certain amount of time, but it is constant

A sequence of operations:

`count = count + 1;`	Cost: c_1
`sum = sum + count;`	Cost: c_2

➔ Total Cost = $c_1 + c_2$

The Execution Time of Algorithms (cont.)

Example: Simple If-Statement

	Cost	**Times**
`if (n < 0)`	c1	1
`absval = -n`	c2	1
`else`		
`absval = n;`	c3	1

Total Cost <= c1 + max(c2,c3)

The Execution Time of Algorithms (cont.)

Example: Simple Loop

Code	Cost	Times
`i = 1;`	c1	1
`sum = 0;`	c2	1
`while (i <= n) {`	c3	n+1
`    i = i + 1;`	c4	n
`    sum = sum + i;`	c5	n
`}`		

Total Cost = c1 + c2 + (n+1)*c3 + n*c4 + n*c5

➔ The time required for this algorithm is proportional to n

The Execution Time of Algorithms (cont.)

Example: Nested Loop

Code	Cost	Times
`i=1;`	c1	1
`sum = 0;`	c2	1
`while (i <= n) {`	c3	n+1
`    j=1;`	c4	n
`    while (j <= n) {`	c5	n*(n+1)
`        sum = sum + i;`	c6	n*n
`        j = j + 1;`	c7	n*n
`    }`		
`    i = i +1;`	c8	n
`}`		

Total Cost = c1 + c2 + (n+1)*c3 + n*c4 + n*(n+1)*c5+n*n*c6+n*n*c7+n*c8

➔ The time required for this algorithm is proportional to n^2

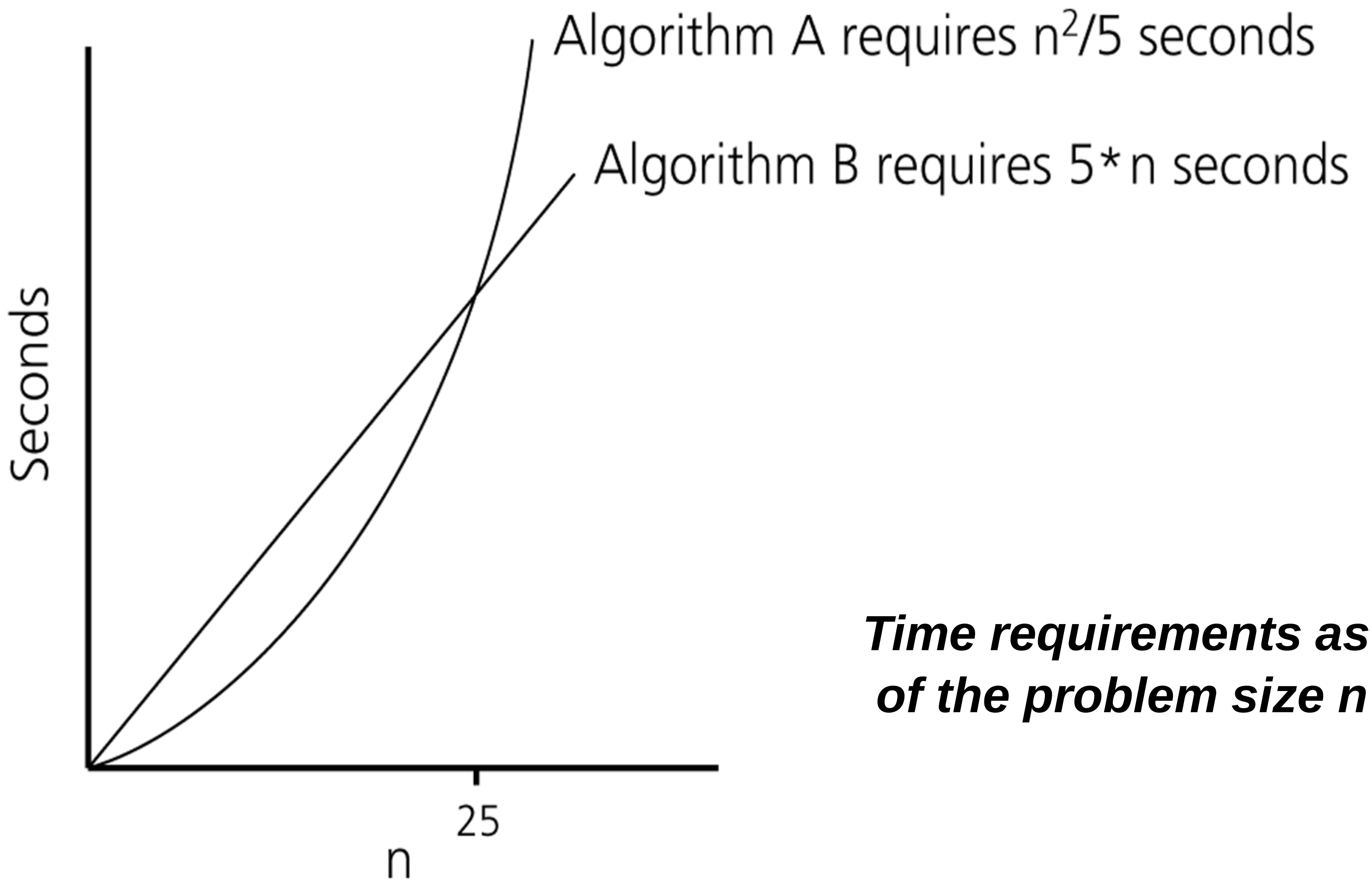
General Rules for Estimation

- **Loops**: The running time of a loop is at most the running time of the statements inside of that loop times the number of iterations.
- **Nested Loops**: Running time of a nested loop containing a statement in the inner most loop is the running time of statement multiplied by the product of the sized of all loops.
- **Consecutive Statements:** Just add the running times of those consecutive statements.
- **If/Else**: Never more than the running time of the test plus the larger of running times of S1 and S2.

Algorithm Growth Rates

- We measure an algorithm's time requirement as a function of the *problem size*.
  - Problem size depends on the application: e.g. number of elements in a list for a sorting algorithm, the number disks for towers of hanoi.
- So, for instance, we say that (if the problem size is n)
  - Algorithm A requires **$5*n^2$** time units to solve a problem of size n.
  - Algorithm B requires **$7*n$** time units to solve a problem of size n.
- The most important thing to learn is how quickly the algorithm's time requirement grows as a function of the problem size.
  - Algorithm A requires time proportional to **n^2**.
  - Algorithm B requires time proportional to **n**.
- An algorithm's proportional time requirement is known as ***growth rate***.
- We can compare the efficiency of two algorithms by comparing their growth rates.

Algorithm Growth Rates (cont.)



Time requirements as a function of the problem size n

Common Growth Rates

Function	Growth Rate Name
c	Constant
$\log N$	Logarithmic
$\log^2 N$	Log-squared
N	Linear
$N \log N$	
N^2	Quadratic
N^3	Cubic
2^N	Exponential

Figure 6.1

Running times for small inputs

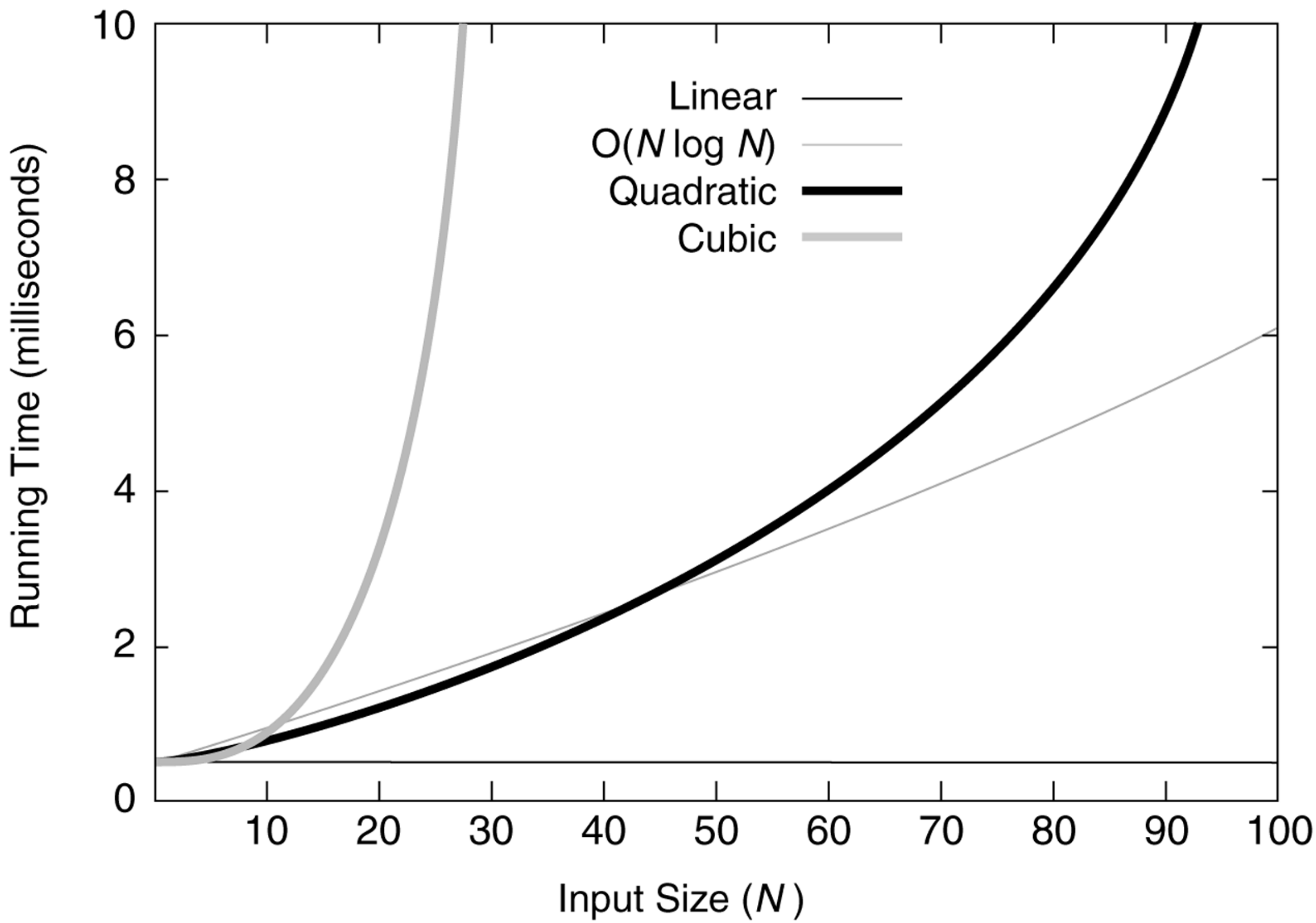
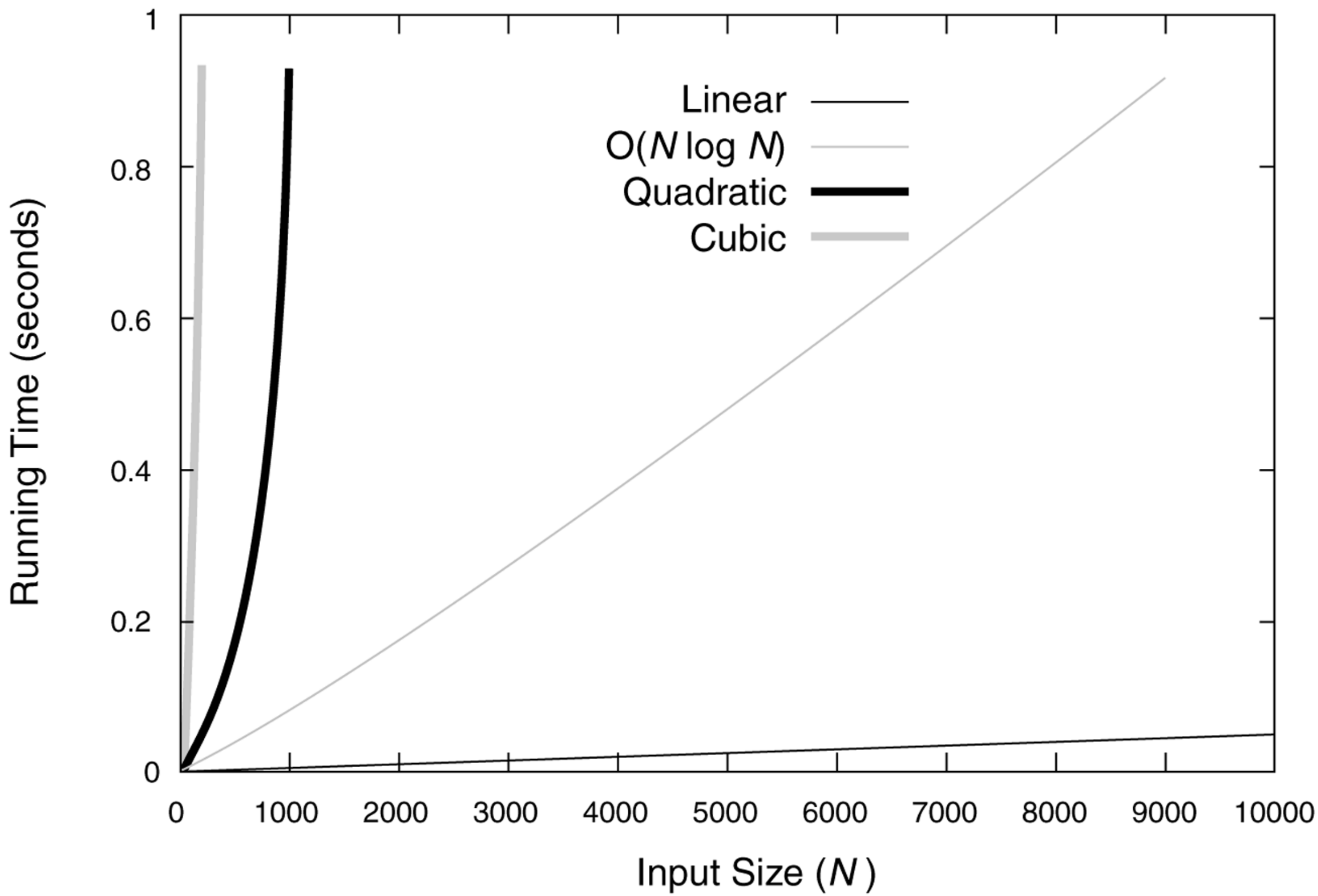


Figure 6.2

Running times for moderate inputs



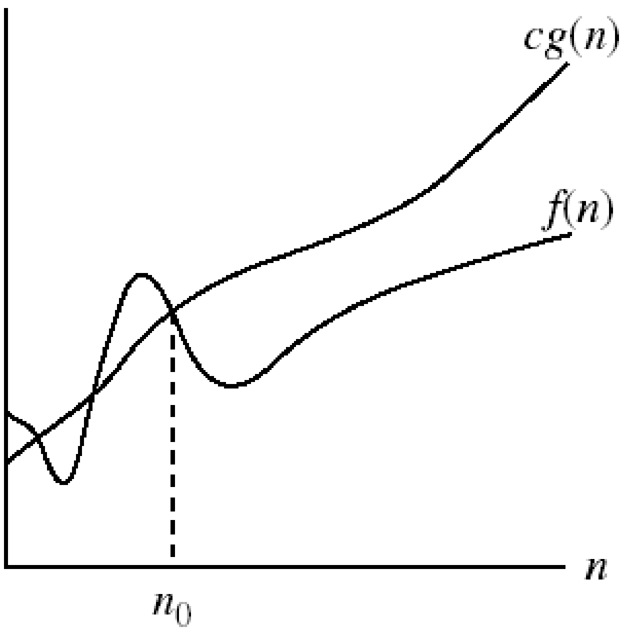
Asymptotic Notation

- O notation :Big-O is the formal method of expressing the upper bound of an algorithm's running time.
- It's a measure of the longest amount of time it could possibly take for the algorithm to complete.
- Formally, for non-negative functions, $f(n)$ and $g(n)$, if there exists an integer n_0 and a constant $c > 0$ such that for all integers $n > n_0$, $f(n) \leq cg(n)$, then $f(n)$ is Big O of $g(n)$.

Asymptotic notations

- *O-notation*

$$O(g(n)) = \{f(n) : \text{ there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \le f(n) \le cg(n) \text{ for all } n \ge n_0\} .$$



$g(n)$ is an ***asymptotic upper bound*** for $f(n)$.

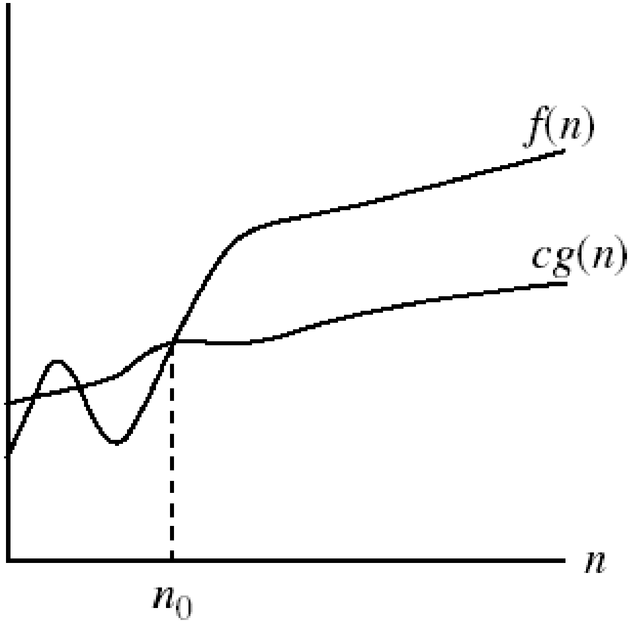
Asymptotic Notation

- ***Big-Omega Notation*** *Ω*
- This is almost the same definition as Big Oh, except that "$f(n) \geq cg(n)$"
- This makes *g(n)* a lower bound function, instead of an upper bound function.
- It describes the **best that can happen** for a given data size.

- For non-negative functions, *f(n)* and *g(n)*, if there exists an integer n_0 and a constant $c > 0$ such that for all integers $n > n_0$, $f(n) \geq cg(n)$, then *f(n)* is omega of *g(n)*. This is denoted as "$f(n) = \Omega(g(n))$".

Asymptotic notations (cont.)

$$\Omega(g(n)) = \{f(n) : \text{ there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq cg(n) \leq f(n) \text{ for all } n \geq n_0\} .$$



Ω(*g*(n)) is the set of functions with larger or same order of growth as *g*(n)

$g(n)$ is an ***asymptotic lower bound*** for $f(n)$.

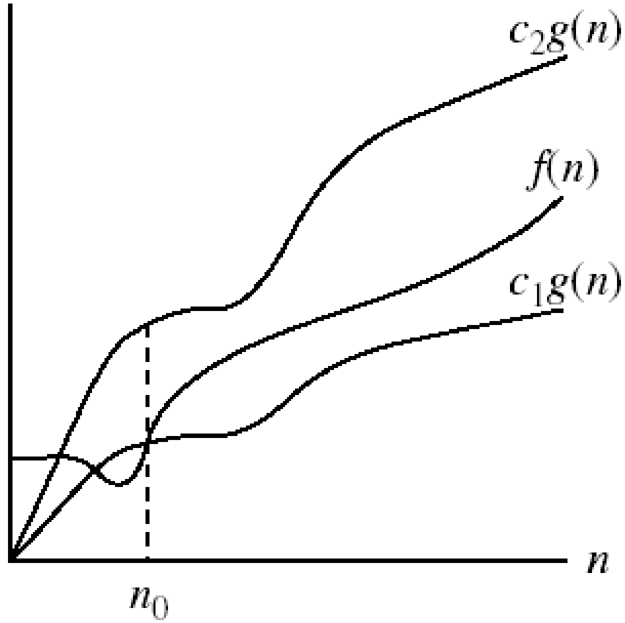
Asymptotic Notation

- **Theta Notation** *Θ*
- Theta Notation For non-negative functions, $f(n)$ and $g(n)$, $f(n)$ is theta of $g(n)$ if and only if $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$. This is denoted as "$f(n) = \Theta(g(n))$".

  This is basically saying that the function, $f(n)$ is bounded both from the top and bottom by the same function, $g(n)$.

Asymptotic notations (cont.)

$$\Theta(g(n)) = \{f(n) : \text{there exist positive constants } c_1, c_2, \text{ and } n_0 \text{ such that } 0 \le c_1 g(n) \le f(n) \le c_2 g(n) \text{ for all } n \ge n_0\} .$$



$\Theta(g(n))$ is the set of functions with the same order of growth as $g(n)$

$g(n)$ is an ***asymptotically tight bound*** for $f(n)$.

Order-of-Magnitude Analysis and Big O Notation

- If *Algorithm A requires time proportional to f(n),* Algorithm A is said to be **order f(n),** and it is denoted as **O(f(n)).**
- The **function f(n)** is called the algorithm's **growth-rate function**.
- Since the capital O is used in the notation, this notation is called the **Big O notation**.
- If Algorithm A requires time proportional to $\mathbf{n^2}$, it is $\mathbf{O(n^2)}$**.**
- If Algorithm A requires time proportional to **n**, it is **O(n).**

Definition of the Order of an Algorithm

Definition:

Algorithm A is order f(n) – denoted as O(f(n)) –

if constants k and n_0 exist such that A requires

no more than k*f(n) time units to solve a problem

of size $n \geq n_0$.

- The requirement of $\mathbf{n \geq n_0}$ in the definition of O(f(n)) formalizes the notion of sufficiently large problems.
    - In general, many values of k and n can satisfy this definition.

Order of an Algorithm

- If an algorithm requires $n^2 – 3*n+10$ seconds to solve a problem size n. If constants k and n_0 exist such that

  $k*n^2 > n^2 – 3*n+10$ for all $n \geq n_0$.

  the algorithm is order n^2 (In fact, k is 3 and n_0 is 2)

  $3*n^2 > n^2 – 3*n+10$ for all $n \geq 2$.

  Thus, the algorithm requires no more than $k*n^2$ time units for $n \geq n_0$,

  So it is **O(n^2)**

Growth-Rate Functions

O(1) Time requirement is **constant**, and it is independent of the problem's size.

O($\log_2 n$) Time requirement for a **logarithmic** algorithm increases increases slowly as the problem size increases.

O(n) Time requirement for a **linear** algorithm increases directly with the size of the problem.

O($n*\log_2 n$) Time requirement for a **$n*\log_2 n$** algorithm increases more rapidly than a linear algorithm.

O(n^2) Time requirement for a **quadratic** algorithm increases rapidly with the size of the problem.

O(n^3) Time requirement for a c**ubic** algorithm increases more rapidly with the size of the problem than the time requirement for a quadratic algorithm.

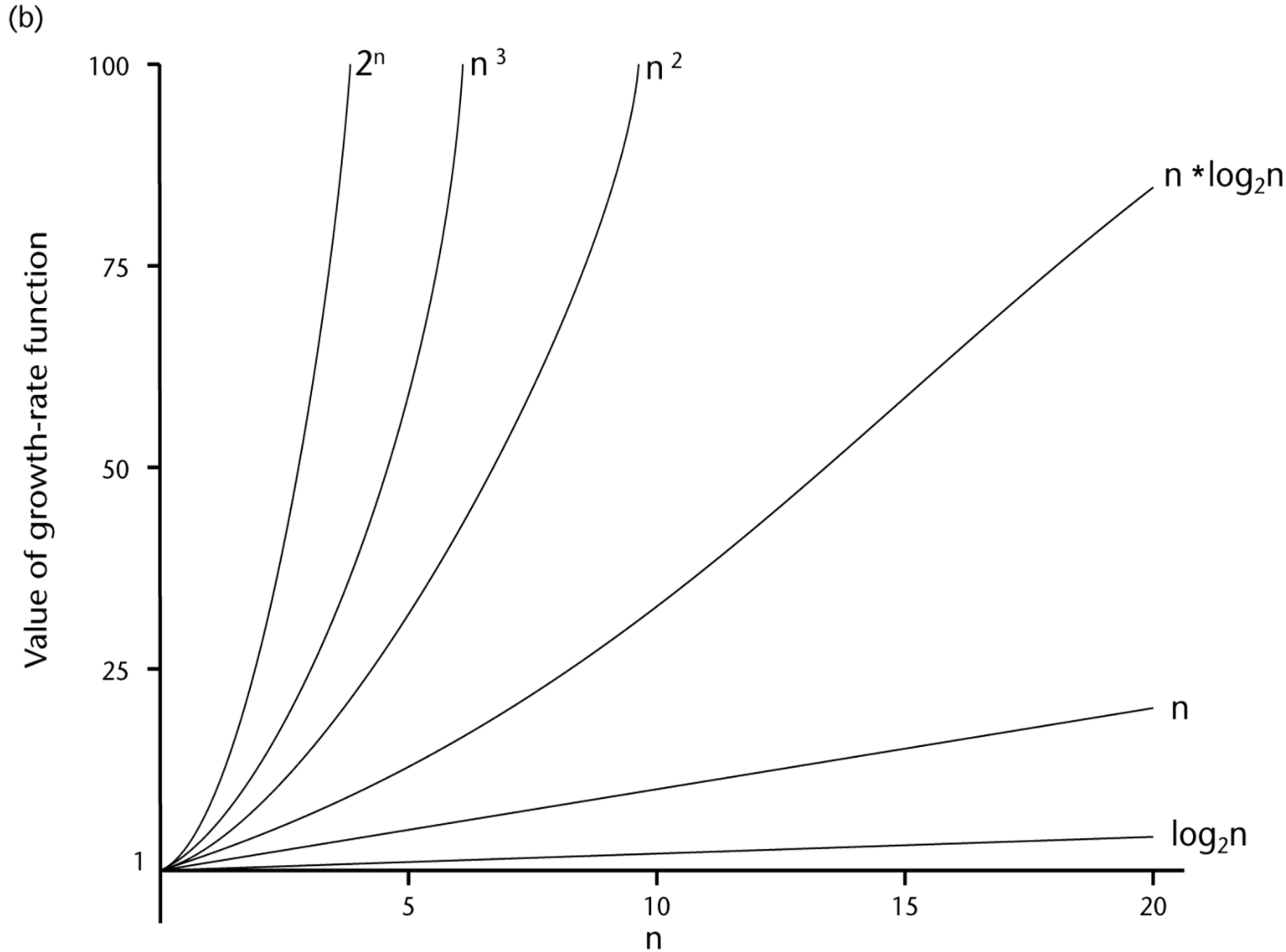
O(2^n) As the size of the problem increases, the time requirement for an **exponential** algorithm increases too rapidly to be practical.

A Comparison of Growth-Rate Functions

(a)

Function	n = 10	100	1,000	10,000	100,000	1,000,000
1	1	1	1	1	1	1
$\log_2 n$	3	6	9	13	16	19
n	10	10^2	10^3	10^4	10^5	10^6
$n * \log_2 n$	30	664	9,965	10^5	10^6	10^7
n^2	10^2	10^4	10^6	10^8	10^{10}	10^{12}
n^3	10^3	10^6	10^9	10^{12}	10^{15}	10^{18}
2^n	10^3	10^{30}	10^{301}	$10^{3,010}$	$10^{30,103}$	$10^{301,030}$

A Comparison of Growth-Rate Functions (cont.)



Growth-Rate Functions

- If an algorithm takes 1 second to run with the problem size 8, what is the time requirement (approximately) for that algorithm with the problem size 16?
- If its order is:

O(1) ➔ T(n) = 1 second

O($\log_2 n$) ➔ T(n) = (1*$\log_2 16$) / $\log_2 8$ = 4/3 seconds

O(n) ➔ T(n) = (1*16) / 8 = 2 seconds

O(n*$\log_2 n$) ➔ T(n) = (1*16*$\log_2 16$) / 8*$\log_2 8$ = 8/3 seconds

O(n^2) ➔ T(n) = (1*16^2) / 8^2 = 4 seconds

O(n^3) ➔ T(n) = (1*16^3) / 8^3 = 8 seconds

O(2^n) ➔ T(n) = (1*2^{16}) / 2^8 = 2^8 seconds = 256 seconds

How much better is $O(log_2 n)$?

n	***$O(log_2 n)$***
16	4
64	6
256	8
1024 (1KB)	10
16,384	14
131,072	17
262,144	18
524,288	19
1,048,576 (1MB)	20
1,073,741,824 (1GB)	30

Properties of Growth-Rate Functions

1. *We can ignore low-order terms in an algorithm's growth-rate function.*
    - If an algorithm is $O(n^3+4n^2+3n)$, it is also $O(n^3)$.
    - We only use the higher-order term as algorithm's growth-rate function.

2. *We can ignore a multiplicative constant in the higher-order term of an algorithm's growth-rate function.*
    - If an algorithm is $O(5n^3)$, it is also $O(n^3)$.

3. *$O(f(n)) + O(g(n)) = O(f(n)+g(n))$*
    - We can combine growth-rate functions.
    - If an algorithm is $O(n^3) + O(4n)$, it is also $O(n^3 +4n^2)$ ➔ So, it is $O(n^3)$.
    - Similar rules hold for multiplication.

Some Mathematical Facts

- Some mathematical equalities are:

$$\sum_{i=1}^{n} 1 = 1 + 1 + \cdots .. + 1 = n$$

$$\sum_{i=1}^{n} i = 1 + 2 + \ldots + n = \frac{n*(n+1)}{2} \approx \frac{n^2}{2}$$

$$\sum_{i=1}^{n} i^2 = 1 + 4 + \ldots + n^2 = \frac{n*(n+1)*(2n+1)}{6} \approx \frac{n^3}{3}$$

$$\sum_{i=1}^{n} i^3 = 1 + 8 + \cdots .. + n^3 = \left(\frac{n(n+1)}{2}\right)^2$$

$$\sum_{i=0}^{n-1} 2^i = 0 + 1 + 2 + \ldots + 2^{n-1} = 2^n - 1$$

Growth-Rate Functions – Example1

	Cost	Times
`i = 1;`	c1	1
`sum = 0;`	c2	1
`while (i <= n) {`	c3	n+1
`    i = i + 1;`	c4	n
`    sum = sum + i;`	c5	n
`}`		

T(n) = c1 + c2 + (n+1)*c3 + n*c4 + n*c5

= (c3+c4+c5)*n + (c1+c2+c3)

= a*n + b

➔ So, the growth-rate function for this algorithm is **O(n)**

Growth-Rate Functions – Example2

	Cost	**Times**
`i=1;`	c1	1
`sum = 0;`	c2	1
`while (i <= n) {`	c3	n+1
`    j=1;`	c4	n
`    while (j <= n) {`	c5	n*(n+1)
`        sum = sum + i;`	c6	n*n
`        j = j + 1;`	c7	n*n
`    }`		
`    i = i +1;`	c8	n
`}`		

$$
\begin{aligned}
T(n) &= c1 + c2 + (n+1)*c3 + n*c4 + n*(n+1)*c5+n*n*c6+n*n*c7+n*c8 \\
&= (c5+c6+c7)*n^2 + (c3+c4+c5+c8)*n + (c1+c2+c3) \\
&= a*n^2 + b*n + c
\end{aligned}
$$

➔ So, the growth-rate function for this algorithm is **$O(n^2)$**

Growth-Rate Functions – Example3

	Cost	**Times**
`for (i=1; i<=n; i++)`	c1	n+1
`for (j=1; j<=i; j++)`	c2	$\sum_{j=1}^{n}(j+1)$
`for (k=1; k<=j; k++)`	c3	$\sum_{j=1}^{n}\sum_{k=1}^{j}(k+1)$
`x=x+1;`	c4	$\sum_{j=1}^{n}\sum_{k=1}^{j}k$

T(n) $= c1*(n+1) + c2*\left(\sum_{j=1}^{n}(j+1)\right) + c3*\left(\sum_{j=1}^{n}\sum_{k=1}^{j}(k+1)\right) + c4*\left(\sum_{j=1}^{n}\sum_{k=1}^{j}k\right)$

$= c1n + c1 + c2n^2 + c2n + c3n^3 + (c3 + c4)2n^2$

$= (c1 + c2)n + (c2 + 2c3 + 2c4)n^2 + c3n^3$

$= a*n^3 + b*n^2 + c*n + d$

➔ So, the growth-rate function for this algorithm is **$O(n^3)$**

Sequential Search

```
int sequentialSearch(const int a[], int item, int n){
  for (i = 0; i < n && a[i]!= item; i++);
  if (i == n)
       return –1;
  return i;
}
```

Unsuccessful Search: ➔ O(n)

Successful Search:

Best-Case: *item* is in the first location of the array ➔O(1)

Worst-Case: *item* is in the last location of the array ➔O(n)

Average-Case: The number of key comparisons 1, 2, ..., n

$$\frac{\sum_{i=1}^{n} i}{n} = \frac{(n^2+n)/2}{n}$$ ➔ O(n)

Binary Search-Iterative

```
int  binarySearch(int a[], int size, int x) {
  int low =0;
  int high = size –1;
  int mid;     // mid will be the index of
                        // target when it’s found.
  while (low <= high) {
    mid = (low + high)/2;
    if (a[mid] == x)
     return mid;
    else if (a[mid] > x)
          high  = mid – 1;
    else
          low = mid + 1;
  }
  return –1;
}
```

Binary Search-Recursive

```
int binarySearch(int []a, int lo, int hi, int x)
{
   int mid= (lo+hi)/2;
   if(a[mid]==x) return mid;
   else if(a[mid]>x)
       return binarySearch(a,lo,mid-1);
   else return binarySearch(a mid+1,hi);
}
```

Binary Search – Analysis

$T(n)=T(n/2)+c$ $T(1)=c$

$$T(\frac{n}{2}) = T(\frac{\frac{n}{2}}{2}) + c$$
$$= T(\frac{n}{2^2}) + c$$

$T(n) = T\left(\frac{n}{2^2}\right)+2c$ $T(n/2^2)=T(n/2^3)+c$

$T(n)=T(n/2^3)+3c$

.

.

.

$T(n)=T(n/2^k)+kc$ suppose $n=2^k \Rightarrow \log(n)=k$

$T(n)=T(n/n)+c\log(n)$

$T(n)=T(1)+c\log(n)$

$T(n)=c+c\log(n)$

$T(n)=c(\log(n)+1)$

$T(n)=O(\log(n))$

Binary Search – Analysis

- For an unsuccessful search:
  - The number of iterations in the loop is $\lfloor \log_2 n \rfloor + 1$

    ➔ $O(\log_2 n)$
- For a successful search:
  - ***Best-Case:*** The number of iterations is 1. ➔ $O(1)$
  - ***Worst-Case:*** The number of iterations is $\lfloor \log_2 n \rfloor + 1$ ➔ $O(\log_2 n)$
  - ***Average-Case:*** The avg. # of iterations < $\log_2 n$ ➔ $O(\log_2 n)$

```
0   1   2   3   4   5   6   7
3   2   3   1   3   2   3   4
```
← an array with size 8

← # of iterations

The average # of iterations = 21/8 < $\log_2 8$

Growth-Rate Functions – Recursive Algorithms

```
void hanoi(int n, char source, char dest, char spare) {  Cost
  if (n > 0) {                                            c1
  hanoi(n-1, source, spare, dest);                        c2
    cout << "Move top disk from pole " << source          c3
       << " to pole " << dest << endl;
  hanoi(n-1, spare, dest, source);                        c4
} }
```

- The time-complexity function T(n) of a recursive algorithm is defined in terms of itself, and this is known as **recurrence equation** for T(n).
- To find the growth-rate function for a recursive algorithm, we have to solve its recurrence relation.

Growth-Rate Functions – Hanoi Towers

- What is the cost of `hanoi(n,'A','B','C')`?

when n=0

$T(0) = c1$

when n>0

$T(n) = c1 + c2 + T(n-1) + c3 + c4 + T(n-1)$

$= 2*T(n-1) + (c1+c2+c3+c4)$

$= \mathbf{2*T(n-1) + c}$ ← recurrence equation for the growth-rate function of hanoi-towers algorithm

- Now, we have to solve this recurrence equation to find the growth-rate function of hanoi-towers algorithm

Growth-Rate Functions – Hanoi Towers (cont.)

- There are many methods to solve recurrence equations, but we will use a simple method known as *repeated substitutions*.

$T(n) = 2*T(n-1) + c$ $\quad$ $T(n-1)= 2*T(n-2)+c$

$\quad = 2 * (2*T(n-2)+c) + c$

$\quad = 2^2*T(n-2)+2c+c$ $\quad$ $T(n-2)= 2*T(n-3)+c$

$\quad = 2^2*(2*T(n-3)+c)+2^1c+2^0c$

$\quad = 2^3 * T(n-3) + (2^2+2^1+2^0)*c$ $\quad$ (assuming n>2)

when substitution repeated k-1[th] times

$\quad = 2^k * T(n-k) + (2^{k-1}+ ... +2^1+2^0)*c$

When k=n

$\quad = 2^n * T(0) + (2^{n-1}+ ... +2^1+2^0)*c$

$$= 2^n * c1 + (\sum_{i=0}^{n-1} 2^i)*c$$

$\quad = 2^n * c1 + (2^n-1)*c = 2^n*(c1+c) – c$ $\quad$ ➔ So, the growth rate function is **$O(2^n)$**

What to Analyze

- An algorithm can require different times to solve different problems of the same size.
  - Eg. Searching an item in a list of n elements using sequential search. ➔ Cost: 1,2,...,n
- ***Worst-Case Analysis*** –The maximum amount of time that an algorithm require to solve a problem of size n.
  - This gives an upper bound for the time complexity of an algorithm.
  - Normally, we try to find worst-case behavior of an algorithm.
- ***Best-Case Analysis*** –The minimum amount of time that an algorithm require to solve a problem of size n.
  - The best case behavior of an algorithm is NOT so useful.
- ***Average-Case Analysis*** –The average amount of time that an algorithm require to solve a problem of size n.
  - Sometimes, it is difficult to find the average-case behavior of an algorithm.
  - We have to look at all possible data organizations of a given size n, and their distribution probabilities of these organizations.
  - ***Worst-case analysis is more common than average-case analysis.***

What is Important?

- An array-based list `retrieve` operation is O(1), a linked-list-based list `retrieve` operation is O(n).
- But insert and delete operations are much easier on a linked-list-based list implementation.

  ➔ When selecting the implementation of an Abstract Data Type (ADT), we have to consider how frequently particular ADT operations occur in a given application.

- If the problem size is always small, we can probably ignore the algorithm's efficiency.
  - In this case, we should choose the simplest algorithm.

What is Important? (cont.)

- We have to weigh the trade-offs between an algorithm’s time requirement and its memory requirements.
- We have to compare algorithms for both style and efficiency.
  - The analysis should focus on gross differences in efficiency and not reward coding tricks that save small amount of time.
  - That is, there is no need for coding tricks if the gain is not too much.
  - Easily understandable program is also important.
- Order-of-magnitude analysis focuses on large problems.