

Figure 1.6 IAS Structure

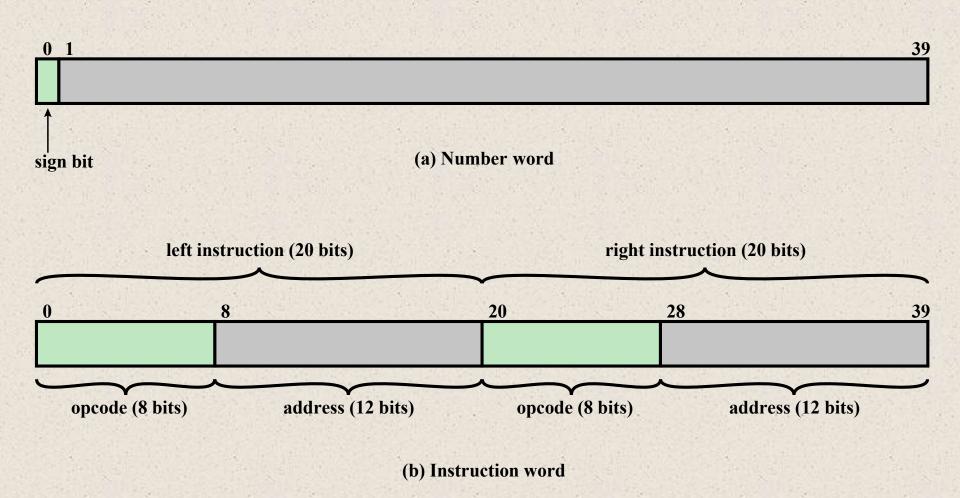


Figure 1.7 IAS Memory Formats

Registers

Memory buffer register (MBR)

- · Contains a word to be stored in memory or sent to the I/O unit
- Or is used to receive a word from memory or from the I/O unit

Memory address register (MAR)

 Specifies the address in memory of the word to be written from or read into the MBR

Instruction register (IR)

Contains the 8-bit opcode instruction being executed

Instruction buffer register (IBR)

 Employed to temporarily hold the right-hand instruction from a word in memory

Program counter (PC)

 Contains the address of the next instruction pair to be fetched from memory

Accumulator (AC) and multiplier quotient (MQ)

Employed to temporarily hold operands and results of ALU operations

Instruction Type Opcode Representation Description O0001010 LOAD MQ Transfer contents of register MQ to the accumulator AC O0001001 LOAD MQ,M(X) Transfer contents of memory location X MQ O0100001 STOR M(X) Transfer contents of accumulator to memory location X Transfer contents of accumulator to memory location X MQ O0100001 LOAD M(X) Transfer M(X) to the accumulator X Transfer M(X) to the accumulator Transfer absolute value of M(X) to the accumulator Transfer absolute value of M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) Transfer - M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) to the accumulator Transfer - M(X) Transfer - M(X) to the accumulator Transfer - M(X) to the	nory I(X) M(X)
$\begin{array}{c} MQ \\ 00100001 & STOR\ M(X) & Transfer\ contents\ of\ accumulator\ to\ meriocation\ X \\ 00000001 & LOAD\ M(X) & Transfer\ M(X)\ to\ the\ accumulator \\ 00000010 & LOAD\ -M(X) & Transfer\ -M(X)\ to\ the\ accumulator \\ 00000011 & LOAD\ M(X) & Transfer\ absolute\ value\ of\ M(X)\ to\ the\ accumulator \\ 00000100 & LOAD\ - M(X) & Transfer\ - M(X) \ to\ the\ accumulator \\ \hline Unconditional & 00001101 & JUMP\ M(X,0:19) & Take\ next\ instruction\ from\ right\ half\ of\ M(X) \\ \hline Data\ transfer & M(X)\ to\ the\ accumulator \\ \hline Unconditional & 00001101 & JUMP\ M(X,0:19) & Take\ next\ instruction\ from\ right\ half\ of\ M(X) \\ \hline Data\ transfer & M(X)\ to\ the\ accumulator \\ \hline Unconditional & 00001101 & JUMP\ M(X,0:19) & Take\ next\ instruction\ from\ right\ half\ of\ M(X) \\ \hline Data\ transfer & M(X)\ to\ the\ accumulator \\ \hline Unconditional & 00001111 & JUMP\ M(X,0:19) & If\ number\ in\ the\ accumulator\ is\ nonneg$	nory I(X) M(X)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	((X) M(X)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	M(X)
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	M(X)
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	M(X)
Unconditional 00001101 JUMP M(X,0:19) Take next instruction from left half of M branch 00001110 JUMP M(X,20:39) Take next instruction from right half of 00001111 JUMP+ M(X,0:19) If number in the accumulator is nonneg	M(X)
branch 00001110 JUMP M(X,20:39) Take next instruction from right half of 00001111 JUMP+ M(X,0:19) If number in the accumulator is nonneg	M(X)
00001111 JUMP+ M(X,0:19) If number in the accumulator is nonneg	
	tive
	uve,
take next instruction from left half of M	(X)
JU If number in the	
MP accumulator is nonnego	
Conditional branch + take next instruction from	m
M(X) right half of $M(X)$, 20:	
39)	
00000101 ADD M(X) Add M(X) to AC; put the result in AC	
1 (7)	. A.C
00001000 SUB $ M(X) $ Subtract $ M(X) $ from AC; put the remain in AC	
Arithmetic 00001011 MUL M(X) Multiply M(X) by MQ; put most significant bits of result in AC, put least significant in MQ	bits
00001100 DIV M(X) Divide AC by M(X); put the quotient in and the remainder in AC	MQ
00010100 LSH Multiply accumulator by 2; i.e., shift let bit position	one
00010101 RSH Divide accumulator by 2; i.e., shift righ position	
00010010 STOR M(X,8:19) Replace left address field at M(X) by 12 rightmost bits of AC	
Address modify 00010011 STOR M(X,28:39) Replace right address field at M(X) by rightmost bits of AC	2

Table 1.1

The IAS Instruction Set

(Table can be found on page 17 in the textbook.)

History of Computers Second Generation: Transistors

- Smaller
- Cheaper
- Dissipates less heat than a vacuum tube
- Is a *solid state device* made from silicon
- Was invented at Bell Labs in 1947
- It was not until the late 1950's that fully transistorized computers were commercially available

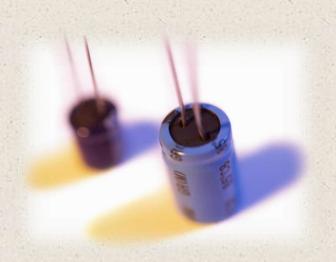


Table 1.2 Computer Generations

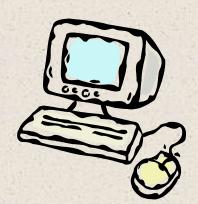
Generation	Approximate Dates	Technology	Typical Speed (operations per second)
1	1946–1957	Vacuum tube	40,000
2	1957–1964	Transistor	200,000
3	1965–1971	Small and medium scale integration	1,000,000
4	1972–1977	Large scale integration	10,000,000
5	1978–1991	Very large scale integration	100,000,000
6	1991-	Ultra large scale integration	>1,000,000,000

+

Second Generation Computers

■Introduced:

- More complex arithmetic and logic units and control units
- The use of high-level programming languages
- Provision of system software which provided the ability to:
 - Load programs
 - Move data to peripherals
 - Libraries perform common computations



History of Computers Third Generation: Integrated Circuits

- 1958 the invention of the integrated circuit
- Discrete component
 - Single, self-contained transistor
 - Manufactured separately, packaged in their own containers, and soldered or wired together onto masonite-like circuit boards
 - Manufacturing process was expensive and cumbersome
- The two most important members of the third generation were the IBM System/360 and the DEC PDP-8



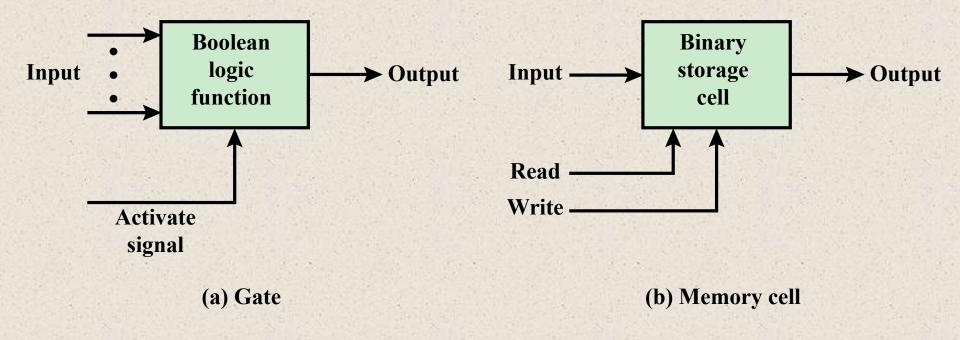


Figure 1.10 Fundamental Computer Elements

Integrated Circuits

- Data storage provided by memory cells
- Data processing provided by gates
- Data movement the paths among components are used to move data from memory to memory and from memory through gates to memory
- Control the paths among components can carry control signals

- A computer consists of gates, memory cells, and interconnections among these elements
- The gates and memory cells are constructed of simple digital electronic components
- Exploits the fact that such components as transistors, resistors, and conductors can be fabricated from a semiconductor such as silicon
- Many transistors can be produced at the same time on a single wafer of silicon
- Transistors can be connected with a processor metallization to form circuits

IBM System/360

- Announced in 1964
- Product line was incompatible with older IBM machines
- Was the success of the decade and cemented IBM as the overwhelmingly dominant computer vendor
- The architecture remains to this day the architecture of IBM's mainframe computers
- Was the industry's first planned family of computers
 - Models were compatible in the sense that a program written for one model should be capable of being executed by another model in the series

+ Family Characteristics

Similar or identical instruction set

Similar or identical operating system

Increasing speed

Increasing number of I/O ports

Increasing memory size

Increasing cost