

## Lecture 8:- Principle of Resolution.

literal:- A variable or its negation.  $P \quad \neg P$

clause:- A disjunction of literals.  $\Rightarrow \left\{ \begin{array}{l} P \vee Q \quad \neg P \vee \neg Q \\ P \vee \neg Q \\ P \vee \neg Q \\ P \vee \neg Q \vee R \end{array} \right.$

" Given a set of clauses a deduction of  $C$  from  $S$  is a finite sequence  $C_1, C_2, \dots, C_k$  of clauses.

A deduction of  $\square$  empty clause is called refutation or a proof of  $S$ .

$$\begin{array}{ll} p_1 & \rightarrow C_1 \\ p_2 & \rightarrow C_2 \\ \vdots & \\ p_n & \rightarrow C_k \\ \hline \therefore C & \neg C_{k+1} \end{array}$$

$$\begin{array}{l} P \vee Q \\ \neg P \vee R \\ \hline \therefore P \vee R \end{array}$$

Ex P65.

$$\begin{array}{ll} p_1 & P \\ p_2 & P \rightarrow Q \\ \hline C & \therefore Q \end{array}$$

$$\begin{array}{ll} (C_1) & P \quad \checkmark \\ (C_2) & \neg P \vee Q \quad \checkmark \\ (C_3) & \neg Q \quad \checkmark \\ \hline & P \vee Q. \\ & \neg P \vee R. \\ & \hline & Q \vee R. \end{array}$$

from  $(C_1, C_3) \quad Q \rightarrow (4) \checkmark$

from  $(4, C_3) \quad \square$

$$P \rightarrow Q = \neg P \vee Q.$$

Ex 11-P65.

$$\begin{array}{ll} p_1 & T \rightarrow M \vee E \\ p_2 & S \rightarrow \neg E \\ p_3 & T \wedge S. \\ \hline \end{array}$$

$$\begin{array}{ll} C_1 & \rightarrow \neg T \vee M \vee E. \checkmark \\ C_2 & \rightarrow \neg S \vee \neg E \checkmark \\ C_3 & \rightarrow T \checkmark \\ \hline \end{array}$$

$$p3 \quad \underline{T \wedge S.}$$

$$C. \quad \therefore M.$$

$$C3 \rightarrow T \checkmark$$

$$C4 \rightarrow S \checkmark$$

$$C5 \rightarrow \neg M \checkmark$$

$$\text{from } C1, C2 \quad \neg T \vee M \vee \neg S. \text{ --- } \checkmark_6$$

$$\text{from } C3, 6. \quad M \vee \neg S. \text{ --- } \checkmark_7$$

$$\text{from } C4, 7 \quad M \text{ --- } \checkmark_8$$

$$\text{from } C5, 8 \quad \square$$

Ex6 P-62.

$$P1 \quad \neg P \wedge q$$

$$P2 \quad r \rightarrow P$$

$$P3 \quad \neg r \rightarrow S$$

$$P4 \quad S \rightarrow t$$

$$C. \quad \therefore t.$$

$$C1 \quad \neg P \checkmark$$

$$C2. \quad q$$

$$C3 \quad \neg r \vee P \checkmark$$

$$C4 \quad r \vee S \checkmark$$

$$C5 \quad \neg S \vee t \checkmark$$

$$C6 \quad \neg t \checkmark$$

$$\text{from } C1, C3 \quad \neg r \quad (7) \checkmark$$

$$\text{from } C4, 7. \quad S \quad (8) \checkmark$$

$$\text{from } C5, 8 \quad t \quad (9) \checkmark$$

$$\text{from } C6, 9 \quad \square \quad (10)$$

Ex7: P 63:-

$$P1 \quad P \rightarrow q$$

$$P2 \quad \neg P \rightarrow r$$

$$P3 \quad \underline{r \rightarrow S}$$

$$C. \quad \therefore \neg q \rightarrow S.$$

$$C1 \quad \neg P \vee q \checkmark$$

$$C2 \quad P \vee r \checkmark$$

$$C3 \quad \neg r \vee S \checkmark$$

$$C4 \quad \neg q \checkmark$$

$$C5 \quad \neg S \checkmark$$

$$\text{from } C1, C2 \quad q \vee r \text{ --- } 6 \checkmark$$

From C1, C2  $q, \vee$  — 6 ✓  
 From C3, 6  $q, \vee$  — 7 ✓  
 From C4, 7  $s$  — 8 ✓  
 From C5, 8  $\square$

Ex 9 P 65.

P1  $L \rightarrow A$

P2  $E \rightarrow \neg I$

P3  $A \rightarrow E$

C.  $\therefore L \rightarrow \neg I$

C1  $\neg L \vee A$  ✓

C2  $\neg E \vee \neg I$  ✓

C3  $\neg A \vee E$  ✓

C4  $L$  ✓

C5  $I$  ✓

From C1, C3  $\neg L \vee E$  — 6 ✓

From C2, 6  $\neg I \vee \neg L$  — 7 ✓

From C4, 7  $\neg I$  — 8 ✓

From C5, 8  $\square$