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Assignment #08

S - AES

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BS-CS - 7B.

## ~~Q~~ Simplified AES example.

Plaintext  $\Rightarrow$  1101 0111 0010 1000

16 bit key  $\Rightarrow$  0100 1010 1111 0101

### (1) Key Generation

(Sol)

The input key, is split into two words  $w_0, w_1$

$$w_0 \Rightarrow 01001010$$

$$w_1 = 11110101$$

Now we have to find  $w_2$

$$w_2 = w_0 \oplus RCON_1 \oplus \text{SubNib}(\text{RotNib}(w_1))$$

~~Now~~ Now  $RCON_1 \Rightarrow 10000000$

$$w_2 \Rightarrow 01001010 \oplus 10000000 \oplus \text{SubNib}(01011111)$$

$$\text{Now, SubNib}(0101) \Rightarrow 0001$$

$$\text{SubNib}(1111) \Rightarrow 0111$$

$$w_2 \Rightarrow 11001010 \oplus 00010111$$

$$w_2 \Rightarrow 11011101$$



~~$w_2 = 1101 1001$~~

~~$w_2 = 1101 1001$~~

$$w_3 = w_2 \oplus w_1$$

$$w_3 = 1101 \ 1101 \oplus 1110 \ 1011$$

$$w_3 = 0010 \ 1000$$

$$w_4 = w_2 \oplus RCON_2 \oplus \text{SubNib}(\text{RotNib}(w_3))$$

$$w_4 \Rightarrow 1101 \ 1101 \oplus 0011 \ 0000 \oplus \text{SubNib}(\text{RotNib}(0010 \ 1000))$$

$$w_4 \Rightarrow 1110 \ 1101 \oplus \text{SubNib}(1000 \ 0010)$$

$$\text{SubNib}(1000) \Rightarrow 0110$$

$$\text{SubNib}(0010) \Rightarrow 1010$$

$$w_4 = 1110 \ 1101 + 0110 \ 1010$$

$$w_4 \Rightarrow 1000 \ 0111$$

$$w_5 \Rightarrow w_4 \oplus w_3$$

$$w_5 = 1000 \ 0111 \oplus 0010 \ 1000$$

$$w_5 \Rightarrow 1010 \ 1111$$

Now, Key0  $\Rightarrow$  wow1

Key0 = 01001010 1111 0101

Key1 = wow3

Key1  $\Rightarrow$  1101 1101 0010 1000

Key2  $\Rightarrow$  w4w5

Key2 = 1000 0111 1010 1111



# S-AES Encryption!

① Add Round 0 Key.

plaintext  $\oplus$  Key

1101 0111 0010 1000  $\oplus$  0100 1010 1111 0101

$\Rightarrow$  1001 1101 1101 1101.

② Round 1 (Nibble Substitution)

$S(1001) \Rightarrow 0016$

$S(1101) \Rightarrow 1110$

$S(1101) \Rightarrow 1110$

$S(1101) \Rightarrow 1110$

$\Rightarrow$  0016 1110 1110 1110

Apply Shift Row, i.e. "Swap and nibble"  
~~nibble~~ with "4th nibble".

$\Rightarrow$  0016 1110 1110 1110.

Mix Columns!

$$M_c = \begin{bmatrix} 1 & 4 \\ 4 & 1 \end{bmatrix}, \Rightarrow \begin{bmatrix} 0010 & 1110 \\ 1110 & 1110 \end{bmatrix}$$

$$S = \begin{bmatrix} 0010 & 1110 \\ 1110 & 1110 \end{bmatrix} \Rightarrow \begin{matrix} S_{00} & S_{01} \\ S_{10} & S_{11} \end{matrix}$$

$$S = M_c \times S$$

~~$$S_{00} = 0010 \oplus (4 \times)$$~~

$$S_{00} \Rightarrow 0010 \times 1 \oplus (1110 \times 4)$$

$$1110 \times 4 \Rightarrow (x^3 + x^2 + x) \times x^2$$

$$\Rightarrow x^5 + x^4 + x^3$$

To be reduced modulo  $(x^4 + x + 1)$

$$\begin{array}{r} x^4 + x + 1 \overline{) x^5 + x^4 + x^3} \\ \underline{+ x^5 + x^4 + x} \phantom{+ 1} \\ x^3 + x^2 + x \phantom{+ 1} \\ \underline{+ x^4 + x + 1} \\ x^3 + x^2 + 1 \end{array}$$

~~$$x^3 + x^2 + 1$$~~  

$$x^3 + x^2 + 1 \Rightarrow 1101$$



$$S_{00} \Rightarrow 0010 \oplus 1101$$

$$S_{00} \Rightarrow 1111$$

$$S_{01} \Rightarrow (1 \times 1110) \oplus (4 \times 1110)$$

$$S_{01} = 1110 \oplus 0011$$

$$S_{01} = 0011$$

$$S_{10} \Rightarrow (4 \times 0010) \oplus (1 \times 1110)$$

$$S_{10} = 1000 \oplus 1110$$

$$S_{10} = 0110$$

~~$$S_{01}$$~~  

$$S_{11} = (4 \times 1110) \oplus (1 \times 1110 \times 1)$$

$$S_{11} = 1101 \oplus 1110$$

$$S_{11} = 0011$$

$$\text{Output} \Rightarrow S_{00} \quad S_{10} \quad S_{01} \quad S_{11}$$

$$1111 \quad 0110 \quad 0011 \quad 0011$$

Now ~~&~~ XOR with key 1.

$$\begin{array}{r} \text{Q} \quad 1111 \ 0110 \ 0011 \ 0011 \quad (+) \\ \text{Q} \quad 1101 \ 1101 \ 0010 \ 1000 \\ \hline \end{array}$$

0010    1011    0001    1011

Final Round Nibble Substitution.

$$\begin{aligned} S(0010) &= 1010 \\ S(1011) &= 0011 \\ S(0001) &= 0100 \\ S(1011) &= 0011 \end{aligned}$$

~~Apply Shift Row (2nd and 4th)~~

~~1010    0100    0011    0011~~

1010    0011    0100    0011

Apply Shift key. (2nd and 4th)

1010    0100    0011    0011

Now, XOR with key 2.



Key  $\oplus$  ~~00~~ 1010 0011 0100 0011

1000 0111 1010 1111  $\oplus$   
1010 0011 0100 0011

00100010011101100

Cipher text  $\Rightarrow$  0010 0100 1110 1100