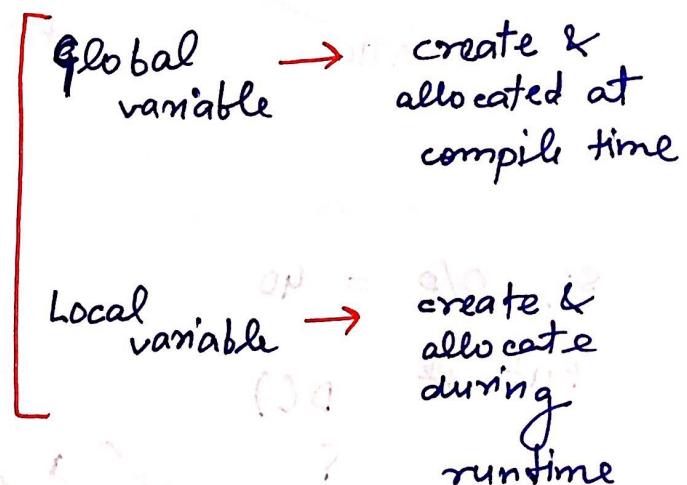
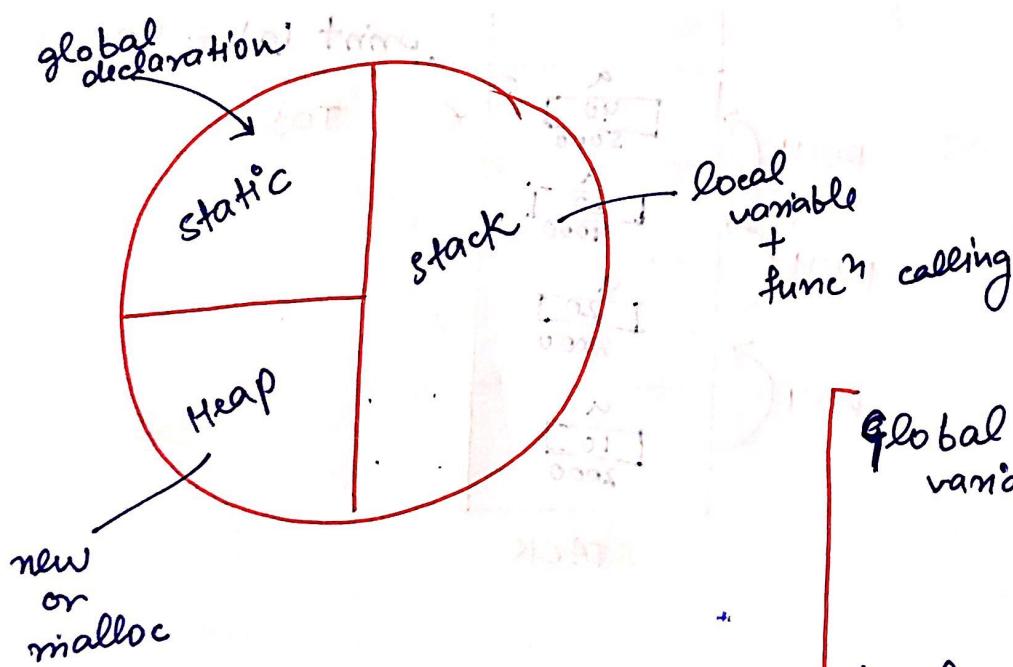
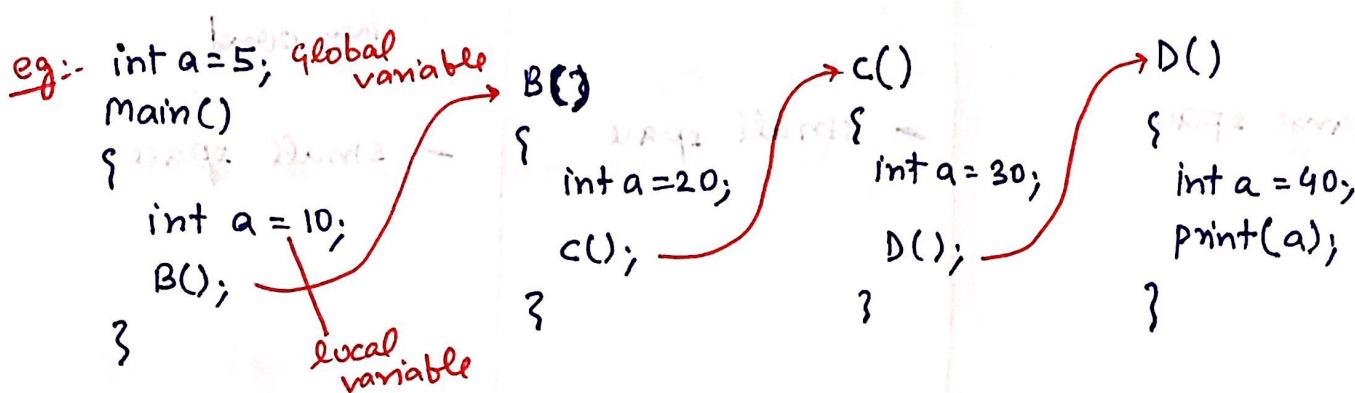


⇒ Programming :-

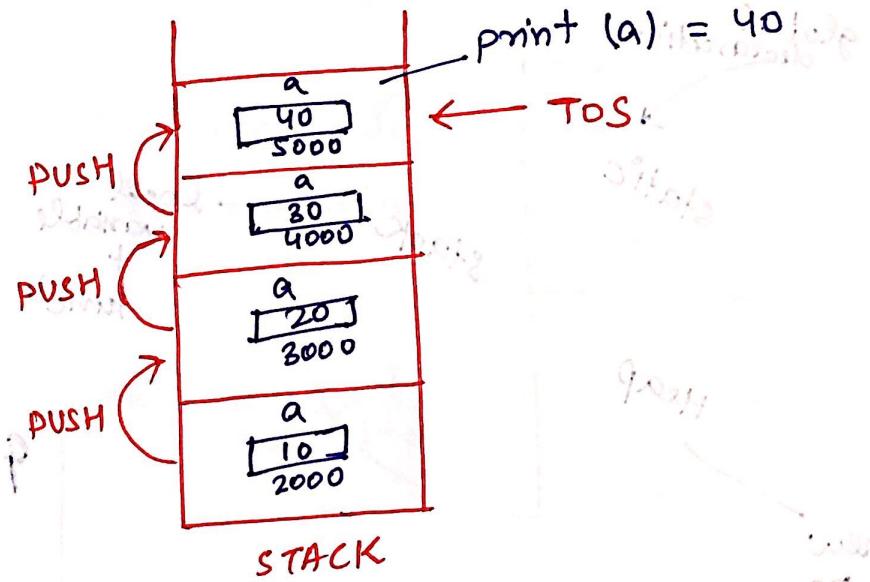
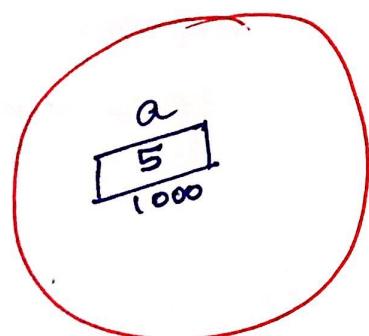
1. Scope of a variable -

(i) Static scoping → resolved by compiler

(ii) Dynamic scoping → resolved by processor



<u>STACK</u>	<u>STATIC</u>	<u>HEAP</u>
- Default value = garbage value	- Default value = 0	- variables allocated using "new" or "malloc" keyword
- Local variable + function calling.	- Global variables	
- Large space	- small space	- small space



so, O/P = 40

But if D()

{ free variable
print(a); // no declaration for a

then scoping problem is occurred.

- If scoping problem is handled by compiler at compile time then it is known as Static scoping!
Compiler checks for static variable ^{in global scope} [and return its value if exist.]
- If scoping problem is handled by processor at runtime then it performs POP opⁿ to find the variable, this is known as Dynamic scoping.

- So, Now O/p is:

static scoping = 5
Dynamic Scoping = 30 from c()

eg:-

```

int a;
main()
{
    int a=10;
    B()
}
c()
print(a);
}
    
```

Ques: static scoping: 0 (Default value is 0)
Dynamic scoping: 10

eg:- int a;
 main() {
 } B() {
 } c() {
 } {
 } } } }

Result :-
 static variable is not shared by both
 B() and c().
 B() has its own local variable a.
 c() has its own local variable a.

O/p: static : 0
 dynamic : 0

Note :-
 static variable will be initialized at compile time and
 dynamic variable will be initialized at run time.

eg:- int a;
 main() {
 } B() {
 } c() {
 } {
 } } } }

Result :-
 static variable is shared by both
 B() and c().
 B() has its own local variable a.
 c() has its own local variable a.

O/p static : 0
 dynamic : garbage value

eg:- main() {
 } B() {
 } c() {
 } {
 } } } }

Result :-
 B() has its own local variable a.
 c() has its own local variable a.

O/p : 10 (No scoping problem)

eg:- int a = 5, b = 10;

main()

```
{
    int a = 20;
    int b = 35;
    c();
    print(a, b);
}
```

D(a)

(~~MAIN ← main → D(a)~~)

(~~D(a) ← D(b) ← E()~~)

C()

```

int a = 7
print(a, b);
D(b);
a = 2, b = 3;

```

int a
D()

```

{
    print(a, b);
    a = 3, b = 7;
    E();
    a = 9, b = 10;
}

```

E()

```

{
    print(a, b);
    a = 11, b = 12;
}

```

STATIC

a = 7, b = 10

a = 10, b = 10

a = 5, b = 7

a = 20, b = 35

a = 20, b = 3

a = 11, b = 7

DYNAMIC

a = 7, b = 35

a = 35, b = 35

a = 35, b = 7

a = 20, b = 3

a = 20, b = 3

a = 3, b = 7

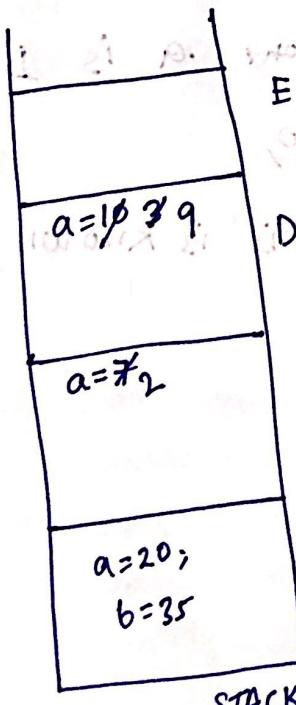
a = 5

b = 10

for STATIC

values

stack



MAIN()

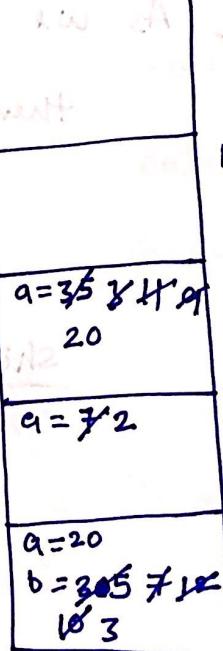
STACK

E()

D()

C()

MAIN()



eg:- main()

{

int i=-1; j=-1; k=-1; l=2, m;

$m = ((i++ \&& j++ \&& k++) \parallel (l++));$

print(i,j,k,l,m);

}

(Non-zero \rightarrow TRUE)
zero \rightarrow FALSE

SIMPLIFICATION
 $m = (\neg (True \& True \& True) \parallel (True))$

or d

(m = 1)

21=d, 22=a

F=d, so, F

E=d, so, E

S=d, so, S

R=d, so, R

i=0
j=0
k=0
l=2
m=1

This condition
is not
checked

so l++
will not
execute

$l=2$

As we know in $a \text{ OR } b$ if a is 1

then no need to calculate b so,

$l++$ will not execute, This is known as

short-circuit

Static Variable :-

```

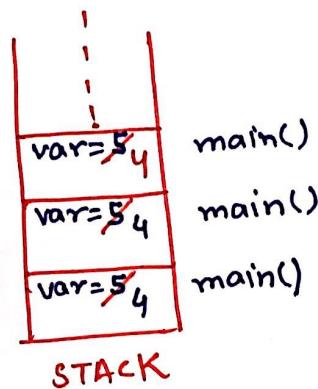
    Shows how static works? . static always prints
main()           main()
    static int var=5;           int var=5;
    if(--var)                 if(--var)
    {                         {
        main();               main();
        print(var);           print(var);
    }
}

```

"var" is a local variable,
but by using static keyword
we can force compiler to
create memory in static
area at compile time.

var=5
4 3 2 1
0
STATIC

var--	main()
var--	main()
var=5	main()
var--	main()
var--	main()



O/p:- STACK overflow

Every time main() called,
new local variable created
with var=5 , so termin-
ation condition i.e var=0
will never reached .

So, PUSH op^h will continue
infinitely until stack
overflow happen.

POP	main() →
POP	main() → 0

O/p :- 0000

NOTE:- Memory creates only once for static variable during compile time. Processor will not create extra stack variable during run time. It skips memory creation by seeing static token.

- For static variable initialization will be done only once.
- Default value for static variable = 0

e.g. main()

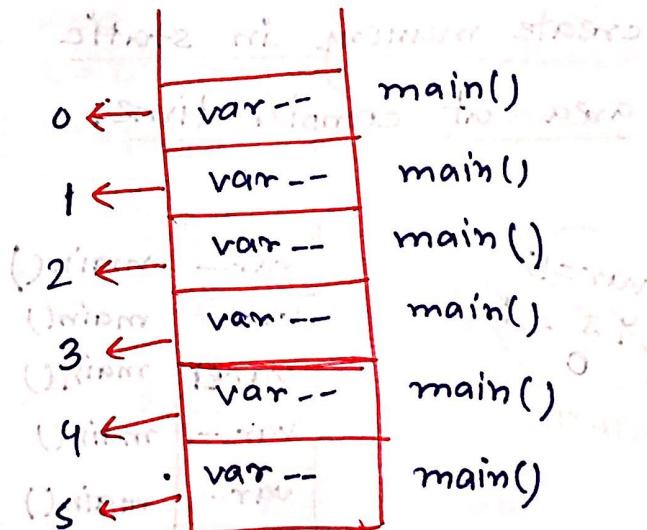
```

    {
        static int var = 5;
        if(var--)
            main();
        printf(var);
    }
}

```

main()	push	main()
main()	print()	pop → -1
main()	print()	pop → -1
main()	print()	pop → -1
main()	print()	pop → -1
main()	print()	pop → -1
main()	print()	pop → -1

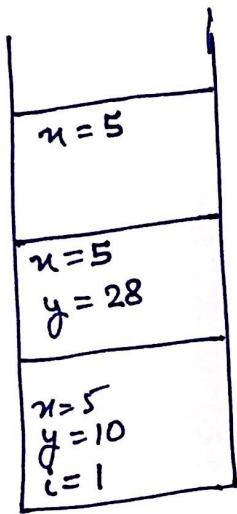
var = 5
X 3 X
X 0 - 1



O/P: - 5 - 4 - 3 - 2 - 1 - 0

eg:- main()

```
int x=5, y=10, i;  
for( i=1; i<=2; i++)  
{  
    y += f(x) + g(x);  
    printf(y);  
}
```



STACK

f(int x)

```
{  
    int y;  
    y = g(x);  
    return(x+y);  
}
```

g(int x)

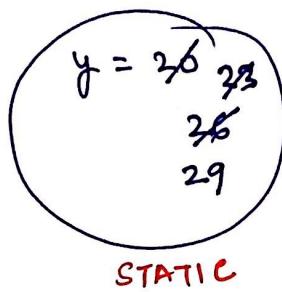
Initialization
modification

```
static int y = 20;  
y += 3  
return(x+y);
```

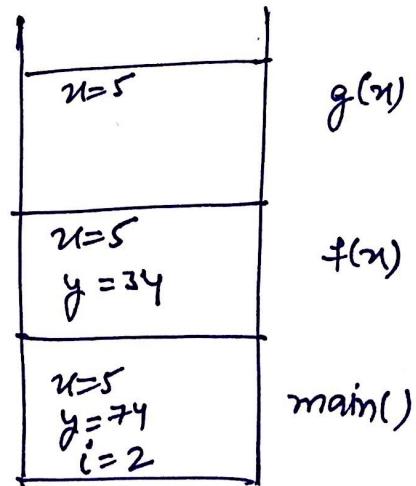
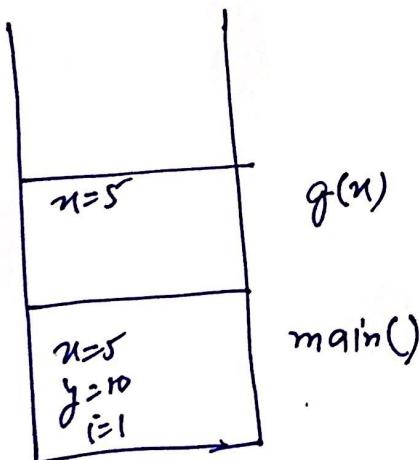
?

$$y += \cancel{28}^{33} + 31$$

$$y = 10 + 64 = 74$$



g(x)
f(x)
main()



g(x)
f(x)
main()

O/p:- 74, 150

$$y += 39 + 37$$

$$y = 150$$

eg:-

$f(\text{int } x)$

{ static int $r = 10$;

if ($n \leq 0$)

{

$r += 7$;

return(3);

}

if ($n > 3$)

return ($r + f(n-2)$);

else

return ($r - f(n-1)$);

}

main()

{

int $n = 6$;

printf("%d", f(n));

}

$$f = f_0 + f_1 + f_2 + f_3 + f_4 + f_5 + f_6$$

$n=0$	$r=r+f_0$
$n=1$	$r-f_0$
$n=2$	$r-f_1$
$n=3$	$r+f_2$
$n=4$	$r-f_3$
$n=5$	$r+f_4$
$n=6$	$r-f_5$
$n=6$	printf(f(6))

$f(0)$

3

$f(1)$

14

$f(2)$

3

$f(3)$

20

$f(4)$

2

$f(5)$

37

main()

37

O/P:- 37

$r=10$
17

Q3:

main()

{
 int $x = 5, y;$

$$y = \underline{++x} * \underline{++x} * \underline{++x}$$

~~4 5 6~~ x

printf(x, y);

}

$$\begin{aligned} O/p &:= (6 \times 6) * 7 \\ &= 36 * 7 = \underline{\underline{252}} \end{aligned}$$

$(y = \underline{++x} * \underline{++x} * \underline{++x})$ will give 252 because increment operator is left associative
 $\hookrightarrow y = (\underline{++x} * \underline{++x}) * \underline{++x}$ $x=5$ $x=6$ $x=7$
 ↓ unary ↓ binary

so, after unary op^h $x=6$

$$\text{so, } x * x = 6 \times 6 = 36$$

~~36~~
temp

Then $++x \rightarrow (x=7)$

so,

$$\boxed{36} * \boxed{7} \quad x$$

~~252~~
y

word
by mistake
(Jodhpur)

Extern Variable :-

eg:- int a=5;
main()

```
    {
        extern int a;
        printf(a);
    }
```

O/p:- 5

```
int a=5;
main()
{
    int a;
    printf(a);
}
```

O/p:- Garbage value

Extern keyword tell compiler
to get memory ^{global} ~~outside~~
with same name

so, no memory created for this
extern variable.

eg:- main()

```
{
    extern int a;
    printf(a);
}
```

}

O/p:- Error
(Undefined symbol)

eg:- main()

```
{
    extern int a;
    printf(a);
    a=10;
    printf(a);
}
```

}

O/p:- 0
10

```
eg:- extern int a;  
main()  
{  
    extern int a;  
    printf(a);  
}
```

```
eg:- extern int a=10;  
main()  
{  
    extern int a;  
    printf(a);  
}
```

O/p:- Error (Undefined symbol a)

O/p:- 10

Note:- If we declare variable with extern then no memory allocated.

If we initialize variable with extern the memory allocated.

```
eg:- extern int a=10;  
main()  
{  
    extern int a = 6;  
    printf(a);  
}
```

GLOBAL
eg:- extern int a=10;
main()
{
 LOCAL {
 extern int a;
 static int a;
 printf(a);
 }
}

O/p:- Error

Extern can't initialize var inside funcn body

O/p:- Error

Ambiguity in a.

```

int a = 5, b = 10;
main()
{
    extern int a;
    register int b = 20;
    c();
    printf(a, b);
    d();
}

```

c()
 static int b;
 printf(a, b); free variable
 d();
 a = 1; b = 2;

{ function won't work
 (a) because

Q: Why printing value of b from main() is not working?

D()
 {
 extern int b;
 static int a = 7;
 printf(a, b);
 printf(a, b);
 a = 4; b = 5;
 }

E() calls printf

static int a = 7;
 auto int b = 23;
 printf(a, b); free variable
 a = 4; b = 5;

{ function won't work

O/P:-

DYNAMIC SCOPING

S, 0
 7, 10
 5, 23
 1, 20
 4, 5
 1, 23

5, 0
 7, 10
 7, 23
 1, 20
 4, 5
 4, 23

STATIC SCOPING

eg:- main()

(SWITCH)

{ int i; i=10;

for(; i<= 25 ; i++)

{
switch(i)
{

case 1 : i+=3;

case 2 : i+=4;

case 3 : i+=5;

default : i+=3;

break;

}

printf("%d", i);

}

}

O/p :-
16
20
24
28

- Switch case statement , break is required after each case otherwise it will execute all cases sequentially .

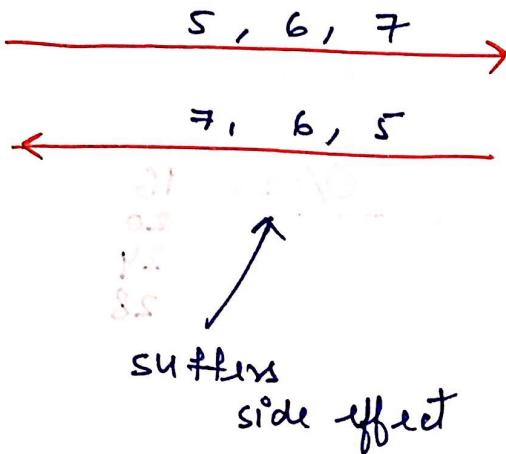
- Default position is not fixed in switch , it can be placed anywhere and it will execute when there is no match there.

- Default (last statement) → no break needed

Default (middle statement) → break needed

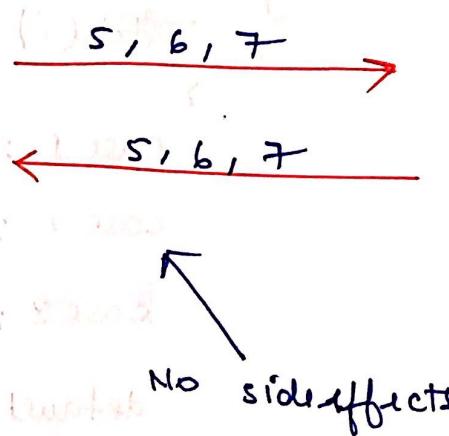
Side Effect :-

```
int i = 4;  
printf( ++i, ++i, ++i);
```



(Answer different
for $L \rightarrow R$ & $R \rightarrow L$)

```
int i=4;  
printf("%d,%d,%d", i+1, i+2, i+3)
```



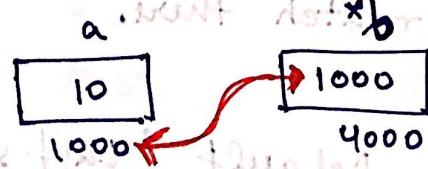
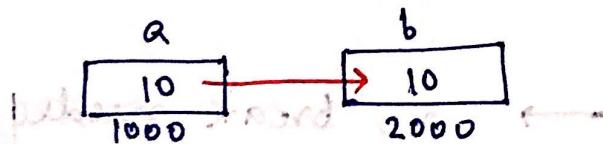
Pointer :-

Als we de berekening in deel 1 door $\frac{a}{1000}$ delen dan blijft de waarde van a onveranderd en zijn de resterende getallen veel kleiner.

printf(a) → 10 (data)

`printf(&a) → 1000 (Address)`

int $b = a$; int $*b = &a$; based



int b = &a



since size of b is 2 Bytes,
but size of address is
depends upon system memory

so, if address size > size of b

it can't store in b.

so, we use pointer,

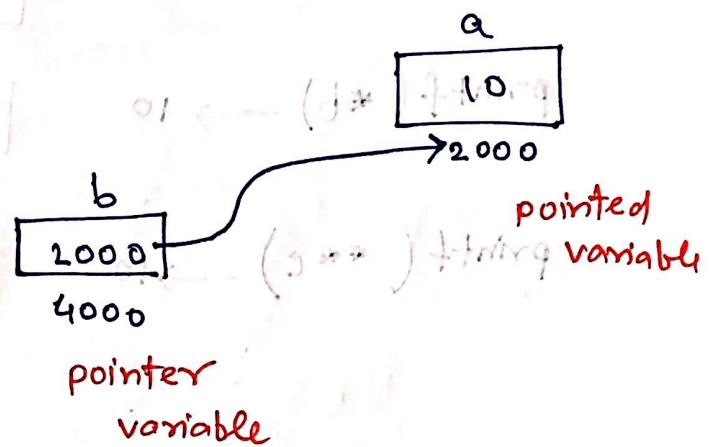
dereference operator

int *b = &a; i.e

so, printf(b) = 2000

printf(&b) = 4000

printf(*b) = 10



So, pointer is a special variable used to store address of other variables.

'*' dereference operator can only be used with variables which address of some other variables.

eg:-

int a = 10;

int *b = &a;

✓ printf(*b);

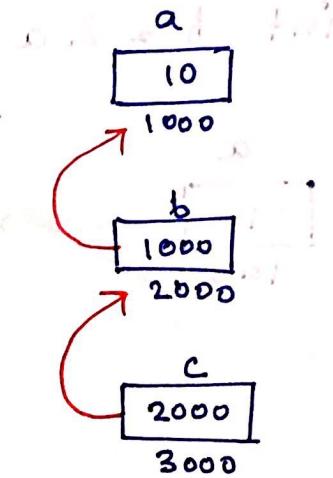
✗ printf(*a);

} a is normal variable which carry data not address

eg:- value of a for variable
int $a = 10;$ — var
in declaration to variable

programme int *b = &a; → pointer

double int **c = &b; → double
pointer
which stores address of pointer



printf(a) → 10 [Immediate mode]

printf(*b) → 10 [Direct addressing]

printf(**c) → 10 [Indirect addressing]

eg:-
int a = 10;
int *b = &a;
printf(a * b)

O/P:- INVALID operation as there is no operator b/w a & *b

After knowing you was storage addressed compiler checks declaration for *b is, and

if it present then * behaves as dereference operator otherwise multiplication.

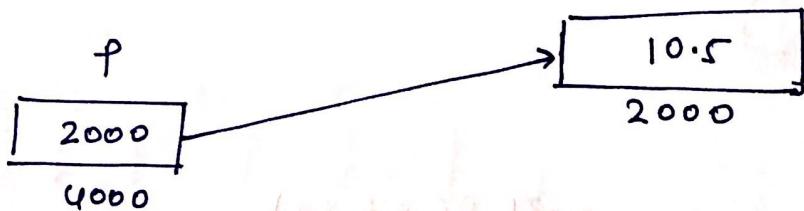
Eg:-

```

int a = 10;
int *b = &a;
int **c = &b;
    
```

size of (c) → 8B
 size of (*c) → 8B
 size of (**c) → 2B.

Eg:- P is float pointer:



size of (p) → 8B

size of (*p) → 4B

print(p+1) → 2000 + 4 = 20004

print((*(p)+1)) → 10.5 + 1 = 11.5

Eg:- int i=5, j=10;

main()

{
printf(i, j);

fun(&i, &j);

printf(i, j);

}

fun(int *i, int *j)

{
*i = 7;

*j = *(i+3);

i = j;

*j = 4;

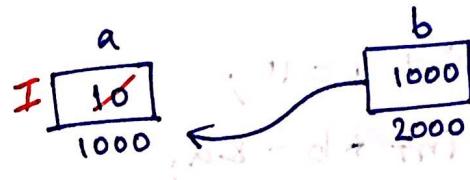
}

[O/p :- 5,10
4,7]

eg:- main()

```
int a=10;
int *b=&a;
scanf("%d", b);
printf("%d", a+25);
```

}



// scanf("%d", b);

address
modified

O/p :- I+25, a + 25



void P(int *y)

{ int x = *y + 2;

Q(x);

*y = x-1;

printf(x);

main() = 1 + 3.01 + 2 + (1+4.01) + 6

main()

{

x=5;

P(&x);

printf(x);

}

O/p :- 12 7 6

eg:- $f(\text{int } n, \text{ int } *p)$: - main() + \rightarrow (10, 20)

```

    {
        int a = 5, b = 6;
        int *p = &a, *q = &b;
        *p = 20;
    }

```

Q. yesterday a person is solving problem $a = *p$; and

\rightarrow $f(a, b) ; \text{ print}(a, b);$

O/P :- 20 20

0001	0001	0001	0001	0001	0001
0001	0001	0001	0001	0001	0001

$*q = 8b;$

$*p = 20;$

$\rightarrow f(a, b);$ direct

0001 print(a, b); q

}

eg:- main()

{ int a[] = { 25, 30, 90, 55, 22, 63 }; }

$\rightarrow \text{printf}("1.d", f(a, 6));$

O/P :- -21

3

$f(\text{int } *a, \text{ int } n)$

{

if ($n \leq 0$)

return (0).

else if ($*a + 2 == 0$) ~~return~~ \rightarrow "20" (H.W.)

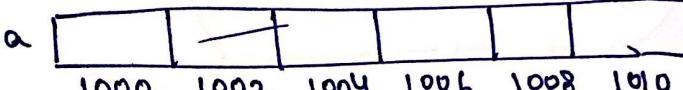
return ($*a - f(a+1, n-1)$);

else

return ($*a + f(a+1, n-1)$);

return;

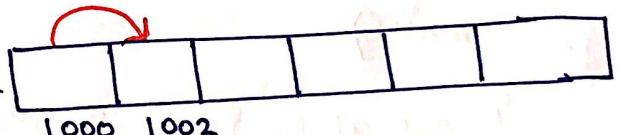
- $a[i] = *(a+i) = *(i+a) = i[a]$
- $\underbrace{a[i][j]}_b = b[j] = *(b+j) = *(a[i]+j) = *(*(a+i) + j)$
- Base address of array which is itself a pointer, can't be modified.

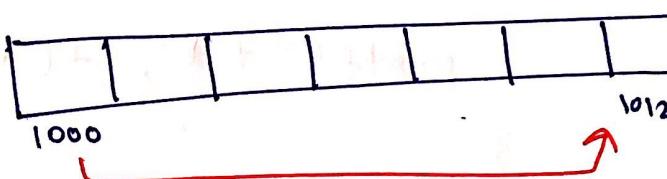
$\text{int } a[6];$ a 

$\text{print}(a) \rightarrow 1000$

$\text{print}(&a) \rightarrow 1000$

But,

skip element $\text{print}(a+1) \rightarrow 1002$ a 

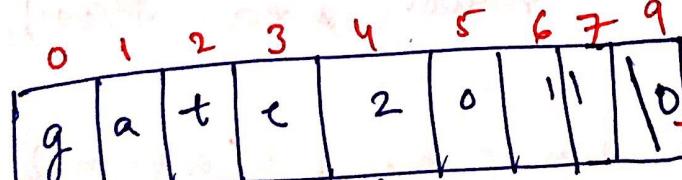
skip array $\text{print}(&a+1) \rightarrow 1012$ 

eg:- main()

{
 char a[] = "gate 2011";

 printf("%s", a + a[3] - a[1]);

3

a 

O/p: 2011

so, $a + e - a = a + 4$

$s = "abcd"$ \rightarrow size of (s) = 5 (Include NULL char)

\downarrow
 $\text{strlen}(s) = 4$

= $\text{printf}("s", \text{base address})$; print all char from BA to null char.

= $\text{printf}("%c", \text{base address})$; print char at BA only

so,

$\text{printf}("%s", "gate2011") \rightarrow \text{gate2011}$

$\text{printf}("%s", a) \rightarrow \text{gate2011}$

$\text{printf}("%s", a+2) \rightarrow \text{te2011}$

$\text{printf}("%c", a+3) \rightarrow e$

$\text{printf}("%d", a+1) \rightarrow 97$

$\text{printf}("%c", *(a+1)+1) \rightarrow b$

Eg:- main() \rightarrow (2) for int. s = "string"; \rightarrow b also \rightarrow s
` char a[10];
char *b = "string"; \rightarrow b also \rightarrow s
int length = strlen(b);
for (i=0; i<length; i++)
{
 a[i] = b[length-i];
}
printf("%s", a);
}

0. 1 2 3 4 5 6
 s t r i n g \0
 1000 01 02 03 04 05 06
 length
 NULL char
 exclude NULL char

$*b \rightarrow s$

$*(\text{b}+1) \rightarrow t$

$*(\text{b}+2) \rightarrow r$

:

:

:

$*(\text{b}+5) \rightarrow g$

a	0	1	2	3	4	5
	\0	g	n	i	r	t

printf("%s", a)

↳ 1. Nothing as

starting char is itself a
 \0 char.

eg:- storing to ~~good~~
main()

{ char a[] = "Hello";
char b[] = "Hello"; O/p:- Error

if ($a == b$)

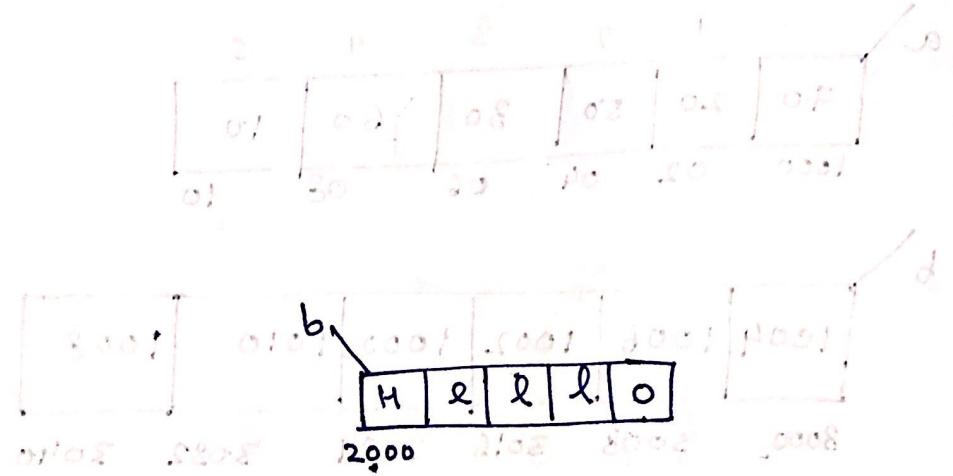
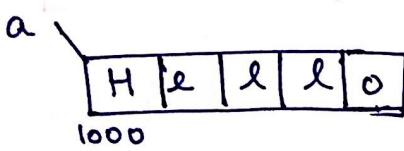
* printf("hi");

else

* printf("bi");

}

(* $a + r$)



since a and b carry base address which can't be changed
 $base = 2 \text{ can't change}$

so, $(a == b)$ can't be performed

if ($a == b$) O/p:- "bi"

as $a = 1000$
 $b = 2000$

if ($*a == *b$)

O/p:- "Hi"

$*a = H$

$*b = H$

Array of pointers

Q:-

main()

{ int a[] = { 70, 20, 50, 20, 60, 10 };

int *b[] = { a+2, a+3, a+1, a, a+5, a+4 }

a	0	1	2	3	4	5
	70	20	50	30	60	10
	1000	1002	1004	1006	1008	1010

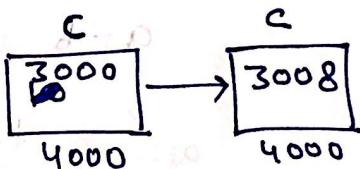
b	1004	1006	1002	1000	1010	1008
	3008	3016	3012	3024	3032	3040

priority:

(++ > *)

mod process of b/w a & a+2
int **c = b;

*c++ ; \rightarrow c = 3008



printf(c-b, *c-a, **c); \rightarrow $\frac{3008-3000}{8}, \frac{1006-1000}{2}, 30$

(1, 3, 30)

printf(c-b, *c-a, **c); \rightarrow (2, 1, 20)
++**c;

printf(c-b, *c-a, **c); \rightarrow (2, 1, 21)
++*++*c

printf(c-b, *c-a, **c); \rightarrow (2, 2, 51)

?

Had a doubt

eg:-
main()

{ float c[] = { 20, 90, 33, 60, 44, 23 };

float * a[] = { c+3, c+2, c+1, c, c+5, c+4 };

float * b = a;

2. $b-a$, $*b-c$, $**b$; : address of first element

printf(b-a, *b-c, **b);

3. $a[2]++ * ++* +b$; : address of fourth element
(because)

printf(b-a, *b-c, **b); $\rightarrow (1, 3, 61)$

4. $a[2]++ * *b + 9$; : address of fifth element
printf(b-a, *b-c, **b); $\rightarrow (1, 3, 62)$

}

c	0	1	2	3	4	5
	20	90, 91	33	60, 61 62	44	23

a	0	1	2	3	4	5
	1012	1008 10012	1004	1000	1020	1016

b	0	1	2	3	4	5
	2000	2008	2004	2000	2020	2016

NOTE:- We can't add pointer ~~with~~ any other Numeric data-type.

- We can subtract two pointers eg:- $\text{ptr}_1 - \text{ptr}_2$
i.e. no. of position ptr_1 ahead of ptr_2
- Adding const to pointer : $\text{ptr} + 5$ (Move 5 steps forward)
Subtract const to pointer : $\text{ptr} - 5$ (Move 5 steps backward)
- We can't perform Addition/ Multiplication/ division because : they had no meaning.

eg:-

main()

```
char *a[] = {"Papa", "gotohell", "break", "ayyooo",
              "bacha", "chachi"};
```

```
char ***b[] = {a+2, a+1, a, a+3, a+4, a+5}
```

```
char ***c = b;
```

```
++*++*c+4;
```

printf("%s", ~~++***(c+2)+2~~); → a

printf("%s", ~~++***c+3~~); → b

printf("%s", ~~*++*++c+2~~); → c

}

or

0	1	2	3	4	5
2000 2001	3000 2001	4000 2001	5000 5001	6000 2001	7000 2001

6	23	000120	0	031	44	548
000120	60	61	62	63	64	65

60

2-D array pointer

eg:-

main()

{

int a[5][3] = {10, 20, 30, ----, 150};

printf("%d", (*a) && (a[0] == a)); → 1

printf("%d", *(a+2)+3); → 1036

printf(*a+10, a[10]+3, **a+30); → 1020, 1066, 40

printf(*(a+3)+10, *(a+3)+60); → 1038, 160

}

a	0	1	2
0	10 1000	20 02	30 04
1	40 06	50 08	60 10
2	70 12	80 14	90 16
3	100 18	110 20	120 22
4	130 24	140 26	150 28
5	22	24	26

since multidimensional array stored as 1-D array inside computer system, so we only need one Base Add.

row-major order

$$a = 1000$$

row selected only

$$*a = a[0]$$

$$**a = 10$$

Both row & column selected

$$\begin{aligned}
 - a+1 &= 1002 \\
 &= 1000 + 2(\text{row size}) \quad (\text{skip-row}) \\
 &= 1006
 \end{aligned}$$

$$\text{Now, } a+1 = 1006$$

$$*(a+1) = 40 \\ = 1006$$

only row selected,
column has to be
selected

$$*\underline{\text{row}} \underline{\text{column}} + 1 = 1008$$

Both row & column selected

$$*\underline{*(a+1)+1} = 50 \quad \text{Data at 1008}$$

$$(a+1)+1 = 1006 + 1006 \quad \begin{array}{l} \text{skip} \\ \text{row 2 times} \end{array} \\ = 1012$$

$$*\underline{(a+1)+1} = 1012$$

$$*\underline{(a+1)+1} + 1 = 1014$$

Row, column
selected

$$*\underline{(*((a+1)+1)+1)} = 80$$

so,

$$a = 1000$$

$$&a = 1000$$

for array
both are
same.

$$a+1 = 1006 \quad \begin{array}{l} \text{skip 1 row} \\ \text{* } a+1 = 1002 \end{array}$$

$$&a+1 = 1030 \quad \begin{array}{l} \text{skip whole} \\ \text{array} \end{array}$$

main()

$$a = \underline{1000}$$

{
 int a[4][5][3];
 printf(*(a+3) +7), → 1132
 printf(*a[4] +20), → 1160
 printf(**a +30, *(*(a+2)+2) +10) → 1060,
 1092
 printf(*(*a+7) +10) → 1062

}

$$= *(a+3) + 7$$

$$\text{matrix size} = 15 \times 2B$$

$$= 1000 + 3(30) + 7(6)$$

$$1701 = (1+1+B) * = \underline{30B}$$

$$= \boxed{1132}$$

$$= \underline{6B}$$

$$= * a[4] + 20$$

$$1001 = (1+1+B) * = \underline{3 \times 2B}$$

$$= * * (a+4) + 20$$

$$= 1000 + 4(30) + 20(2)$$

$$= \boxed{1160}$$

$$= **a + 30 , * (*(a+2)+2) +10$$

$$= \boxed{1060}$$

$$1000 + 2(30) + 2(6) + 10(2)$$

$$\boxed{1092}$$

$$= *(*a+7) +10$$

$$= 1000 +$$

$$7(6) + 10(2) = \boxed{1062}$$

eg:- ~~char s = "1-0 \n";~~ ~~int a = 10; printf("%d", a);~~

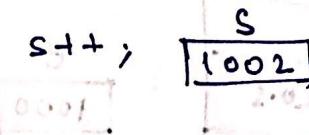
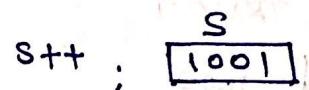
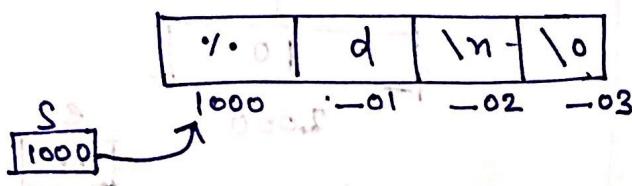
main()

```
{ char *s = "1-0 \n";  
    s++;  
    s++;
```

printf(s-2, 300);

}

$$\begin{array}{r} 1002 \\ - 1002 \\ \hline 0 \end{array}$$



↳ printf("1-0 \n", 300)

= 300 → 300 shows standard output

(S. 10.1) H/w

Structure & Union :-

eg:- ~~struct node~~

```
int a; - 2B  
float b; - 4B } 7B  
char c; - 1B
```

}

User data type is created

d (definition)

10
20.5
g

User defined datatype

struct node d = {10, 20.5, 'g'};

int d = 10;

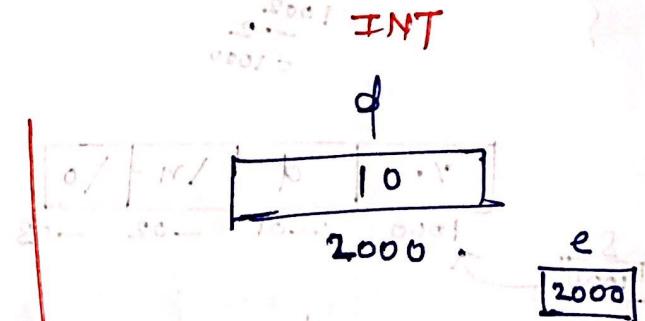
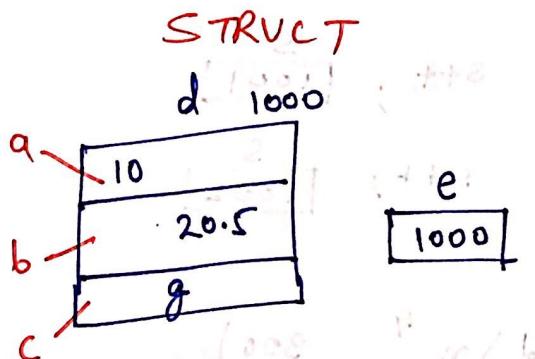
2B	d	10
----	---	----

pre-defined datatype

`printf("%d", d) → 10` (int d is accrued)

`printf("%d", d.a) → 10` (node d is accrued)

member access operator



`struct node *e = &d;`

`int *e = &d;`

`printf("%u", e);`
↳ 1000

`printf("%u", e);`
↳ 2000

`printf("%d", (*e).a);`

`printf("%d", *e);`
↳ 10

`printf("%f", (*e).b);`
↳ 20.5

`printf("%c", (*e).c);`
↳ g

`printf("%c", e->c);`
↳ g

NOTE:-
Priority:-

(. > *)

eg:-

struct s_1

{

int a; 2B

float b; 4B

};

struct s_2

{

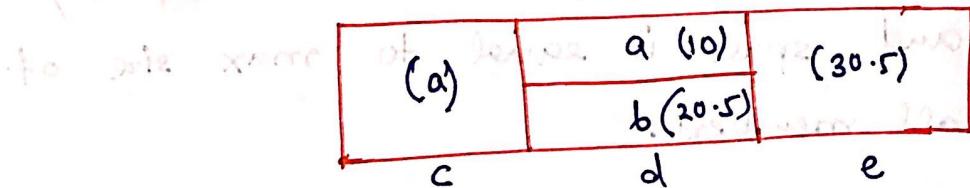
char c; 1B

struct s_1 d; 6B

float e; 4B

};

for (int i = 0; i < 10; i++) {
 f.a = 'a'; f.b = 20.5; f.d.a = i; f.d.b = 10 + i; f.d.c = 'c' + i;
 f.e = 30.5 + i;
}

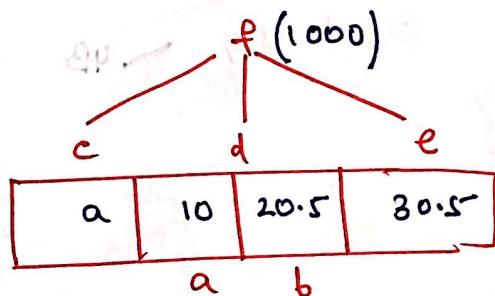


so the value of f is instance will be like this
 $\text{struct } s_2 \Rightarrow f = \{ 'a', \{ 10, 20.5 \}, 30.5 \} ;$

$\text{printf}(\cdot \cdot \cdot , f.d.b) \Rightarrow \underline{\underline{20.5}}$

$\text{printf}(\cdot \cdot \cdot , f.e) \Rightarrow \underline{\underline{30.5}}$

struct s_2 * g = &f;



$\text{printf}(\cdot \cdot \cdot , (*g).d.b) \Rightarrow 20.5$



$\text{printf}(\cdot \cdot \cdot , (g \rightarrow d).b) \Rightarrow 20.5$

$\text{printf}(\cdot \cdot \cdot , g \rightarrow e) \Rightarrow 30.5$

`printf(g) ⇒ 1000`

`printf(g+1) ⇒ 1000 + 11 = 1011`

(SKIP whole struct node)

Union :- All member share common space, and space is equal to max size of all members.

In this one of the member is active at a time.

`union u,`

{

`char c; → 1B`

`struct node d; → 6B`

`float e; → 4B`

so, size of (u) = 6

}

eg:-

`int a[] = { 6, 7, 8, 18, 34, 67 };`

`int a2[] = { 25, 56, 28, 29 };`

`int a3[] = { -12, 27, -31 };`

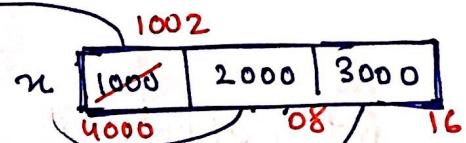
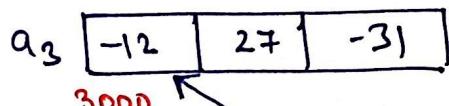
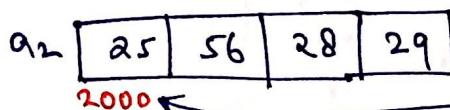
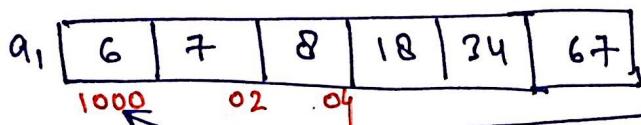
int *x[] = { a₁, a₂, a₃ }; relation to main()

f(int **a)

```

{
    printf( a[0][2] ); ⇒ 108 → f(a)
    printf( *a[2] ); ⇒ -12 → f(a);
    printf( *++a[0] ); ⇒ 7 → f(a);
    printf( *(++a)[0] ); ⇒ 25 → f(a);
    printf( a[-1][1] );
}

```



$$a[0][2] = *(*(\underline{a+0}) + 2) = *\underbrace{(*a + 2)}_{1000} = *1004 = \underline{\underline{8}}$$

$$*\underline{a[2]} = *\underline{(*(\underline{a+2}))} = -12$$

$$*\underline{++a[0]} = *\underline{++*\underline{a}} = \underline{\underline{7}}$$

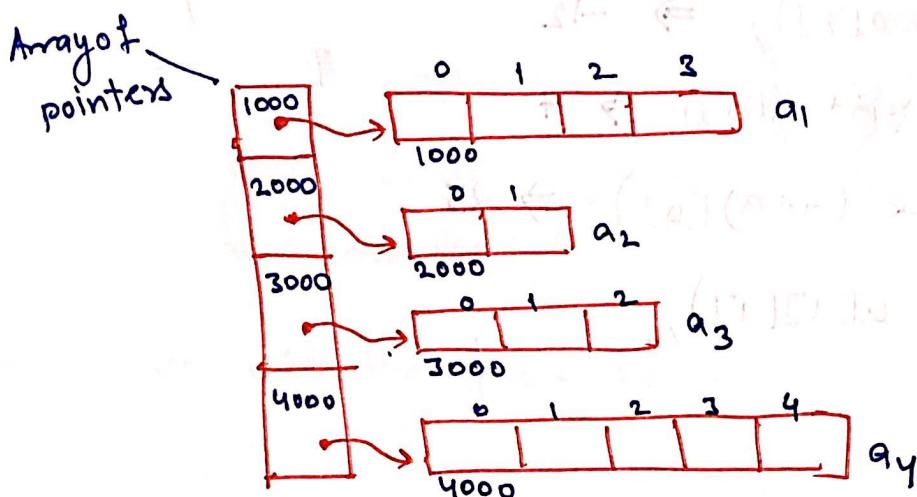
$$*\underline{*(++a)[0]} = *\underline{2000} = \underline{\underline{25}}$$

$$*\underline{a[-1][1]} = *\underline{(\underline{a-1} + 1)} = \underline{\underline{8}}$$

$$\underline{a = 2000}$$

Note:- Array of pointers can be used as } 2D array format.

- We can create an ~~non~~ 2-D array of non-uniform columns using array of pointers.



Q:
 int $a[10][5]$;
 int $*b[10]$;

which option can't possible:

(a) $a[5][3] = \text{data}$

(b) $a[5] = \text{address}$

(c) $b[5] = \text{address}$

(d) $b[5][2] = \text{data}$

$a[5] = *(\&a+5)$ → address of row of

2D array
which can't be

changed.

Type Casting:- Conversion of one data-type to another during some operations.

Implicit

Explicit

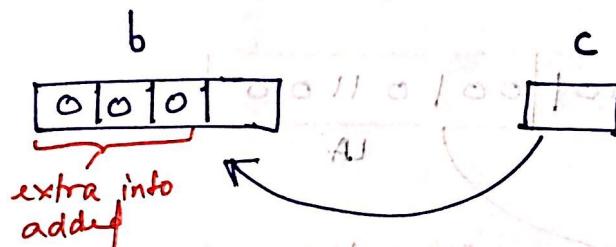
Implicit :-

```
int a;  
float b;  
char c;
```

b = c;

b = c
float 4B char 1B
small size var large size var

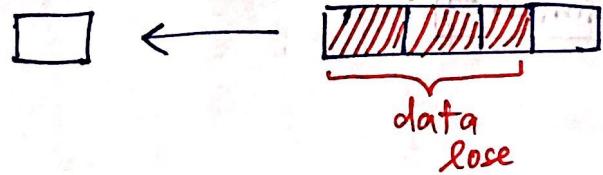
(small size → large size)



storing small size info into bigger storage space then compiler cast it by itself.

Explicit :-

c = b;
char 1B float 4B



small space

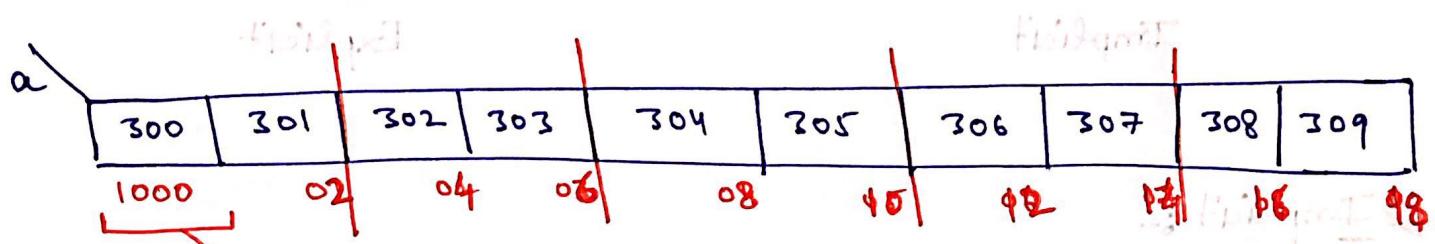
in this will compiler generates error.

~~c = (char) b;~~
target type

so, user has to done it explicitly.

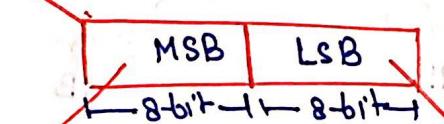
Type casting in pointers :-

int a[10] = { 300, 301, 302, ..., 309 };



2B

MSB LSB



Higher address

Lower address

300

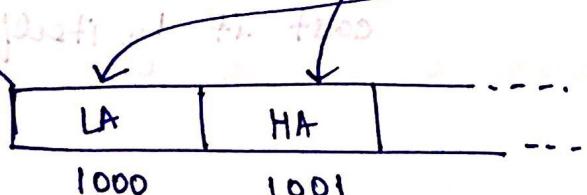
\Rightarrow 00000001 00101100

HA

LA

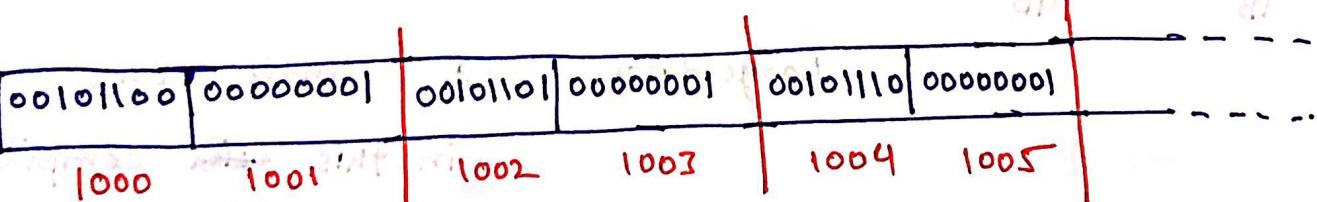
Little endian

a

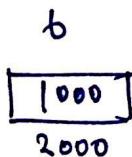


similarly :-

a

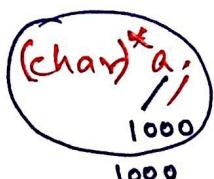


`int *b = a;`



No type casting required
(same type).

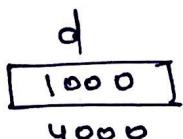
`char *c = (char *) a;`



int
a is address which
and after type casting it is an
char address of 8B

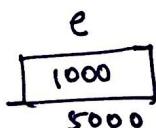
so, same size.

`float *d = (float *) a;`



INT add a → FLOAT address
1000 1000
(8B) (8B)

`void *e = a;`



INT
`b → 1000`
`*b → 300`
Read 2B

CHAR
`c → 1000`
`*c → 44`
Read 1B only

FLOAT
`d → 1000`
`*d → data`
Read 4B

VOID
`e → 1000`
`*e → X`
can't read \$
as it don't
know how
many B have
to read

`b+1 → 1002`
MOVE 2B

`c+1 → 1001`
MOVE 1B

`d+1 → 1004`
MOVE 4B

`e+1 → X`
can't mov
`c = (char *) e`

so a
void pointer
can be
cast into
any type