# Quiz-2 (ADA-2024)

## February 14, 2024

# Roll Number: Section:

- 1. Which of the following is true?
- (A) While soliving original problem Dynamic programming divides problem into non-overlapping subproblems.
- (B) Divide and conquer divides a problem into subproblems while dynamic programming does not divide into subproblems.
- (C) Dynamic programming does not always guarantee an optimal solution.
- (D) While soliving original problem Dynamic programming divides problem into overlapping subproblems

## Option (D) is correct.

- **2.** Let T be a node weighted rooted tree that is not necessarily a binary tree. The weight function is  $\operatorname{wt}:V(T)\to\mathbb{R}^+$ . Moreover, let the root of the tree be r. Also, for every node x, let  $T_x$  denotes the subtree of T rooted at x. Let  $\mathsf{MVC}(x,0)$  denotes the total weight of a minimum weighted vertex cover of  $T_x$  that does not contain x and  $\mathsf{MVC}(x,1)$  denotes the total weight of a minimum weighted vertex cover of  $T_x$  that does contains x. Let Y denotes the set of all children of x. Clearly,  $Y=\emptyset$  if x is a leaf and  $Y\neq\emptyset$  otherwise. Which of the following statements is correct?
- (A) The subproblem MVC(r, 1) returns the value of an optimal solution for the tree T.

(B) If 
$$Y \neq \emptyset$$
, then  $\mathsf{MVC}(x,1) = \mathsf{wt}(x) + \sum_{y \in Y} \mathsf{MVC}(y,0)$ .

(C) If 
$$Y \neq \emptyset$$
, then  $\mathsf{MVC}(x,0) = \sum_{y \in Y} \mathsf{MVC}(y,1)$ .

(D) If 
$$Y \neq \emptyset$$
, then  $\mathsf{MVC}(x,0) = \mathsf{wt}(x) + \sum_{y \in Y} \mathsf{MVC}(y,0)$ .

#### Option (C) is correct

3. Suppose that there are n weeks, a plan is specified by a 'low stress job', a 'medium stress job', and a 'high stress job' or 'none' in each of the weeks. For any  $i \geq 3$ , if a high stress job is chosen for the week i, then no job can be chosen in the week i-1 and i-2. If a medium stress job is chosen for the week  $i \geq 2$ , then no job can be chosen at the week i-1. However, if a low stress job is chosen for the week i, then both the high stress, medium stress, and low stress job can be chosen for the week i-1. In the week 1, a high stress job, or a medium stress job, or a low stress job both can be chosen. Similarly, in the second week, a high stress job can be chosen if no job was chosen in the first week. A low stress job gives revenue  $a_i$ , a medium stress job gives revenue  $b_i$ , and a high stress job gives a revenue of  $c_i$ . The objective is to design an algorithm that maximizes the revenue.

Let  $\mathsf{REVENUE}(k)$  denotes the maximum possible revenue in the weeks  $\{1, 2, \dots, k\}$ . Which of the following options is correct?

$$(A) \ \, \text{For every } k \geq 4, \, \text{the value REVENUE}(k) = \max \begin{cases} b_k + \mathsf{REVENUE}(k-1) \\ c_k + \mathsf{REVENUE}(k-2) \\ a_k + \mathsf{REVENUE}(k-3) \end{cases}$$
 
$$(B) \ \, \text{For every } k \geq 4, \, \text{the value REVENUE}(k) = \max \begin{cases} c_k + \mathsf{REVENUE}(k-1) \\ a_k + \mathsf{REVENUE}(k-2) \\ b_k + \mathsf{REVENUE}(k-3) \end{cases}$$

(B) For every 
$$k \ge 4$$
, the value  $\mathsf{REVENUE}(k) = \max \begin{cases} c_k + \mathsf{REVENUE}(k-1) \\ a_k + \mathsf{REVENUE}(k-2) \\ b_k + \mathsf{REVENUE}(k-3) \end{cases}$ 

(C) For every 
$$k \ge 4$$
, the value  $\mathsf{REVENUE}(k) = \max \begin{cases} b_k + \mathsf{REVENUE}(k-1) \\ a_k + \mathsf{REVENUE}(k-2) \\ c_k + \mathsf{REVENUE}(k-3) \end{cases}$ 

(D) For every 
$$k \geq 4$$
, the value  $\mathsf{REVENUE}(k) = \max \begin{cases} a_k + \mathsf{REVENUE}(k-1) \\ b_k + \mathsf{REVENUE}(k-2) \\ c_k + \mathsf{REVENUE}(k-3) \end{cases}$ 

#### Option (D) is correct.

- **4.** Which of the following is not true?
- (A) Memoization makes recursive function calls to solve subproblems.
- (B) Table-Filling starts from the simplest subproblem and gradually reaches up to the original problem.
- (C) Unlike memoization, table-Filling method does not use solutions of subproblem while solving the original problem.
- (D) Memoization starts from the original problem and gradually breaks it down into subproblems.

#### Option (C) is correct

- 5. Suppose that an equipment manufacturing company manufactures  $s_i$  units in the *i*-th week. Each week's production has to be shipped by the end of that week. Every week, one of the two shipping agents A and B are involved in shipping that week's production and they charge in the following:
  - Company A charges a rupees per week but will only ship for a block of three consecutive weeks.
  - Company B charges b rupees per unit.

The objective is to find the minimum cost of a schedule. Let Cost(k) denotes the minimum cost of any schedule for the weeks  $\{1, 2, \dots, k\}$ . Which of the following options is correct?

(A) For all 
$$k \ge 4$$
,  $Cost(k) = min \begin{cases} 3a + Cost(k-3) \\ bs_k + Cost(k-1) \end{cases}$ 

(A) For all 
$$k \ge 4$$
,  $\operatorname{Cost}(k) = \min \begin{cases} 3a + \operatorname{Cost}(k-3) \\ bs_k + \operatorname{Cost}(k-1) \end{cases}$   
(B) For all  $k \ge 4$ ,  $\operatorname{Cost}(k) = \min \begin{cases} 3a + \operatorname{Cost}(k-2) \\ bs_k + \operatorname{Cost}(k-1) \end{cases}$ 

(C) For all 
$$k \ge 4$$
,  $\operatorname{Cost}(k) = \min \begin{cases} as_k + \operatorname{Cost}(k-1) \\ 3b + \operatorname{Cost}(k-3) \end{cases}$ 

(D) For all 
$$k \ge 4$$
,  $\mathsf{Cost}(k) = \min \begin{cases} as_k + \mathsf{Cost}(k-1) \\ 4b + \mathsf{Cost}(k-3) \end{cases}$ 

## Option (A) is correct