

# Assignment-1

1. Suppose hypothetical processor has 2 I/O Regs:

0011 = Load AC from I/O

0111 = Store AC to I/O

In these case, 12 bit address identifies an external device. Show prog. execution

(a) Load AC from device 6

(b) Add content of memory location 880

(c) Store AC to device 7

Assume next value retrieved from device 6 is 5 2

Memory 880 contains a value 4. i.e.,  $4 + 5 = 9$

Memory				Memory			
300	3006	300	PC	300	3006	301	PC
301	5880	301	AC	301	5880	5	AC
302	7007	302	IR	302	7007	3006	IR

880	4	006	5	device G	880	4	006	5	device 6
881		007		device 7	881		007		device 7

Memory				Memory			
200	2006	301	PC	900	3006	302	PC
301	5880	301	AC	301	5880	9	AC
302	7007	302	IR	302	7007	5880	IR

Memory				Memory			
200	2006	302	PC	300	3006	303	PC
301	5880	301	AC	301	5880	9	AC
302	7007	302	IR	302	7007	7007	IR

(Q2) Consider a hypothetical 64-bit microprocessor having 64-bit ins. composed of 2 fields. The first 4 bytes contain op-code & next 4 contain operand address. What is max directly addressable memory capacity?

Sol:  $4 \text{ bytes} = 32 \text{ bits}$

So,  $64 - 32 = 32 \text{ bits}$  contain operand address.

So,

$$\begin{aligned} \text{Maximum Memory Capacity} &= 2^{32} \text{ bytes} \\ &= 4 \text{ Gbytes} \end{aligned}$$

(Q3) Cal. the system's avg. access time. The key data includes cache access time (50ns), main memory access time (500ns), the fraction of read operation (0.9), & cache hit ratio for read operation (0.9).

Sol.

Fraction of read operation = 0.8

Fraction of write operation = 0.2

MISS ratio = 0.1

$$\begin{aligned} \text{Time to read} &= (0.9 \times 50 \text{ ns}) + 0.1 \times (500 \text{ ns} + 50 \text{ ns}) \\ &= (45 + 55) \text{ ns} \\ &= 100 \text{ ns} \end{aligned}$$

Time to write = 500 ns.

$$\begin{aligned} \text{Avg. access time} &= 0.8 \times \text{time to read} + 0.2 \times \text{time write} \\ &= (0.8 \times 100 \text{ ns}) + (0.2 \times 500 \text{ ns}) \\ &= 80 \text{ ns} + 100 \text{ ns} \\ &= 180 \text{ ns} \end{aligned}$$

(Q4) Consider a computer with 400 Mbytes of available memory. Three prog. Job1, Job2 & Job3 are submitted for execution at same time with a time limit 15 sec.

Type of Job	Job1	Job2	Job3
Heavy Computation	(90%)	(10%)	(10%)
Duration	10 min	20 min	15 min.
Memory req.	100M	150M	125M

executed in sequence Job1, Job2 & Job3. Find out CPU utilization, memory utilization & throughput in uniprogramming & multiprogramming

Sol: CPU utilization -

#### Uniprogramming:

$$\text{Job1} \rightarrow 90\% \text{ for } 10 \text{ mins} = 9 \text{ mins. of CPU time}$$

$$\text{Job2} \rightarrow 10\% \text{ for } 20 \text{ mins} = 2 \text{ mins of CPU time}$$

$$\text{Job3} \rightarrow 10\% \text{ for } 15 \text{ mins} = 1.5 \text{ mins of CPU time}$$

$$\text{CPU Utilization} = \frac{12.5}{45} = 27.78\%$$

#### Multiprogramming:

CPU utilization is same as max time of uniprogramming, i.e., 90% of 10 min. = average effective time.

#### Memory Utilization:

Uniprogramming: The total memory used by all jobs is 400 Mbytes.

Multiprogramming: The total memory used by all jobs is  $100 + 150 + 125 = 375 \text{ Mbytes}$ .

$$\text{So, } 375 \text{ Mbytes is utilized.}$$

#### Throughput:

Uniprogramming: 1 job in 20 min.

Multiprogramming: 3 jobs in 20 min.



Q5. In a batch OS, 3 jobs are submitted for execution. Each job involves an I/O activity, & CPU time & another I/O activity of same time & spans 1st job.

Job 1 takes total 23 ms to complete, with 3ms for CPU

Job 2 takes total 29 ms to complete with 5ms for CPU

Job 3 takes total 14 ms to complete with 4ms for CPU.

Solu:

Uniprogramming:

$$\text{Total CPU time} = 3 + 5 + 4 = 12 \text{ ms}$$

$$\text{Total time} = 23 + 29 + 14 = 66 \text{ ms}$$

$$\text{CPU utilization} = \frac{12}{66} \times 100 = 18.18\%$$

Multiprogramming

$$\text{Total CPU time} = \max(3, 5, 4) = 12 \text{ ms}$$

$$\text{Total time} = \text{longest time of Job i.e., Job 2} = 29 \text{ ms}$$

$$\text{CPU utilization} = \frac{12}{29} \times 100 = 41.38\%$$

Q6. Trace following progr. segment & determine how many processes created

a. int main() {  
pid\_t c1, c2;  
c2=0;  
c1=fork();  
if(c1==0)  
c2=fork();  
if(c2>0)  
fork();  
printf("1");  
return 0;

b. int main() {  
pid\_t c1=1, c2=1;  
c1=fork();  
if(c1!=0)  
c2=fork();  
if(c2==0)  
fork();  
printf("1");  
return 0;

3

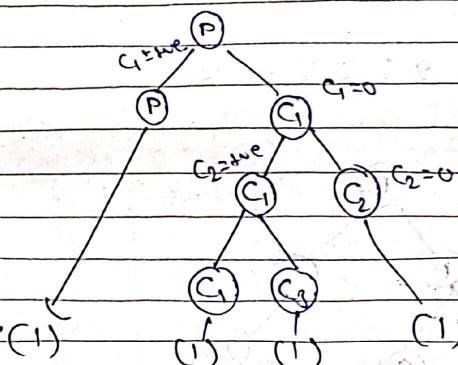
c. int main() {  
if(fork() || fork())  
fork();  
printf("1");  
return 0;

3

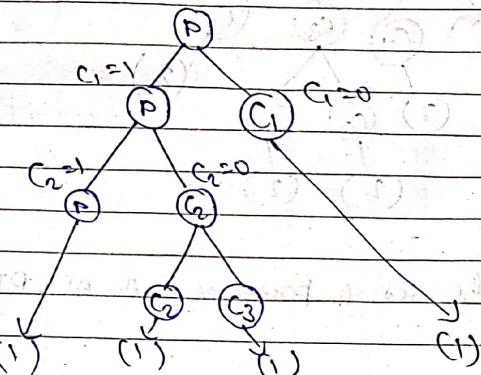
d. int main() {  
if(fork() > 0 || !fork())  
if(fork() || fork())  
fork();  
printf("2");  
return 0;

3

(a) Output: 1 1 1 1



(b) Output: 1 1 1 1



Q3. In a batch OS, 3 jobs are submitted for execution.  
Each job involves an I/O activity, & CPU time &  
overall I/O activity of same time spent on 1st job.

Job 1 takes total 29 ms to complete, with 18 ms for CPU  
Job 2 takes total 29 ms to complete with 14 ms for CPU  
Job 3 takes total 14 ms to complete with 8 ms for CPU.

Sol:

### Uniprogramming.

$$\text{Total CPU time} = 3 + 5 + 4 = 12 \text{ ms.}$$

$$\text{Total time} = 23 + 29 + 14 = 66 \text{ ms}$$

$$\text{CPU utilisation} = \frac{12}{66} \times 100 = 18.18\%$$

### Multiprogramming

$$\text{Total CPU time} = \text{longest time} = 12 \text{ ms}$$

$$\text{Total time} = \text{longest time of Job i.e., Job 2} \\ = 29 \text{ ms.}$$

$$\text{CPU utilisation} = \frac{12}{29} \times 100 = 41.38\%.$$

Q4. Trace following Prgm. segment & determine how many processes created

- a. int main() {

pid\_t c1, c2;

c1=0;

c1 = fork();

if (c1==0)

c2 = fork();

if (c2>0)

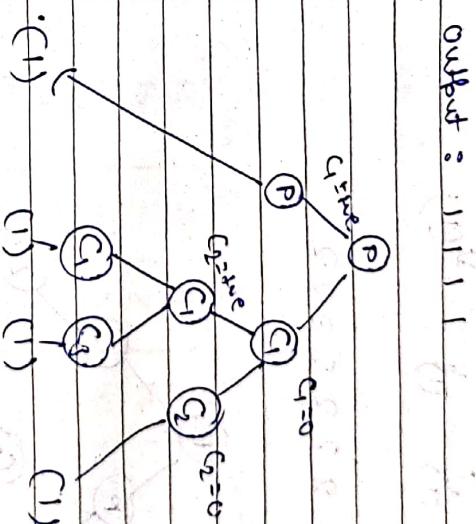
fork();

printf("%d",

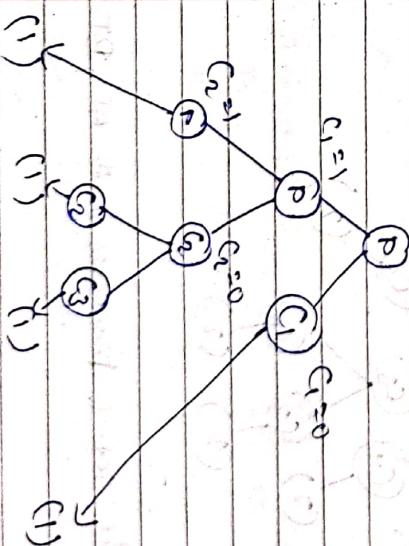
return 0;

3

(a) Output : 1 1 1 1



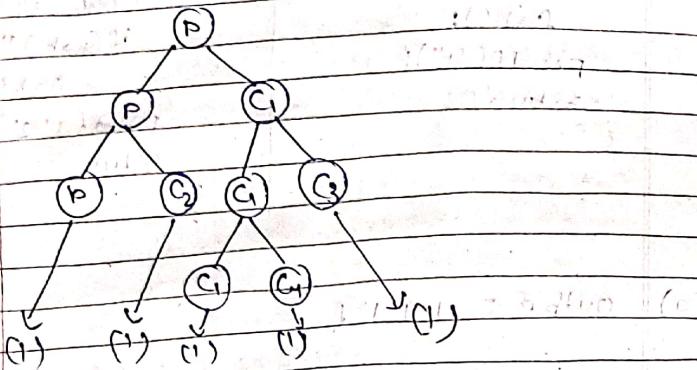
(b) Output: 1 1 1 1



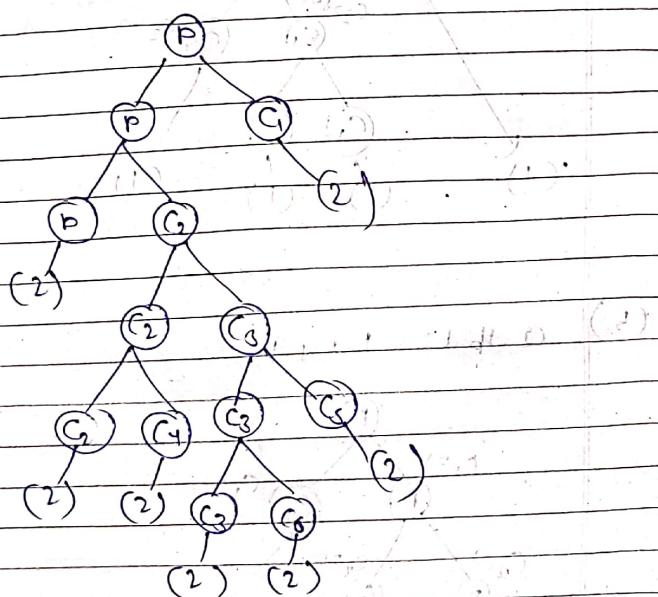
```

int main()
{
    if (fork() == 0)
        if (fork() == 0)
            if (fork() == 0)
                printf("1");
            else
                printf("2");
        else
            printf("3");
    else
        if (fork() == 0)
            if (fork() == 0)
                printf("1");
            else
                printf("2");
        else
            printf("3");
}
  
```

(c) Output: 11111



(d) Output: 2222222



Q7. Consider the set of processes with AT, BT &amp; Priority

Process	A.T.	B.T.	Priority
P <sub>1</sub>	0	11	5
P <sub>2</sub>	0	8	3
P <sub>3</sub>	12	2	4
P <sub>4</sub>	2	6	2
P <sub>5</sub>	9	16	1

Finding avg. TAT, avg. BWT &amp; avg. RT.

- FCFS
- SJF
- SRTF
- Priority (non-preemptive)
- Round robin (RR) - highest priority time

• FCRS:

Process	A.T.	B.T.	C.T.	T.A.T	W.T	R.T
P <sub>1</sub>	0	11	11	11	0	0
P <sub>2</sub>	0	8	19	19	8	11
P <sub>3</sub>	12	2	43	31	29	29
P <sub>4</sub>	2	6	25	23	17	17
P <sub>5</sub>	9	16	41	32	16	16

P<sub>1</sub> P<sub>2</sub> P<sub>4</sub> P<sub>5</sub> P<sub>3</sub>      avg TAT = 28.2 ms

avg W.T = 14.6 ms

avg R.T = 14.6 ms

• SJF

Process	A.T.	B.T.	C.T.	T.A.T	W.T.	R.S.
P <sub>1</sub>	0	12	27	27	16	16
P <sub>2</sub>	0	8	8	8	0	0
P <sub>3</sub>	12	2	16	4	2	2
P <sub>4</sub>	2	6	14	12	6	6
P <sub>5</sub>	9	16	43	34	18	18

P<sub>2</sub> P<sub>4</sub> P<sub>3</sub> P<sub>1</sub> P<sub>5</sub>      0 8 14 16 27 + 43 = 102 ms

avg TAT = 17 ms

avg W.T = 8.4 ms

avg R.T = 9.4 ms

## SRTF

Process	A.T.	W.T.	C.T.	T.A.T	W.T	R.T
P <sub>1</sub>	0	11	27	27	16	16
P <sub>2</sub>	0	8	8	8	0	0
P <sub>3</sub>	12	2	16	4	2	2
P <sub>4</sub>	2	6	14	12	6	6
P <sub>5</sub>	9	16	43	34	18	18

P <sub>2</sub>	P <sub>2</sub>	P <sub>4</sub>	P <sub>4</sub>	P <sub>3</sub>	P <sub>1</sub>	P <sub>3</sub>
0	2	8	9	12	14	16

$P_2=11$     $P_2=11$     $P_2=11$     $P_2=11$     $P_2=11$   
 $P_3=6$     $P_3=6$     $P_3=2$     $P_3=16$     $P_3=16$   
 $P_4=6$     $P_4=35$     $P_4=2$     $P_4=16$     $P_4=16$   
 $P_5=6$     $P_5=16$     $P_5=2$     $P_5=16$     $P_5=16$   
 $P_3=12$

avg TAT = 17 ms

avg WT = 8.4 ms

avg RT = 9.4 ms

## Non preemptive priority

Priority	Process	A.T.	W.T	C.T.	TAT	WT	R.T
1	P <sub>1</sub>	0	11	11	11	0	0
0	P <sub>2</sub>	0	8	43	43	35	
3	P <sub>3</sub>	12	2	29	29	15	15
2	P <sub>4</sub>	2	6	35	35	27	27
4	P <sub>5</sub>	9	16	27	18	2	2

P <sub>1</sub>	P <sub>5</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>2</sub>
0	11	27	29	35

avg TAT = 24.4 ms

avg WT = 15.8 ms

avg RT = 15.8 ms

## Preemptive Priority

Priority	Process	A.T	W.T	C.T	TAT	WT	RT
1	P <sub>1</sub>	0	11	35	35	24	0
0	P <sub>2</sub>	0	8	43	43	35	35
3	P <sub>3</sub>	12	2	27	15	13	13
2	P <sub>4</sub>	2	6	8	6	0	0
4	P <sub>5</sub>	9	16	25	16	0	0

P <sub>1</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>5</sub>	P <sub>3</sub>	P <sub>1</sub>	P <sub>2</sub>
0	2	8	9	12	25	35

avg TAT = 23 ms

avg WT = 14.4 ms

avg RT = 9.6 ms

## Round Robin

Process	Priority	A.T	W.T	C.T	TAT	WT	RT
P <sub>1</sub>	1	0	11	32	32	21	0
P <sub>2</sub>	0	0	8	28	28	20	5
P <sub>3</sub>	3	12	2	30	18	16	28
P <sub>4</sub>	2	2	6	31	29	23	10
P <sub>5</sub>	4	9	16	43	34	18	20

P <sub>1</sub>	P <sub>2</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>5</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>5</sub>
0	5	10	15	20	25	28	30	31	32

avg TAT = 128.2 ms

avg WT = 19.6 ms

avg RT = 12.6 ms

## Highest response ratio

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>5</sub>
0	10	10	10	10
1	10	10	10	10
2	2	3	3	3
3	1	1	1	1



Processes	A.T.	B.T.	C.T.	T.A.T	W.T	R.T
P <sub>1</sub>	0	4	4	4	0	0
P <sub>2</sub>	0	2	6	6	4	4
P <sub>3</sub>	1	3	9	8	5	5
P <sub>4</sub>	2	2	11	9	7	7
<b>B</b>						

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>5</sub>
0	11	17	19	27 : 43

$$P_2 = \frac{11+8}{8} = \frac{19}{8}$$

$$P_3 = \frac{9+2}{6} = \frac{11}{6}$$

$$P_4 = \frac{2+16}{16} = \frac{18}{16}$$

$$\text{avg TAT} = 18.8$$

$$\text{avg W.T.} = 10.2$$

- Q8. Consider the set of processes with arrival time, B.T, Priority (lowest no. implies high priority)
- | Processes      | A.T. | B.T. | Priority |
|----------------|------|------|----------|
| P <sub>1</sub> | 0    | 4    | 3        |
| P <sub>2</sub> | 0    | 2    | 1        |
| P <sub>3</sub> | 1    | 3    | 2        |
| P <sub>4</sub> | 2    | 2    | 4        |

Find avg TAT, avg WT & avg RT

- FCFS • SJF • SRTF • Non preemptive priority
- Preemptive priority • Highest response ratio
- Round robin (q=2 ms)

→ FCFS

P	A.T.	B.T.	C.T.	T.A.T	W.T	R.T
P <sub>1</sub>	0	4	4	4	0	0
P <sub>2</sub>	0	2	6	6	4	4
P <sub>3</sub>	1	3	9	8	5	5
P <sub>4</sub>	2	2	11	9	7	7

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>
0	4	6	9

$$\text{avg TAT} = 6.75 \text{ ms}$$

$$\text{avg WT} = 4 \text{ ms}$$

$$\text{avg RT} = 4 \text{ ms}$$

→ SJF

P	C.T.	T.A.T	W.T	R.T
P <sub>1</sub>	11	11	7	7
P <sub>2</sub>	2	2	0	0
P <sub>3</sub>	7	6	3	3
P <sub>4</sub>	4	2	0	0

$$\text{avg TAT} = 5.25 \text{ ms}$$

$$\text{avg WT} = 2.5 \text{ ms}$$

$$\text{avg RT} = 2.5 \text{ ms}$$

→ SRTF

P	C.T.	TAT	W.T	R.T
P <sub>1</sub>	11	11	7	7
P <sub>2</sub>	2	2	0	0
P <sub>3</sub>	7	6	3	3
P <sub>4</sub>	4	2	0	0

$$\text{avg TAT} = 5.25 \text{ ms}$$

$$\text{avg WT} = 2.5 \text{ ms}$$

$$\text{avg RT} = 2.5 \text{ ms}$$

→ Non Preemptive Priority

P <sub>2</sub>	P <sub>3</sub>	P <sub>1</sub>	P <sub>4</sub>
0	2	5	9

$$\text{avg TAT} = 6 \text{ ms}$$

$$\text{avg WT} = 2.25 \text{ ms}$$

$$\text{avg RT} = 3.25 \text{ ms}$$

Priority	P	C.T.	T.A.T	W.T.	R.T.
3	P <sub>1</sub>	9	9	5	5
1	P <sub>2</sub>	2	2	0	0
2	P <sub>3</sub>	5	4	1	1
4	P <sub>4</sub>	11	9	7	7

→ Preemptive Priority.

P	C.T.	T.A.T	W.T.	R.T.
P <sub>1</sub>	9	9	5	5
P <sub>2</sub>	2	2	0	0
P <sub>3</sub>	5	4	1	1
P <sub>4</sub>	11	9	7	7

$$\text{avg TAT} = 6 \text{ ms}$$

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	A	P <sub>4</sub>
0	1	2	4	5	9 11

$$\text{avg WT} = 3.25 \text{ ms}$$

$$\text{avg RT} = 3.25 \text{ ms}$$

→ Highest response ratio

P	A.T.	B.T.	TAT	T.A.T	W.T.	R.T.
P <sub>1</sub>	0	4	4	4	0	
P <sub>2</sub>	0	2	6	6	4	
P <sub>3</sub>	1	3	11	10	7	
P <sub>4</sub>	2	2	8	6	4	

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>
0	4	6	8 11

$$\text{avg TAT} = 6.5 \text{ ms}$$

$$\text{avg WT} = 3.75 \text{ ms}$$

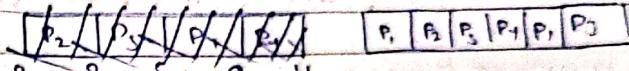
$$P_1 = \frac{4+2}{2} = \frac{6}{2} = 3$$

$$P_2 = \frac{3+2}{2} = \frac{5}{2}$$

$$P_3 = \frac{2+2}{2} = \frac{4}{2} = 2$$

→ Round robin ( $q=2 \text{ ms}$ )

Priority	P	A.T.	B.T.	C.I.	T.A.T.	W.T.	R.T.
3	P <sub>1</sub>	0	4	10	10	6	
1	P <sub>2</sub>	0	2	4	4	2	
2	P <sub>3</sub>	1	3	11	10	7	
4	P <sub>4</sub>	2	2	8	6	4	



P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>
0	2	4	6	8	10	11	

$$P_1 = 2, P_2 = 2, P_3 = 2, P_4 = 2$$

$$P_2 = 2, P_3 = 3, P_4 = 1, P_2 = 1$$

$$P_3 = 3, P_4 = 2, P_4 = 2$$

$$P_4 = 2$$

Q9. Consider the following 4 processes with arrival time & their length of CPU bursts

P	A.T.	B.T.
P <sub>1</sub>	0	3
P <sub>2</sub>	1	1
P <sub>3</sub>	3	3
P <sub>4</sub>	4	2

Find n. Such

avg WT of

processes 1 ms

using SRTF.

Sol:

Let  $n=1$

P	A.T.	B.T.	C.T.	TAT	W.T.	R.T.
P <sub>1</sub>	0	3	4	4	1	0
P <sub>2</sub>	1	1	2	1	0	0
P <sub>3</sub>	3	3	5.8	5	2	2
P <sub>4</sub>	4	1	5	1	0	0

$$\text{avg WT} = \frac{3}{4} = 0.75$$

Let  $n=2$ .

P	A.T.	B.T.	C.T.	T.A.T	WT	n.t.
P <sub>1</sub>	0	3	4	4	1	0
P <sub>2</sub>	1	1	2	1	0	0



Scanned with OKEN Scanner

P	A.T	B.T	C.T	TAT	W.T	R.T
P <sub>3</sub>	3	3	9	6	3	3
P <sub>4</sub>	4	2	6	2	0	0
	4	4	8	14	10	10

avg. WT =  $\frac{10}{4} = 2.5 \text{ ms}$   
when n=2, avg. waiting time is 1ms.

Q10. Multilevel feedback queue scheduling (MLFQ) with q<sub>1</sub>, q<sub>2</sub>, q<sub>3</sub>. q<sub>1</sub> has highest priority & q<sub>3</sub> has lowest priority. q<sub>1</sub> & q<sub>2</sub> use round robin with time = 8 ms  
ms. q<sub>3</sub> follows FCFS. Find avg. W.T. & avg T.A.T.

Process	Burst T.	A.T	C.T	T.A.T	W.T	R.T
P <sub>1</sub>	8	0	17	17	9	0
P <sub>2</sub>	22	0	42	42	20	3
P <sub>3</sub>	4	0	23	23	19	6
P <sub>4</sub>	12	0	46	46	34	9

Q <sub>1</sub>	P <sub>1</sub>   P <sub>2</sub>   P <sub>3</sub>   P <sub>4</sub>
Q <sub>2</sub>	P <sub>1</sub>   P <sub>2</sub>   P <sub>3</sub>   P <sub>4</sub>

avg T.A.T = 24.5 ms

Q <sub>3</sub>	P <sub>2</sub>   P <sub>4</sub>
	avg WT = 20.5 ms

avg RT = 4.5 ms

P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>	P <sub>1</sub>	P <sub>2</sub>	P <sub>3</sub>	P <sub>4</sub>
0	3	6	9	12	17	22	23

avg T.A.T = 24.5 ms

avg WT = 20.5 ms

avg RT = 4.5 ms

Q11. Two process P & Q. Synchronize the execution process P & Q with suitable Semaphore operation using 2 boolean semaphores S & T. Sequence will be...

```

Process P          Process Q
while(true){       while(true){
    wait(S);        wait(T);
    print("1");     print("1");
    print("1");     print("0");
    signal(T);      signal("0");
}}                  signal("S");

```

initial value S = false  
T = true

Q12. Two concurrent Process P & Q are executing following inst. Synchronize execution of P & Q with suitable Semaphore operation using 2 boolean semaphore S & T.

Process P	Process Q
print("1")	print("0")
print("2")	process("0")

(a) 1 3 2 4

Process P	Process Q
wait(S)	wait(T);
print("1");	print("3");
signal(T);	signal(Q);
wait(S);	wait(T);
print("2");	print("4");
signal(T);	signal(S);

S = True, T = False.

(b) 3 1 2 4

Process P	Process Q
wait(S);	wait(T);
print("1");	print("3");
print("2");	print("2");
signal(T);	signal(Q);
wait(S);	wait(T);
print("4");	print("1");
signal(T);	signal(S);

S = False, T = True.