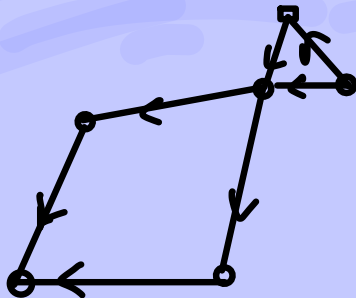


Digraph immersions

Shoham Letzter

UCL

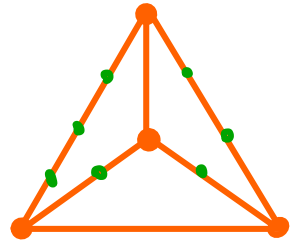


Joint work with António Girão

Clique subdivisions

Bollobás-Thomason / Komlós-Szemerédi '96:

$\exists c > 0$ s.t. if G has average ^{min} deg $\geq ct^2$ then G has a subdivision of K_t .



Tight (up to a constant factor): $G = K_{t,t}$

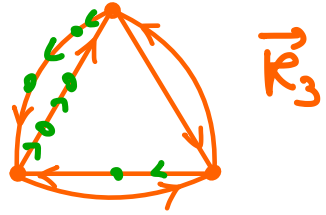
$$\left. \begin{array}{c} t/2 \leq \\ \Rightarrow c \gtrsim \frac{1}{8} \end{array} \right\} \geq \binom{t/2}{2} \sim \frac{t^2}{8}$$

Best known bounds on least c :

$$\begin{array}{ll} \geq \frac{9}{64} & (\text{\u0179uczak}) \\ \leq \frac{10}{23} & (\text{K\u00fch\u00f1-Osthus '06}) \end{array}$$

What about digraphs?

\vec{K}_t complete digraph on t vertices.



① Is there $f(t)$ st.: if G has $\min \text{ in \& out-deg} \geq f(t)$
then G contains a subdivision of \vec{K}_t ?

- Yes for $t=2$. $\min \text{ out-deg} \geq 1 \Rightarrow \text{subdivision of } \vec{K}_2$. $f(2)=1$
- No for $t \geq 3$

(Mader '85 using construction of Thomassen 85'
DeVos-McDonald-Mohar-Scheidt '12).

$\mathcal{T}\mathcal{T}_t$ transitive tournament on t vertices.

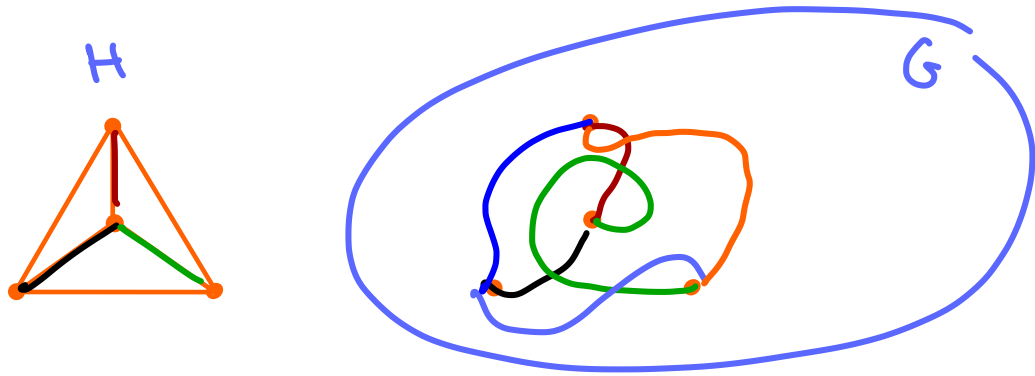


- ② Is there $f(t)$ s.t.: if G has min in & out $\deg \geq f(t)$ then G contains a subdivision of $\mathcal{T}\mathcal{T}_t$?

This is (almost) a conjecture of Mader '96.

Immersions

G immerses H if \exists injection $f: V(H) \rightarrow V(G)$ and edge-disjoint paths P_{uv} , for $uv \in E(H)$, s.t. P_{uv} starts at $f(u)$ and ends at $f(v)$.



DeVos - Dvořák - Fox - McDonald - Mohar - Scheide '14:

If G has average $\deg \geq 200t$ then G immerses K_t . K_{t-1} does not immerse K_t .

Dvořák - Yepremyan '17: $\min \deg \geq 11t + 7 \Rightarrow$ immersion of K_t .

Hong - Wang - Yang '20: average $\deg \geq (1+\varepsilon)t$ & H -free for H bipartite \Rightarrow immersion of K_t .

Conjecture (Lescure - Meyniel '89): G immerses $K_{\chi(G)}$.

Immersion in digraphs

③ Is there $f(t)$ s.t.: if G has min in & out-deg $\geq f(t)$ then G immerses \vec{K}_t ?

No for $t \geq 3$ (DMMS '12).

Lochet '19: $\exists f(t)$ s.t.: if G has min out-deg $\geq f(t)$ then it immerses \vec{K}_t .

Conceivable that can take $f(t) = \underbrace{ct}_3$.

G is Eulerian if $d^+(u) = d^-(u)$ for every vertex u .

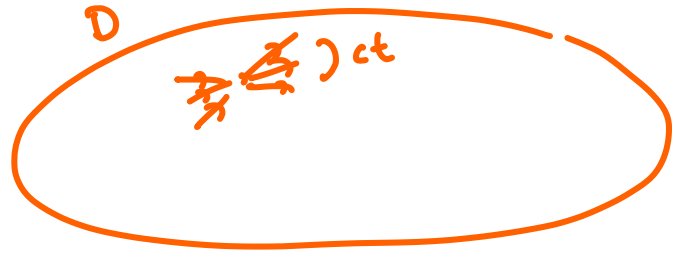


DeVos-McDonald-Mohar-Scheide '12: If G is Eulerian with $\min \text{out-deg} \geq t^2$ then it immerses \vec{K}_t .

Thm (Girão-L. 22+). $\exists c > 0$ s.t. if G is Eulerian with $\min \text{out-deg}$ at least ct then G immerses \vec{K}_t .

Overview of the proof.

I) Let c be a large constant.



Lemma. D Eulerian with $\min \text{out-deg} \geq ct \Rightarrow D$ immerses a digraph G with $\Theta(t)$ vertices and $\Omega(t^2)$ edges.

We use 'sparse expanders'.

* Introduced by Rónyai-Szemerédi '96.

* Can be found in graphs with average deg at least a large constant.

$$G \quad d(G) \geq 100$$



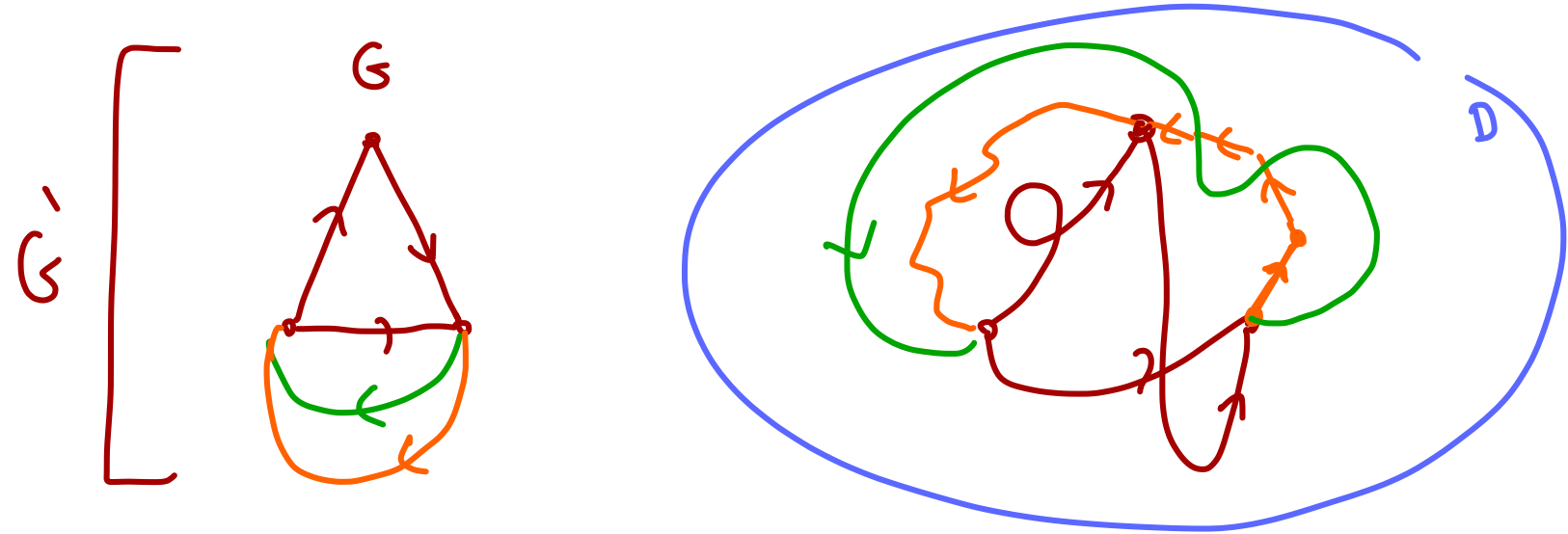
* many recent applications:

- odd cycle problem (Liu-Montgomery '20+)
- clique subdivisions in C_4 -free graphs (Liu-Montgomery '17)
- Kozlós conjecture on Hamiltonian sets
(Kim-Liu-Sharifzadeh-Staden '17)



Our proof is a rare use of expanders in digraphs
(Eulerian + immersion help).

II) Observation: If D is Eulerian and immerses G then it immerses an Eulerian multidigraph $G' \geq G$ with $V(G') = V(G)$.



III) Lemma. If G' is an Eulerian multidigraph on n vertices whose underlying graph (obtained by removing directions and multiplicities) has $\min \deg \geq \alpha n$, then it immerses \vec{K}_s , where $s = c'\alpha^{-4}n$.

I) + II) + III) \Rightarrow theorem.

D Eulerian, $\delta^+ \geq ct$.

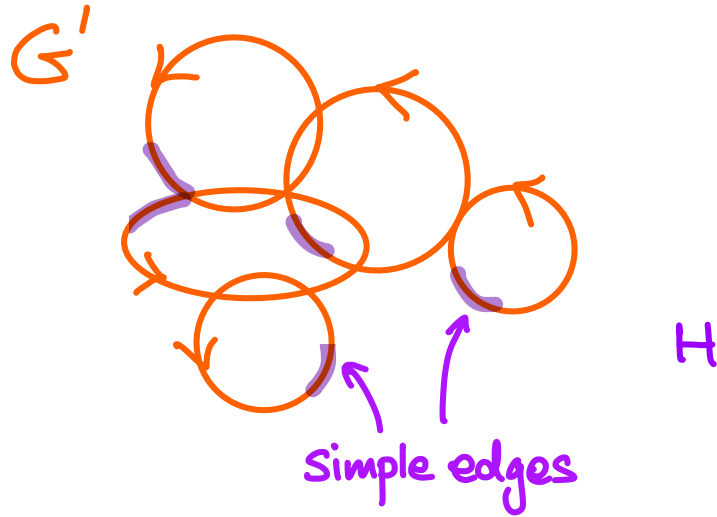
I) $D \rightsquigarrow G$, G has $\Theta(t)$ vertices, $\geq c't^2$ edges.
 \downarrow
 immersion

II) $D \rightsquigarrow G'$, G' Eulerian, multidigraph $G' \supseteq G$, $V(G') = V(G)$

III) $G' \rightsquigarrow \vec{K}_t$. $\Rightarrow D \rightsquigarrow \vec{K}_t$.

More about III

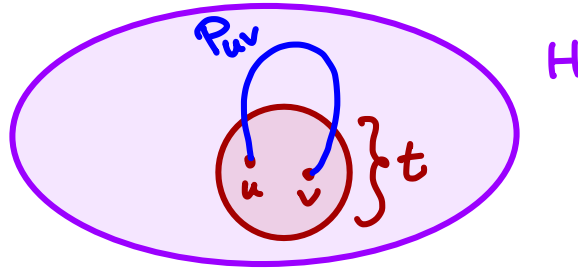
We will find a collection \mathcal{C} of $\Omega(n^2)$ edge-disjoint dicycles in G' each containing an edge which is simple in $\bigcup_{C \in \mathcal{C}} C$.



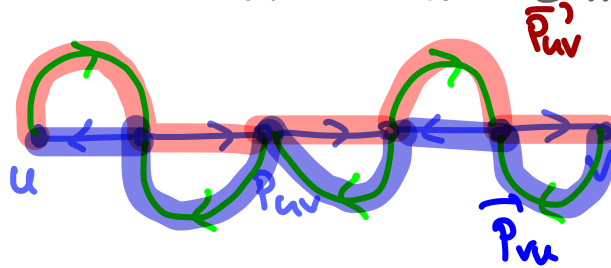
Let H be the undirected graph formed by the **purple** (simple) edges.

Then H has average $\deg \Omega(n)$.

Thus, by DDFMMS '14: H immerses K_t , where $t = \Omega(n)$.



Each P_{uv} corresponds to paths $u \rightarrow v$ and $v \rightarrow u$ in G' that are edge-disjoint.



$\Rightarrow G'$ immerses \vec{K}_t .

Open problems

① What is $\min f(t)$ s.t. if $\delta^+(G) \geq f(t)$ then G immerses \mathcal{H}_t ?

Lochet '19: $f(t) = O(t^3)$.

Maybe $f(t) = O(t)$?

② Mader '96: Is there $g(t)$ s.t. if $\delta^+(G) \geq g(t)$ then G contains a subdivision of \mathcal{H}_t ?

③ Conjecture (Lescure-Meyniel '89): G immerses $\mathcal{K}_{\chi(G)}$.

Finding \mathcal{C}

Lemma 1. \mathcal{D} multigraph on n vertices with $\min \text{out-deg} \geq \alpha n$.
Then \exists cycle with $\leq \frac{4}{\alpha}$ simple edges.

Apply Lemma 1 repeatedly to find \mathcal{C} .

preprocessing of G' that ensures that we don't have a dipath of length 2 with two multiple edges.



immersion .

G immerses H iff H can be obtained from G by:

* delete edge / vx

* or replace a path uvw by uw .



Proof of Lemma 1

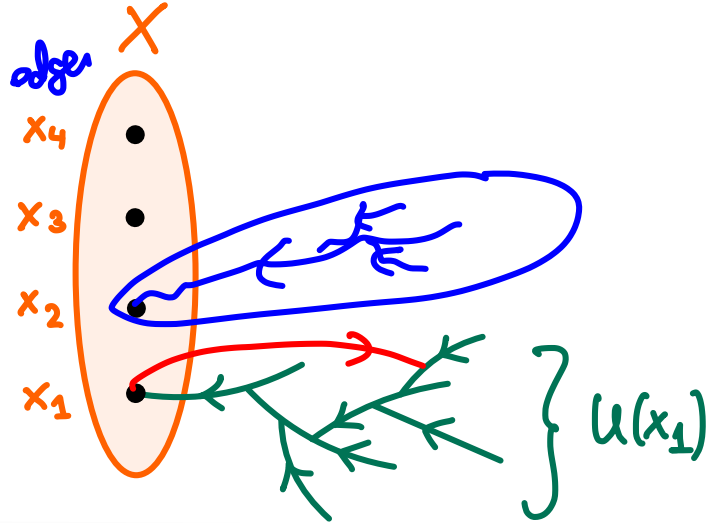
$$\delta^+(u) \geq 2n$$

Lemma 2. D digraph, $\omega: V(D) \rightarrow \mathbb{R}^+$. If $\omega(N^+(u)) \geq \alpha \cdot \omega(V(D))$ then \exists dicycle of length $\leq \frac{4}{\alpha}$.

$D' \subseteq D$ simple subdigraph, $E(D') = \{xy : xy \text{ is a multiple edge in } D\}$.

$X = \{x \in V(D') : d_{D'}^+(x) = 0\}$. $X = \emptyset \Rightarrow \exists$ cycle w/ no simple

Find sets $U(x)$, $x \in X$, as follows:



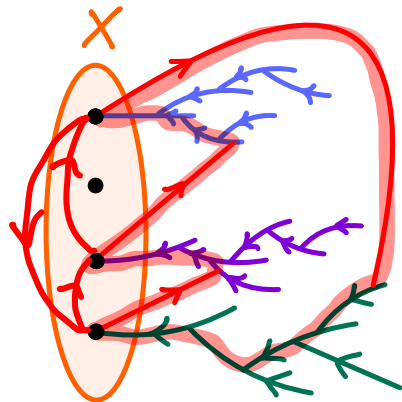
If $\exists \text{ edge } x \xrightarrow{\mathcal{D}} u(x)$, then \exists dicycle in \mathcal{D} with ≤ 1 simple edge.
 Suppose \nexists no such edges.

\mathcal{D}_0 digraph on X with $x \rightarrow y$ iff $\exists \text{ edge } x \xrightarrow{\mathcal{D}} u(y)$.

Can check: $w(N_{\mathcal{D}_0}^+(x)) \geq \alpha w(X) \quad \forall x \in X$, where $w(x) = |u(x)|$.

By Lemma 2: \exists dicycle of length $\leq \frac{4}{\alpha}$ in \mathcal{D}_0 .

$\Rightarrow \exists$ dicycle in \mathcal{D} with $\leq \frac{4}{\alpha}$ simple edges.



$$\delta^+ \geq ct$$

$$c_1 t \leq d^+(u) \leq c_2 t$$

