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What is Cryptography?

Cryptography derived its name from a Greek word called "Kryptos" which means "Hidden Secrets".

1. Cryptography is the practice and study of hiding information.
2. It provides Confidentiality, Integrity, Accuracy

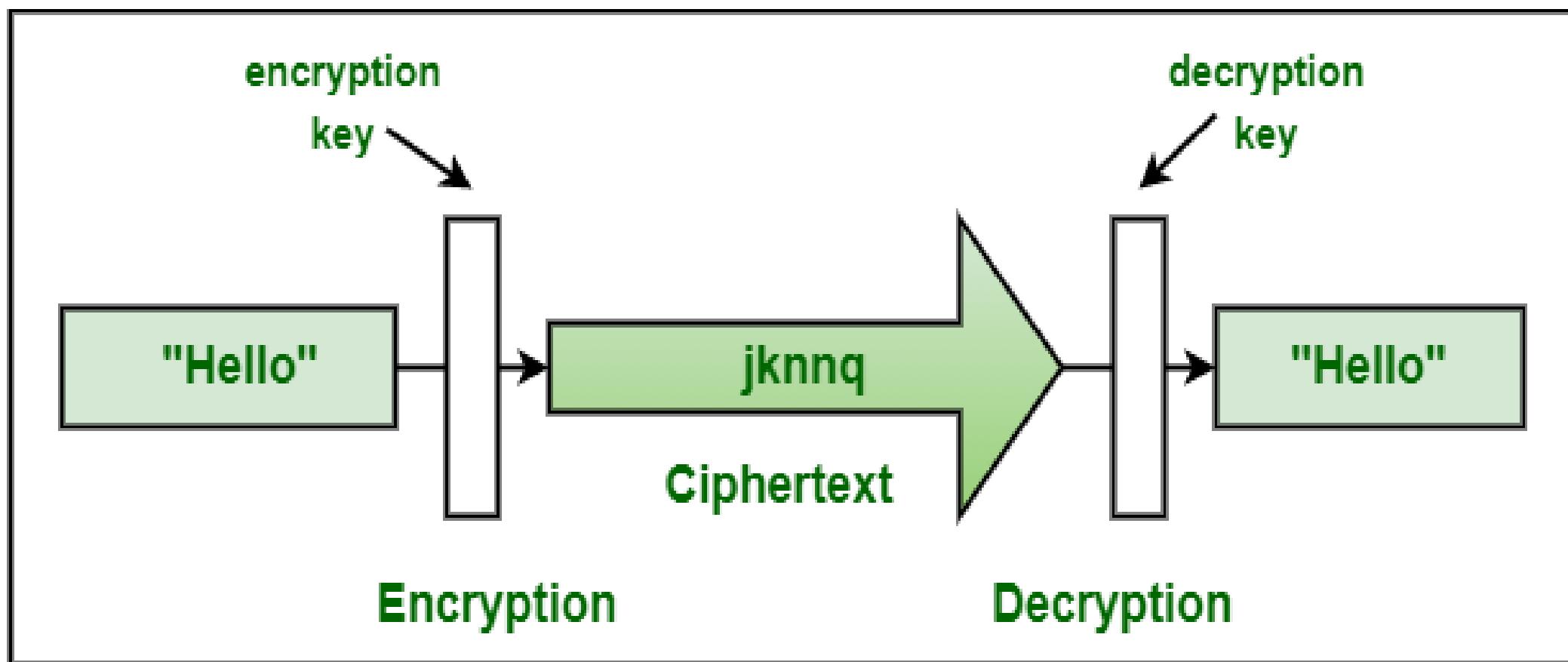
Cryptography

Cryptography is process to secure your information from unauthorise access.

Encrypted **501393c26075af84b9e612b5d90b4cff790ce10c**

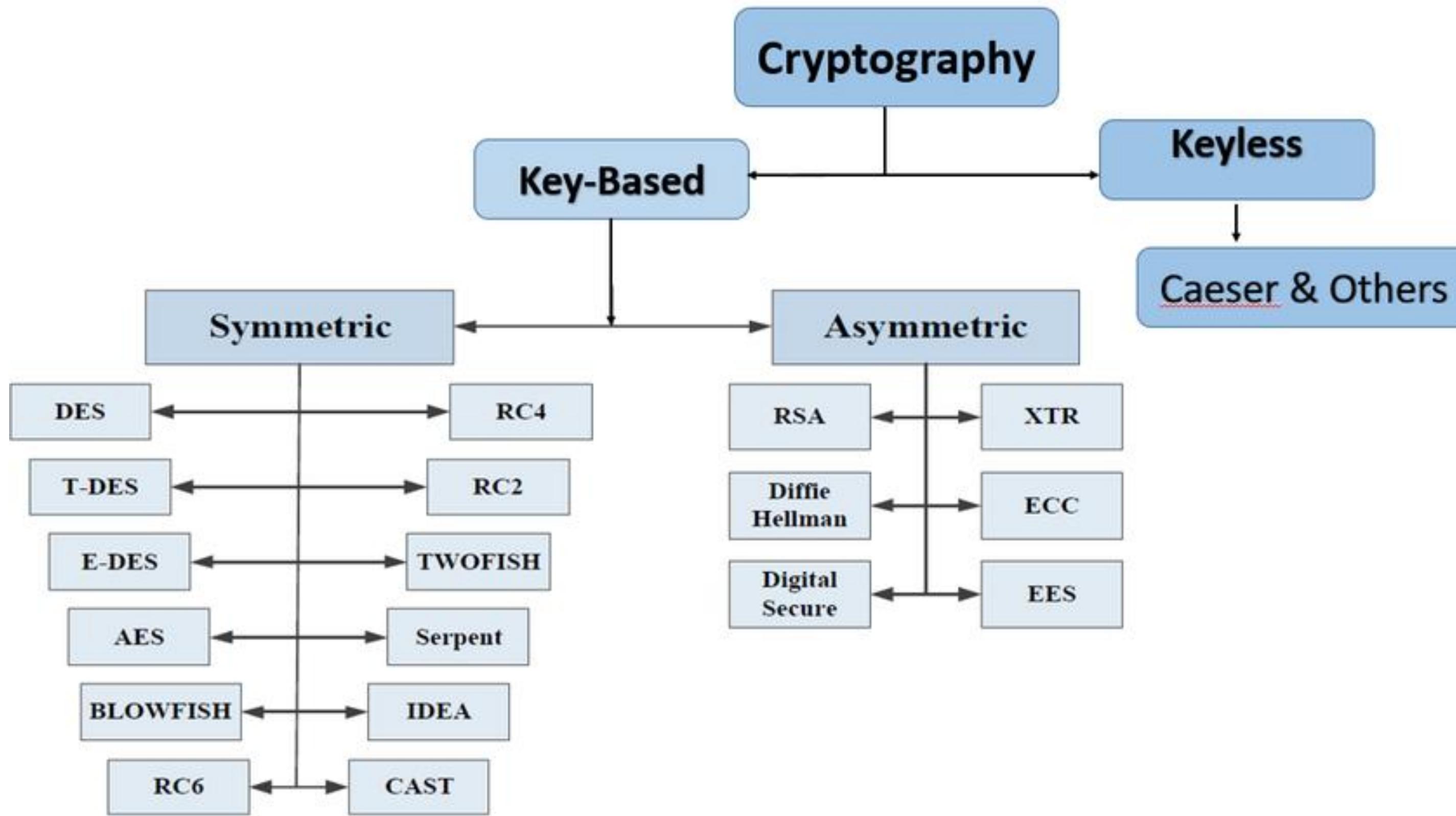
Decrypted **Hello_NSU_Hackers**

Encryption & Decryption



Cryptography

What are the Types of Cryptography



Some Common Cryptography

Example

- Caesar Cipher
- Rot 13
- Vigenere Cipher
- Morse code
- Bacon Cipher
- Alphabetical substitution

Tools

- Cyberchef
- dcode.fr
- boxentriq.com
- cryptii.com

Symmetric Cryptography

Symmetric cryptography is type of cryptography where both sender and receiver use same private key.

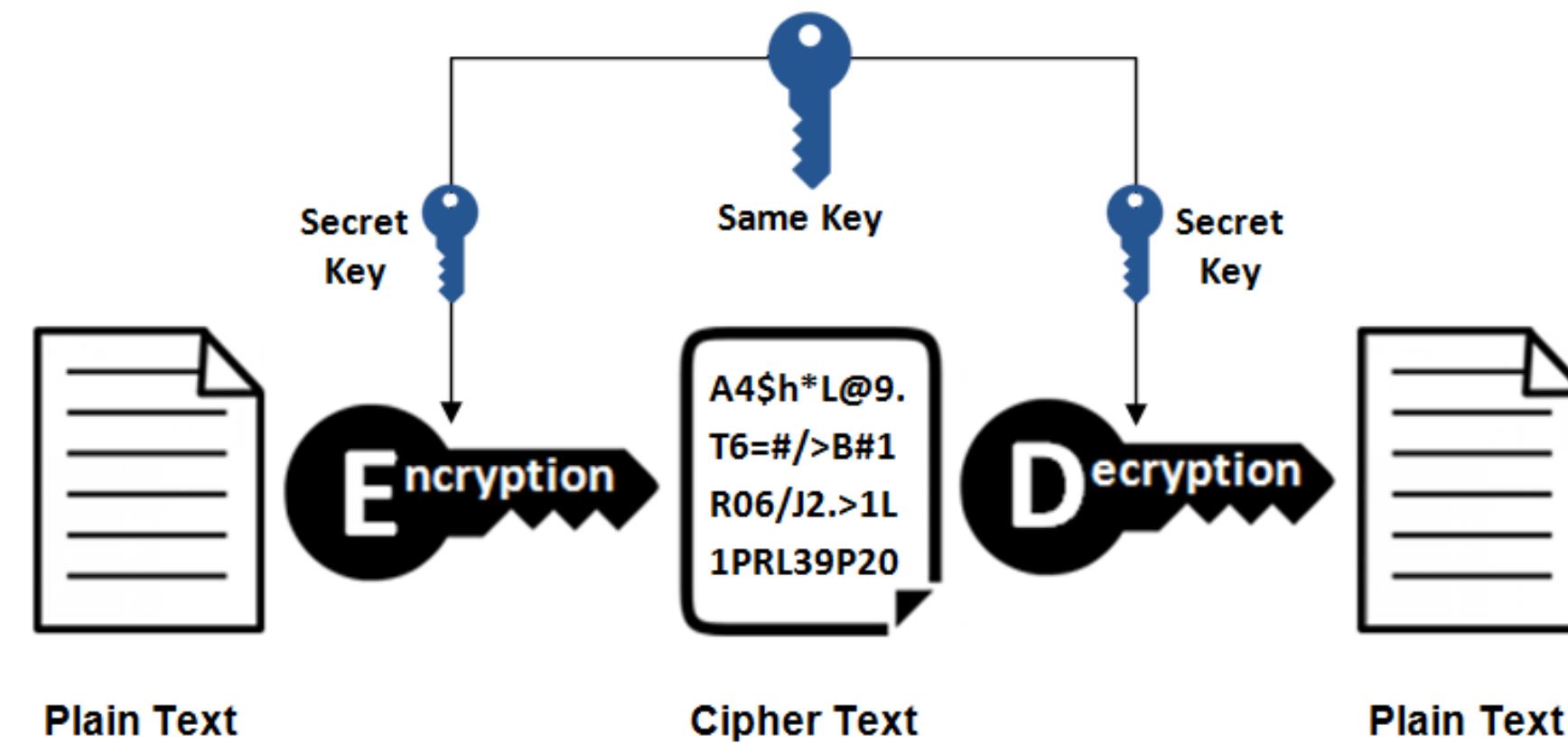
Symmetric Encryption

Advantage:

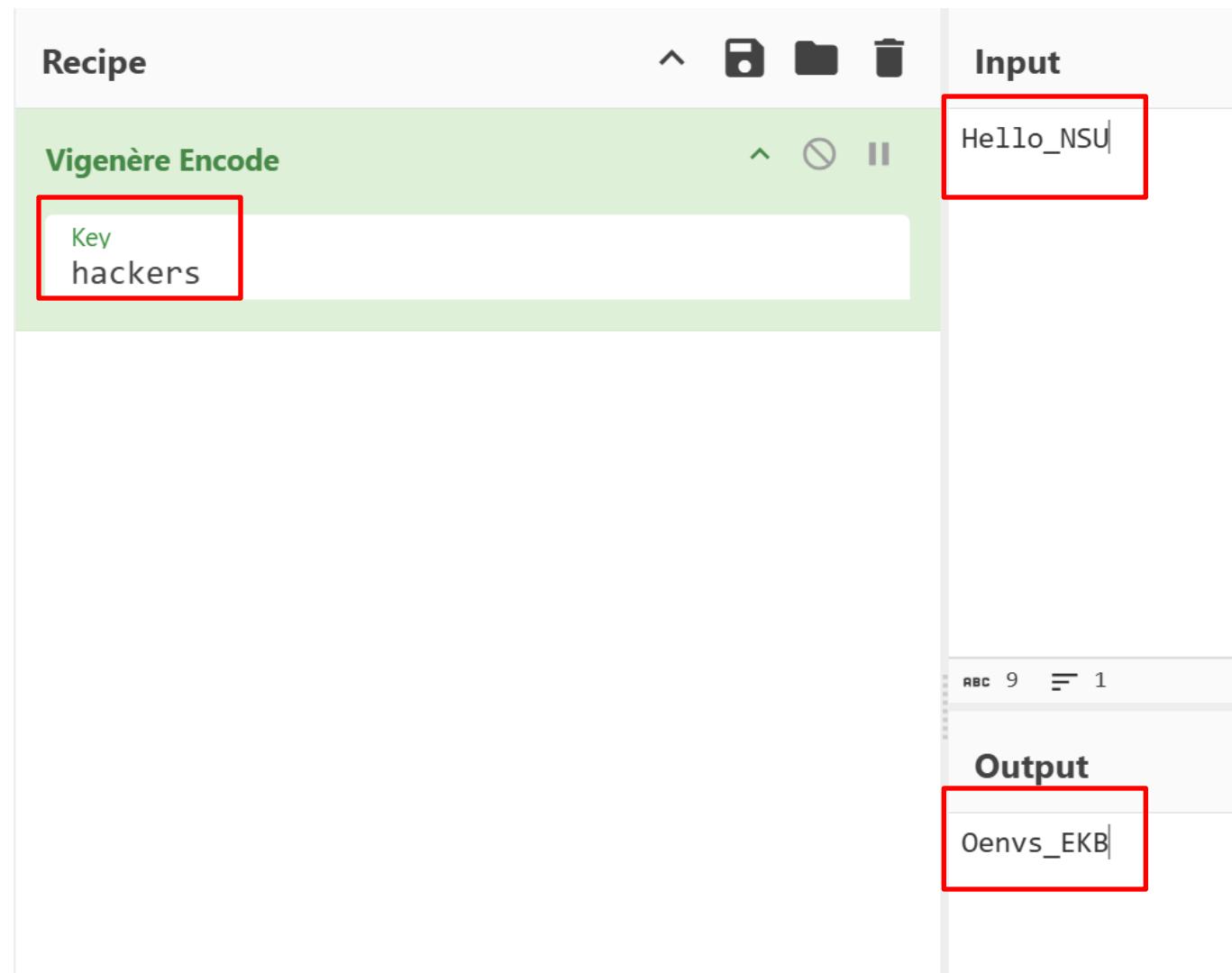
Simpler and Faster

Disadvantage:

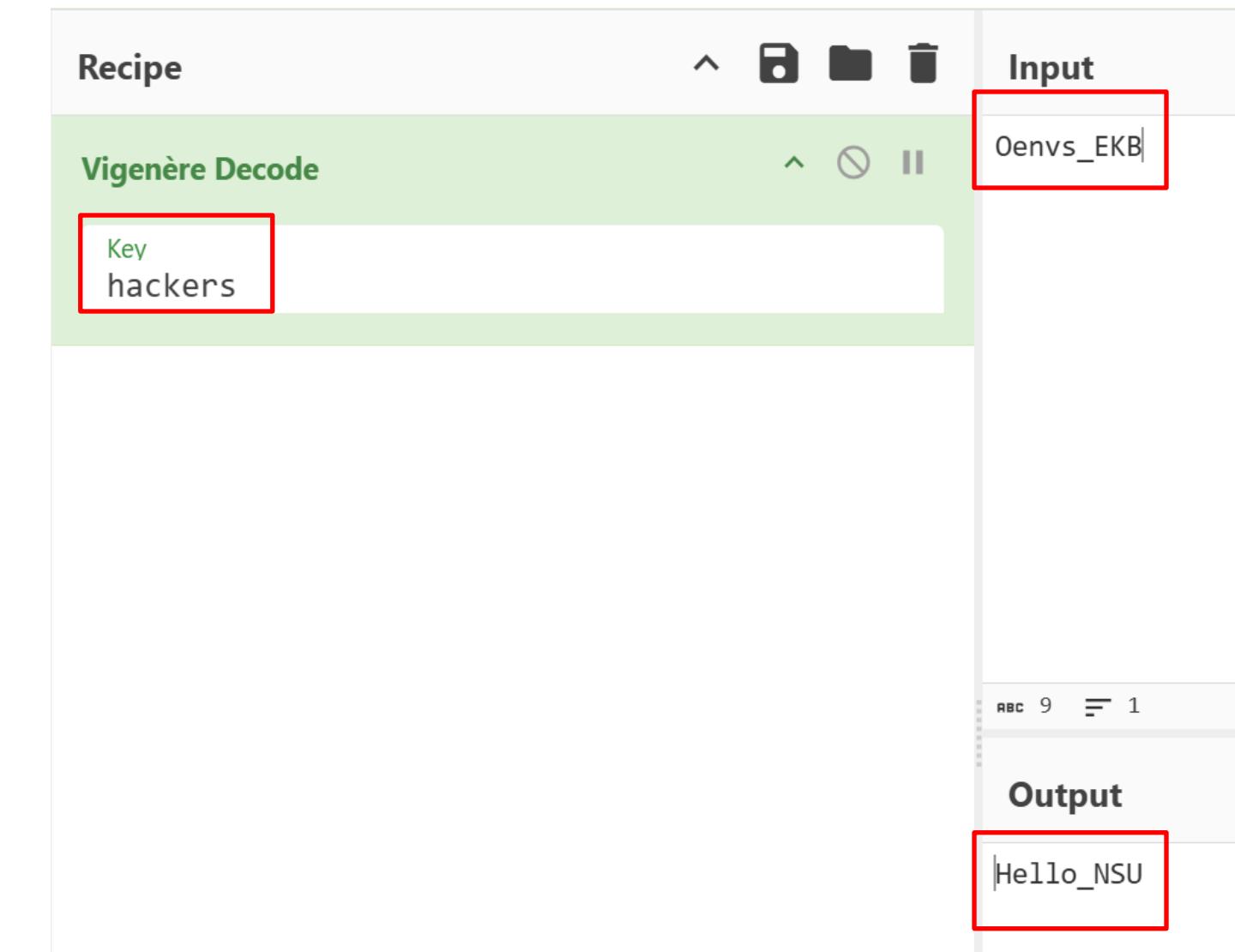
Less Secured



Symmetric Cryptography Example



Encrypt



Decrypt

Asymmetric Cryptography

Use public and private key:

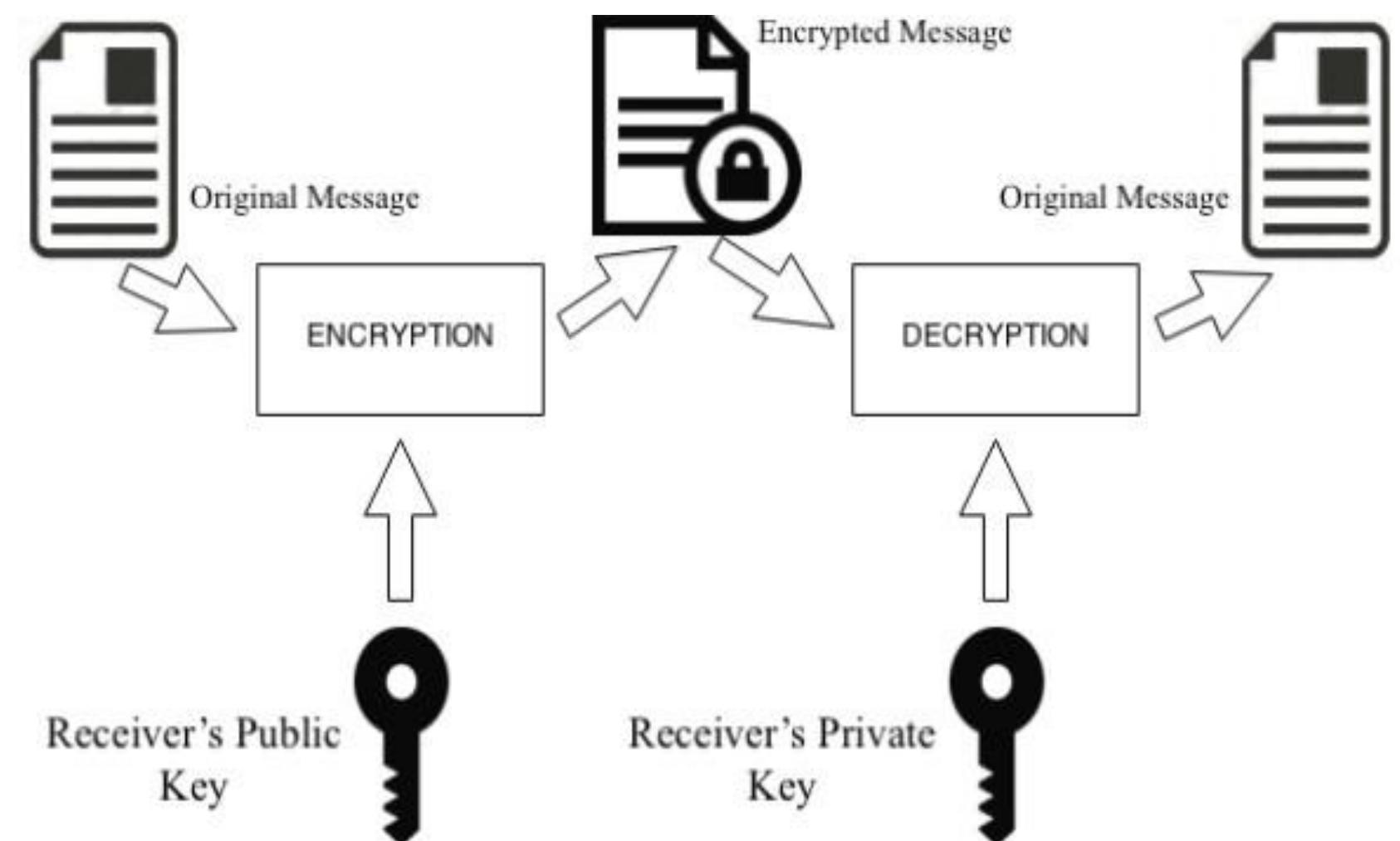
- Public key use for encryption
- Private key use for decryption

Advantage:

1. More Secured
2. Authentication

Disadvantage:

1. Relatively Complex



Asymmetric Cryptography

Example

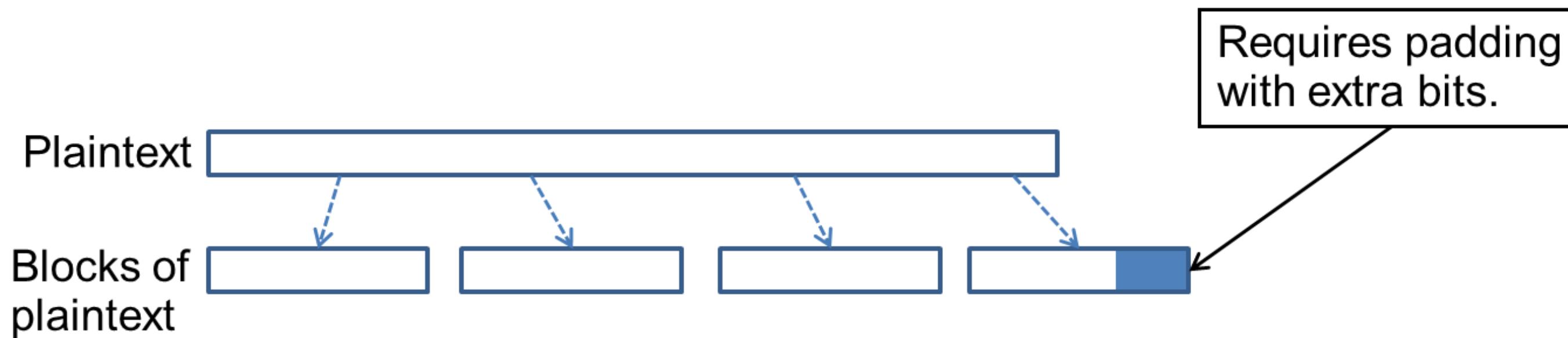
- RSA Algorithm
- Diffie-Hellman
- Elliptic Curve (ECC)

Tools

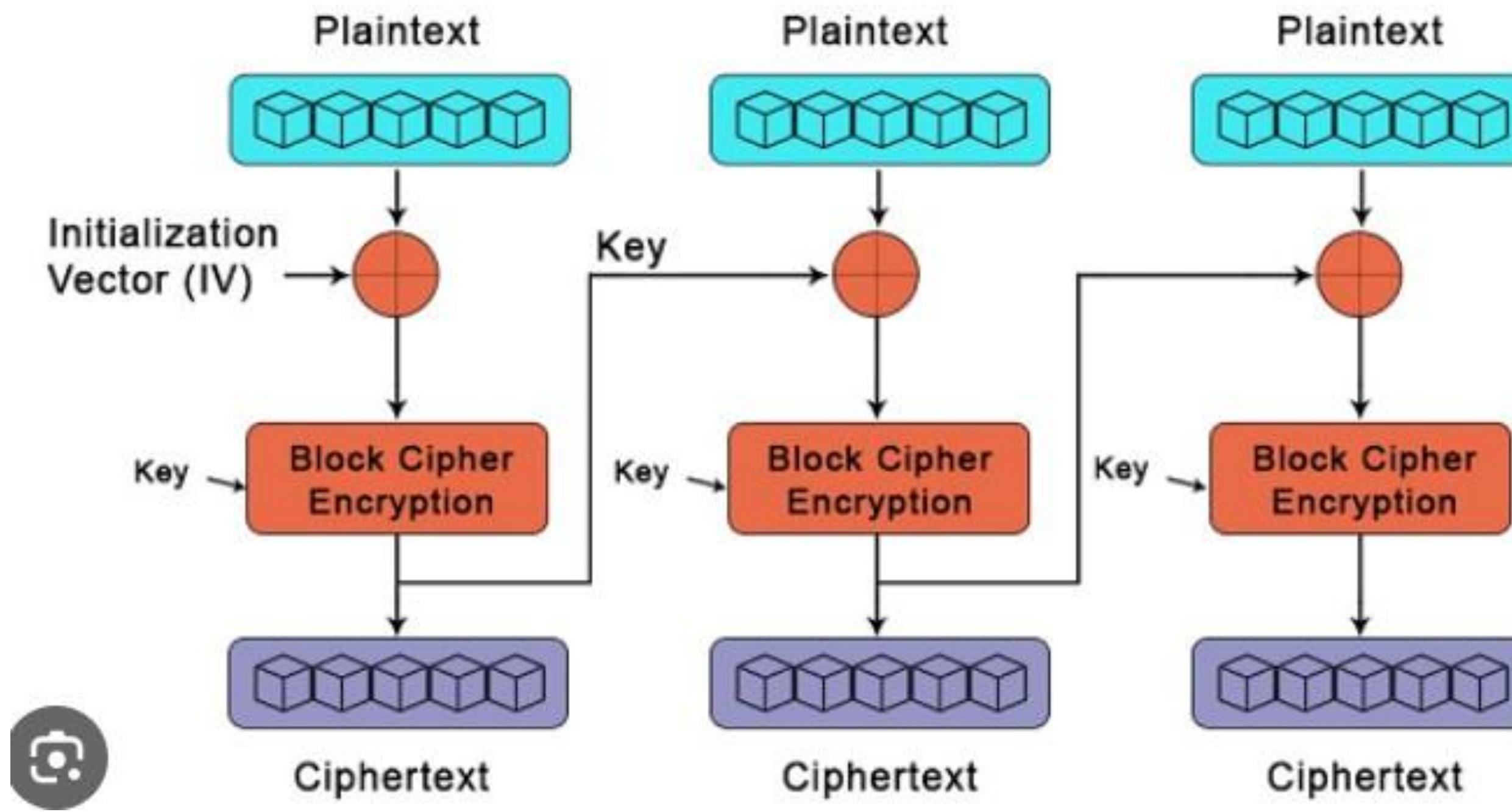
- Python code
- openssl (Default on Kali Linux OS)
- Online Tools

Block Ciphers

- In a **block cipher**:
 - Plaintext and ciphertext have fixed length b (e.g., 128 bits)
 - A plaintext of length n is partitioned into a sequence of m **blocks**, $P[0], \dots, P[m-1]$, where $n \leq bm < n + b$
- Each message is divided into a sequence of blocks and encrypted or decrypted in terms of its blocks.



Block Cipher



Paddings

- Block ciphers require the length n of the plaintext to be a multiple of the block size b
- Padding the last block needs to be unambiguous (cannot just add zeroes)

Where are block ciphers used

Data Encryption Standard (DES) →

64-bit blocks &
56-bit keys

Triple DES (3DES)

- Nested application of DES with three different keys KA, KB, and KC
- Effective key length is 168 bits, making exhaustive search attacks unfeasible

Advanced Encryption Standard (AES) →

128-bit blocks
key lengths: 128, 192 & 256
bits

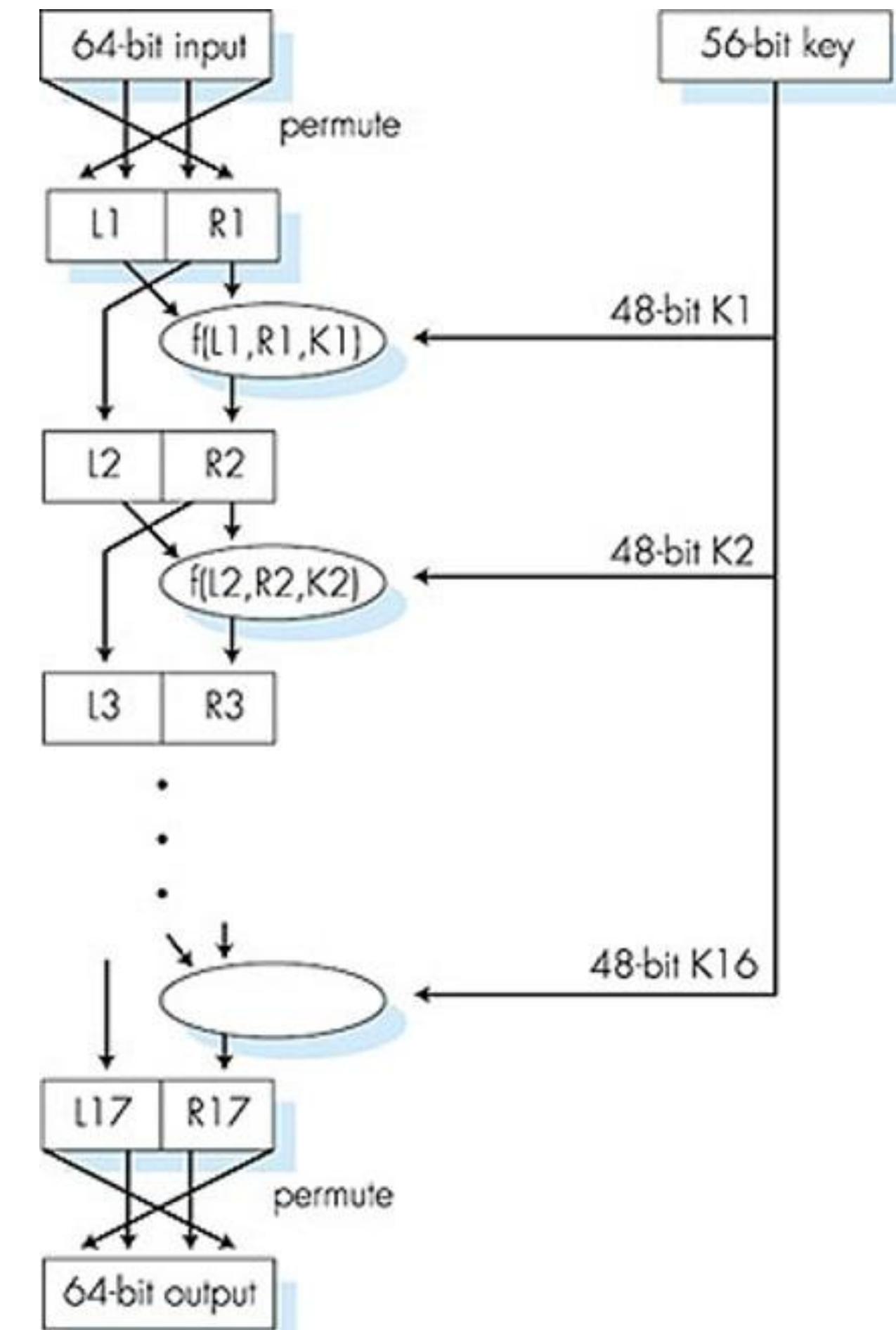
Symmetric key crypto: DES

DES operation

initial permutation

16 identical “rounds” of function application,
each using different 48 bits of key

final permutation

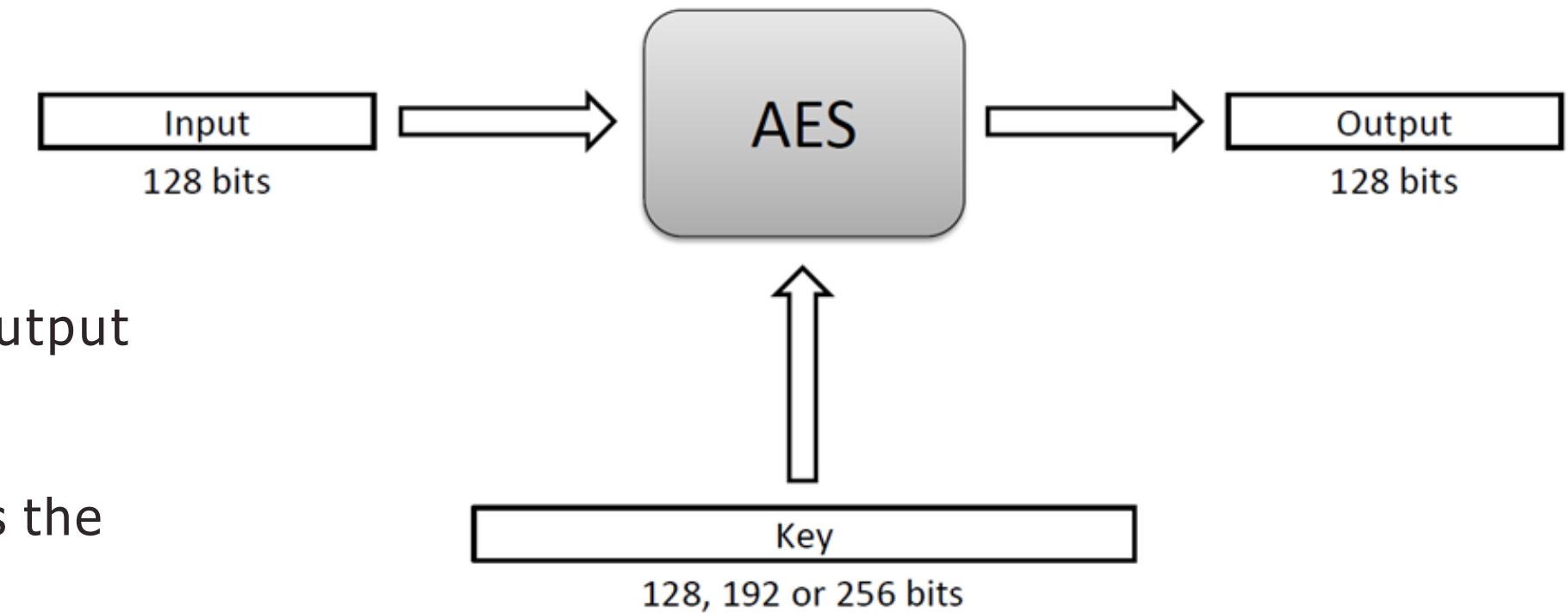


Advanced Encryption Standard (AES)

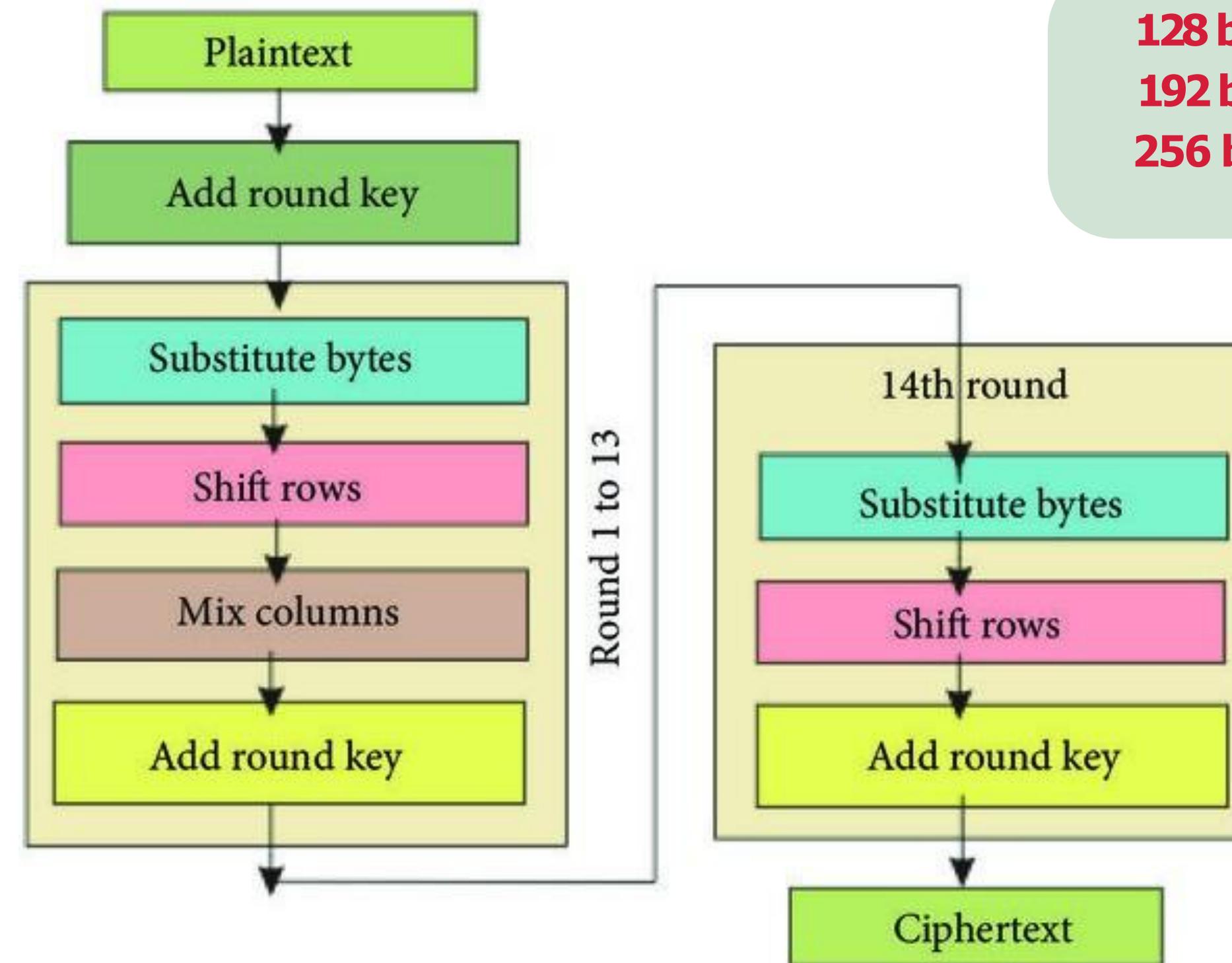
- Block-based symmetric encryption
- input is chunked into blocks of 16, 24, or 32 bytes
(128, 192, 256 bits)

Multiple modes

- ECB
 - **Flaw:** same input block results in the same output block
- CBC
 - **Flaw:** malleable. XOR on the ciphertext XORs the plaintext when decrypted
- CTR
- GCM
- ...many more

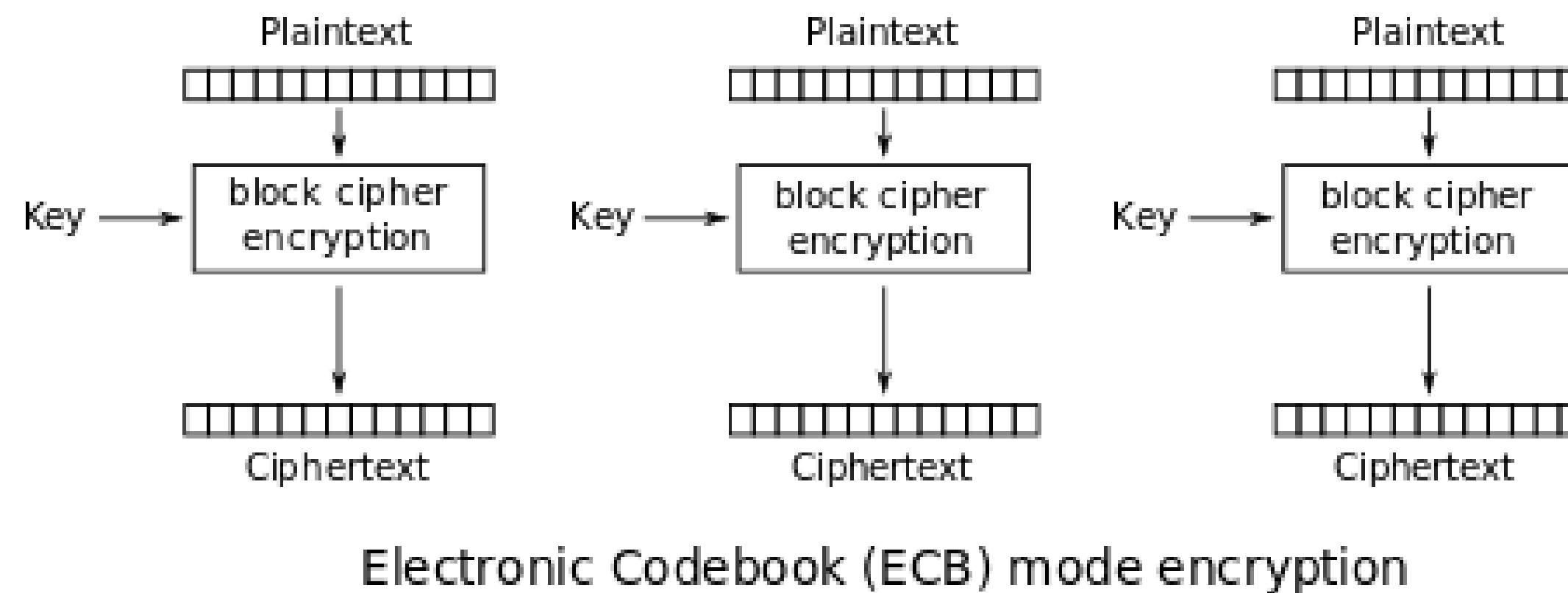


AES Round Structure



ECB Mode

- Electronic Code Book (ECB) Mode (is the simplest):
 - Block $P[i]$ encrypted into ciphertext block $C[i] = E_K(P[i])$
 - Block $C[i]$ decrypted into plaintext block $M[i] = D_K(C[i])$

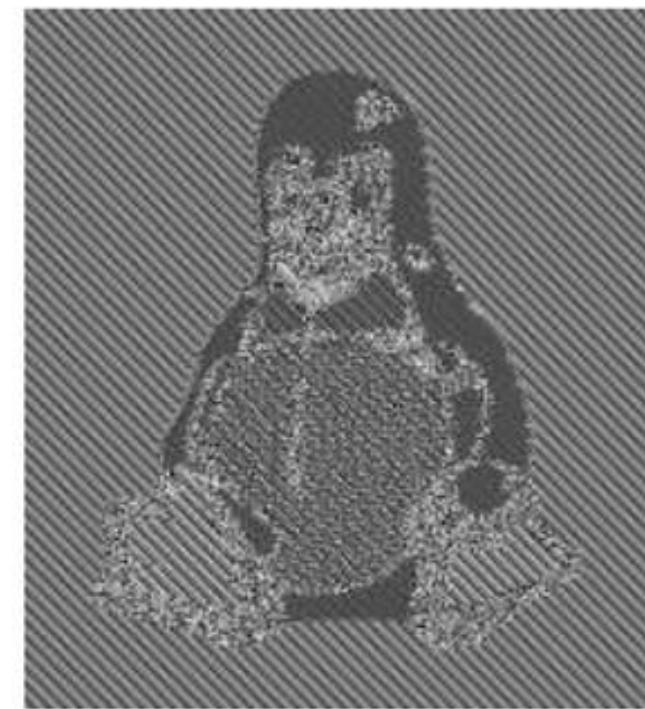


Strengths & Weaknesses of ECB

- Strengths:
 - Is very simple
 - Allows for parallel encryptions of the blocks of a plaintext
 - Can tolerate the loss or damage of a block
- Weakness:
 - Documents and images are not suitable for ECB encryption since patterns in the plaintext are repeated in the ciphertext:



(a)



(b)

RSA: Choosing keys

RSA

RSA (Rivest-Shamir-Adleman) is a public-key cryptosystem that is widely used for secure data transmission.

1. Choose two large prime numbers p, q .
(e.g., 1024 bits each)
 2. Compute $n = pq$, $z = (p-1)(q-1)$
 3. Choose e (with $e < n$) that has no common factors with z . (e, z are "relatively prime").
 4. Choose d such that $ed-1$ is exactly divisible by z .
(in other words: $ed \bmod z = 1$).
 5. Public key is (n, e) . Private key is (n, d) .
- $\overbrace{K_B^+}^{(n,e)}$ $\overbrace{K_B^-}^{(n,d)}$

RSA: Encryption & Decryption

0. Given (n,e) and (n,d) as computed above

1. To encrypt bit pattern, m , compute

$$c = m^e \text{mod } n \quad (\text{i.e., remainder when } m^e \text{ is divided by } n)$$

2. To decrypt received bit pattern, c , compute

$$m = c^d \text{mod } n \quad (\text{i.e., remainder when } c^d \text{ is divided by } n)$$

Magic happens!

$$m = \underbrace{(m^e \text{mod } n)}_c^d \text{mod } n$$

RSA Cryptography Example

Bob chooses $p=5, q=7$. Then $n=35, z=24$.

$e=5$ (so e, z relatively prime).

$d=29$ (so $ed-1$ exactly divisible by z).

| | <u>letter</u> | <u>m</u> | <u>m^e</u> | <u>$c = m^e \bmod n$</u> |
|----------|---------------|--------------------------------------|-------------------------------------|-------------------------------------|
| encrypt: | I | 12 | 1524832 | 17 |
| decrypt: | <u>c</u> | <u>c^d</u> | <u>$m = c^d \bmod n$</u> | <u>letter</u> |
| | 17 | 481968572106750915091411825223071697 | 12 | I |

RSA Algorithm

RSA Basic Term:

1. Choose **two prime numbers p, q**
2. Multiply, $n = p * q$
3. $\phi(n) = (p-1)(q-1)$
4. Generate Public key, Choose an integer number e --- where $1 < e < \phi(n)$ and e is co-prime of $\phi(n)$
- 5. Generate Private key, Compute d to satisfy, $(d * e) \text{ MOD } \phi(n) = 1$
 $d = \text{inverse}(e, \phi)$ -----python code

Note :

For Public key:

- 1. Must be prime number
- 2. Must be less than $\phi(n)$
- 3. NOT be a factor of $\phi(n)$

RSA Cryptography

Encrypt Message :

Public key is ($n = 3233$, $e = 17$)

suppose, $m = 123$

So, Cipher $c = m^e \text{ mod } n$

$$\begin{aligned} &= 123^{17} \text{ mod } 3233 \\ &\quad = 855 \end{aligned}$$

Encrypt:

$$C = m^e \text{ mod } n$$

Here, $C \rightarrow$ Cipher text

$m \rightarrow$ Message

Decrypt:

$$M = C^d \text{ mod } n$$

Decrypt Message :

Private Key is ($n=3233$, $d = 2753$)

$m = C^d \text{ mod } n$

$$= 855^{2753} \text{ mod } 3233$$

$$= 123$$

THANK YOU