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# What is Cryptography?

Cryptography derived its name from a Greek word called "**Kryptos**" which means "**Hidden Secrets**".

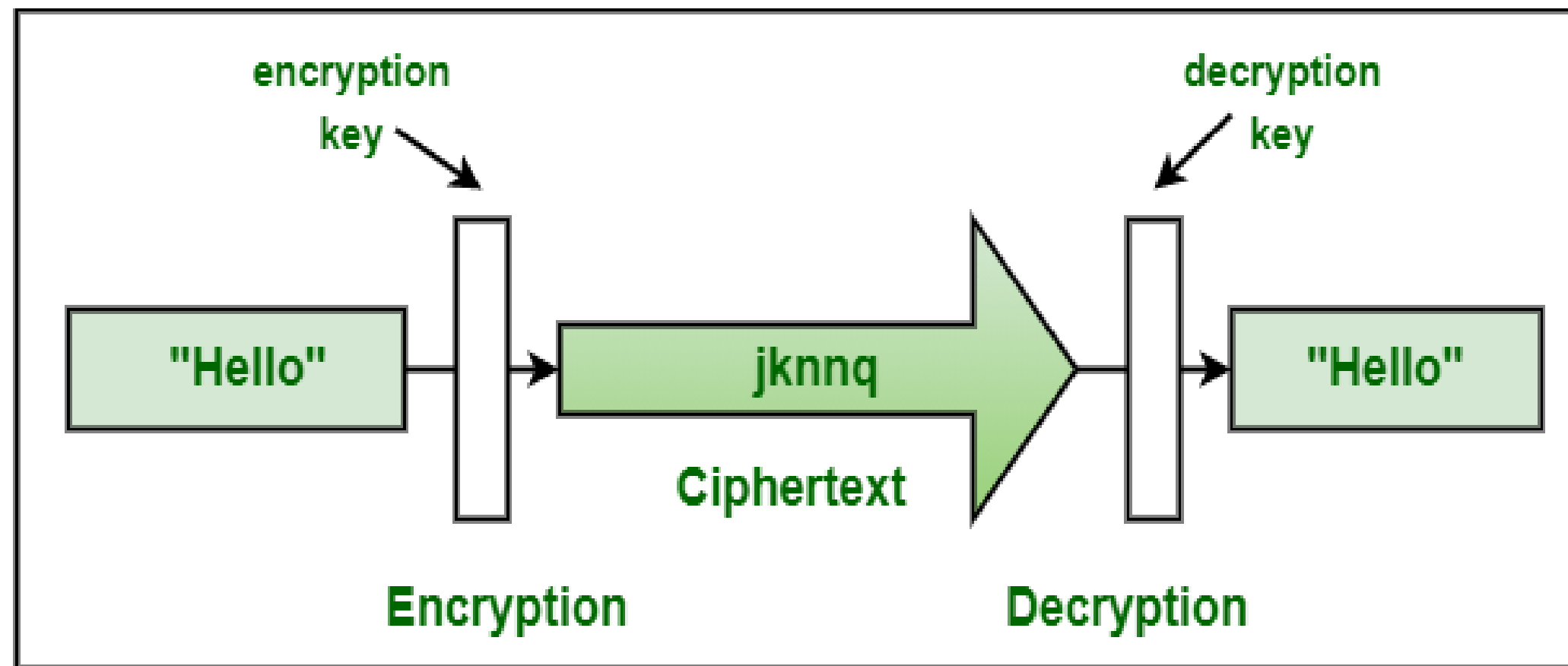
1. Cryptography is the practice and study of hiding information.
2. It provides Confidentiality, Integrity, Accuracy

# Cryptography

Cryptography is process to secure your information from unauthorise access.

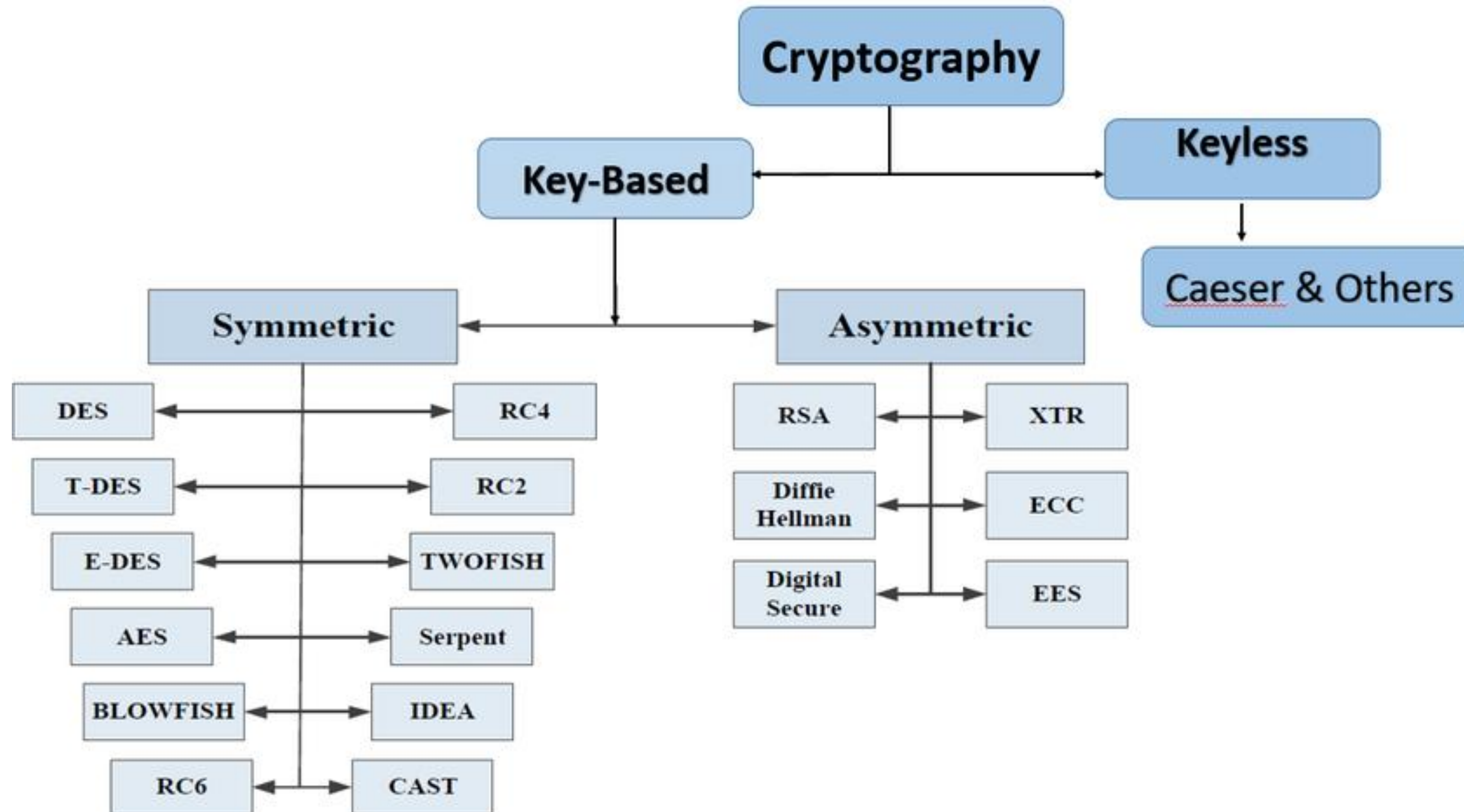
<b>Encrypted</b>	<b>501393c26075af84b9e612b5d90b4cff790ce10c</b>
<b>Decrypted</b>	<b>Hello_NSU_Hackers</b>

# Encryption & Decryption



Cryptography

# What are the Types of Cryptography



# Some Common Cryptography

## Example

- Caesar Cipher
- Rot 13
- Vigenere Cipher
- Morse code
- Bacon Cipher
- Alphabetical substitution

Tools
<ul style="list-style-type: none"><li>• <a href="#">Cyberchef</a></li><li>• <a href="#">dcode.fr</a></li><li>• <a href="#">boxentriq.com</a></li><li>• <a href="#">cryptii.com</a></li></ul>

# Symmetric Cryptography

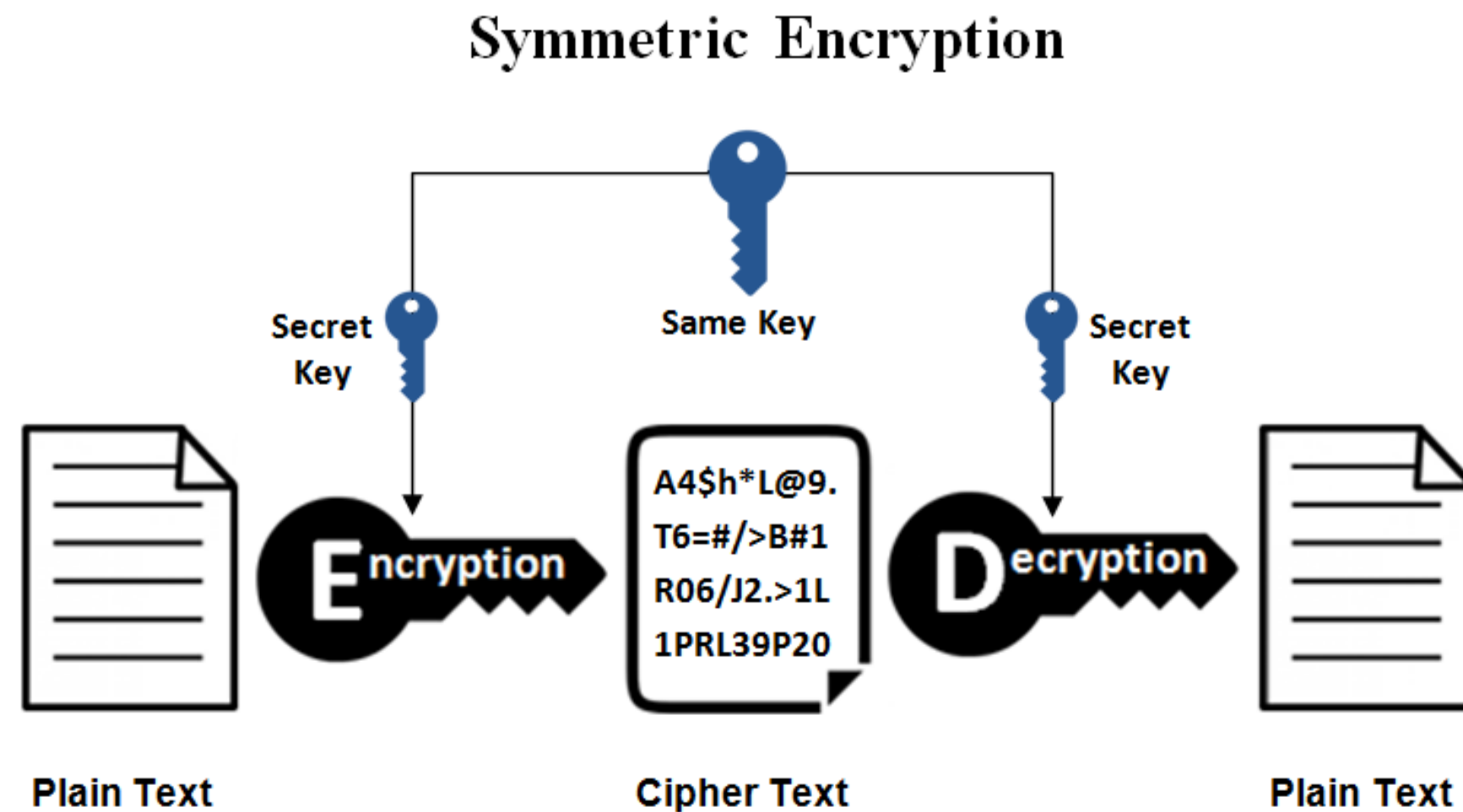
Symmetric cryptography is type of cryptography where both sender and receiver use same private key.

**Advantage:**

Simpler and Faster

**Disadvantage:**

Less Secured



# Symmetric Cryptography

## Example

Recipe

^

📁

🗑

Input

Vigenère Encode

^

🚫

⏸

Key

hackers

ABC 9

≡ 1

Output

Oenvs\_EKB

Encrypt

Recipe

^

📁

🗑

Input

Vigenère Decode

^

🚫

⏸

Key

hackers

ABC 9

≡ 1

Output

Hello\_NSU

Decrypt



# Asymmetric Cryptography

Use public and private key:

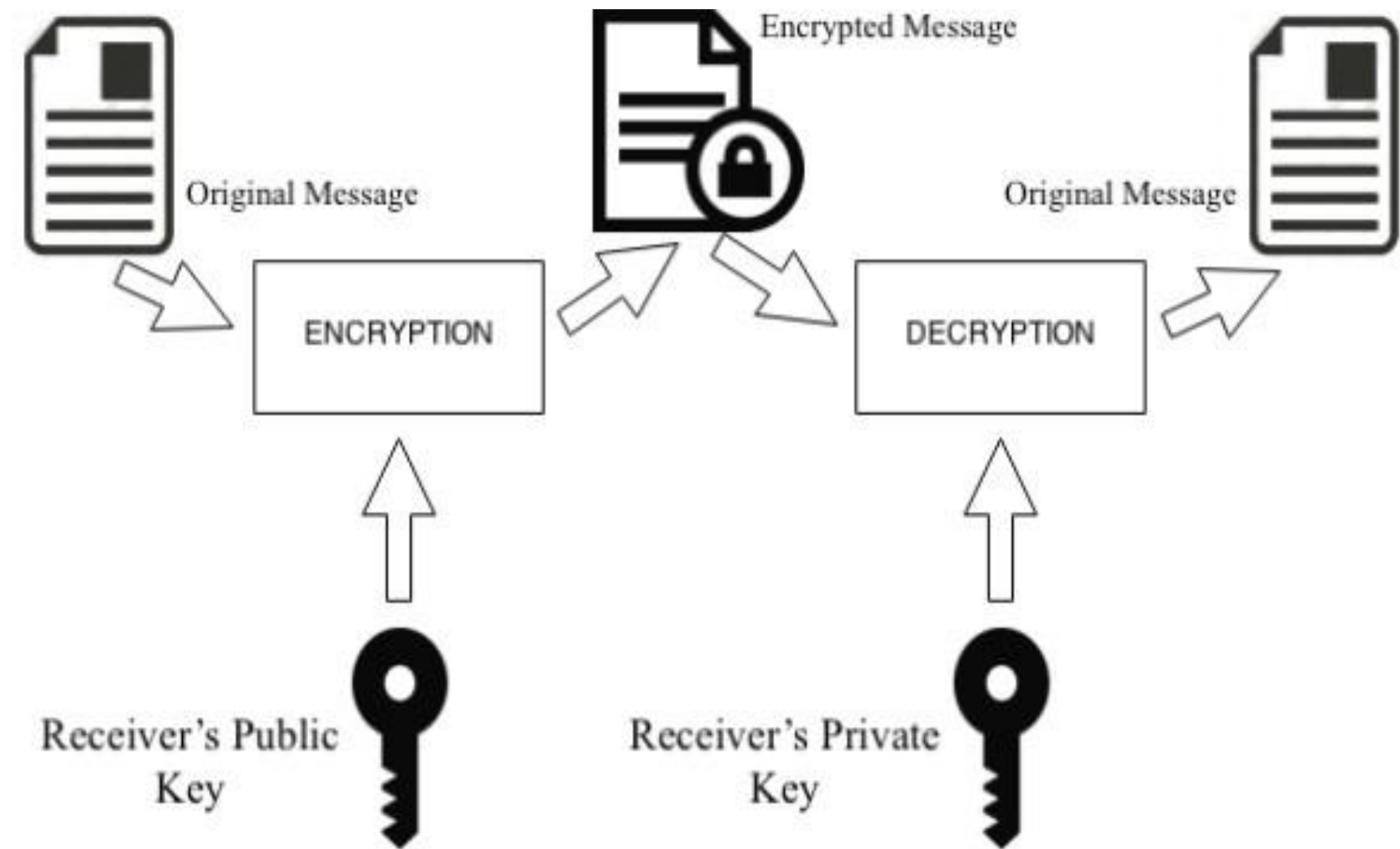
- **Public key** use for encryption
- **Private key** use for decryption

## Advantage:

1. More Secured
2. Authentication

## Disadvantage:

1. Relatively Complex



# Asymmetric Cryptography

## Example

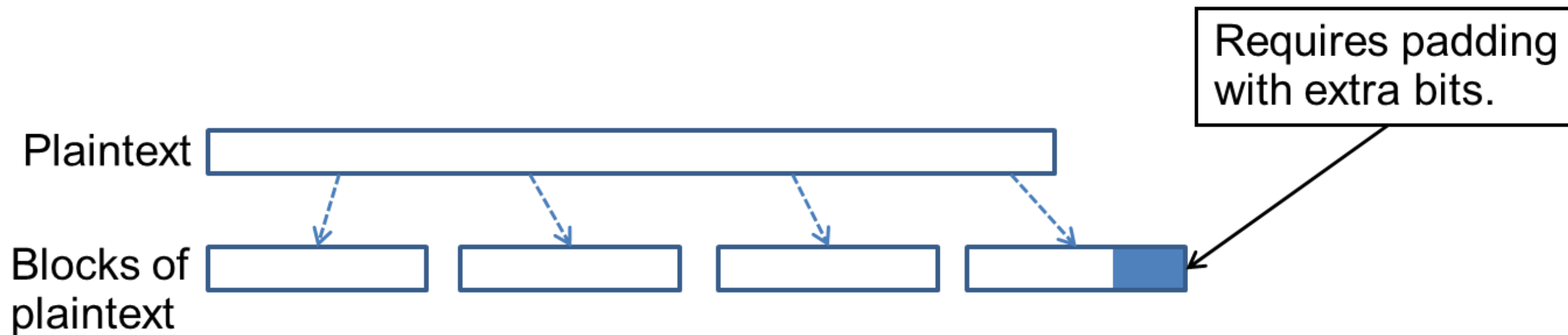
- RSA Algorithm
- Diffie-Hellman
- Elliptic Curve (ECC)

## Tools

- Python code
- openssl (Default on Kali Linux OS)
- Online Tools

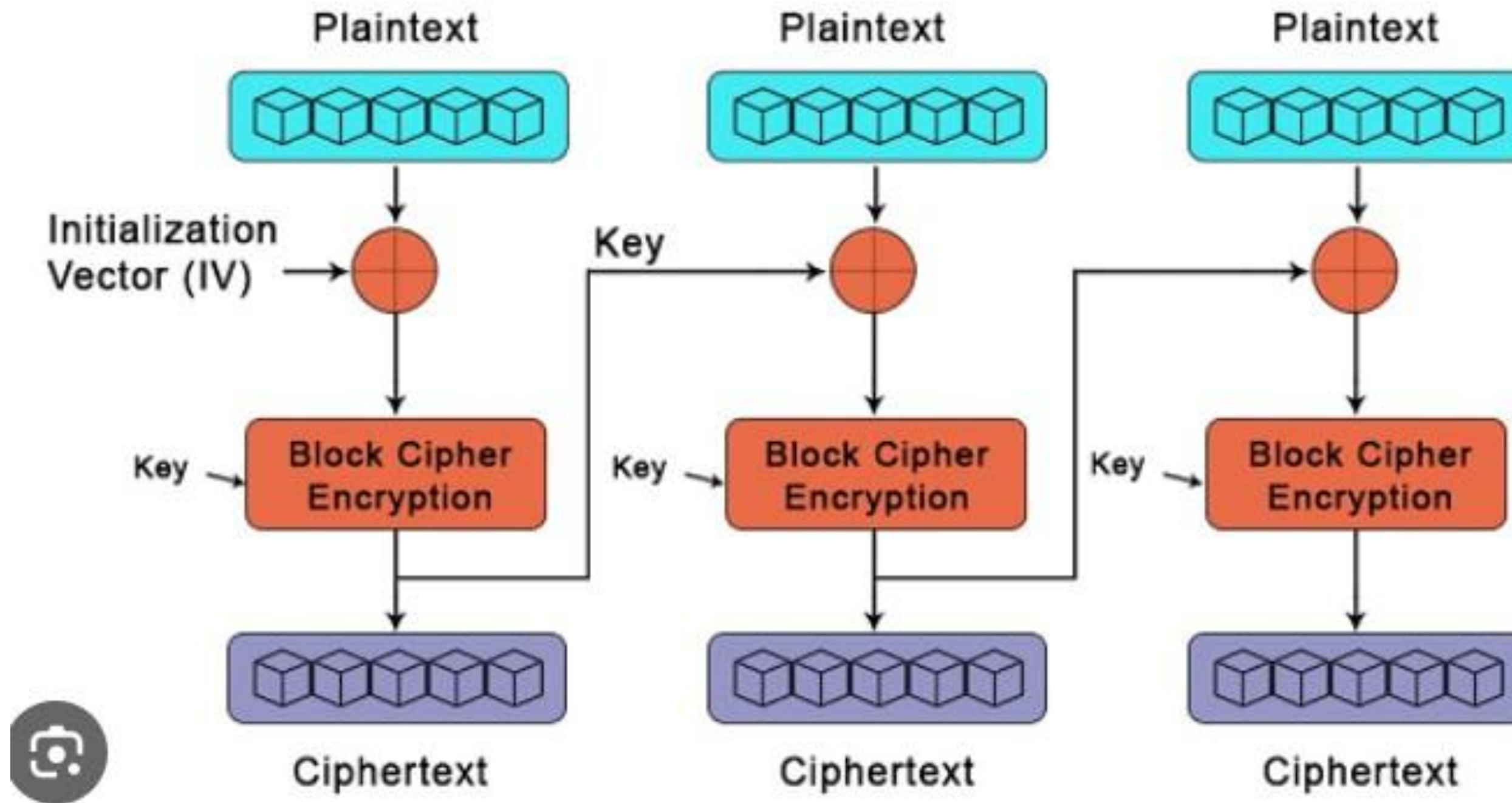
# Block Ciphers

- In a **block cipher**:
  - Plaintext and ciphertext have fixed length  $b$  (e.g., 128 bits)
  - A plaintext of length  $n$  is partitioned into a sequence of  $m$  **blocks**,  $P[0], \dots, P[m-1]$ , where  $n \leq bm < n + b$
- Each message is divided into a sequence of blocks and encrypted or decrypted in terms of its blocks.





# Block Cipher

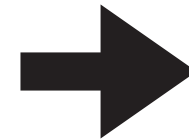


# Padding

- **Block ciphers** require the length  $n$  of the plaintext to be a multiple of the block size  $b$
- Padding the **last block needs to be unambiguous** (cannot just add zeroes)

# Where are block ciphers used

**Data Encryption Standard (DES)**

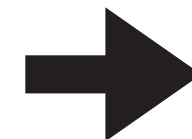


64-bit blocks &  
56-bit keys

## Triple DES (3DES)

- Nested application of DES with **three different keys KA, KB, and KC**
- Effective **key length is 168 bits**, making exhaustive search attacks unfeasible

**Advanced Encryption Standard (AES)**



128-bit blocks  
key lengths: 128, 192 & 256  
bits

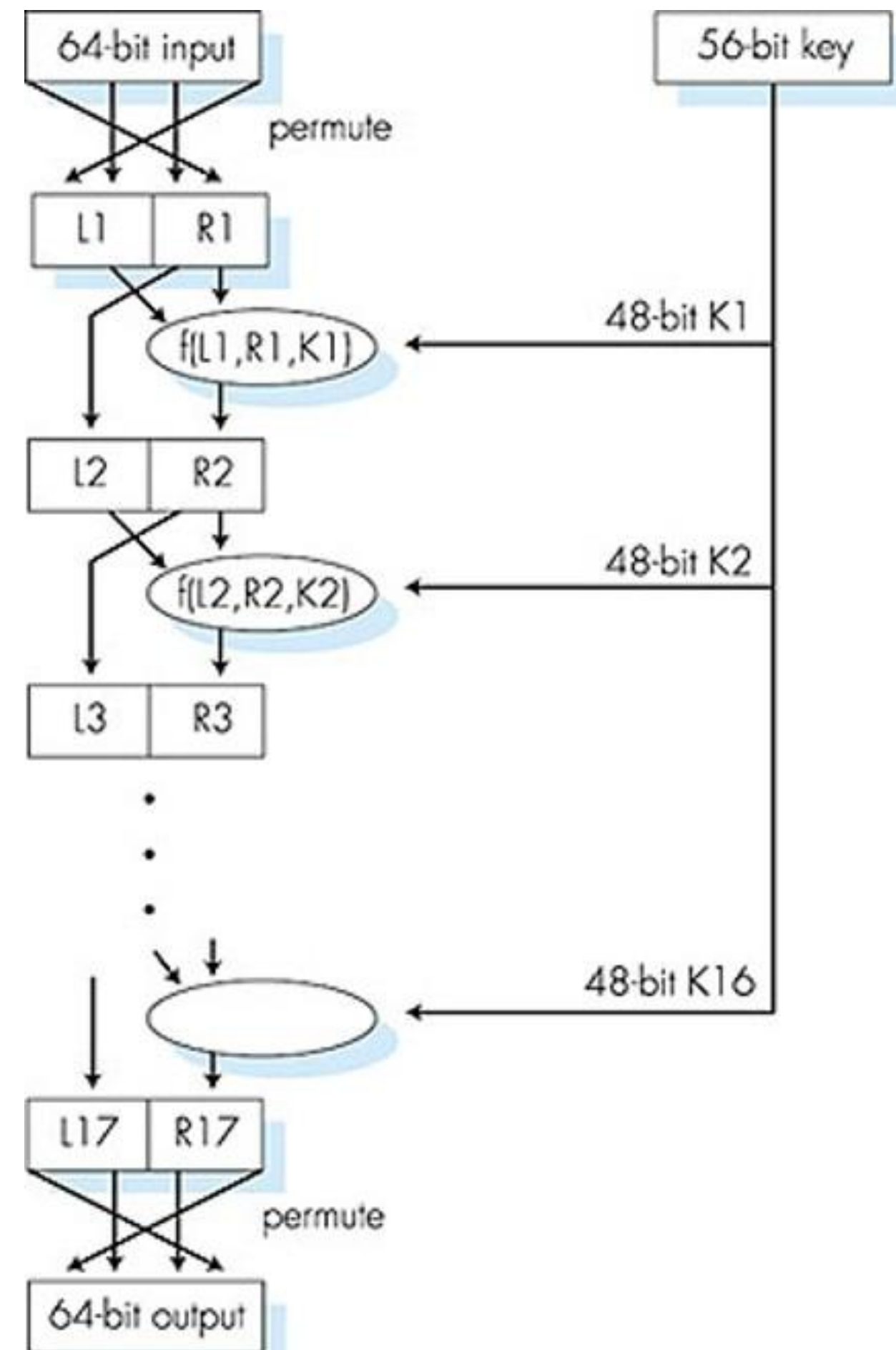
# Symmetric key crypto: DES

## DES operation

initial permutation

16 identical “rounds” of  
function application,  
each using different 48  
bits of key

final permutation

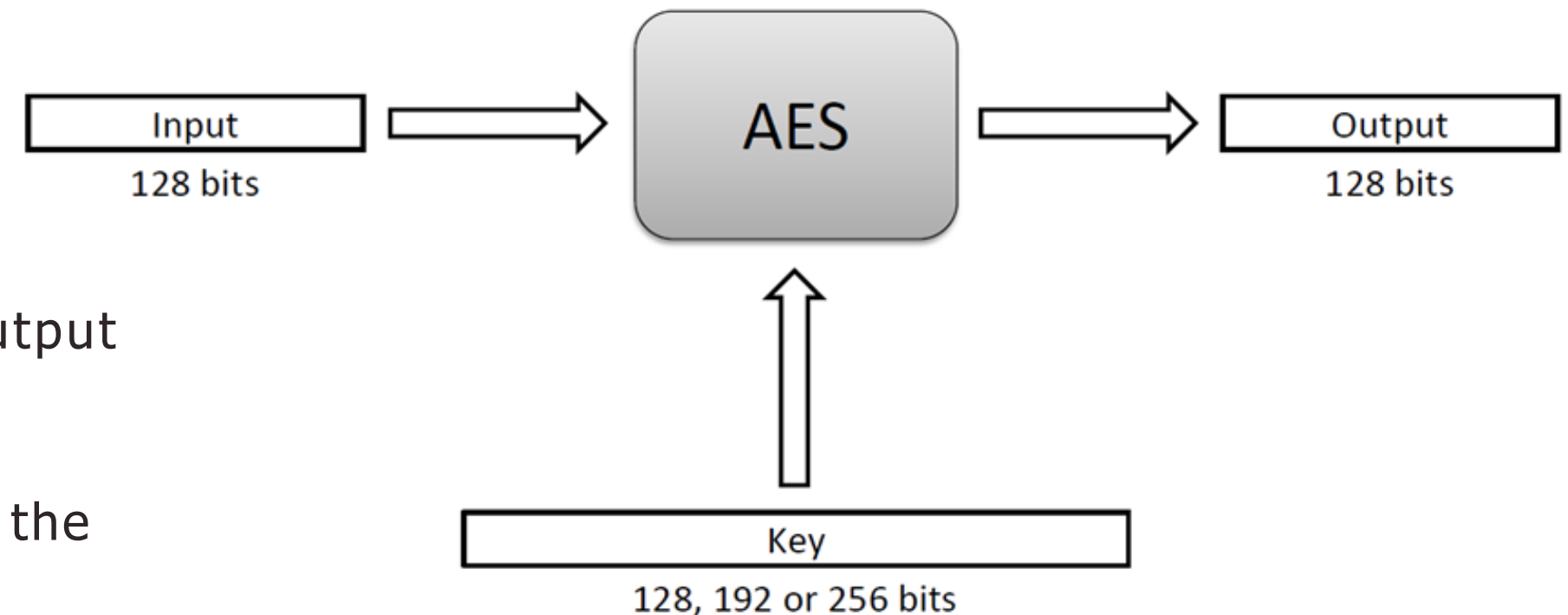


# Advanced Encryption Standard (AES)

- Block-based symmetric encryption
- input is chunked into blocks of 16, 24, or 32 bytes (128, 192, 256 bits)

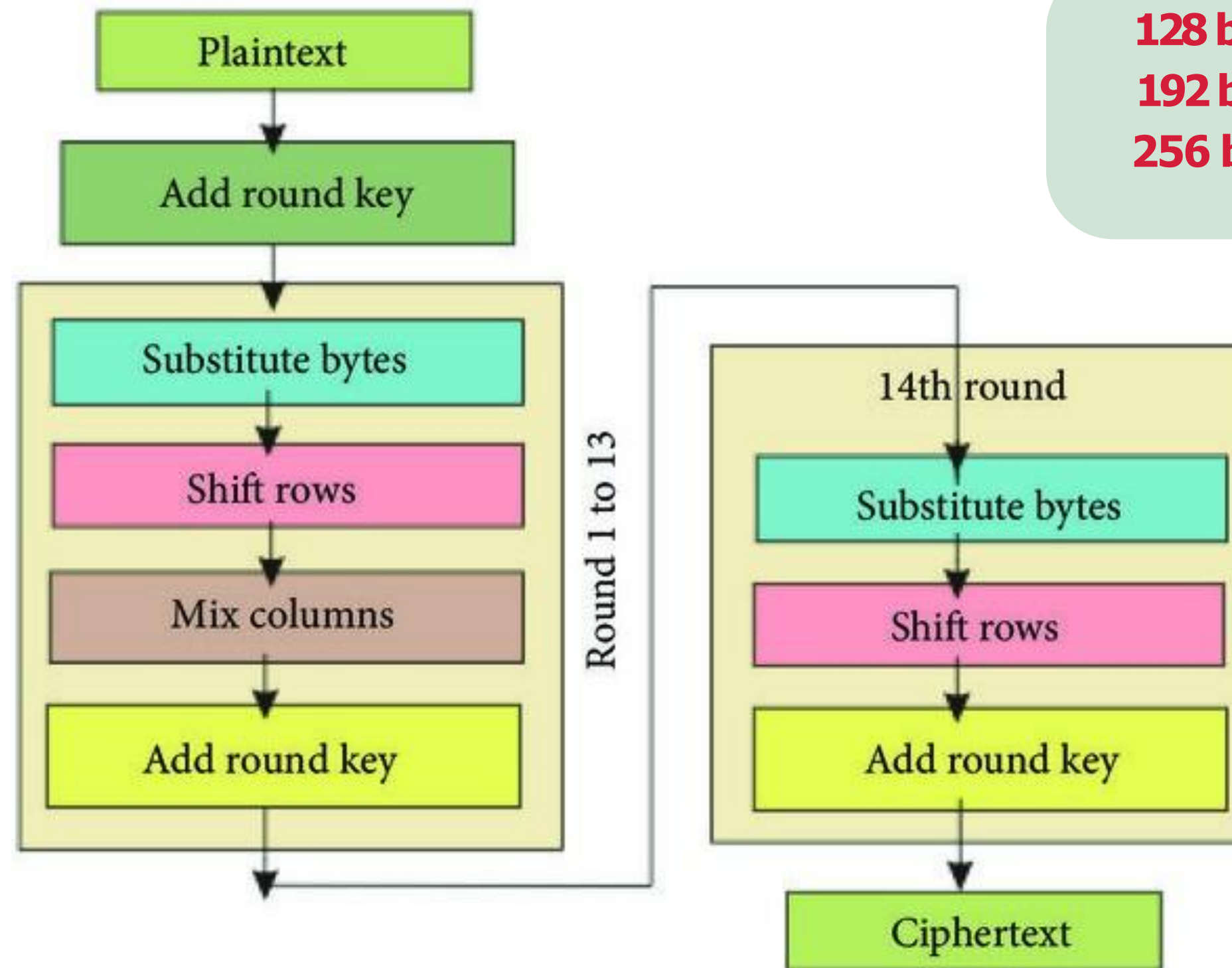
## Multiple modes

- ECB
  - **Flaw**: same input block results in the same output block
- CBC
  - **Flaw**: malleable. XOR on the ciphertext XORs the plaintext when decrypted
- CTR
- GCM
- ...many more





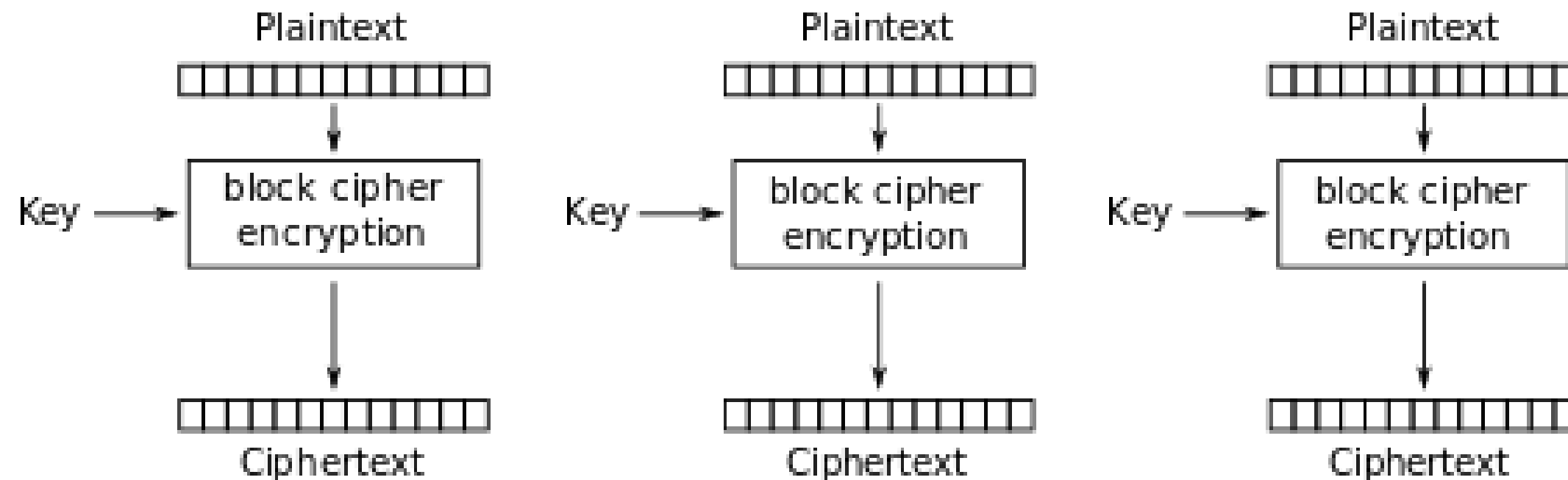
# AES Round Structure



**128 bit key –10 rounds**  
**192 bit key –12 rounds**  
**256 bit key –14 rounds**

# ECB Mode

- Electronic Code Book (ECB) Mode (is the simplest):
  - Block  $P[i]$  encrypted into ciphertext block  $C[i] = E_K(P[i])$
  - Block  $C[i]$  decrypted into plaintext block  $M[i] = D_K(C[i])$



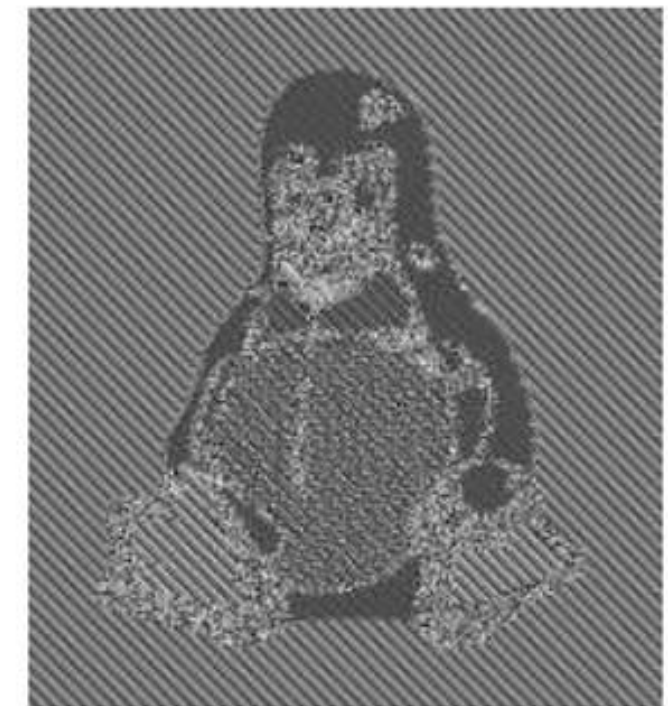
Electronic Codebook (ECB) mode encryption

# Strengths & Weaknesses of ECB

- Strengths:
  - Is very simple
  - Allows for parallel encryptions of the blocks of a plaintext
  - Can tolerate the loss or damage of a block
- Weakness:
  - Documents and images are not suitable for ECB encryption since patterns in the plaintext are repeated in the ciphertext:



(a)



(b)

# RSA

RSA (Rivest-Shamir-Adleman) is a public-key cryptosystem that is widely used for secure data transmission.

## RSA: Choosing keys

1. Choose two large prime numbers  $p, q$ .  
(e.g., 1024 bits each)
2. Compute  $n = pq$ ,  $z = (p-1)(q-1)$
3. Choose  $e$  (with  $e < n$ ) that has no common factors with  $z$ . ( $e, z$  are "relatively prime").
4. Choose  $d$  such that  $ed-1$  is exactly divisible by  $z$ .  
(in other words:  $ed \bmod z = 1$ ).
5. Public key is  $(n, e)$ . Private key is  $(n, d)$ .  

$\underbrace{\hspace{1.5cm}}_{K_B^+}$

$\underbrace{\hspace{1.5cm}}_{K_B^-}$

## RSA: Encryption & Decryption

0. Given  $(n,e)$  and  $(n,d)$  as computed above
1. To encrypt bit pattern,  $m$ , compute
$$c = m^e \bmod n \quad (\text{i.e., remainder when } m^e \text{ is divided by } n)$$
2. To decrypt received bit pattern,  $c$ , compute
$$m = c^d \bmod n \quad (\text{i.e., remainder when } c^d \text{ is divided by } n)$$

Magic happens!

$$m = \underbrace{(m^e \bmod n)}_c^d \bmod n$$

# RSA Cryptography

## Example

Bob chooses  $p=5$ ,  $q=7$ . Then  $n=35$ ,  $z=24$ .

$e=5$  (so  $e$ ,  $z$  relatively prime).

$d=29$  (so  $ed-1$  exactly divisible by  $z$ ).

	<u>letter</u>	<u>m</u>	<u><math>m^e</math></u>	<u><math>c = m^e \bmod n</math></u>
encrypt:	I	12	1524832	17
	<u>c</u>	<u><math>c^d</math></u>	<u><math>m = c^d \bmod n</math></u>	<u>letter</u>
decrypt:	17	481968572106750915091411825223071697	12	I



# RSA Algorithm

## RSA Basic Term:

1. Choose **two prime numbers**  $p, q$
2. Multiply,  $n = p * q$
3.  $\phi(n) = (p-1)(q-1)$
4. Generate Public key, Choose an integer number  $e$  --- where  $1 < e < \phi(n)$  and  $e$  is co-prime of  $\phi(n)$
- 5. Generate Private key, Compute  $d$  to satisfy,  $(d * e) \text{ MOD } \phi(n) = 1$   
 $d = \text{inverse}(e, \phi)$  -----python code

## Note :

### For Public key:

- 1. Must be prime number
- 2. Must be less than  $\phi(n)$
- 3. NOT be a factor of  $\phi(n)$

# RSA Cryptography

## Encrypt Message :

Public key is (  $n = 3233$  ,  $e = 17$  )

suppose,  $m = 123$

So, Cipher  $c = m^e \bmod n$

$$\begin{aligned} &= 123^{17} \bmod 3233 \\ &= 855 \end{aligned}$$

## Decrypt Message :

Private Key is (  $n=3233$  ,  $d = 2753$  )

$$\begin{aligned} m &= C^d \bmod n \\ &= 855^{2753} \bmod 3233 \\ &= 123 \end{aligned}$$

Encrypt:

$$C = m^e \bmod n$$

Here,  $C \rightarrow$  Cipher text

$m \rightarrow$  Message

Decrypt:

$$M = C^d \bmod n$$



**THANK YOU**