

# CMPUT366: Assignment 1

(Fall 2021)

**Material Coverage:** Lecture 1 to 4

**Algorithms Involved:** Dijkstra and A\*

**Assignment Weight:** 12%

**Mark Distribution:** 50% for correct code implementation; 50% for answering questions about your results

**Goal of this assignment:**

Your goal is to implement Dijkstra's algorithm and A\* for solving pathfinding problems

- We will consider a grid environment where each action in the four cardinal directions (north, south, east, and west) has the cost of 1.0 and each action in one of the four diagonal directions has the cost of 1.5.
- Each search problem is defined by a video game map, a start location and a goal location.
- You will use a single map in our experiments; feel free to explore other maps if you want.

Gridded\_map = Map("dao-map/brc000d.map") in main.py

**Python Environment Setup:**

- Use of virtual environment is encouraged
  1. Download the assignment zip, unzip it
  2. Open a terminal at the unzipped folder
  3. `virtualenv -p python3 venv` (choose a name for your virtual environment)
  4. `source venv/bin/activate`
  5. `pip install -r requirements.txt`

Now the assignment environment is set up.

## Implementing Dijkstra:

- $S_o$ : start node
- $S_g$ : goal node
- $T(n)$ : returns the children of node  $n$
- $g(n)$ : cost from  $S_{init}$  to  $n$
- $c(n, n')$ : cost from  $n$  to  $n'$
- $g(n) + c(n, n')$ : cost of  $n'$
- OPEN(priority queue)
- CLOSED(python dictionary (hash\_value: node))

```
Def Dijkstra( $S_o, S_g, T$ ):  
    OPEN.insert( $S_o, 0$ )  
    CLOSED.add( $S_o$ )  
    # node_expanded will be used for explaining the results later  
    node_expanded = 0  
    while OPEN is not empty:  
         $n = \text{OPEN.pop}()$   
        node_expanded += 1  
        # goal found  
        If  $n == S_g$ :  
            Return Path  
  
    for  $n'$  in  $T(n)$ :  
        #  $n'$  is not expanded before  
        If  $n'$  not in CLOSED:  
            OPEN.insert( $n', g(n) + c(n, n')$ )  
            CLOSED.Add( $n'$ )  
        #  $n'$  is expanded before, but a cheaper path is found to  $n'$   
        If  $n'$  in CLOSED & ( $g(n) + c(n, n') < g(n')$ ):  
            Update  $g(n')$  to  $g(n) + c(n, n')$  in OPEN  
            Update  $g(n')$  to  $g(n) + c(n, n')$  in CLOSED  
  
    # no solution  
    Return -1
```

**FAQ:**

1. This Assignment did not ask for a returning solution path. So you do not need to update parent of  $n'$  if a cheaper path to  $n'$  is found
2. Hash\_value is a unique identification for a node. For example, you can use `self.CLOSED[hash_value]` to access a specific node in the CLOSED dictionary.
3. `map.successors(n)` can be used to access the children( or neighbors) of a node  $n$ .
4. You don't need to apply heap on the CLOSED dictionary since you don't need to pop out any value from.

In the OPEN list. Nodes can be sorted by their cost. (See the `__lt__` method in State class, algorithms.py). You can easily use the python heap queue algorithm to turn the OPEN list to a priority queue.

Reference: <https://docs.python.org/3/library/heapq.html>

### **heapq (you may need to use it for OPEN):**

This module provides an implementation of the heap queue algorithm, also known as the priority queue algorithm.

heapq.**heappush**(*heap*, *item*)

Push the value *item* onto the *heap*, maintaining the heap invariant.

heapq.**heappop**(*heap*)

Pop and return the smallest item from the *heap*, maintaining the heap invariant. If the heap is empty, `IndexError` is raised. To access the smallest item without popping it, use `heap[0]`.

heapq.**heapify**(*x*)

Transform list *x* into a heap, in-place, in linear time.

## Implementing A\* Algorithm:

- $S_o$ : start node
- $S_g$ : goal node
- $h(n, S_g)$ : heuristic function; estimates the cost to reach goal from node  $n$
- current: current node ( $n$ )
- $T(n)$ : returns the children of node  $n$
- $g(n)$ : cost from  $S_o$  to  $n$
- $c(n, n')$ : cost from  $n$  to  $n'$
- $g(n) + c(n, n')$ : cost of  $n'$
- OPEN(priority queue)
- CLOSED(python dictionary)

Def A\*( $S_o, S_g, h$ ):

```
OPEN.insert( $S_o, 0$ )
while OPEN is not empty:
     $n = \text{OPEN.pop}()$ 
    # goal found
    If  $n == S_g$ :
        Return Path
    CLOSED.Add( $n$ )
    # Loop through the node's children
    for  $n'$  in  $T(n)$ :
        # If it is already in CLOSED, skip it
        if  $n'$  in CLOSED:
            continue
        if  $n'$  in OPEN:
            if  $g(n') > g(n) + c(n, n')$ :
                G score of node  $n' = g(n) + c(n, n')$ 
            else:
                # If it isn't in the open set, calculate the G and H score for  $n'$ 
                G score of node  $n' = g(n) + c(n, n')$ 
                H score of node  $n' = h(n', S_g)$ 
                # Add it to OPEN
                OPEN.insert( $n', \text{G score of } n' + \text{H score of } n'$ )
        # no solution
    Return -1
```