

Outline

ARP

• ARP Package

RARP

Logical Addresses

- The hosts and routers are recognized at the network level by their *logical addresses*
 - n A logical address is an internet address
 - n Called a *logical* address because it is usually implemented in software
 - n The logical addresses in the TCP/IP are called **IP** address and are 32 bits long

Physical Address

- However, hosts/routers are recognized at the physical layer by their *physical address*
 - n A physical address is an local address
 - n Called a *physical* address because it is usually implemented in hardware
 - n Examples
 - 48-bit MAC addresses in Ethernet

Translation

- We need both the physical address and the logical address for packet delivery
- Thus, we need to be able to map a logical address to its corresponding physical address and vice versa
- Solutions
 - n Static mapping
 - n Dynamic mapping

Static Mapping

- Create a table that associates a logical address with a physical address and store in each machine
- However, physical addresses may change
 - n A machine could change its NIC resulting in a new physical address
 - n In some LANs, such as LocalTalk, the physical address changes every time the computer is turned on
 - n A mobile station can move from one physical network to another, resulting in a change in its physical address

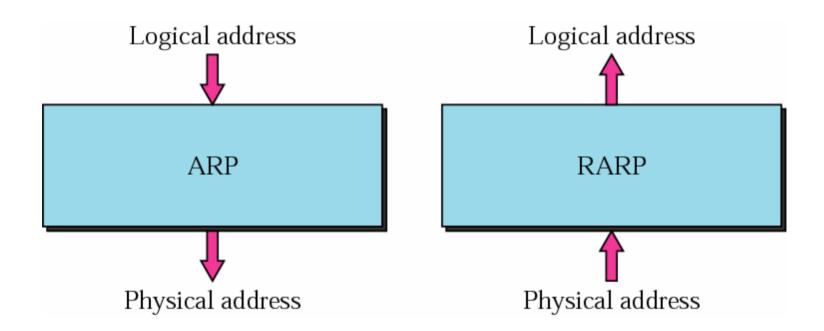
Dynamic Mapping

Use a protocol to find another address

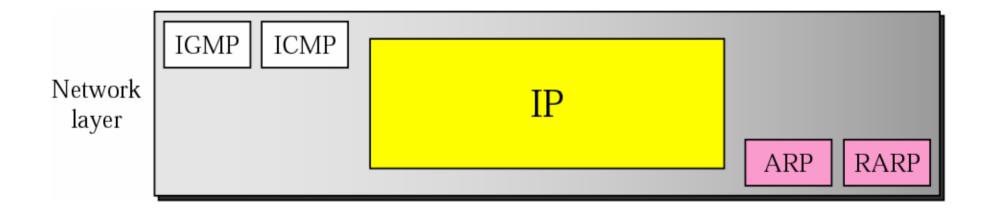
- ARP: Address Resolution Protocol
 - n Map a logical address to a physical address

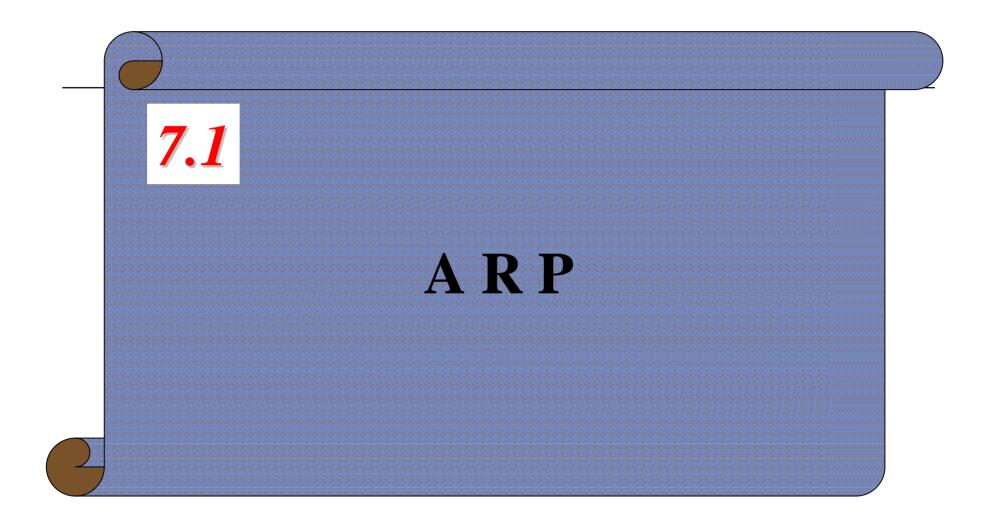
- RARP: Reverse Address Resolution Protocol
 - n Map a physical address to a logical address

ARP and RARP



Position of ARP and RARP in TCP/IP Protocol Suite





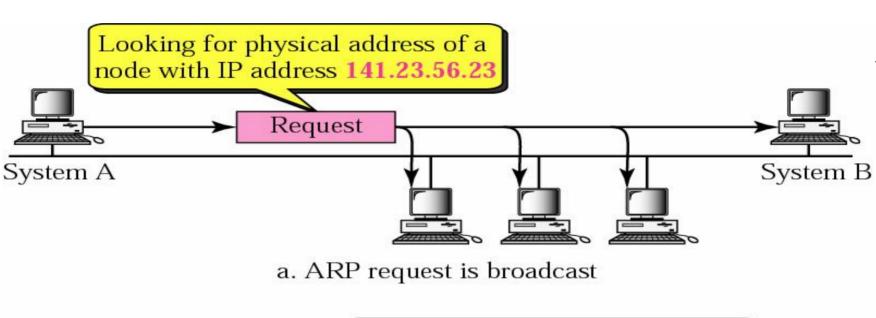
ARP Operation

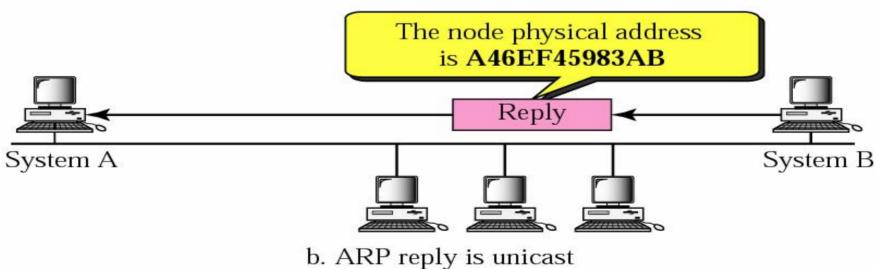
- To find the physical address of another host or router on its network
 - n Send an ARP request message
- ARP request message
 - n The physical address of the sender
 - n The IP address of the sender
 - n The physical address of the receiver is θs
 - n The IP address of the receiver

ARP Operation (Cont.)

- Then, ARP request message is broadcast by the physical layer
 - n For example: in Ethernet, MAC header's destination address is all *1s* (broadcast address)
 - n Received by every station on the physical network
- The intended recipient send back an ARP reply message
 - n ARP reply message packet is *unicast*

ARP Operation





ARP Packet

Hardware Type		Protocol Type
Hardware length	Protocol length	Operation Request 1, Reply 2
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled in a request)		
Target protocol address (For example, 4 bytes for IP)		

Packet Format

- HTYPE (Hardware type)
 - n 16-bit field defining the underlying type of the network
 - Ethernet is given the type 1
 - ARP can be used on any physical network
- PTYPE (Protocol type)
 - n 16-bit field defining the protocol
 - IPv4 is 0800_{16}
 - ARP can be used with any higher-level protocol

Packet Format (Cont.)

- HLEN (Hardware length)
 - 8-bit field defining the length of the physical address in bytes
 - Ethernet has the value of 6
- PLEN (Protocol length)
 - n 8-bit field defining the length of the logical address in bytes
 - IPv4 has the value of 4
- OPER (Operation)
 - n 16-bit field defining the type of packet
 - n (1) = ARP request, (2) = ARP reply

Packet Format (Cont.)

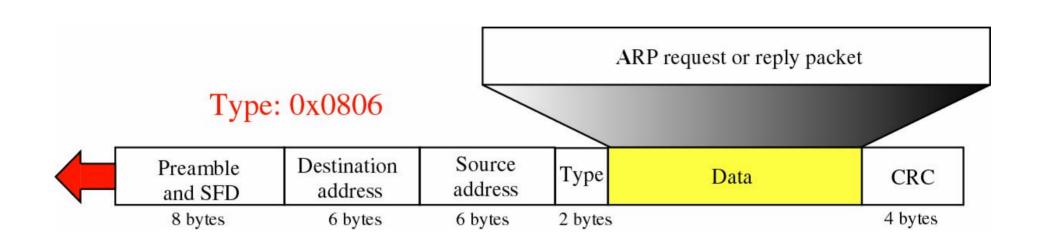
- SHA (Sender hardware address)
 - n A variable-length field defining the physical address of the sender

- SPA (Sender protocol address)
 - n A variable-length field defining the logical address of the sender

Packet Format (Cont.)

- THA (Target hardware address)
 - n A variable-length field defining the physical address of the target
 - n For an ARP request operation packet
 - This field is all 0s
- TPA (Target protocol address)
 - n A variable-length field defining the logical address of the target

Encapsulation of ARP Packet



- An ARP packet is encapsulated directly into a data link frame
- Type field indicates that the data carried by the frame is an ARP packet

Operations

- The sender knows the target's IP address
- IP asks ARP to create an ARP request message
 - n The sender physical address
 - n The sender IP address
 - n The target physical address field is filled with 0s
 - n The target IP address
- The message is passed to the data link layer to encapsulate in a data link frame
 - n Physical destination address is broadcast address

Operations (Cont.)

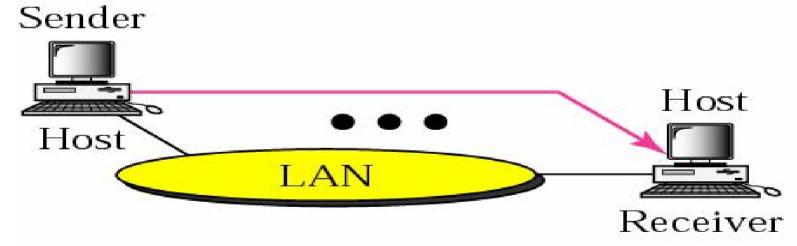
- Every host or routers receives the frame and since the destination address is broadcast, pass it to the ARP
 - n All machines' ARP except the one targeted drop the packet
- The target reply with an ARP reply message that contains its physical address and is unicast
- The sender receives the reply message and knows the target's physical address

Four Cases to Use ARP

- Case 1: The sender is a host and wants to send a packet to another host on the same network
 - n Use ARP to find another host's physical address
- Case 2: The sender is a host and wants to send a packet to another host on another network
 - n Sender looks at its routing table
 - n Find the IP address of the next hop (router) for this destination
 - n Use ARP to find the router's physical address

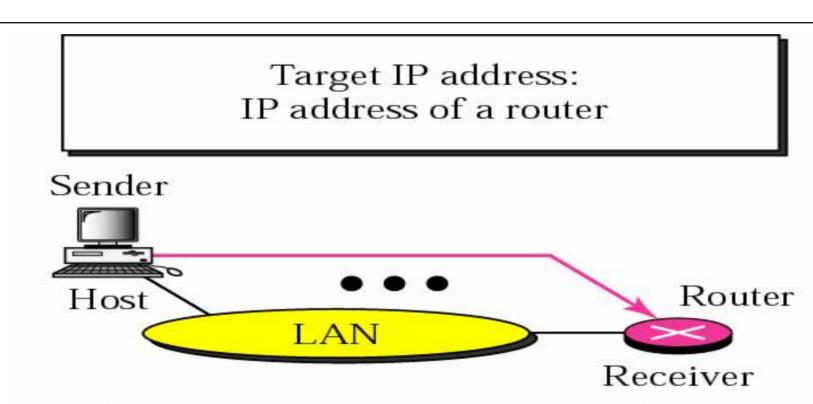
Four Cases Using ARP: Case 1

Target IP address: Destination address in the IP datagram



Case 1. A host has a packet to send to another host on the same network.

Four Cases Using ARP: Case 2



Case 2. A host wants to send a packet to another host on another network.

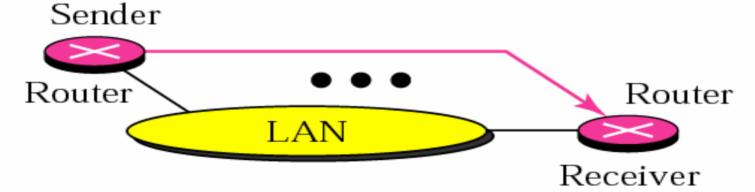
It must first be delivered to a router.

Four Cases to Use ARP (Cont.)

- Case 3: the sender is a router and received a datagram destined for a host on another network
 - n Router check its routing table
 - n Find the IP address of the next router
 - n Use ARP to find the next router's physical address
- Case 4: the sender is a router that has received a datagram destined for a host in the same network
 - n Use ARP to find this host's physical address

Four Cases Using ARP: Case 3

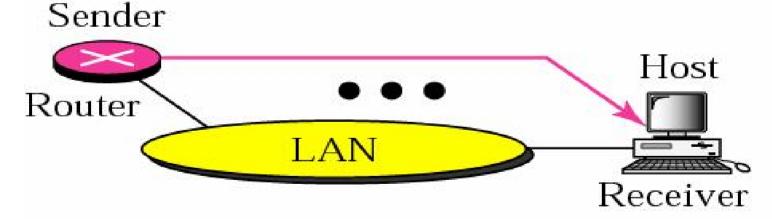
Target IP address: IP address of the appropriate router found in the routing table



Case 3. A router receives a packet to be sent to a host on another network. It must first be delivered to the appropriate router.

Four Cases Using ARP: Case 4

Target IP address: Destination address in the IP datagram



Case 4. A router receives a packet to be sent to a host on the same network.

Note

An ARP request is broadcast; an ARP reply is unicast.

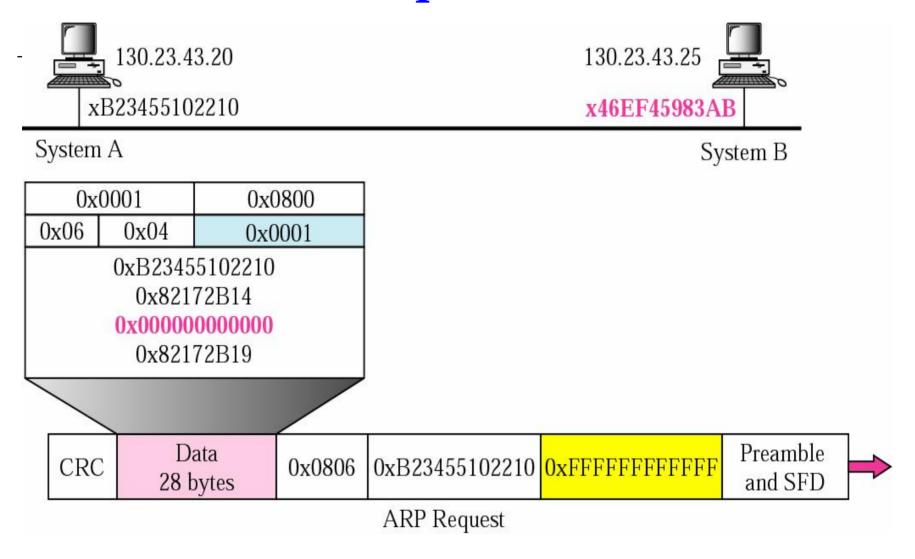
Example 1

- A host with IP address 130.23.43.20 and physical address 0xB23455102210
- Another host with IP address 130.23.43.25 and physical address 0xA46EF45983AB.
- The two hosts are on the same Ethernet network
- Show the ARP request and reply packets encapsulated in Ethernet frames

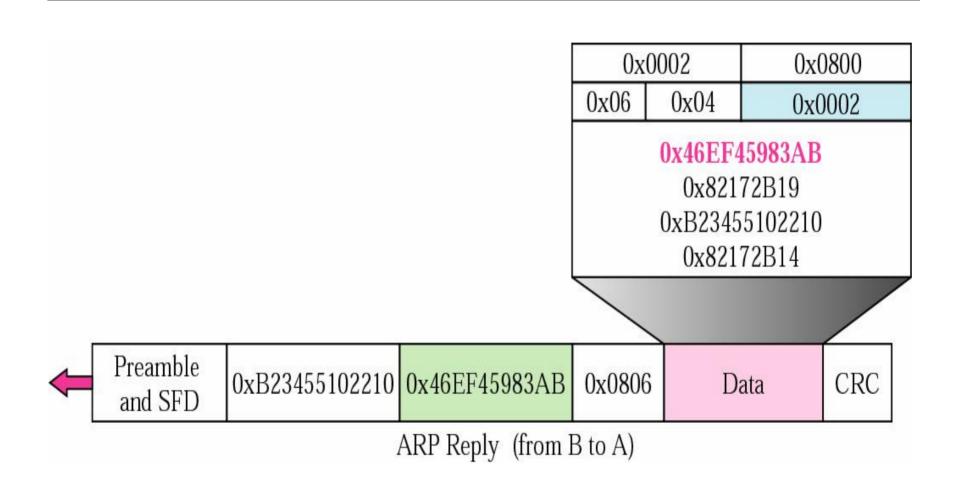
Solution

- Figure 7.7 shows the ARP request and reply packets
- Note that
 - n The IP addresses are shown in hexadecimal

Example 1



Example 1 (Continued)



Proxy ARP

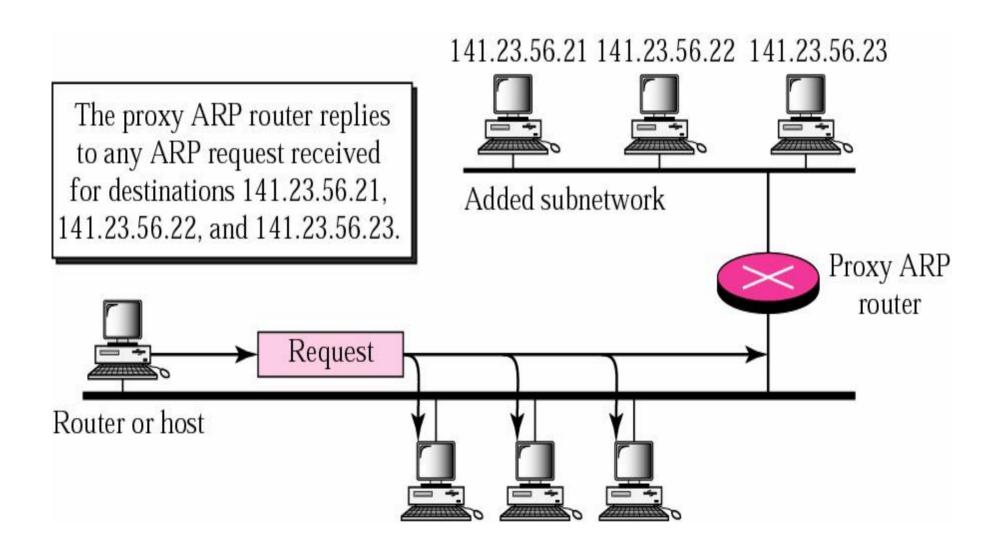
- Used to create a subnetting effect
- A router running a proxy ARP
 - n Its ARP acts on behalf of a set of hosts
 - n If it receives an ARP request message looking for the address of one of these host
 - The router sends an ARP reply announcing its own hardware (physical) address
 - n After the router receives the actual IP packet
 - It sends the packet to the appropriate host or router

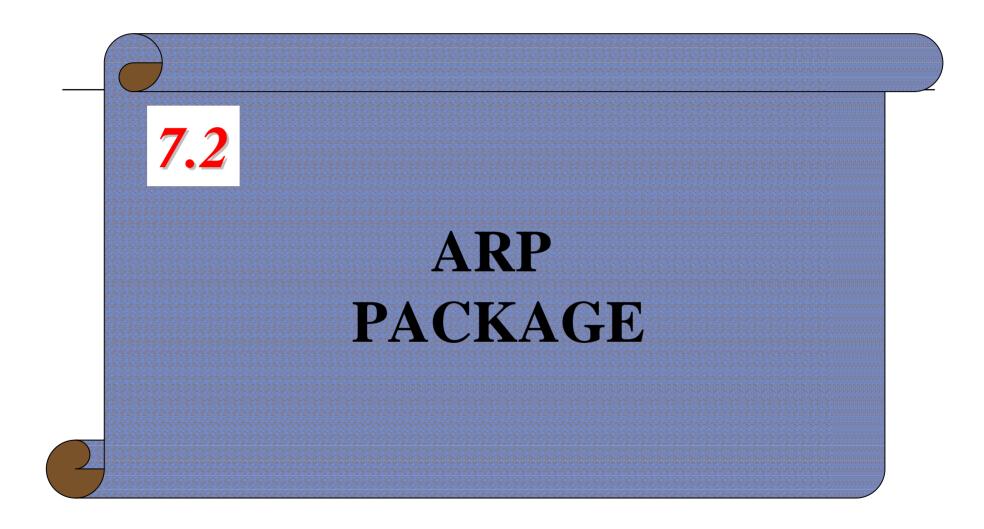
Example

• Administrator need to create a subnet without changing the whole system

Add a router running a proxy ARP

Proxy ARP

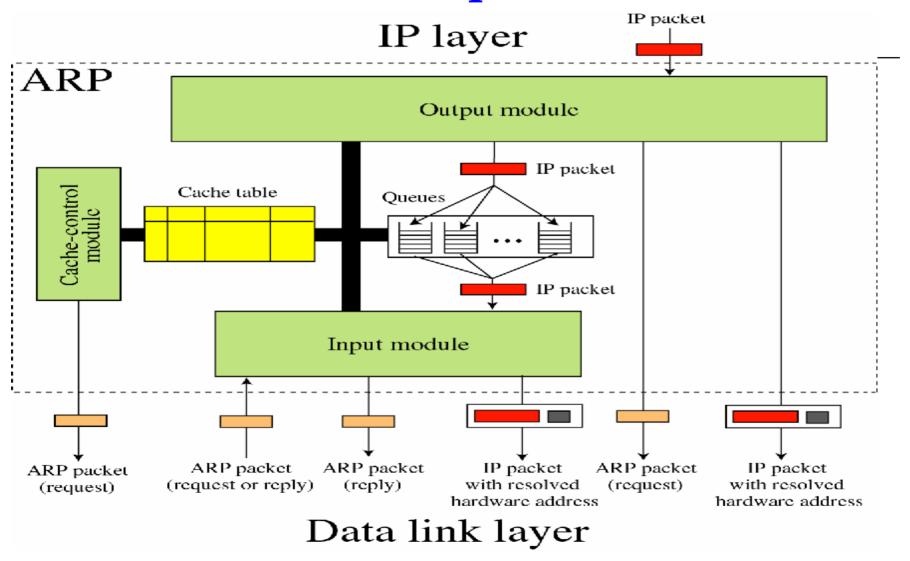




ARP Package

- Five components in an ARP package
 - n A cache table
 - n Queues
 - n An output module
 - n An input module
 - n A cache-control module

ARP Components



Cache Table

- Inefficient to use ARP to each datagram destined for the same host or router
 - n Introduce the cache table

• Cache table: an array of entries that contains the following's entries

Content of a Cache Table Entry

- State:
 - n FREE: the lime-to-live for this entry has expired
 - n PENDING: a request for this entry has been sent, but the reply has not yet been received
 - n RESOLVED: the entry is complete and valid
- Hardware type
- Protocol type
- Hardware length
- Protocol length
 - n Above fields are all the same as in the ARP packet

Content of a Cache Table Entry (Cont.)

- Interface number
- Queue number: ARP uses numbered queues to enqueue the packet waiting for address resolution
- Attempts: the number of times an ARP request is sent out for this entry
- Time-out: the lifetime of an entry in seconds
- Hardware address: the destination hardware address
- Protocol address: the destination IP address

Queues

- ARP package maintains a set of queues to hold the IP packets while ARP tries to resolve the hardware address
- Packets for the same destination are usually enqueued in the same queue
- The output module sends unsolved packets into the queue
- The input module removes a packet from the queue and sends it, with the resolved physical address, to data link layer for transmission

Output Module

- Wait until an IP packet from the IP software
- Check the cache table if receiving a IP packet
 - n If found and state = RESOLVED
 - Passed to the data link layer for transmission
 - n If found and state = PENDING
 - Send packet to this queue and wait
 - n If not found
 - Create an entry with state = PENDING
 - Create a queue and enqueue this packet
 - Send an ARP request

Input Module

- Wait until an ARP packet (request or reply) arrives and check the cache table
 - n If found state = PENDING
 - Copy the target hardware address in the packet
 - Change the state to RESOLVED
 - Set the value of TIME-OUT for this entry
 - Dequeue the packets from the corresponding queue and set them to the data link layer

Input Module (Cont.)

- n If found and state = RESOLVED
 - Copy the target hardware address in the packet
 - Set the value of TIME-OUT for this entry
 - This is because the target hardware address could have been changed
- n If not found
 - Create a new entry and adds it to the table
- If the packet is a request
 - n Send an ARP reply

Cache-Control Module

- Maintain the cache table by periodically check the cache table, entry by entry
- If state is PENDING
 - n Increment the value of attempts by 1
 - n If (attempts greater than maximum)
 - Change the state to FREE and Destroy the corresponding queue
 - n Else
 - Send an ARP request

Cache-Control Module (Cont.)

- If state is RESOLVED
 - n Decrement the value of time-out by the value of elapsed time
 - n If (time-out ≤ 0)
 - Change the state to FREE
 - Destroy the corresponding queue
- If state is FREE
 - n Continue to the next entry

Original Cache Table

S <u>tate</u> (Queue A	Attempt	Time-out	Protocol Addr.	Hardware Addr.
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
F					
R	9		60	19.1.7.82	4573E3242ACA
P	18	3		188.11.8.71	

Example 2

- The ARP output module receives an IP datagram from the IP layer with the destination address 114.5.7.89
- It checks the cache table and finds that an entry exists for this destination with the RESOLVED state
- It extracts the hardware address, which is 457342ACAE32, and sends the packet and the address to the data link layer

Example 3

- Twenty seconds later, the ARP output module receives an IP datagram from the IP layer with the destination address 116.1.7.22.
- It checks the cache table and does not find this destination in the table
- The module adds an entry to the table with the state PENDING and the Attempt value 1
- It also creates a new queue for this destination and enqueues the packet
- It then sends an ARP request to the data link layer for this destination

Cache table for Example 3

State	e Queue	? Attem	pt Time-o	ut Protocol Addr.	Hardware Addr.
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
P	23	1		116.1.7.22	
R	9		60	19.1.7.82	4573E3242ACA
P	18	3		188.11.8.71	

Example 4

- Fifteen seconds later, the ARP input module receives an ARP packet with target protocol (IP) address 188.11.8.71
- The module checks the table and finds this address
- It changes the state of the entry to RESOLVED and sets the time-out value to 900
- The module then adds the target hardware address (E34573242ACA) to the entry
- Now it accesses queue 18 and sends all the packets in this queue, one by one, to the data link layer

Cache table for Example 4

State	e Queue	? Attem	pt Time-o	ut Protocol Addr.	Hardware Addr.
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
P	23	1		116.1.7.22	
R	9		60	19.1.7.82	4573E3242ACA
R	18		900	188.11.8.71	E34573242ACA

Example 5

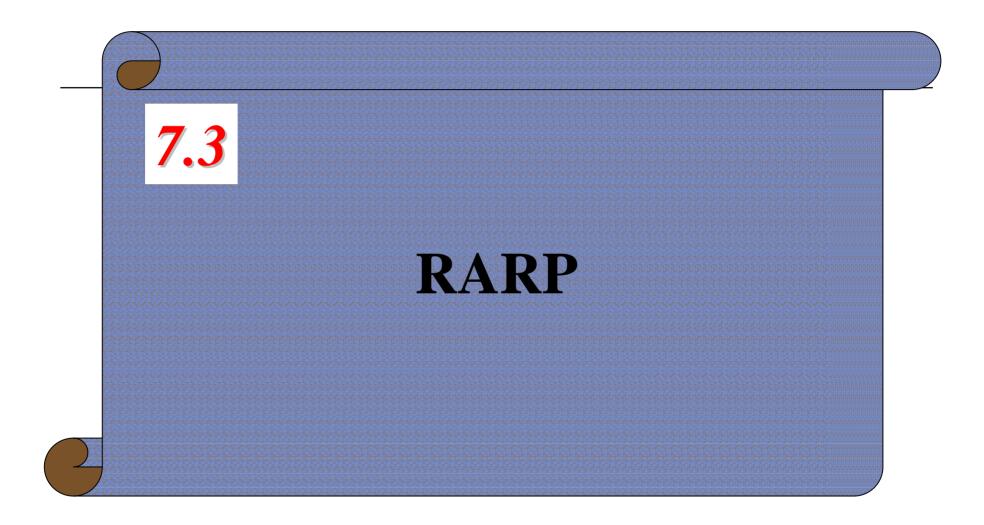
- Twenty-five seconds later, the cache-control module waits up
- The time-out values for the first three resolved entries are decremented by 60
- The time-out value for the last resolved entry is decremented by 25
- The state of the next-to-the last entry is changed to FREE because the time-out is zero

Example 5 (Cont.)

- For each of the three pending entries, the value of the attempts field is incremented by one
- o Then, the attempts value for one entry (the one with IP protocol address 201.11.56.7) is more than the maximum
 - n the state is changed to FREE, the queue is deleted
 - n An ICMP message is sent to the original destination

Cache Table for Example 5

State	Queue	Attem	pt Time-o	ut Protocol Addr.	Hardware Addr.
R	5		840	180.3.6.1	ACAE32457342
P	2	3		129.34.4.8	
F					
R	8		390	114.5.7.89	457342ACAE32
P	12	2		220.55.5.7	
P	23	2		116.1.7.22	
F					
R	18		875	188.11.8.71	E34573242ACA

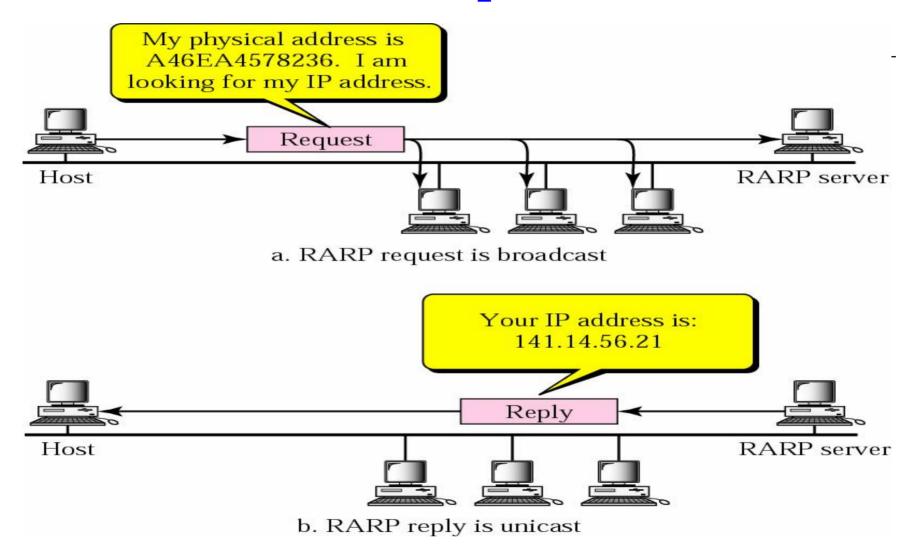


RARP

- A diskless machine is usually booted from ROM
- It cannot include the IP address
 - n IP address are assigned by the network administrator
- Obtain its logical address by the physical address using the RARP protocol

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RARP Operation



Note

The RARP request packets are broadcast; the RARP reply packets are unicast.

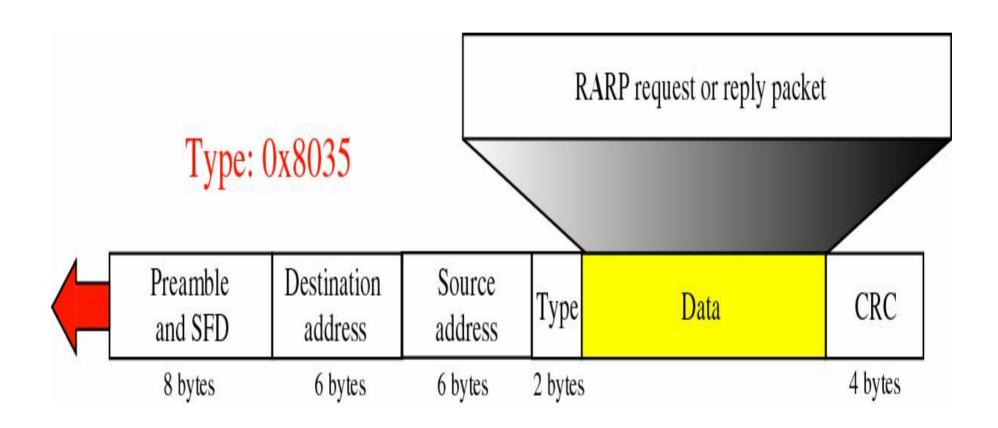
Packet Format

- The format of the RARP packet is the same as the ARP packet
- Except that the operation field is
 - n Three for RARP request message
 - n Four for RARP reply message

RARP Packet

Hardwa	are type	Protocol type			
Hardware length	Protocol length	Operation Request 3, Reply 4			
	Sender hardware address (For example, 6 bytes for Ethernet)				
Sender protocol address (For example, 4 bytes for IP) (It is not filled for request)					
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled for request)					
Target protocol address (For example, 4 bytes for IP) (It is not filled for request)					

Encapsulation of RARP Packet



Alternative Solutions to RARP

- When a diskless computer is booted, it needs more information in addition to its IP address
 - n The subnet mask
 - n The IP address of a router
 - n The IP address of a name server
- RARP cannot provide this extra information
- Two protocols, BOOTP and DHCP, can be used instead of RARP