



COMPUTER ENGINEERING



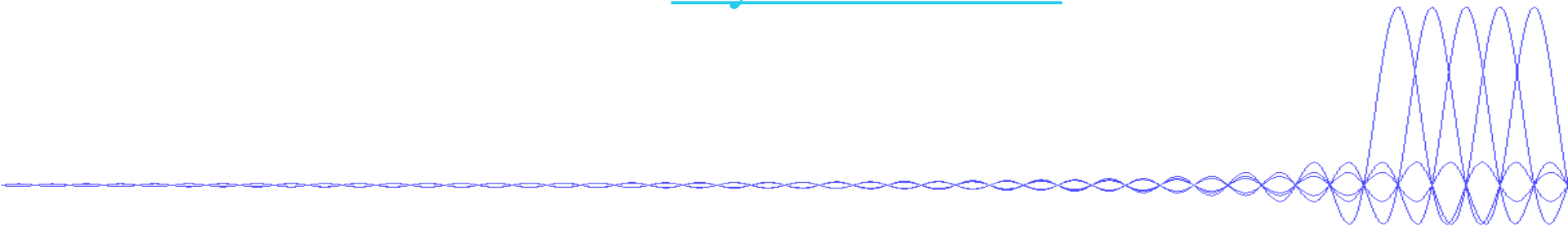
**UIT**  
TRƯỜNG ĐẠI HỌC  
CÔNG NGHỆ THÔNG TIN

# EMBEDDED SYSTEM DESIGN

## INPUT AND OUTPUT

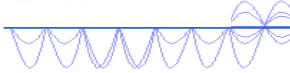
Doan Duy, Ph. D.

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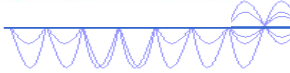




# Objective and Content



- Overview of IO interface
- Hardware and Software IO
- Interrupt



- Overview of IO interface
- Hardware and Software IO
- Interrupt



# Connecting the Analog and Digital Worlds

## Cyber:

- Digital
- Discrete in time
- Sequential

## Physical:

- Continuum
- Continuous in time
- Concurrent

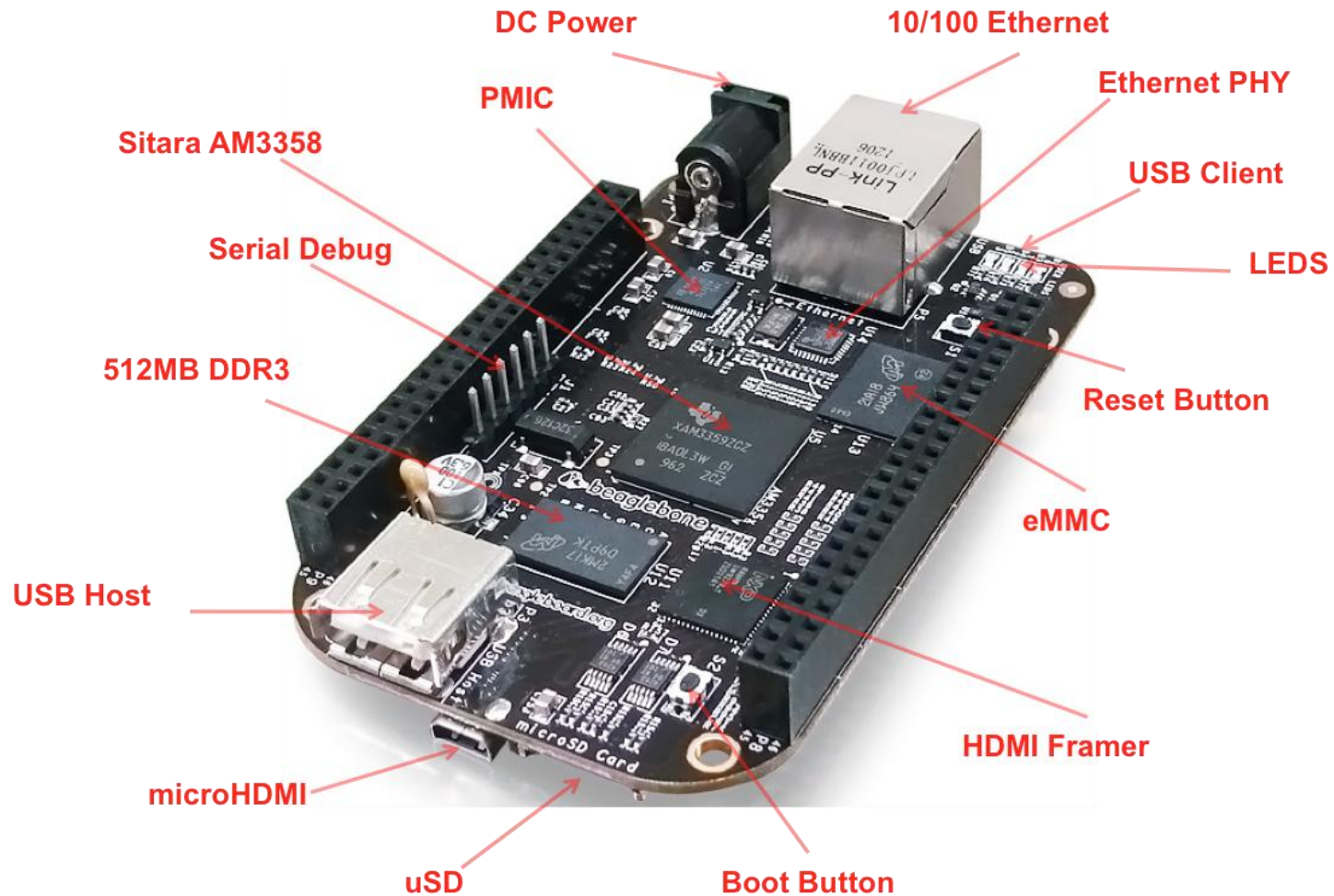


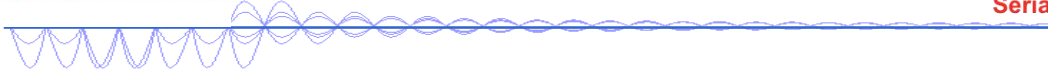
# Practical Issues

- Analog vs. digital
- Wired vs. wireless
- Serial vs. parallel
- Sampled or event triggered
- Bit rates
- Access control, security, authentication
- Physical connectors
- Electrical requirements (voltages and currents)

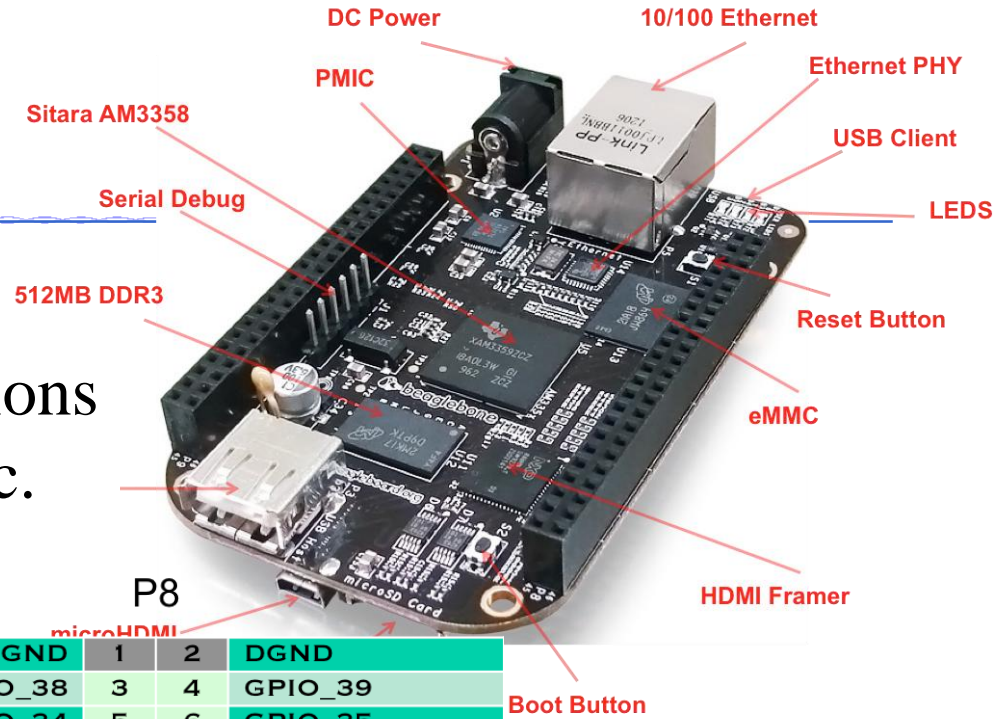


# A Typical Microcomputer Board





■ One of eight configurations with SPI buses, analog I/O, etc.



P9

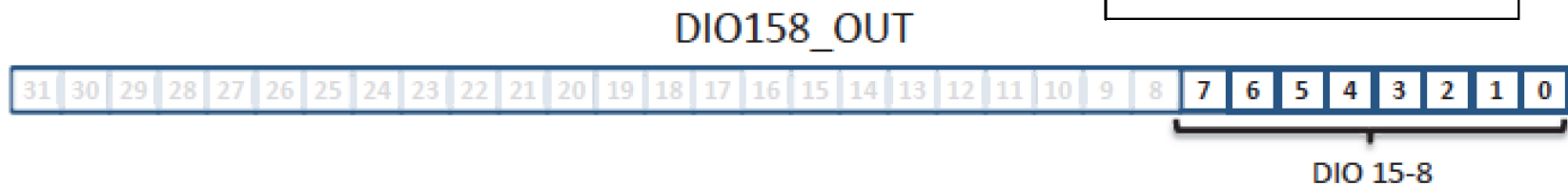
DGND	1	2	DGND
VDD_3V3	3	4	VDD_3V3
VDD_5V	5	6	VDD_5V
SYS_5V	7	8	SYS_5V
PWR_BUT	9	10	SYS_RESETh
GPIO_30	11	12	GPIO_60
GPIO_31	13	14	GPIO_40
GPIO_48	15	16	GPIO_51
SPIO_CS0	17	18	SPIO_D1
SPI1_CS1	19	20	SPI1_CS0
SPIO_DO	21	22	SPIO_SCLK
GPIO_49	23	24	GPIO_15
GPIO_117	25	26	GPIO_14
GPIO_125	27	28	SPI1_CS0
SPI1_DO	29	30	SPI1_D1
SPI1_SCLK	31	32	VDD_ADC
AIN4	33	34	GNDA_ADC
AIN6	35	36	AIN5
AIN2	37	38	AIN3
AIN0	39	40	AIN1
GPIO_20	41	42	SPI1_CS1
DGND	43	44	DGND
DGND	45	46	DGND

P8

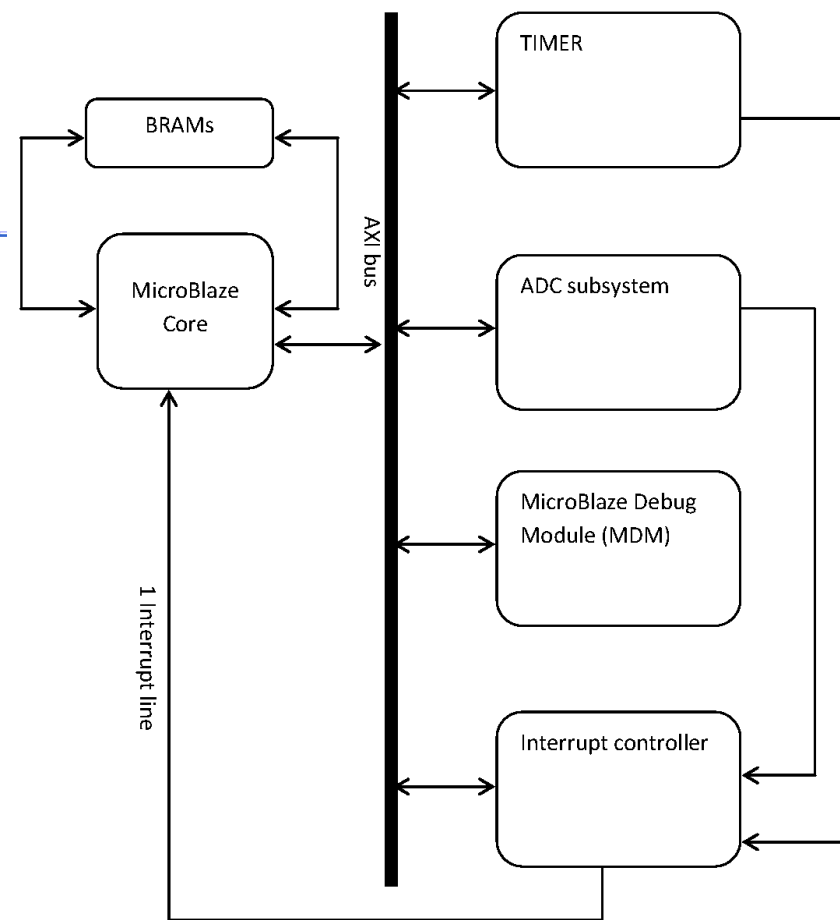
DGND	1	2	DGND
GPIO_38	3	4	GPIO_39
GPIO_34	5	6	GPIO_35
GPIO_66	7	8	GPIO_67
GPIO_69	9	10	GPIO_68
GPIO_45	11	12	GPIO_44
GPIO_23	13	14	GPIO_26
GPIO_47	15	16	GPIO_46
GPIO_27	17	18	GPIO_65
GPIO_22	19	20	GPIO_63
GPIO_62	21	22	GPIO_37
GPIO_36	23	24	GPIO_33
GPIO_32	25	26	GPIO_61
GPIO_86	27	28	GPIO_88
GPIO_87	29	30	GPIO_89
GPIO_10	31	32	GPIO_11
GPIO_9	33	34	GPIO_81
GPIO_8	35	36	GPIO_80
GPIO_78	37	38	GPIO_79
GPIO_76	39	40	GPIO_77
GPIO_74	41	42	GPIO_75
GPIO_72	43	44	GPIO_73
GPIO_70	45	46	GPIO_71

Many GPIO pins can be reconfigured to be PWM drivers, timers, etc.

DIO158\_OUT is a C preprocessor macro defined in a header file in your IDE project. It defines the memory address of this register.



Bit	Field	Access	Values
7 – 0	Digital Output	Write	myRIO DIO 15-8 output state



**Figure 1.10:** DIO158\_OUT (DIO 15-8 Out) register. myRIO MXP Connector B pins 8-15. Bit 0 writes MXP Connector B DIO8, and bit 7 writes DIO15.





- Overview of IO interface
- **Hardware and Software IO**
- Interrupt



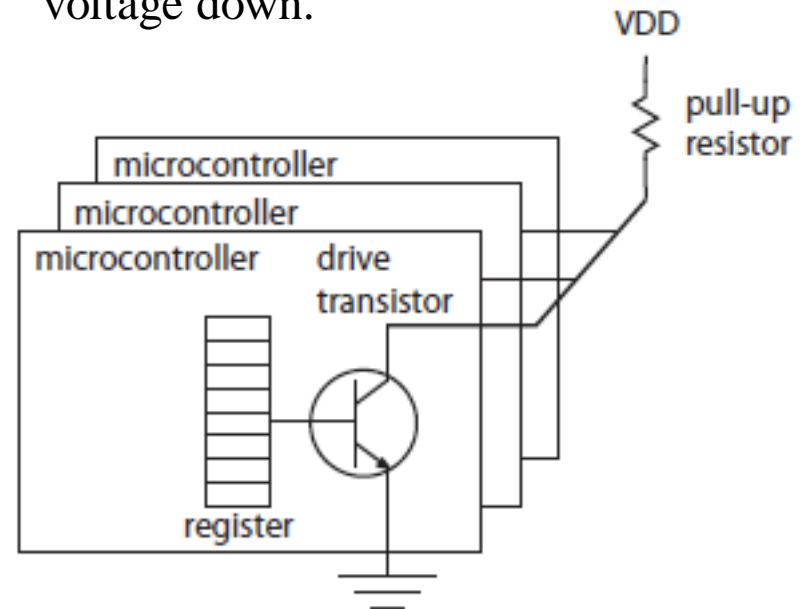
# Typical GPIO

- **Open collector** circuits are often used on GPIO (general-purpose I/O) pins of a microcontroller.

- The same pin can be used for input and output. And multiple users can connect to the same bus.

- The current is limited!?

GPIO pins configured for bus output. Any one controller can pull the bus voltage down.

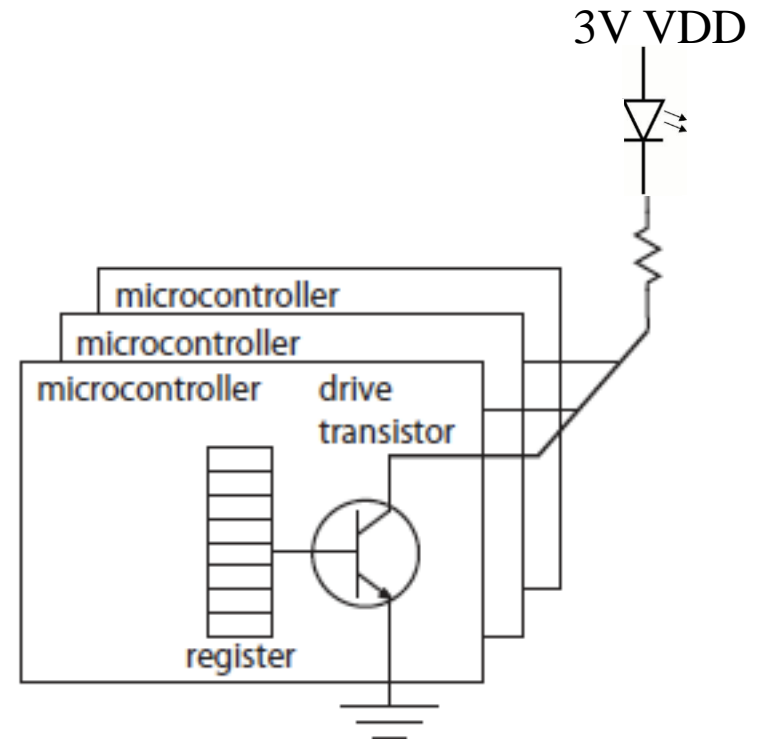




# Example of LED connection

■ Assume GPIO pins can sink up to 18 mA. Assume the LED, when forward biased (turned on), has a voltage drop of 2 volts.

■ What resistor should you use?





# Example of LED connection

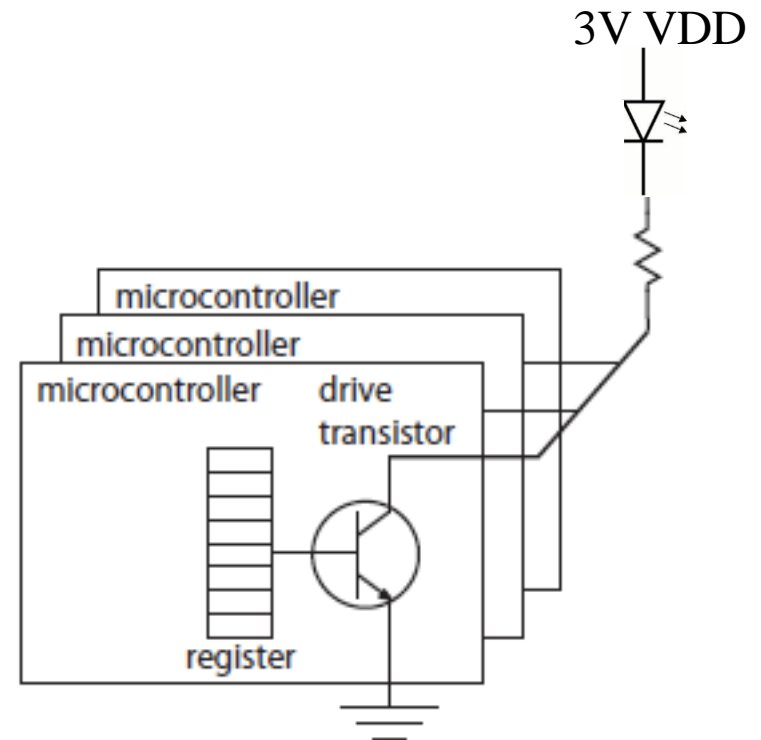
Ohm's law:

$$V = IR$$

When LED is on,  $V = 1$  volt.

To limit to 18mA,

$$R \geq 1/0.018 \approx 56 \text{ ohms}$$



# Wired Connections

## Parallel vs. Serial Digital Interfaces

### ○ Parallel (one wire per bit)

- ▣ ATA: Advanced Technology Attachment
- ▣ PCI: Peripheral Component Interface
- ▣ SCSI: Small Computer System Interface
- ▣ ...

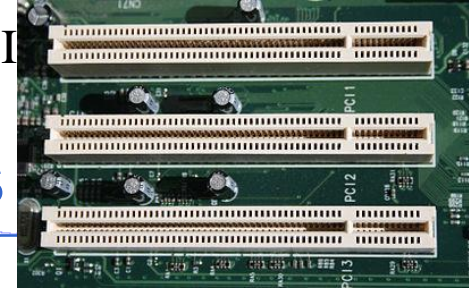
### ○ Serial (one wire per direction)

- ▣ RS-232
- ▣ SPI: Serial Peripheral Interface bus
- ▣ I<sup>2</sup>C: Inter-Integrated Circuit
- ▣ USB: Universal Serial Bus
- ▣ SATA: Serial ATA
- ▣ ...

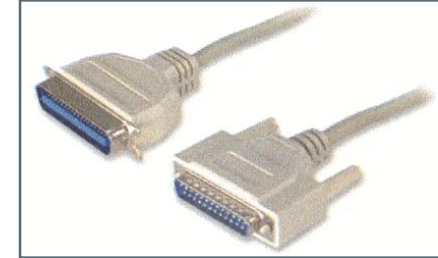
### ○ Mixed (one or more “lanes”)

- ▣ PCIe: PCI Express

PCI



SCSI



USB



RS-232

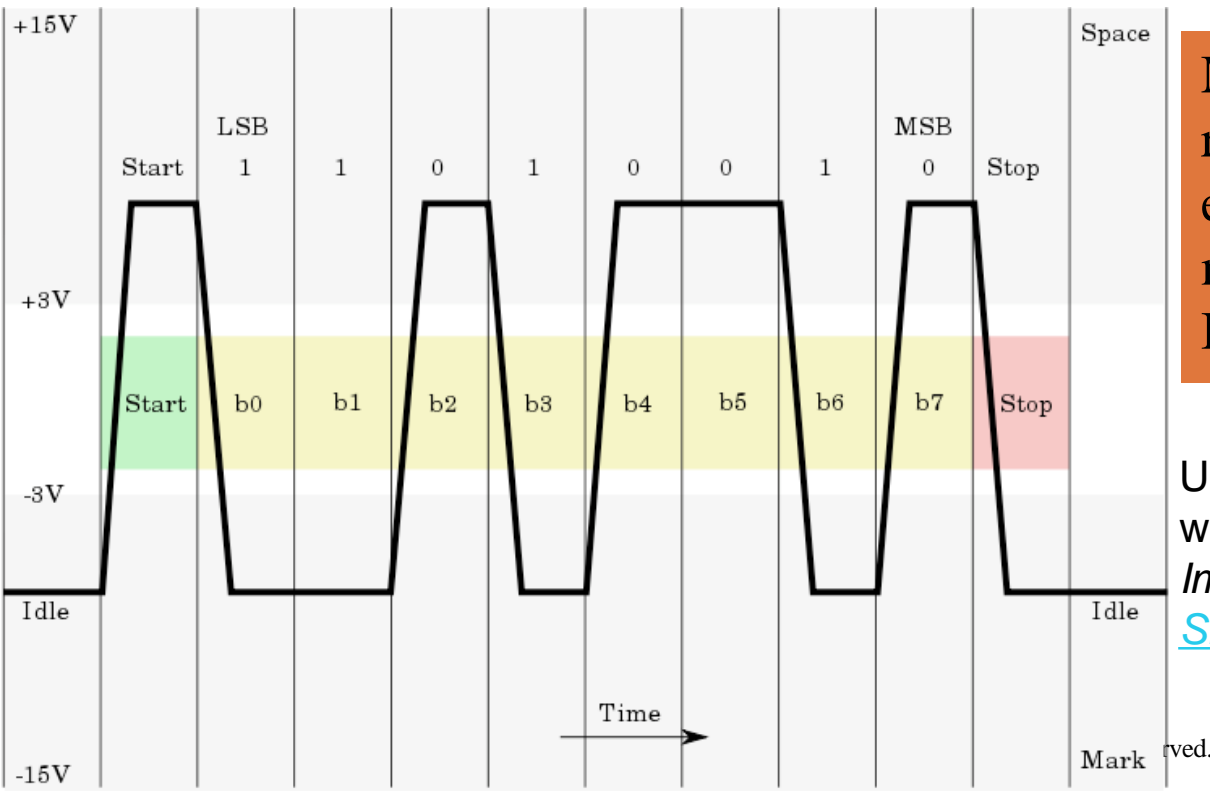




# Serial Interfaces

■ The old but persistent RS-232 standard supports asynchronous serial connections (no common clock).

■ How does it work?



Many uses of RS-232 are being replaced by USB, which is electrically simpler but with a more complex protocol, or Bluetooth, which is wireless.

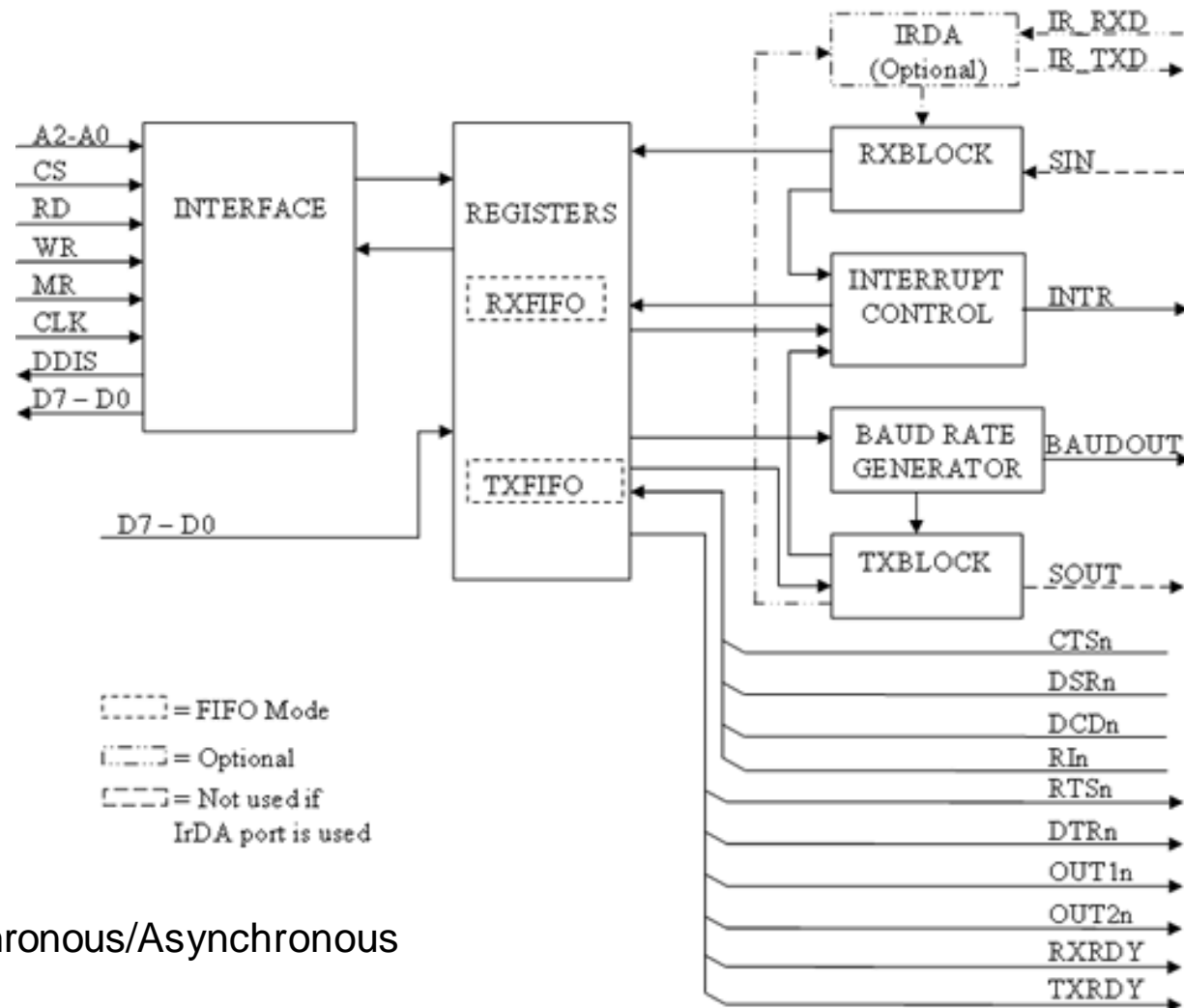
Uppercase ASCII "K" character (0x4b) with 1 start bit, 8 data bits, 1 stop bit.

Image license: [Creative Commons ShareAlike 1.0 License](#)



# UART: Universal Asynchronous Receiver-Transmitter

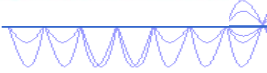
- Convert serial data to parallel data, and vice versa.
- Uses shift registers to load/store data
- Can raise interrupt when data is ready
- Commonly used with RS-232 interface



Variant: USART: Universal Synchronous/Asynchronous Receiver-Transmitter



# Input/Output Mechanisms in Software



## ○ Polling

- Main loop uses each I/O device periodically.
- If output is to be produced, produce it.
- If input is ready, read it.

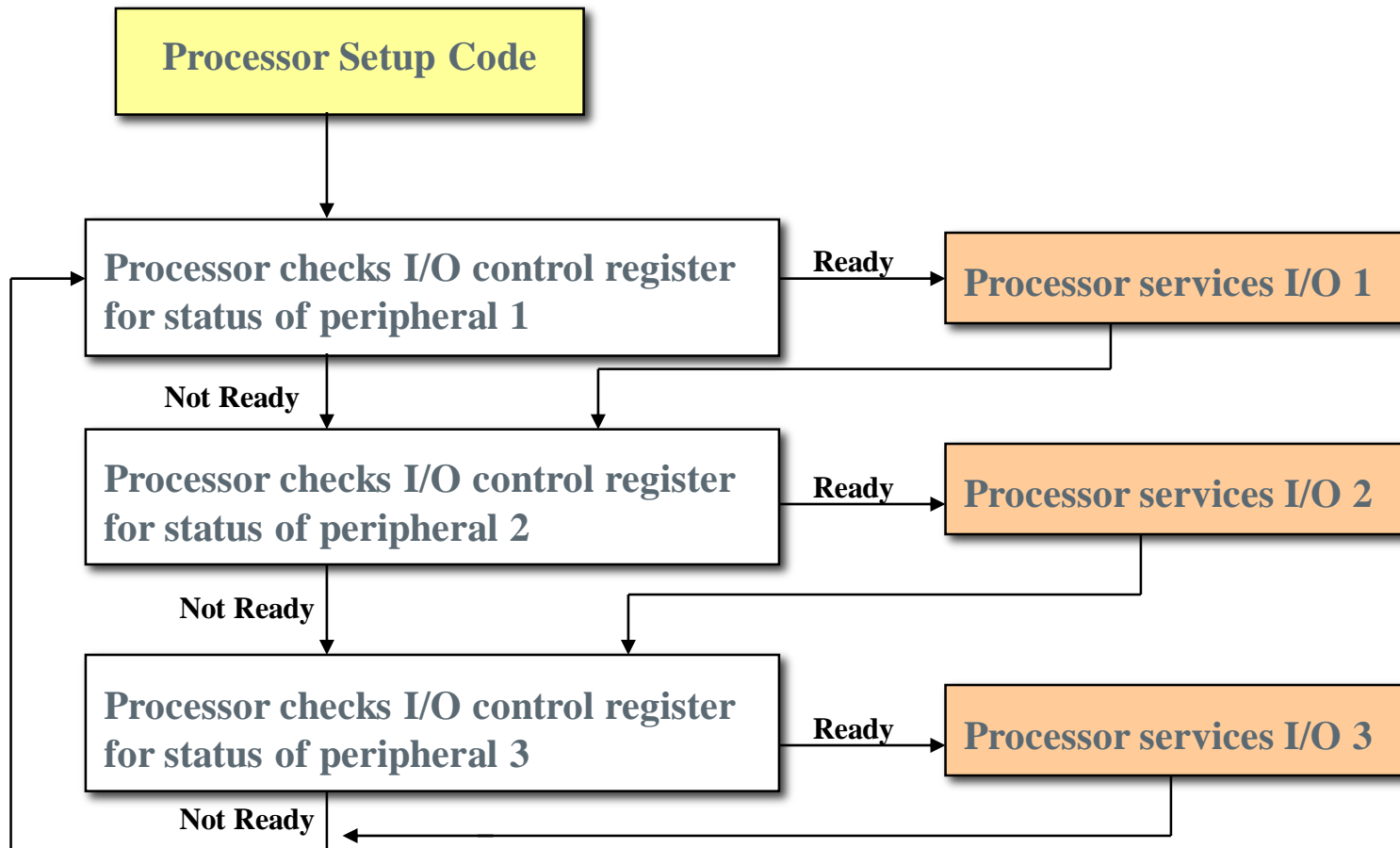
## ○ Interrupts

- External hardware alerts the processor that input is ready.
- Processor suspends what it is doing.
- Processor invokes an interrupt service routine (ISR).
- ISR interacts with the application concurrently.



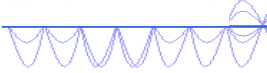


# Polling





# Using a Serial Interface



In an Atmel AVR 8-bit microcontroller, to send a byte over a serial port, the following C code will do:

```
while( ! (UCSR0A & 0x20) ) ;  
UDR0 = x;
```

- x is a variable of type uint8.
- UCSR0A and UDR0 are variables defined in a header.
- They refer to memory-mapped registers in the UART (Universal Asynchronous Receiver-Transmitter)



# Send a Sequence of Bytes

```
for(i = 0; i < 8; i++) {  
    while(!(UCSR0A & 0x20));  
    UDR0 = x[i];  
}
```

How long will this take to execute? Assume:

- 57600 baud serial speed.
- $8/57600 = 139$  microseconds.
- Processor operates at 18 MHz.

Each for loop iteration will consume about 2502 cycles.



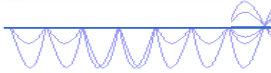
# Receiving via UART

Again, on an Atmel AVR:

```
while ( ! (UCSR0A & 0x80) ) ;  
return UDR0 ;
```

- Wait until the UART has received an incoming byte.
- *The programmer must ensure there will be one!*
- If reading a sequence of bytes, how long will this take?

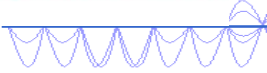
Under the same assumptions as before, it will take about 2502 cycles to receive each byte.



- Overview of IO interface
- Hardware and Software IO
- **Interrupt**



# Input/Output Mechanisms in Software

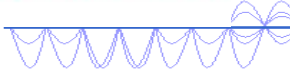


## ○ Polling

- Main loop uses each I/O device periodically.
- If output is to be produced, produce it.
- If input is ready, read it.

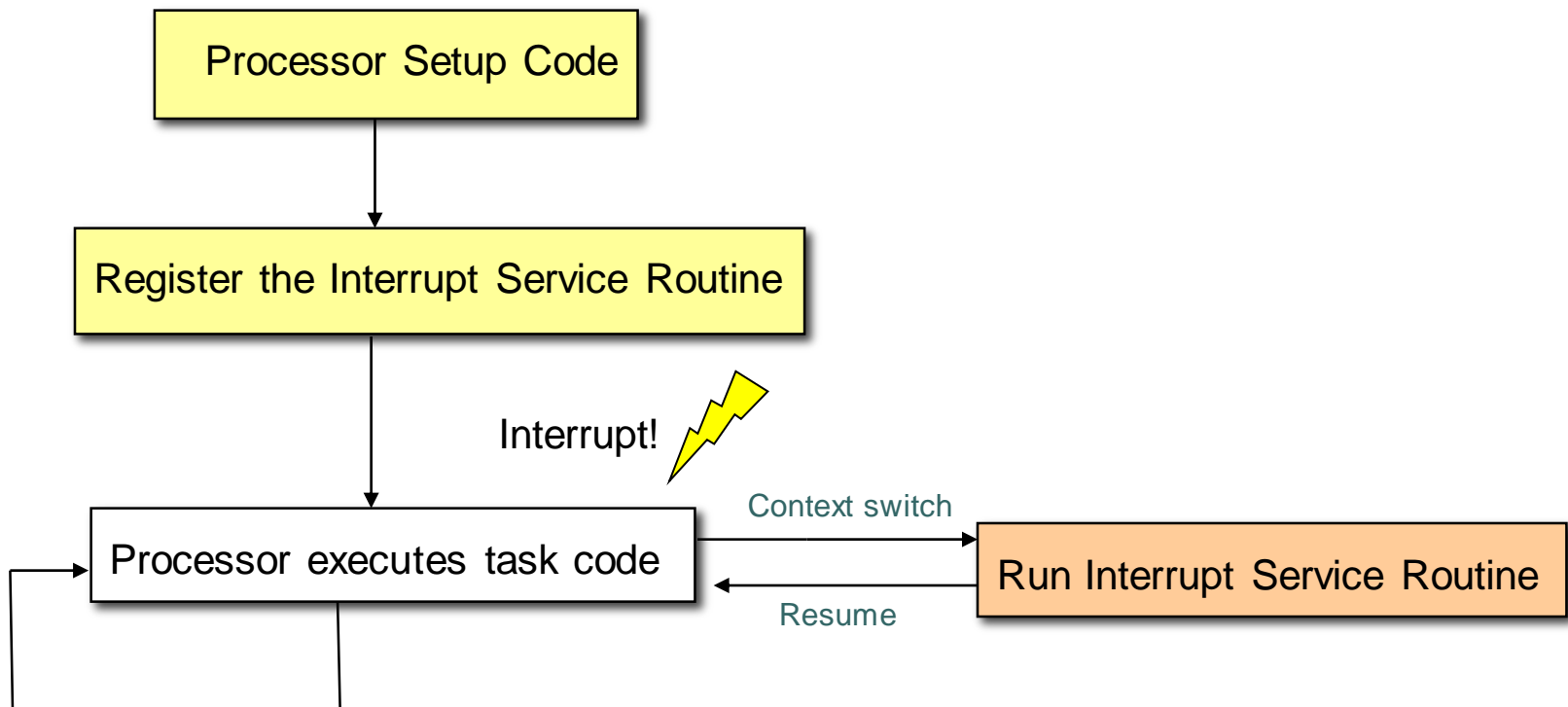
## ○ Interrupts

- External hardware alerts the processor that input is ready.
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- Processor invokes an interrupt service routine (ISR).
- ISR interacts with the application concurrently.



## ○ Interrupt Service Routine

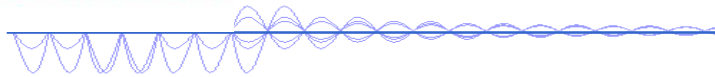
Short subroutine that handles the interrupt





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# Interrupts



Program memory addresses,  
not data memory addresses.

Address	Labels	Code	Comments
0x0000	jmp	RESET	; Reset Handler
0x0002	jmp	EXT_INT0	; IRQ0 Handler
0x0004	jmp	EXT_INT1	; IRQ1 Handler
0x0006	jmp	PCINT0	; PCINT0 Handler
0x0008	jmp	PCINT1	; PCINT1 Handler
0x000A	jmp	PCINT2	; PCINT2 Handler
0x000C	jmp	WDT	; Watchdog Timer Handler
0x000E	jmp	TIM2_COMPA	; Timer2 Compare A Handler
0x0010	jmp	TIM2_COMPB	; Timer2 Compare B Handler
0x0012	jmp	TIM2_OVF	; Timer2 Overflow Handler
0x0014	jmp	TIM1_CAPT	; Timer1 Capture Handler

## ■ Triggers:

- A level change on an interrupt request pin
- Writing to an interrupt pin configured as an output ( “ software interrupt”) or executing special instruction

Source: ATmega168 Reference Manual

## ■ Responses:

- Disable interrupts.
- Push the current program counter onto the stack.
- Execute the instruction at a designated address in program memory.

## ■ Design of interrupt service routine:

- Save and restore any registers it uses.
- Re-enable interrupts before returning from interrupt.





# Memory Map

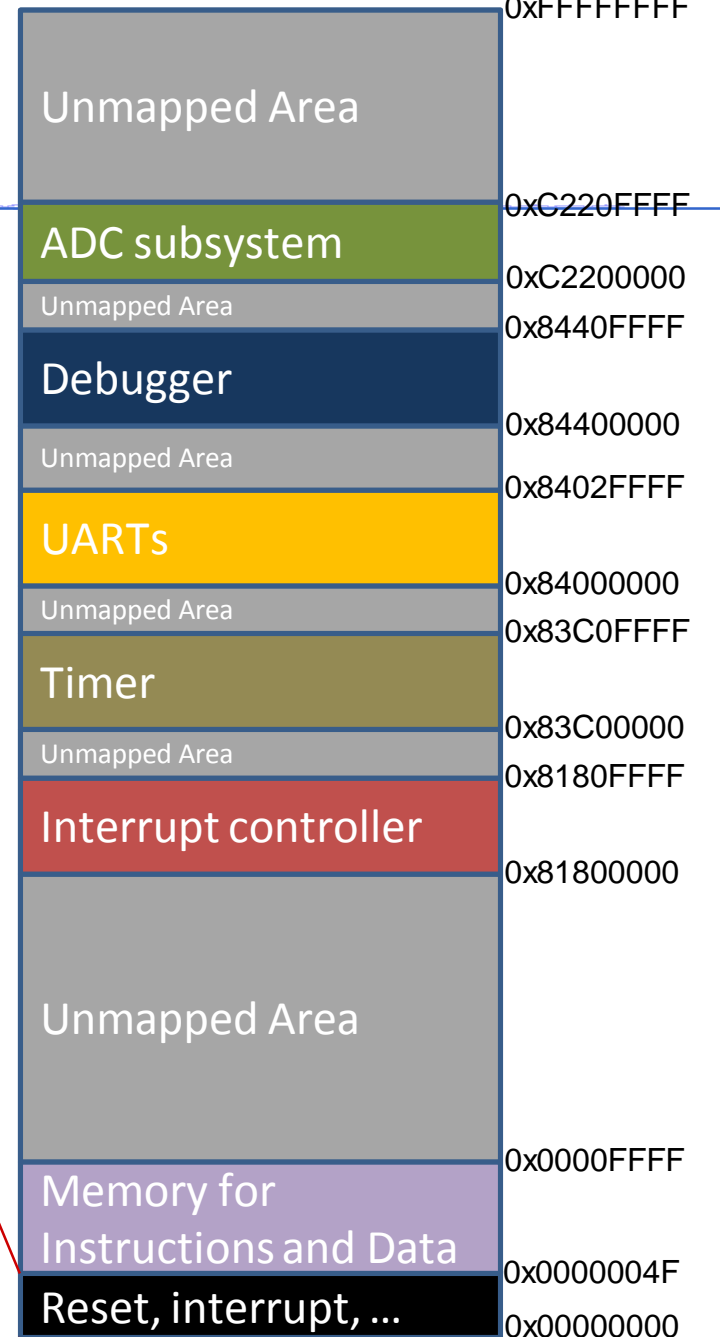
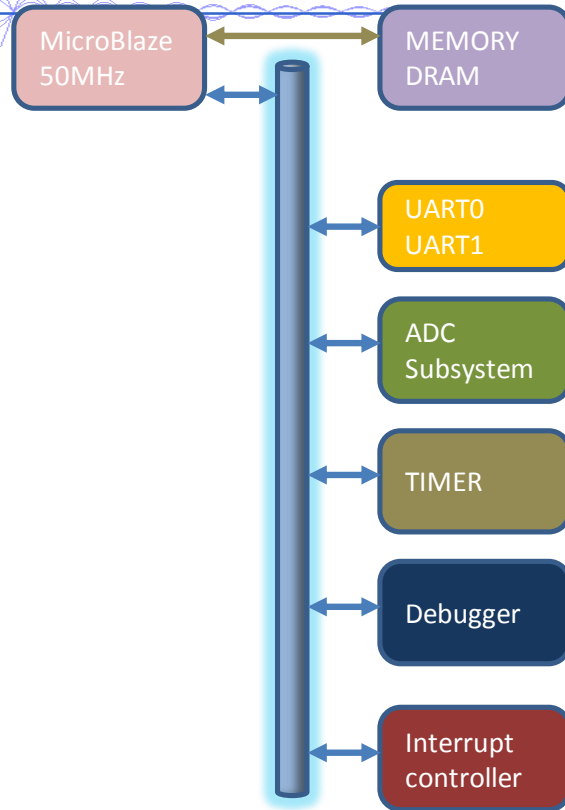
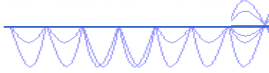


Table 3-4: Interrupt and Exception Handling

On	Hardware jumps to	Software Labels
Start / Reset	0x0	_start
User exception	0x8	_exception_handler
Interrupt	0x10	_interrupt_handler
Break (HW/SW)	0x18	-
Hardware exception	0x20	_hw_exception_handler
Reserved by Xilinx for future use	0x28 - 0x4F	-



# Interrupts are Evil

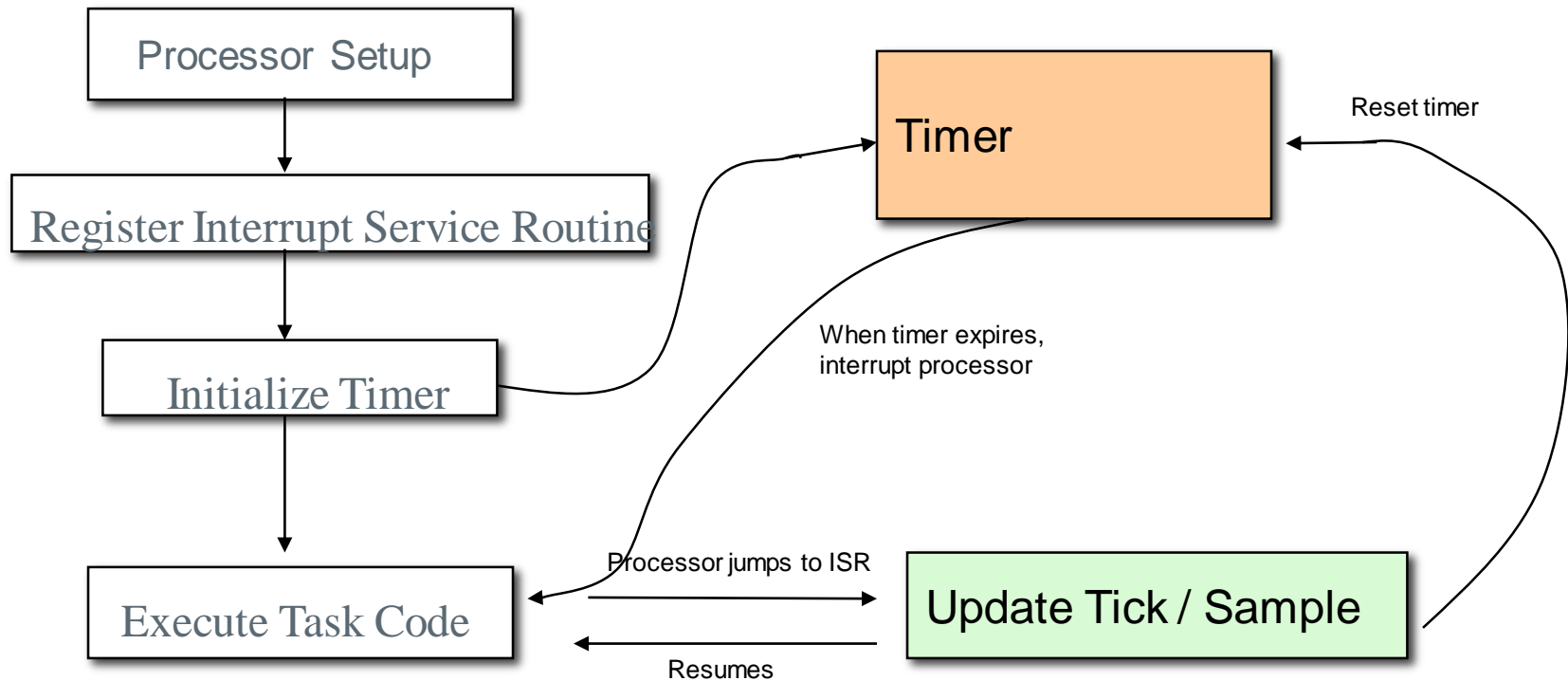


In one or two respects modern machinery is basically more difficult to handle than the old machinery. Firstly, we have got the interrupts, occurring at unpredictable and irreproducible moments; compared with the old sequential machine that pretended to be a fully deterministic automaton, this has been a dramatic change, and many a systems programmer's grey hair bears witness to the fact that we should not talk lightly about the logical problems created by that feature.

(Dijkstra, “The humble programmer” 1972)



# Timed Interrupt





# Concurrency

```
volatile uint timer_count;
void ISR(void) {
    timer_count--;
}

int main(void) {
    // initialization code
    SysTickIntRegister(&ISR);
    ... // other init
    timer_count = 2000;
    while(timer_count != 0) {
        ... code to run for 2 seconds
    }
}
```

concurrent code:  
logically runs at the  
same time. In this case,  
between any two  
machine instructions in  
main() an interrupt can  
occur and the upper  
code can execute.

## What could go wrong?



# Concurrency

```
volatile uint timer_count;
void ISR(void) {
    timer_count--;
}

int main(void) {
    // initialization code
    SysTickIntRegister(&ISR);
    ... // other init
    timer_count = 2000;
    while(timer_count != 0) {
        ... code to run for 2 seconds
    }
}
```

what if the interrupt  
occurs twice during  
the execution of this  
code?

What could go wrong?

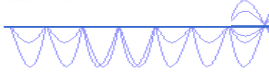


What' s the difference between

Concurrency  
and  
Parallelism



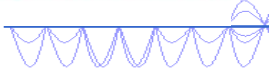
# Concurrency and Parallelism



- A program is said to be **concurrent** if different parts of the program conceptually execute simultaneously.
- A program is said to be **parallel** if different parts of the program physically execute simultaneously on distinct hardware.
- A parallel program is concurrent, but a concurrent program need not be parallel.



# Concurrency and Parallelism



- Interrupt Handling
  - ▣ Reacting to external events (interrupts)
  - ▣ Exception handling (software interrupts)
- Processes
  - ▣ Creating the illusion of simultaneously running different programs (multitasking)
- Threads
  - ▣ How is a thread different from a process?
- Multiple processors (multi-cores)
- . . .





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# Q&A





# Example: Set up a timer on an ATmega168 to trigger an interrupt every 1ms

The frequency of the processor in the command module is 18.432 MHz.

1. Set up an interrupt to occur once every millisecond. Toward the beginning of your program, set up and enable the timer1 interrupt with the following code:

```
TCCR1A = 0x00;  
TCCR1B = 0x0C;  
OCR1A = 71;  
TIMSK1 = 0x02;
```

The first two lines of the code put the timer in a mode in which it generates an interrupt and resets a counter when the timer value reaches the value of OCR1A, and select a prescaler value of 256, meaning that the timer runs at 1/256th the speed of the processor. The third line sets the reset value of the timer. To generate an interrupt every 1ms, the interrupt frequency will be 1000 Hz. To calculate the value for OCR1A, use the following formula:

```
OCR1A = (processor_frequency / (prescaler *  
interrupt_frequency)) - 1  
  
OCR1A = (18432000 / (256 * 1000)) - 1 = 71
```

The fourth line of the code enables the timer interrupt. See the ATmega168 datasheet for more information on these control registers.

- TCCR: Timer/Counter Control Register
- OCR: output compare register
- TIMSK: Timer Interrupt Mask

■ The “prescaler” value divides the system clock to drive the timer.

■ Setting a non-zero bit in the timer interrupt mask causes an interrupt to occur when the timer resets.

Source: iRobot Command Module Reference Manual v6



# Setting up the timer interrupt hardware in C

```
#include <avr/io.h>
```

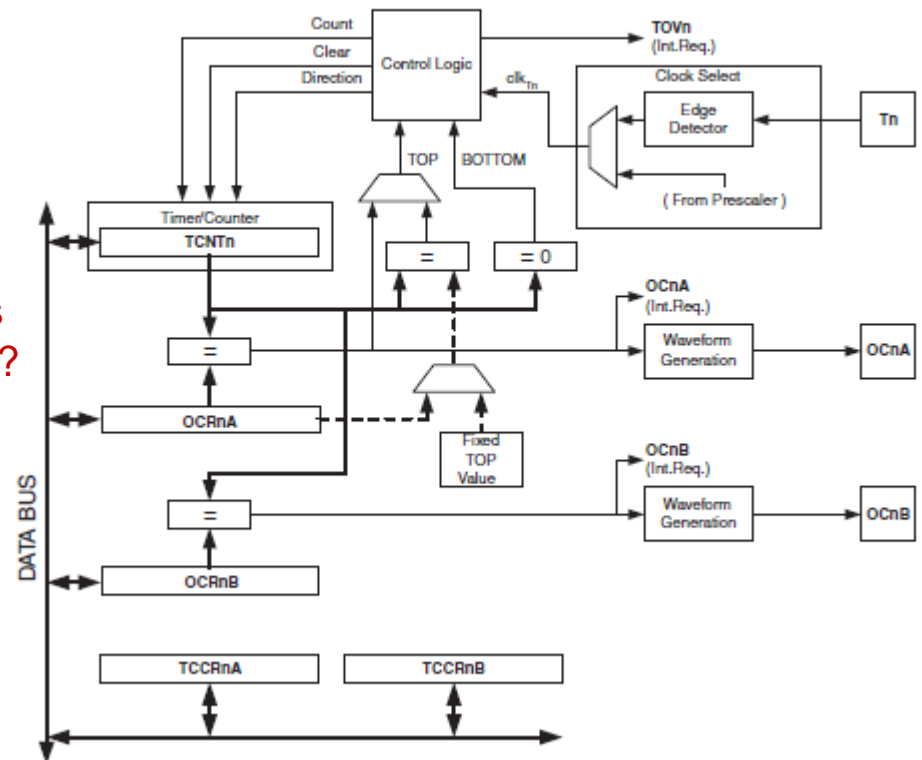
```
int main (void) {  
    TCCR1A = 0x00;  
    TCCR1B = 0x0C;  
    OCR1A = 71;  
    TIMSK1 = 0x02;  
    ...  
}
```

memory-  
mapped  
register.

But how is this  
proper C code?

This code sets the hardware up  
to trigger an interrupt every 1ms.  
How do we handle the interrupt?

Figure 16-1. 8-bit Timer/Counter Block Diagram



Source: ATmega168 Reference Manual



```
#define _MMIO_BYTE(mem_addr) (*(volatile uint8_t*)(mem_addr))
#define _SFR_IO8(io_addr) _MMIO_BYTE((io_addr) + 0x20)
#define _SFR_MEM8(mem_addr) _MMIO_BYTE(mem_addr)
#define BV(bit) (1 << (bit))
```

```
//Timer defines (iomx8.h)
#define TCCR1A _SFR_MEM8(0x80)
#define TCCR1B _SFR_MEM8(0x81)
/* TCCR1B */
#define WGM12 3
#define CS12 2
```

```
//Enable interrupts (interrupt.h)
#define sei() __asm__ __volatile__ ("sei" ::)
//Disable interrupts (interrupt.h)
#define cli() __asm__ __volatile__ ("cli" ::)
#define SIGNAL(signame) \
void signame(void) __attribute__((signal)); \
void signame(void)
```

SEI		Global Interrupt Enable
CLI		Global Interrupt Disable

```
void initialize(void) {
    cli();

    // Set I/O pins
    DDRB = 0x10;
    PORTB = 0xCF;
    .....

    // Set up timer 1 to generate an interrupt every 1 ms
    TCCR1A = 0x00;
    TCCR1B = (_BV(WGM12) | _BV(CS12));
    OCR1A = 71;
    TIMSK1 = _BV(OCIE1A);

    // Set up the serial port with rx interrupt
    .....

    // Turn on interrupts
    sei();
}
```

```
// Global variables
volatile uint16_t timer_cnt = 0;
volatile uint8_t timer_on = 0;

// Timer 1 interrupt to time delays in ms
SIGNAL(SIG_OUTPUT_COMPARE1A) {
    if(timer_cnt) {
        timer_cnt--;
    } else {
        timer_on = 0;
    }
}
```

```
void delayMs(uint16_t time_ms) {
    timer_on = 1;
    timer_cnt = time_ms;
    while(timer_on);
}
```



```
#define MMIO_BYTE(mem_addr) (*(volatile uint8_t*)(mem_addr))
#define SFR_IO8(io_addr) MMIO_BYTE((io_addr) + 0x20)
#define _SFR_MEM8(mem_addr) MMIO_BYTE(mem_addr)
#define BV(bit) (1 << (bit))
```

```
//Timer defines (iomx8.h)
```

```
#define TCCR1A _SFR_MEM8(0x80)
```

```
#define TCCR1B _SFR_MEM8(0x81)
```

```
/* TCCR1B */
```

```
#define WGM12 3
```

```
#define CS12 2
```

```
//Enable interrupts (interrupt.h)
```

```
# define sei() __asm__ __volatile__ ("sei" ::)
```

```
//Disable interrupts (interrupt.h)
```

```
# define cli() __asm__ __volatile__ ("cli" ::)
```

```
#define SIGNAL(signame)
```

```
void signame (void) __attribute__((signal)); \
```

```
void signame (void)
```

SEI		Global Interrupt Enable
CLI		Global Interrupt Disable

```
void initialize(void) {
```

```
cli();
```

```
// Set I/O pins
```

```
DDRB = 0x10;
```

```
PORTB = 0xCF;
```

```
.....
```

```
// Set up timer 1 to generate an interrupt every 1 ms
```

```
TCCR1A = 0x00;
```

```
TCCR1B = (_BV(WGM12) | _BV(CS12));
```

```
OCR1A = 71;
```

```
TIMSK1 = _BV(OCIE1A);
```

```
// Set up the serial port with rx interrupt
```

```
.....
```

```
// Turn on interrupts
```

```
sei();
```

```
}
```

```
// Global variables
```

```
volatile uint16_t timer_cnt = 0;
```

```
volatile uint8_t timer_on = 0;
```

```
// Timer 1 interrupt to time delays in ms
```

```
SIGNAL(SIG_OUTPUT_COMPARE1A) {
```

```
if(timer_cnt) {
```

```
timer_cnt--;
```

```
} else {
```

```
timer_on = 0;
```

```
}
```

```
}
```

```
void delayMs(uint16_t time_ms) {
```

```
timer_on = 1;
```

```
timer_cnt = time_ms;
```

```
while(timer_on) ;
```

```
}
```

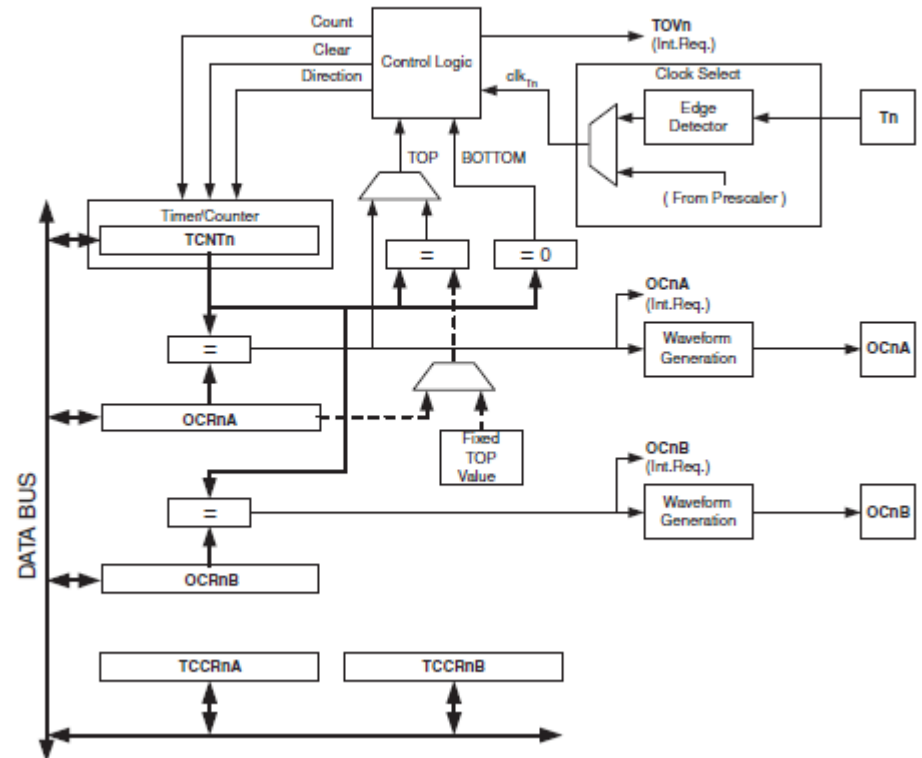


# Setting up the timer interrupt hardware in C

```
#include <avr/io.h>
```

```
int main (void) {  
    TCCR1A = 0x00;  
    TCCR1B = 0x0C;  
    OCR1A = 71;  
    TIMSK1 = 0x02;  
    ...  
}
```

Figure 16-1. 8-bit Timer/Counter Block Diagram



Source: ATmega168 Reference Manual



## Example 2: Set up a timer on a Luminary Micro board to trigger an interrupt every 1ms

```
■ // Setup and enable SysTick with interrupt every
  1ms
■ void initTimer(void) {
■     SysTickPeriodSet(SysCtlClockGet() / 1000);
■     SysTickEnable();
■     SysTickIntEnable();
■ }
■ // Disable SysTick
■ void disableTimer(void) {
■     SysTickIntDisable();
■     SysTickDisable();
■ }
```

Number of cycles per sec.

Start SysTick counter

Enable SysTick timer interrupt

Source: Stellaris Peripheral Driver Library User's Guide