

CS-210 Homework 1

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1. (a) $\exists x P(x)$
 (b) $\exists x [S(x) \wedge P(x)]$
 (c) $\forall x [P(x) \rightarrow \neg R(x)]$
 (d) $\forall x [\neg P(x) \wedge R(x)]$
 (e) $\exists x [P(x) \wedge R(x)]$
 (f) $\forall x [(P(x) \wedge T(x)) \rightarrow Q(x)]$

4. (a) Statement: "If P is a square, then P is a rectangle" (True)
 (b) Converse: "If P is a rectangle, then P is a square" (False)
 (c) Contrapositive: "If P is not a rectangle, then P is not a square" (True)
 (d) Inverse : "If P is not a square, then P is not a rectangle" (False)

All statements are not true.

5.
$$\begin{array}{ll}
 [(P \vee Q) \wedge (\neg P \vee R)] \rightarrow (Q \vee R) & \text{(Implication Law)} \\
 \neg[(P \vee Q) \wedge (\neg P \vee R)] \vee (Q \vee R) & \text{(De Morgan's Law)} \\
 (\neg(P \vee Q) \vee \neg(\neg P \vee R)) \vee (Q \vee R) & \text{(Associate Law)} \\
 (\neg(P \vee \neg P) \vee \neg(Q \vee R)) \vee (Q \vee R) & \text{(Negation Law)} \\
 (\neg(T) \vee \neg(Q \vee R)) \vee (Q \vee R) & \text{(Negation Law)} \\
 (F \vee \neg(Q \vee R)) \vee (Q \vee R) & \text{(Identity Law)} \\
 \neg(Q \vee R) \vee (Q \vee R) & \text{(Negation Law)} \\
 T & \text{(Hence, a tautology)}
 \end{array}$$

7.

P	Q	R	$\neg R$	$(Q \vee R)$	$[P \rightarrow (Q \vee R)]$	$(P \rightarrow Q)$	$[\neg R \rightarrow (P \rightarrow Q)]$
T	T	T	F	T	T	T	T
T	T	F	T	T	T	T	T
T	F	T	F	T	T	F	T
T	F	F	T	F	F	F	F
F	T	T	F	T	T	T	T
F	T	F	T	T	T	T	T
F	F	T	F	T	T	T	T
F	F	F	T	F	T	T	T

The truth values for $[P \rightarrow (Q \vee R)]$ and $[\neg R \rightarrow (P \rightarrow Q)]$ are same , hence the statement is true.

8. (a) $(P \rightarrow R) \vee (Q \rightarrow R) \equiv [(P \wedge Q) \rightarrow R]$

Starting with L.H.S

$$\begin{aligned}
 & (P \rightarrow R) \vee (Q \rightarrow R) && \text{(Implication Law)} \\
 & (\neg P \vee R) \vee (\neg Q \vee R) && \text{(Associative Law)} \\
 & (\neg P \vee \neg Q) \vee (R \vee R) && \text{(Idempotent Law)} \\
 & (\neg P \vee \neg Q) \vee (R) && \text{(De Morgan's Law)} \\
 & \neg(P \wedge Q) \vee (R) && \text{(Implication Law)} \\
 & (P \wedge Q) \rightarrow R && \text{(L.H.S = R.H.S hence, equivalent)}
 \end{aligned}$$

(b) $(P \wedge (Q \vee R)) \equiv [(P \wedge Q) \vee (P \wedge R)]$

Starting with L.H.S

$$\begin{aligned}
 & (P \wedge (Q \vee R)) && \text{(Distributive Law)} \\
 & (P \wedge Q) \vee (P \wedge R) && \text{(L.H.S = R.H.S hence, equivalent)}
 \end{aligned}$$

$$(c) \neg[\neg[(P \vee Q) \wedge R] \vee \neg Q] \equiv (Q \wedge R)$$

Starting with L.H.S

$$\begin{aligned}
 & \neg[\neg[(P \vee Q) \wedge R] \vee \neg Q] && \text{(De Morgan's Law)} \\
 & \neg(\neg[(P \vee Q) \wedge R] \wedge \neg(\neg Q)) && \text{(Double Negation Law)} \\
 & ((P \vee Q) \wedge R) \wedge Q && \text{(Commutative Law)} \\
 & (R \wedge (P \vee Q)) \wedge Q && \text{(Distributive Law)} \\
 & ((R \wedge P) \vee (R \wedge Q)) \wedge Q && \text{(Commutative Law)} \\
 & Q \wedge ((R \wedge P) \vee (R \wedge Q)) && \text{(Distributive Law)} \\
 & (Q \wedge R \wedge P) \vee (R \wedge Q \wedge Q) && \text{(Idempotent Law)} \\
 & (Q \wedge R \wedge P) \vee (R \wedge Q) && \text{(Commutative Law)} \\
 & (R \wedge Q) \vee (R \wedge Q \wedge P) && \text{(Domination Law)} \\
 & (Q \wedge R) && \text{(L.H.S = R.H.S hence, equivalent)}
 \end{aligned}$$

$$(d) (P \vee Q \vee R) \wedge (P \vee T \vee \neg Q) \wedge (P \vee \neg T \vee R) \equiv P \vee [R \wedge (T \vee \neg Q)]$$

Starting from L.H.S

$$\begin{aligned}
 & (P \vee Q \vee R) \wedge (P \vee T \vee \neg Q) \wedge (P \vee \neg T \vee R) && \text{(Distributive Law)} \\
 & P \vee [(Q \vee R) \wedge (T \vee \neg Q) \wedge (\neg T \vee R)] && \text{(Associative Law)} \\
 & P \vee [(R \vee Q) \wedge (R \vee \neg T) \wedge (T \vee \neg Q)] && \text{(Distributive Law)} \\
 & P \vee [R \vee (Q \wedge F) \wedge (T \vee \neg Q)] && \text{(Identity Law)} \\
 & P \vee [R \vee (F) \wedge (T \vee \neg Q)] && \text{(Identity Law)} \\
 & P \vee [R \wedge (T \vee \neg Q)] && \text{(L.H.S = R.H.S)}
 \end{aligned}$$

12. (a) F when $x = 1$ and $y = -1$

x is not equal to y because $-1 \neq 1$.

$(x = y)$ is (F) hence, $\neg(x = y)$ is (T) .

Similarly, x^2 is equal to y^2 because $(-1)^2 = 1^2$

$(x^2 = y^2)$ is (T) hence $\neg(x^2 = y^2)$ is (F)

Hence, the whole is expression is going to be (F)

(b) F when $x = 7$ and $y = 1.9129$

If we take $x = 7$ then for $(x = y^3)$ to be True, $y = \sqrt[3]{7}$.

However since $\sqrt[3]{7}$ is not an integer this statement is (F) .

(c) F when $x = 2$ and $y = 0.5$

If we take $x = 2$ then $(\frac{1}{2} = y)$, since y must be an integer and can't be a fraction therefore it can never have a value that can equal $\frac{1}{2}$.

(d) T

(e) F when $x = 2$ and $y = 3$

$(xy = y)$ is (F) when $x = 2$ and $y = 3$.

(f) F when $x = 0$ and $y = 0$

$(x^3 \neq y^2)$ is (F) when $x = 0$ and $y = 0$.

13. (a) $\neg \forall x \forall y P(x, y)$
 $\exists x \exists y \neg p(x, y)$

(b) $\neg \forall x \exists y (P(x, y) \vee Q(x, y))$
 $\exists x \forall y \neg (P(x, y) \vee Q(x, y))$ (De Morgan's Law)
 $\exists x \forall y (\neg P(x, y) \wedge \neg Q(x, y))$

(c) $\neg \forall x (\forall y P(x, y) \wedge \exists y Q(x, y))$
 $\exists x \neg (\forall y P(x, y) \wedge \exists y Q(x, y))$ (De Morgan's Law)
 $\exists x (\neg \forall P(x, y) \vee \neg \exists y Q(x, y))$ (De Morgan's Law)
 $\exists x (\exists y \neg P(x, y) \vee \forall y \neg Q(x, y))$

(d) $\neg (\exists x \exists y \neg P(x, y) \wedge \forall x \forall y Q(x, y))$ (De Morgan's Law)
 $\neg (\exists x \exists y \neg P(x, y)) \vee \neg (\forall x \forall y Q(x, y))$
 $\neg \exists x \exists y \neg P(x, y) \vee \neg \forall x \forall y Q(x, y)$
 $\forall x \forall y P(x, y) \vee \exists x \exists y \neg Q(x, y)$