

# CS-310, Fall 2021

## Assignment 3

Assigned: Nov. 3, Due: Monday, Nov. 15, 11:55PM

November 2, 2021

**Total Points = 195**

**Problem 1. (Points: 15)** Suppose you are choosing between the following 3 algorithms:

- Algorithm A solves the problem of size  $n$  by dividing it into 8 subproblems of size  $n/4$ , recursively solving each subproblem, and then combining the solutions in linear time.
- Algorithm B solves the problem of size  $n$  by recursively solving two subproblems of size  $n - 1$  and then combining the solutions in constant time
- Algorithm C solves the problem of size  $n$  by dividing it into nine subproblems of size  $n/3$ , recursively solving each subproblem, and then combining the solutions in  $O(n^2)$  time.

What are the running times of each algorithm and which would you choose and why?

**Problem 2. (Points: 10)** Assume that  $n = 2^k$  for some positive integer  $k$ . Using induction prove that if  $T(n)$  is given as follows, then  $T(n) = n \log n$ .

$$T(n) = \begin{cases} 2 & \text{if } n = 2 \\ 2T\left(\frac{n}{2}\right) + n & \text{if } n > 2 \end{cases}$$

**Problem 3. (Points: 20)** Assume you have an array  $A[1..n]$  of  $n$  elements. A majority element of  $A$  is any element occurring in more than  $n/2$  positions (so if  $n = 6$  or  $n = 7$ , any majority element will occur in at least 4 positions). Assume that elements cannot be ordered or sorted, but can be compared for equality. (You might think of the elements as chips, and there is a tester that can be used to determine whether or not two chips are identical.) Design an efficient divide and conquer algorithm to find a majority element in  $A$  (or determine that no majority element exists). Aim for an algorithm that does  $O(n \log n)$  equality comparisons between the elements. A more difficult  $O(n)$  algorithm is possible, but may be difficult to find.

**Problem 4. (Points: 20)** Given a binary string  $S$  of type  $\{1^m 0^n\}$ , devise an algorithm that finds the number of zeroes in  $O(\log k)$  time. Let  $m + n = k$

**Problem 5. (Points: 20) [K-way Merge]** Suppose you have  $k$  sorted arrays  $A_1, A_2, \dots, A_k$  each with  $n$  elements. You want to combine them into a single sorted array of size  $kn$ . One way to do this would be to use the merge operation we discussed in class. First merge arrays  $A_1, A_2$  then merge the result with  $A_3$  and so on

- Figure out how many steps this algorithm would take.
- Design a better algorithm for this problem.

**Problem 6. (Points: 20) [Fibonacci Numbers]**

Given below is a recursive algorithm for finding the  $n_{th}$  Fibonacci number

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**Algorithm 1 :**

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function FIB(n)
  if  $n == 0$  or  $n == 1$  then
    return 1
  else
    return FIB(n-1) + FIB(n-2)
```

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- Analyze the running time of this algorithm
- A lot of computation seems to be repeated over and over in this algorithm. For example  $fib(4)$  would compute  $fib(3)$  and  $fib(2)$  and then the called  $fib(3)$  would also compute  $fib(2)$ . Is there a way we could avoid computing the same thing over and over again while still using a recursive approach?

**Problem 7. (Points: 30)** Given the following recurrence relations, find the running time (in big-O notation) using the recursion tree method and the Master Theorem.

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$$T(n) = \begin{cases} 8T\left(\frac{n}{2}\right) + \Theta(n^2) & \text{if } n \geq 2 \\ 1 & \text{else} \end{cases}$$

•

$$T(n) = \begin{cases} 3T\left(\frac{n}{4}\right) + n \log n & \text{if } n \geq 2 \\ 1 & \text{else} \end{cases}$$

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$$T(n) = \begin{cases} 2T\left(\frac{n}{4}\right) + \sqrt{n} & \text{if } n \geq 2 \\ 1 & \text{else} \end{cases}$$

**Problem 8. (Points: 20)**

A bank confiscates  $n$  bank cards suspecting them to be involved in fraud. Each card corresponds to a unique account but an account can have many cards corresponding to it. Therefore, two bank cards are equivalent if they belong to the same account. The bank has a testing machine for finding out if two cards are equivalent.

They want to know that among the collection of  $n$  cards, are there more than  $n/2$  cards that are equivalent. The only operation that the machine can do is to select two cards and tests them for equivalence. You are required to come up with a solution to this using only  $O(n \log n)$  invocations of the testing machine.

**Problem 9. (Points: 20)** Prove the correctness of the following algorithm for incrementing natural numbers. Analyze the algorithm to find out how many times is line 3 executed in the worst case.

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**Algorithm 2 :** Increment Natural Numbers

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function INCREMENT( $x$ )                                ▷ Returns  $x + 1$ 
    if  $x \bmod 2 == 1$  then
        return  $2 \times \text{INCREMENT}(\lfloor \frac{x}{2} \rfloor)$ 
    else
        return  $x + 1$ 

```

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**Problem 10. (Points: 20)** Prove the correctness of the following recursive algorithm for the multiplication of two natural numbers is correct,  $\forall$  integer constants  $\geq 2$ . Analyze this algorithm to find out how many times line 5 executed in the worst case.

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**Algorithm 3 :** Multiply Two Numbers

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function MULTIPLY( $x, y$ )                                ▷ Returns product  $xy$ 
    if  $y == 0$  then
        return 0
    else
        return MULTIPLY( $cx, \lfloor y/c \rfloor$ ) +  $x(y \bmod c)$ 

```

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