```
CHANGE(n, V, D)
FOR v = 0 TO V
        M[0, v] \leftarrow 0
FOR i = 1 TO n
        FOR v = 0 TO V
               IF (Di > v)
                        M[i, v] \leftarrow M[i-1, v].
               ELSE
                        M[i, v] \leftarrow min \{ M[i-1, v], 1 + M[i-1, v-Di] \}.
combinations = []
start = V
i = n
WHILE (start != 0)
        IF M[i, start] == M[i-1, start]
               i = i - 1
        ELSE
               combinations ← Di
               start = start - Di
return combinations
Discussing optimality:
Assistance taken from slides
M[i, v] = optimal number of coins with denomination 1,....i subjected to limit V
Goal = M[i, V]
Case 1: item i is not selected then the best of {1,2,3..., i-1} are selected subject to limit V
```

Case 2: item i selected then,

V = V - Di

number of coins ++

then the best of {1,2,3..., i-1} are selected subject to limit V

So due to exchange argument, these substructures are optimal, hence, making the whole solution optimal.