Handson 12

Ch# 17

1.

a) Each insert costs $O(1)$, we have n inserts
Hence, for ~~each~~ all inserts we get $O(n)$.
Each complex operation would cost $O(2n)$
because it doubles in size

$$\frac{O(n) + O(2n)}{n} = \frac{O(n)}{n}$$

which is $O(1)$

b) Normal : $O(1)$
Complex : $O(i)$
Each insert costs you \$3
\$1 cost of insertion
\$2 cost of future copying
So, its always \$3 per operation
which is $O(1)$