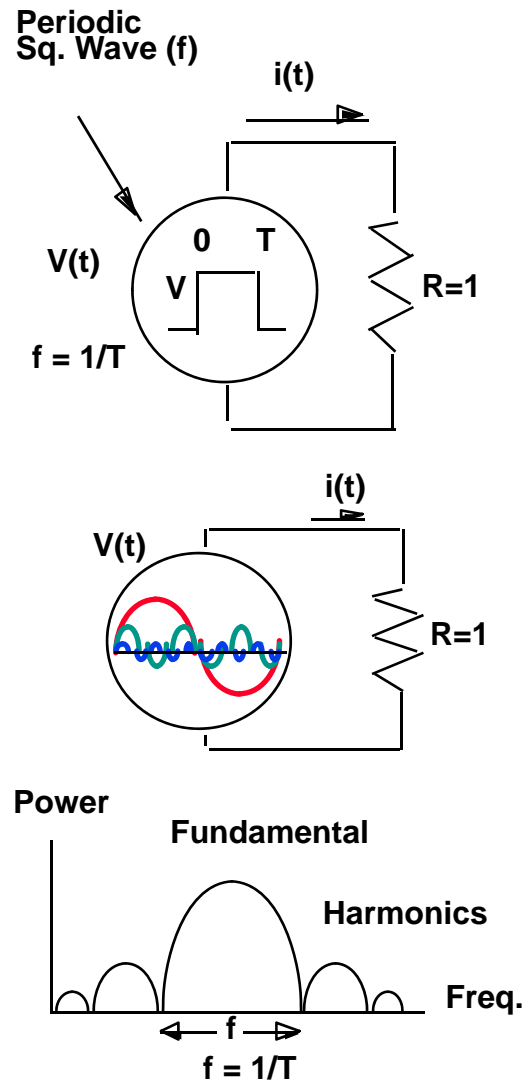


# **Lecture 23**

## **Code Division Multiple Access (CDMA)**

### **TOPICS**

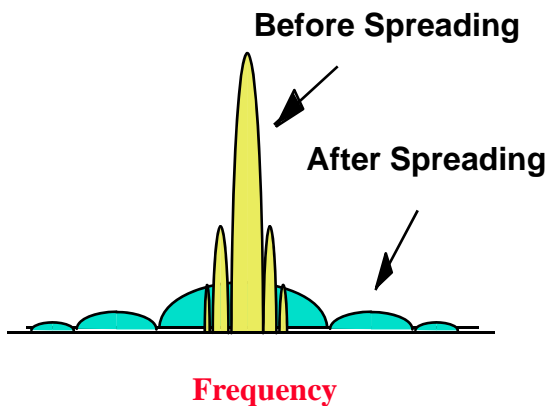
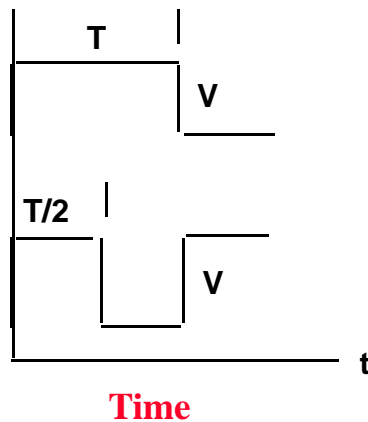
- **Describe Spread Spectrum**
- **Provide DS-CDMA concept**
- **Recognize Orthogonal codes**
- **Illustrate how it is implemented in IS-95 standard**
- **Examine forward and reverse link waveform**
- **Describe soft hand-off & power control issues**



## Spectrum

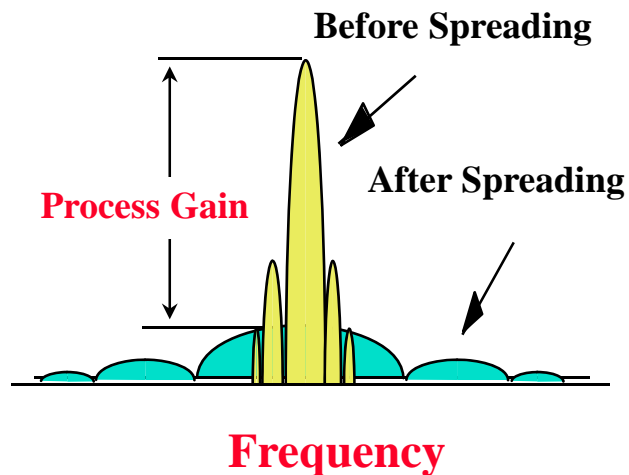
- A discrete time signal is composed of an infinite number of harmonically related sinusoidal waves  
- *Spectrum*
- The associated power is due to infinite number of sinusoidal components:
  - main lobe corresponds to the fundamental frequency
  - side lobes correspond to harmonic components
- The bandwidth of the power spectrum is proportional to the frequency

# Spectrum Spreading



- Spectrum spreading is accomplished by increasing the frequency of the discrete time signal
  - A high frequency signal has wider spectrum than a low frequency signal
  - A low frequency signal has higher power amplitude than a low high power signal
  - This is due to conservation of energy
- This is one of the key element in DS-CDMA
  - It also leads to the concept of “Process Gain”

## Process Gain



- Process gain is due to spectrum spreading

- Defined as

$$G = 10\log\left(\frac{BW}{R_b}\right)$$

**BW=Bandwidth after spreading**

**$R_b$ = Original bit rate (Before spreading)**

## CDMA Process Gain: *Illustration*

- Spreading gain gives CDMA a big advantage over traditional, non-spread-spectrum systems

- Example: CDMA vs TDMA Cellular

- TDMA Cellular

- $BW = 30 \text{ KHz}$     $R_b = 10 \text{ KHz}$

- $\text{Process Gain} = 10 \log (30/10) = 4.7 \text{ db}$

- CDMA

- $BW = 1.23 \text{ MHz}$     $R_b = 10 \text{ KHz}$

- $\text{Process Gain} = 10 \log (1230/10) = 20.9 \text{ db}$

- CDMA advantage:  $20.9 - 4.7 = 16.2 \text{ dB}$

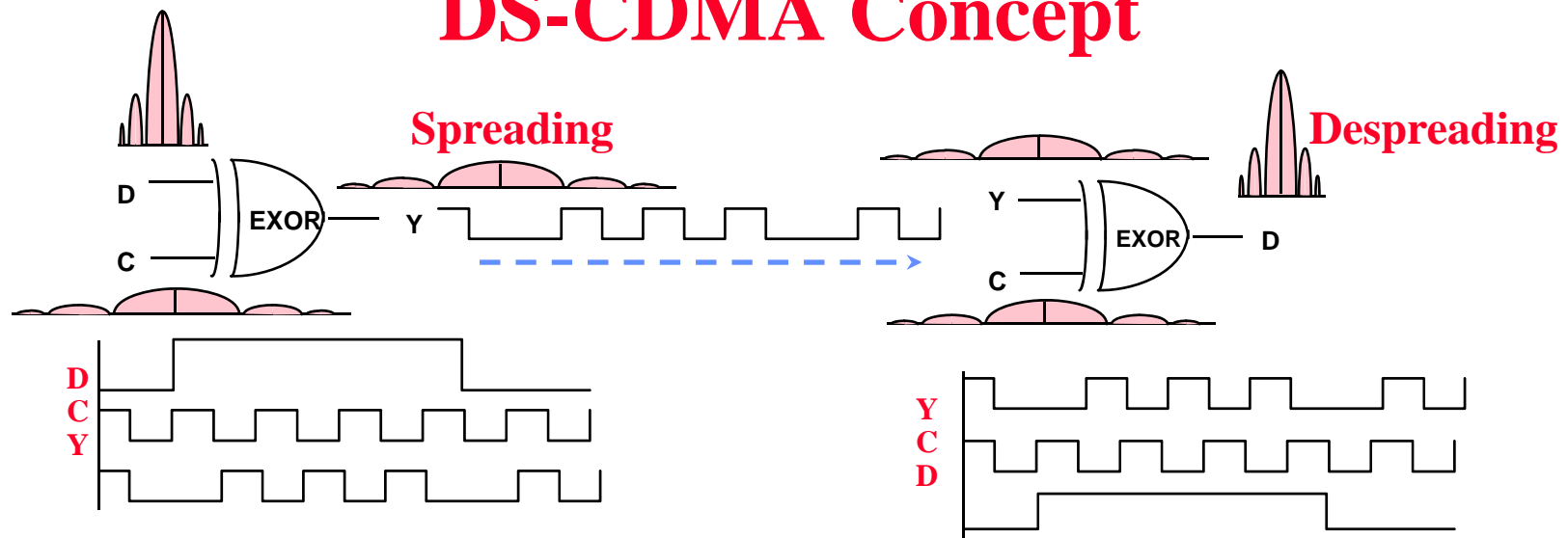
- this advantage assumes only one user on CDMA system

- doubling number of users reduces advantage 3 db

- with 16 users, CDMA still has 4 db advantage over TDMA

- How can we spread the bandwidth of information ?

## DS-CDMA Concept



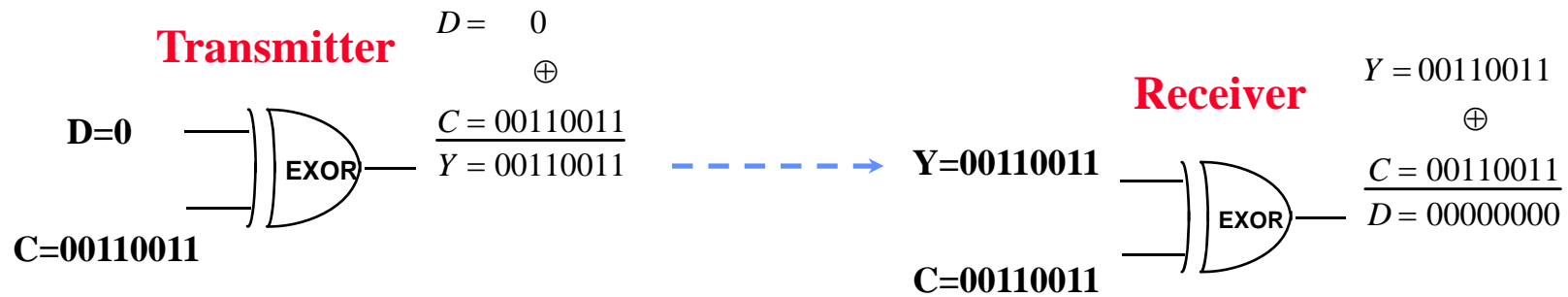
### Spreading:

- $D$  = low speed data
- $C$  = high speed PN Code
- $Y = D \text{ EXOR } C$

### Despreading:

- $Y$  = high speed received data
- $C$  = same PN Code
- $D = Y \text{ EXOR } C$   
 $= D \text{ EXOR } C \text{ EXOR } C$

## DS-CDMA: Example-1



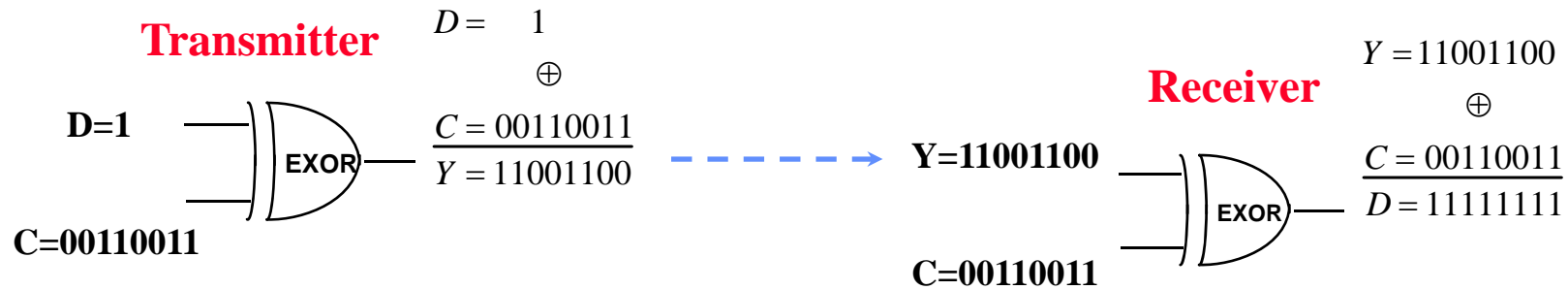
### Spreading:

- $D = 0$
- $C = 00110011$  (8-bit code)
- $Y = D \text{ EXOR } C$   
 $= 0 \text{ EXOR } 00110011$   
 $= 00110011$

### Despreading:

- $Y = 00110011$
- $C = 00110011$  (same code)
- $D = Y \text{ EXOR } C$   
 $= 00110011 \text{ EXOR } 00110011$   
 $= 00000000$

## DS-CDMA: Example-2



### Spreading:

- $D = 1$
- $C = 00110011$  (8-bit code)
- $Y = D \text{ EXOR } C$   
 $= 1 \text{ EXOR } 00110011$   
 $= 11001100$

### Despreading:

- $Y = 11001100$
- $C = 00110011$  (same code)
- $D = Y \text{ EXOR } C$   
 $= 11001100 \text{ EXOR } 00110011$   
 $= 11111111$



## Our Observations

### Spreading Process:

- A PN code multiplied by BIN-0 (binary zero) produces the same PN code:

Bin-0 EXOR PN code = PN code

0 EXOR 00110011 = 00110011

- A PN code, multiplied by BIN-1 produces the inverse code, also known as an “Antipodal code”

BIN-1 EXOR PN code = PN code

1 EXOR 00110011 = 11001100

### Despreading Process:

- A PN code multiplied by the same PN code, produces “Zeros”

PN code EXOR PN code = 0

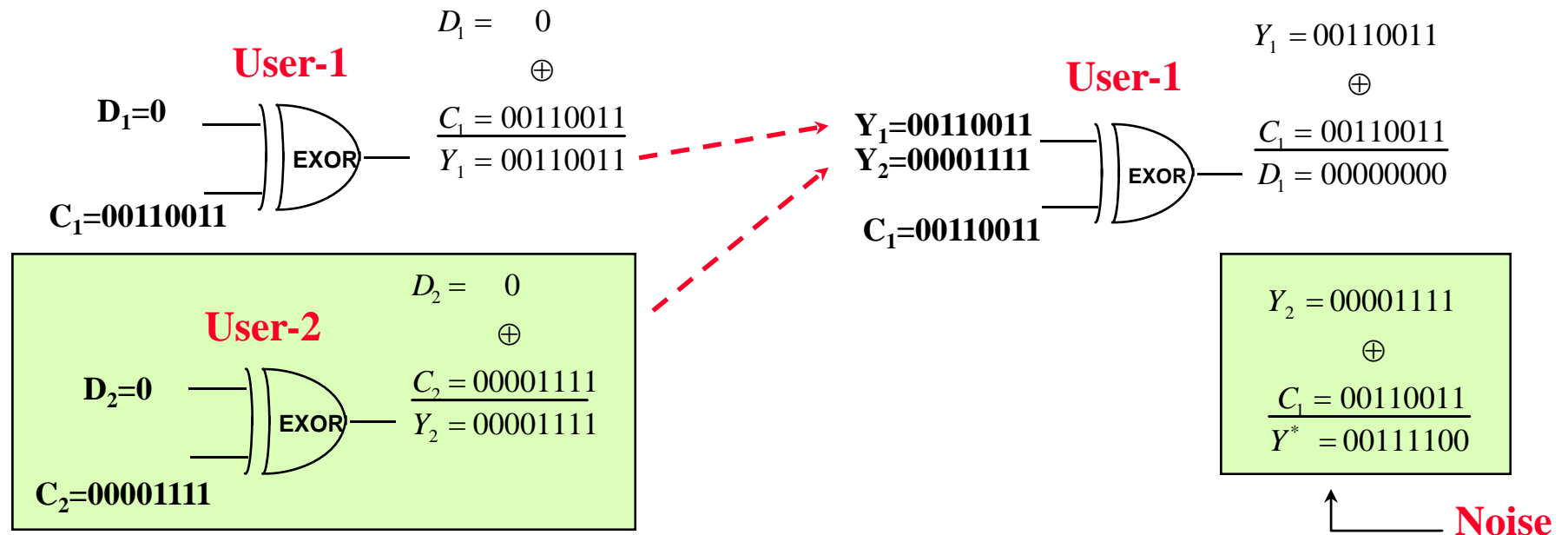
00110011 EXOR 00110011 = 0

- A PN code multiplied by the antipodal code, produces “Ones”

PN code EXOR Antipodal code = 1

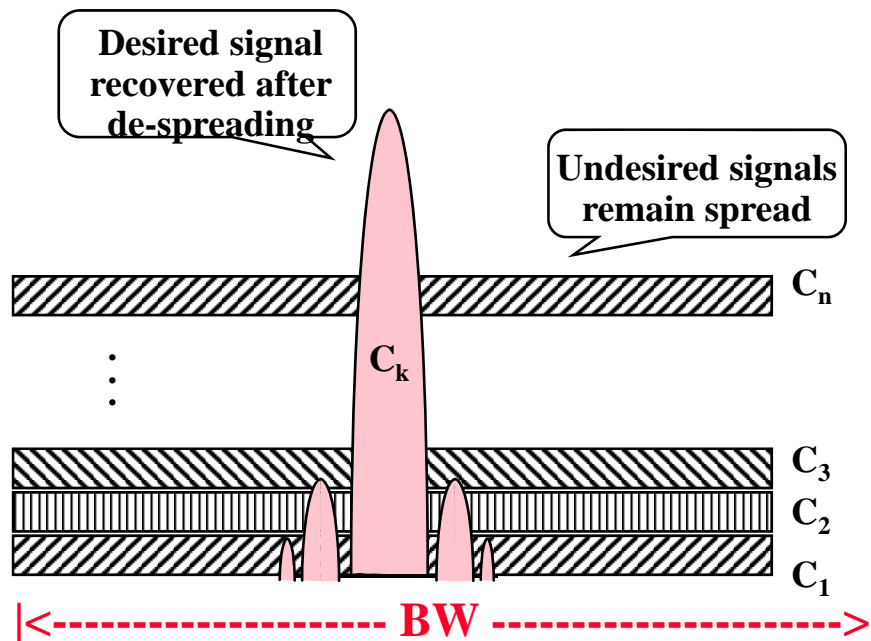
00110011 EXOR 11001100 = 1

## Multi User DS-CDMA



- **Data from User-1:**
  - Tx Data  $D_1=0$ , Code  $C_1=00110011$ , Output  $Y_1=D_1 \text{ EXOR } C_1= 00110011$
  - Rx Data  $D_1=Y_1 \text{ EXOR } C_1= 00110011 \text{ EXOR } 00110011 = 0$  (O.K)
- **Data from User-2:**
  - Tx Data  $D_2=0$ , Code  $C_2=00001111$ , Output  $Y_2=D_2 \text{ EXOR } C_2 = 00001111$
  - Rx Data  $D_2 = Y_2 \text{ EXOR } C_1= 00001111 \text{ EXOR } 00110011 = 00111100$  (spread, noise)
- **All users, other than the desired signal, appears as noise**

## How Many users can Share the Same Spectrum ?



- Each user contributes noise to the system
- Therefore the number of users that can share the same bandwidth would be given by the Process gain minus Total noise contributions
- Also, an additional

- This is CDMA capacity issue
- We shall explore this later when it is appropriate

performance

## PN Codes and Orthogonal Codes

**Example:**

**$n = 64$  (IS-95 CDMA)**

**Number of PN Sequences**

$$= 2^{64} - 1$$

$$= 1.8447 \times 10^{19}$$

**Out of  $1.8447 \times 10^{19}$  PN sequences,  
only 64 of them are orthogonal  
codes**

- **An  $n$ -bit PN code has  $2^n - 1$  PN sequences**
- **Out of these PN sequences, only  $n$ -codes are orthogonal**
- **A 64-bit PN code has  $1.8447 \times 10^{19}$  PN sequences**
- **Only 64 of them are Orthogonal codes**
- **Orthogonal codes have zero cross-correlation properties**

## Properties of Orthogonal Codes

**Example-1:**

**x = 0 0 1 1**

**y = 0 1 1 0**

$$R_{xy} = \begin{array}{cccc} -1 & -1 & 1 & 1 \\ -1 & 1 & 1 & -1 \\ \hline 1 & -1 & 1 & -1 \end{array} = 0$$

**Example-2:**

**x = 0 0 1 1**

**y = 1 1 0 0**

$$R_{xy} = \begin{array}{cccc} -1 & -1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ \hline -1 & -1 & 1 & 1 \end{array} = -4$$

- A pair of code:

$x_1, x_2, \dots, x_n$

$y_1, y_2, \dots, y_n$

is Orthogonal if the cross-correlation is zero.

- This is given by:  $R_{xy}(0) = \sum_{i=1}^n x_i y_i = 0$

- An Orthogonal code has equal number of 1's and 0's. But not all codes having equal number of 1's and 0's are orthogonal

## Properties of Orthogonal Codes: *Cont.*

**Example-1:**

$x = 0\ 0\ 1\ 1$   
 $y = 0\ 1\ 1\ 0$

$$R_{xy} = \frac{-1\ -1\ 1\ 1}{1\ -1\ 1\ -1} = 0$$

**Example-2:**

$x = 0\ 0\ 1\ 1$   
 $y = 1\ 1\ 0\ 0$

$$R_{xy} = \frac{-1\ -1\ 1\ 1}{1\ 1\ -1\ -1} = -4$$

**Antipodal**

- Orthogonal codes are binary numbers ( $2^n$ )
- An Orthogonal code has equal number of 1's and 0's
- A pair of code:  
 $x_1, x_2, \dots, x_n$   
 $y_1, y_2, \dots, y_n$   
 is Orthogonal if the cross-correlation is zero.  
 $R_{xy}(0) = \sum_{i=1}^n x_i y_i = 0$
- This is given by:

## Properties of Orthogonal Codes: *Cont.*

8-bit Orthogonal code has  
16 Bi-Orthogonal codes

	1	2	3	4	5	6	7	8
0	0	0	0	0	0	0	0	0
1	0	1	0	1	0	1	0	1
2	0	0	1	1	0	0	1	1
3	0	1	1	0	0	1	1	0
4	0	0	0	0	1	1	1	1
5	0	1	0	1	1	0	1	0
6	0	0	1	1	1	1	0	0
7	0	1	1	0	1	0	0	1
8	1	1	1	1	1	1	1	1
9	1	0	1	0	1	0	1	0
10	1	1	0	0	1	1	0	0
11	1	0	0	1	1	0	0	1
12	1	1	1	1	0	0	0	0
13	1	0	1	0	0	1	0	1
14	1	1	0	0	0	0	1	1
15	1	0	0	1	0	1	1	0

- An n-bit Orthogonal code has:
  - n Orthogonal codes and
  - n antipodal codes
- Total number of codes =  $2n$
- The  $2n$  code-set is known as Bi-Orthogonal code-set
  - An 8-bit Orthogonal code has 16 Bi-Orthogonal codes
  - A 64-bit Orthogonal code has 128 Bi-Orthogonal code-set

## Generation of Orthogonal Codes

C1	1st Quadrant	2nd Quadrant		b	b
	3rd Quadrant	4th Quadrant		b	$\overline{b}$
C2					

- Step-1: Divide N x N matrix as FOUR Quadrants
- Step-2: Make 1st, 2nd and 3rd quadrant identical & Invert the 4th

➤ b is the bit value 0 or 1

- Orthogonal Code: C1= 0 0

C2= 0 1

- Antipodal Code: C1= 1 1

C2= 1 0

Repeat the process for longer codes



## Generation of Long Orthogonal Codes

		4 x 4			
2 x 2					
C1	0 0	C1	0 0	0 0	
C2	0 1	C2	0 1	0 1	
C3	0 0	C3	0 0	1 1	
C4	0 1	C4	0 1	1 0	
		8 x 8			
C1	0 0	0 0	0 0	0 0	
C2	0 1	0 1	0 1	0 1	
C3	0 0	1 1	0 0	1 1	
C4	0 1	1 0	0 1	1 0	
C5					
C6	0 0	0 0	1 1	1 1	
C7	0 1	0 1	1 0	1 0	
C8	0 0	1 1	1 1	0 0	
	0 1	1 0	1 0	0 1	

- Begin with a 2 x 2 matrix as before
- Generate a 4 x 4 by repeating 2 x 2
- Generate a 8 x 8 by repeating 4 x 4

so on . . .

This process was developed by J. L Walsh in 1923  
- known as Walsh Code

codes

## IS-95A Codes & Their Usage

### ■ Walsh Code:

- 64 bits long, 64 Orthogonal codes, Code rate = 1.2288 Mb/s
- Unique identifier for mobiles
- Walsh Code Assignment:
  - » W0 = Walsh code 0, used in Pilot Channel
  - » W1-W7: Used in paging channels (7-paging channels)
  - » W32: Used in Sync. channel
  - » W8 to W31 and W33 to W63: Used in Traffic channels

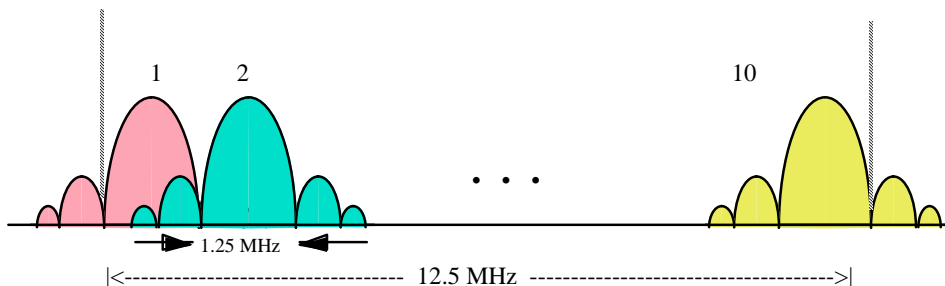
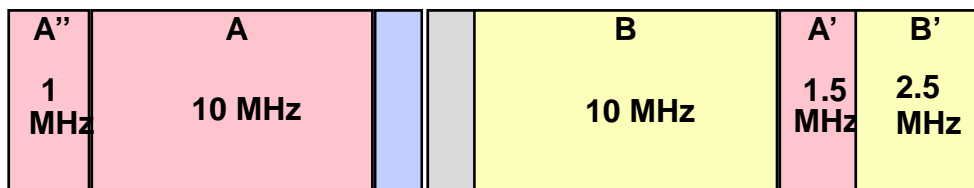
### ■ Long PN Code:

- 42-bits long,  $2^{42}-1=4.398 \times 10^{12}$  PN codes, Code rate=1.2288Mb/s
- Used for reverse channel spreading and data scrambling

### ■ Short PN-Code:

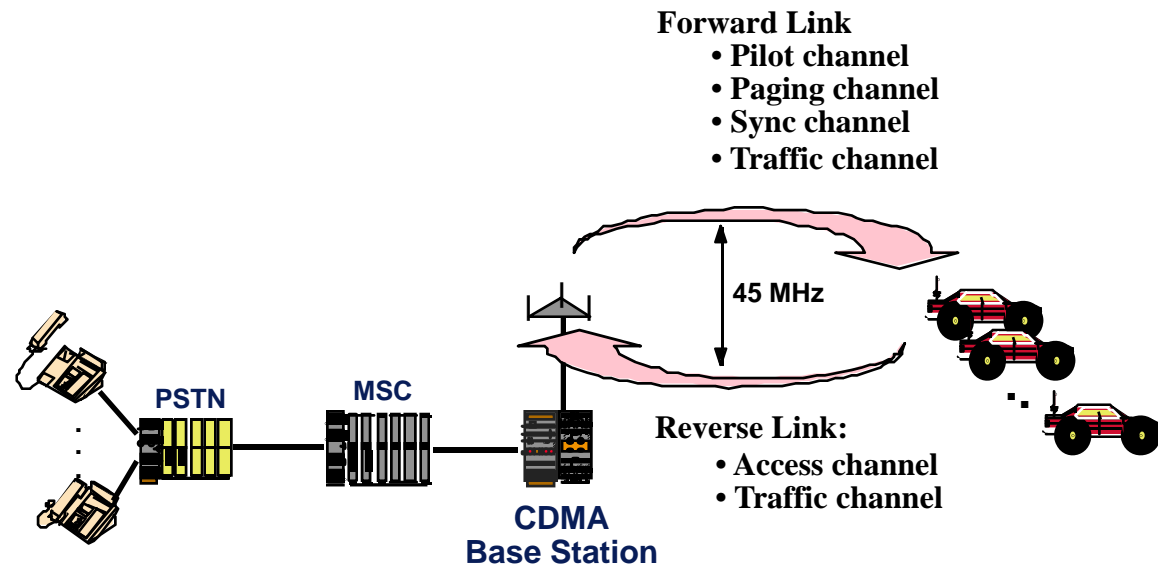
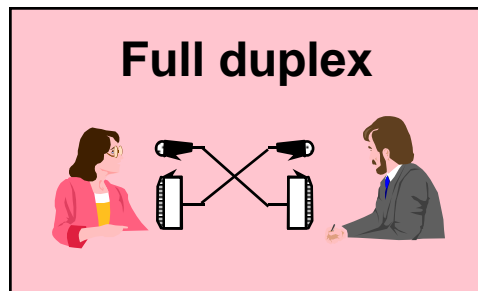
- 15-bits long,  $2^{15}-1=32,767$  PN codes, Code rate=1.2288 Mb/s
- Used for cell identification in reused cells (Similar to SAT in

## IS-95A DS-CDMA *Spectrum*



- The existing 12.5 MHz cellular spectrum ( A or B) is divided into TEN CDMA bands
- Each band is 1.25 MHz
- Each 1.25 MHz band supports 64 Walsh codes

## CDMA Architecture



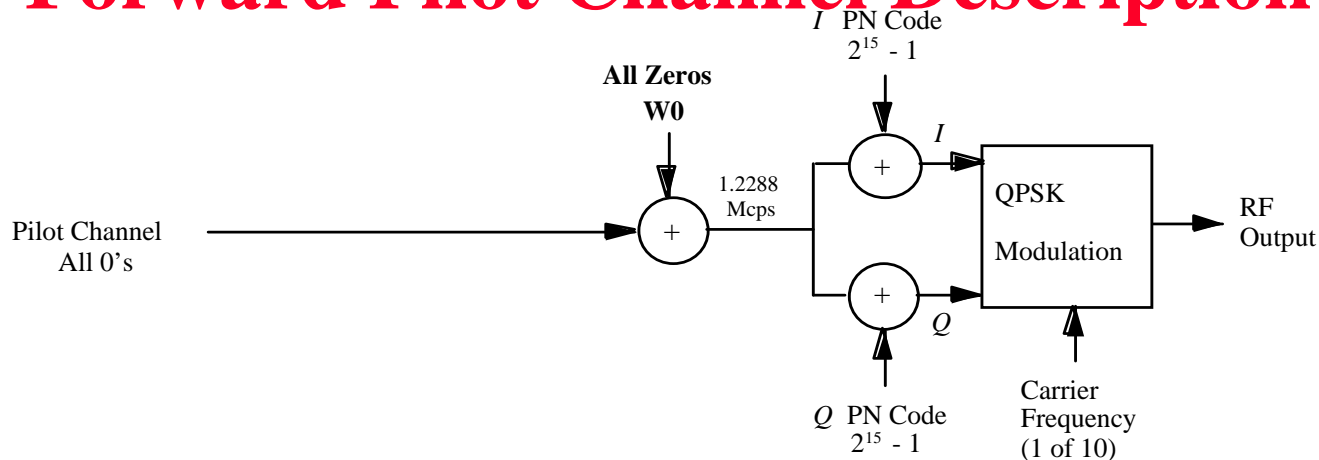
### ■ Forward link -- *Base to mobile*

- Pilot channel
- Paging channel
- Sync. channel
- Traffic channel

### ■ Reverse link -- *Mobile to base*

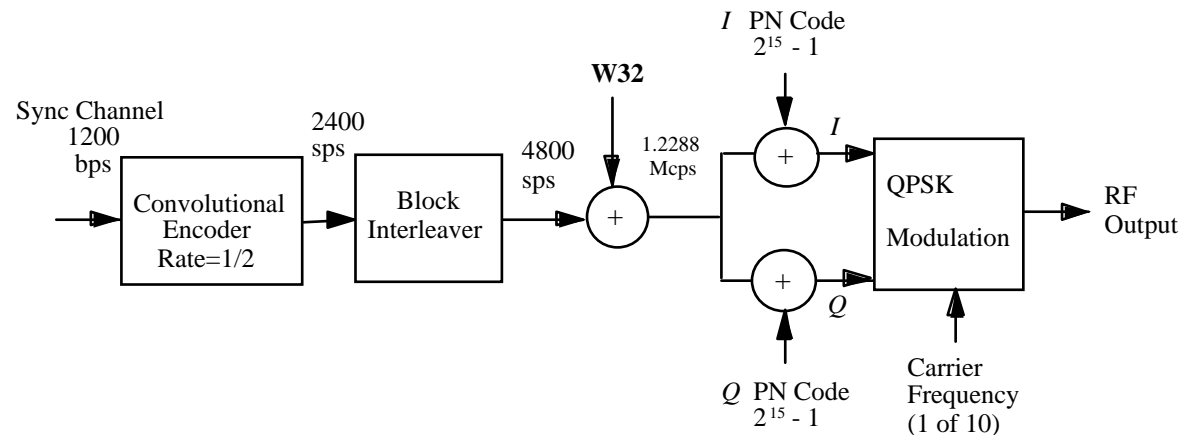
Reverse access channel  
Reverse traffic channel

## Forward Pilot Channel Description



- All zero, uncoded spread spectrum channel
- Spreading Code: W0, 64-zeros
- Continuously transmitted from the base station
- Each base station uses one of  $2^{15}-1$  PN sequences to identify a forward pilot channel (there are 512 PN offsets)
- Provides phase & Timing reference
- Provides signal strength to the mobile

## Forward Sync. Channel Description

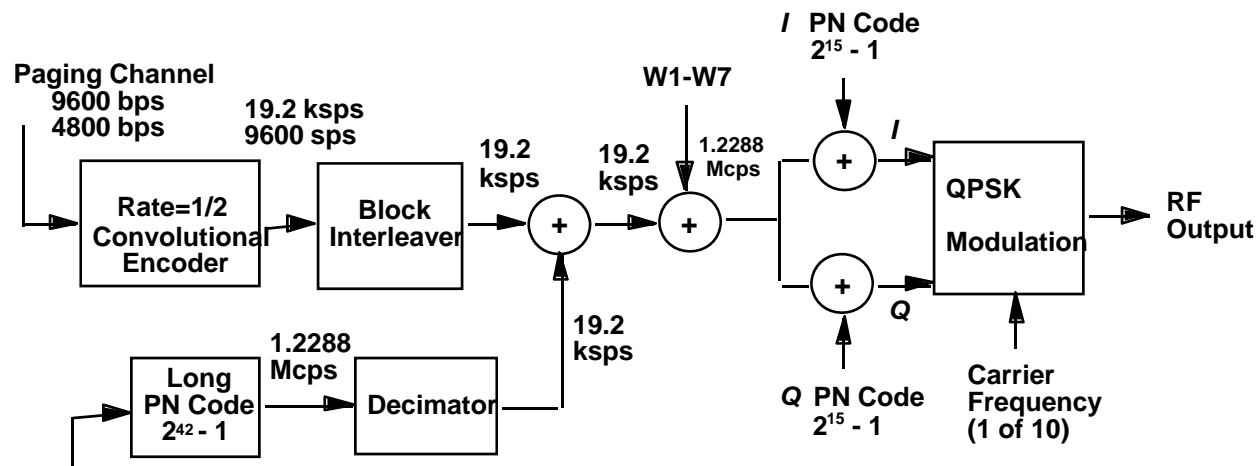


- **Rate 1/2 convolutionally encoded spread spectrum channel**
- **Spread by means of W32 and then MOD2 added (EXOR) with a unique 15-bit PN code (same PN code as the pilot)**
- **Modulated and transmitted**
- **Provides:**

➤ **Unique timing for synchronization**

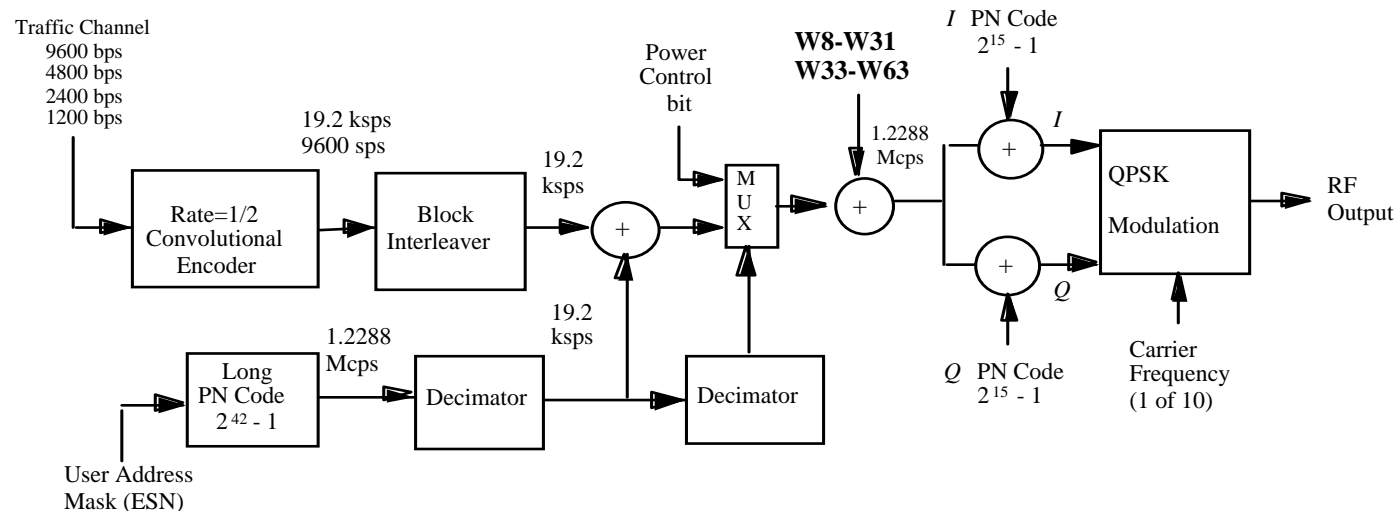
22

## Forward Paging Channel Description



- Spread by means of W1-W7
- Uses the same PN code (offset )
- Pages a mobile
- Provides system parameter MSG
- Assigns traffic channel
- Overhead information
- Neighbor list

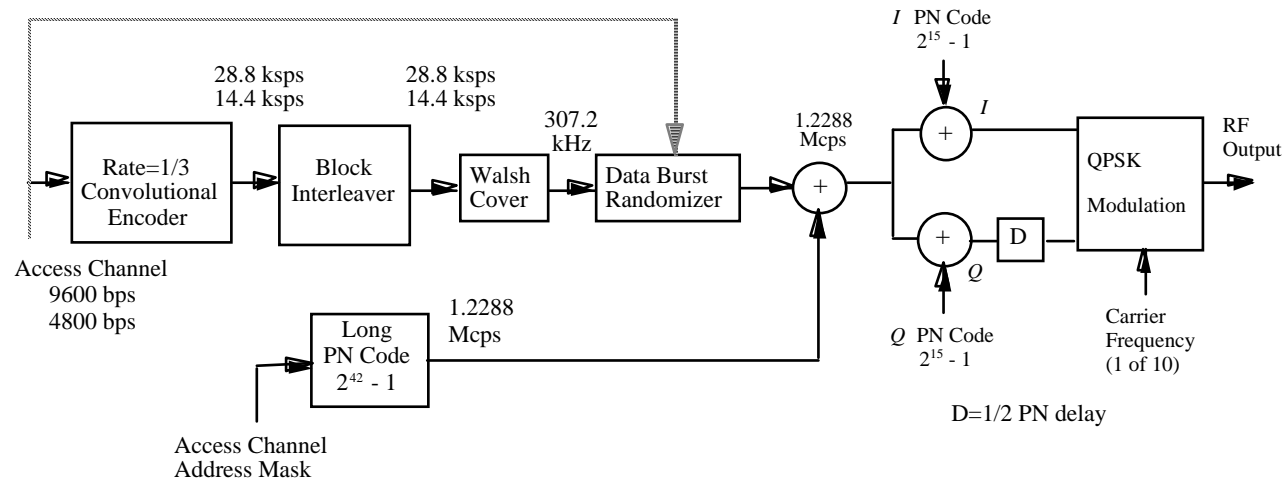
# Forward Traffic Channel Description



- Spread by means of W8-W31 and W33-W63
- Uses the same PN code (offset )
- Used for voice communications, Signaling
- Provides power control
- Assigns traffic channel
- Overhead information
- Neighbor list

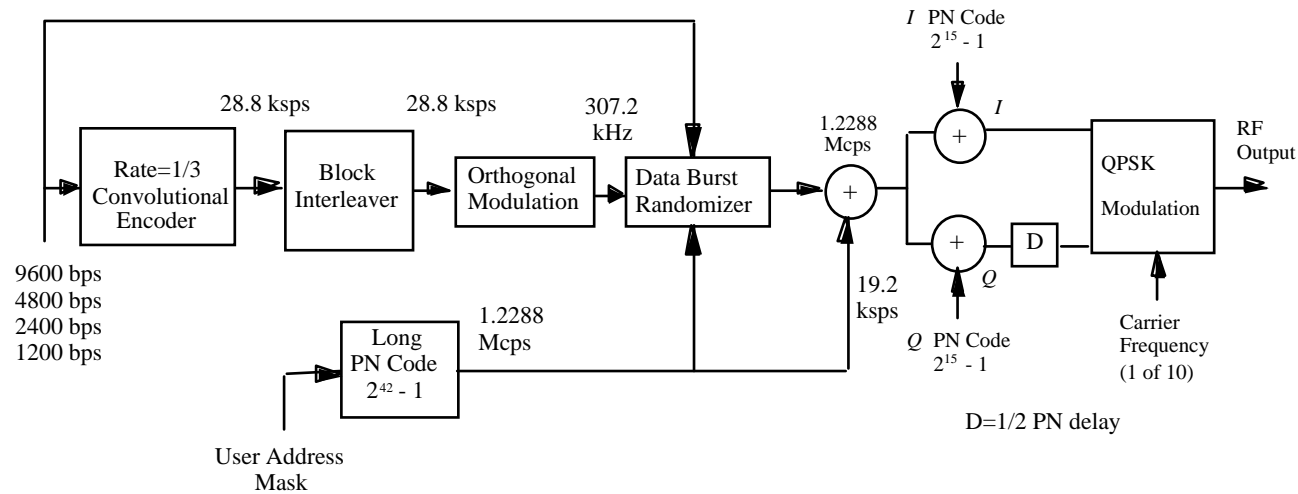


## Reverse Access Channel Description



- Works in association with Forward Paging Channel
- Spread by means of Long PN code (42-bits, 1.2288Mb/s)
- Used for Call origination, system access
- Channel request
- Page response
- Registration

## Reverse Traffic Channel Description



- **Used for voice communications**
- **Spread by means of Long PN code (42-bits, 1.2288Mb/s)**
- **Same short PN offset**
- **Response to command**
- **Seeks information from the base**

## CDMA Capacity *WITHOUT* Cellular Features

- Carrier to interference ratio is given by

$$\frac{C}{I} = \frac{R_b \times E_b}{N_o \times W} \quad (\text{A})$$

where

$E_b$  = energy per bit

$R_b$  = bit rate

$N_o$  = thermal noise

$W$  = transmission  
bandwidth

- In CDMA, the interference is due to all users except one

$$I = C(N - 1)$$

$$\frac{C}{I} = \frac{1}{N - 1} \quad (\text{B})$$

- Equating (A) and (B) we get

$$\frac{1}{N - 1} = \frac{R_b \times E_b}{N_o \times W}$$

$$N = 1 + \frac{W / R_b}{E_b / N_o}$$

$W/R_b = \text{Process gain}$

## CDMA Capacity *WITH* Cellular Features

- In the cellular environment, the maximum achievable capacity is given by

$$N = 1 + \frac{W / R_b}{E_b / N_o} \bullet \frac{F}{D} \bullet S \bullet H$$

F=frequency reuse factor (1)

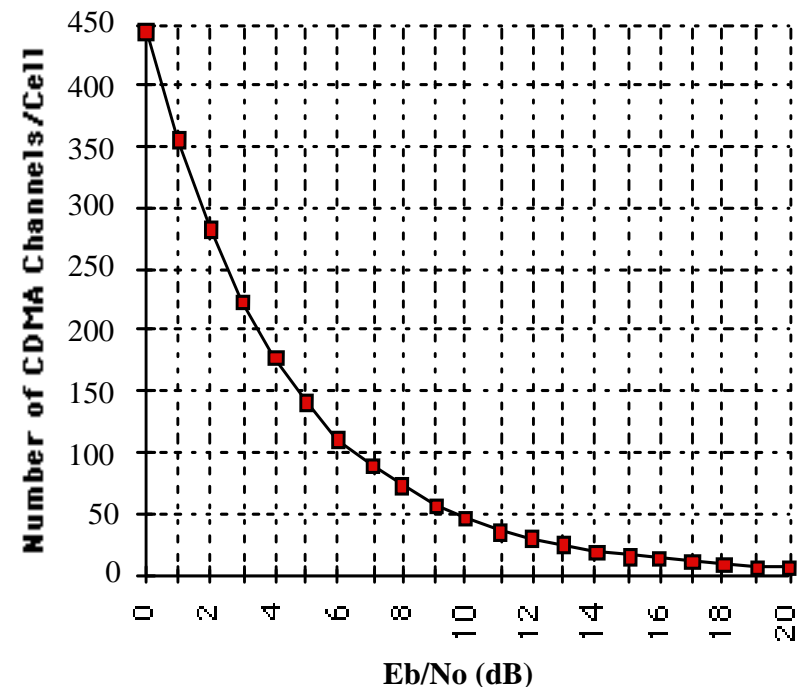
d=voice duty cycle (0.45)

S=sectorization factor (3)

H=soft hand-off factor

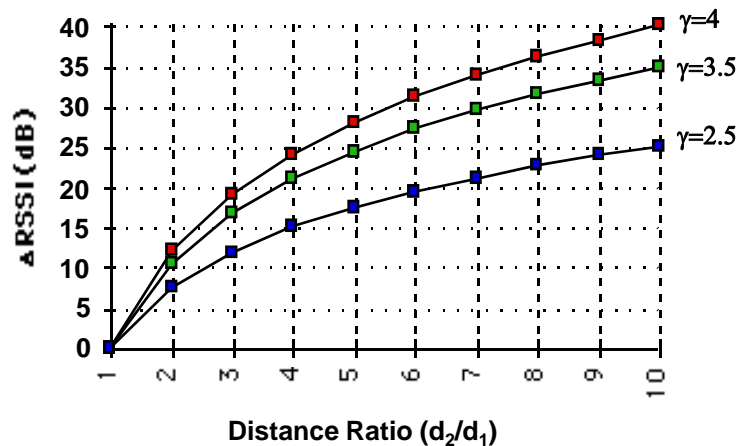
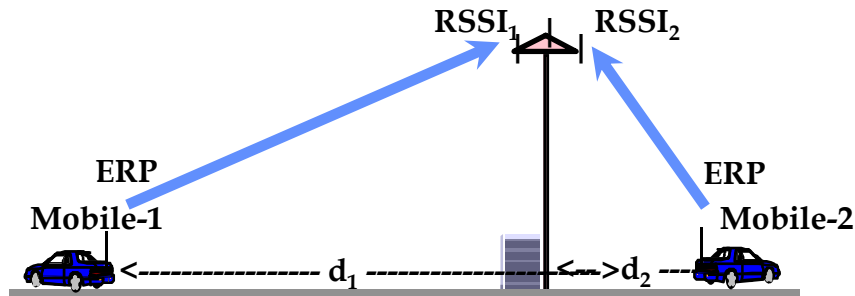
W=1.25MHz

R<sub>b</sub>=9600bps



Low Eb/No is the key to enhance CDMA capacity

## Why CDMA Power Control



$\gamma$  = Propagation Constant

### Near-Far Problem:

- If the mobiles are permitted to transmit the same power from different distances, the ratio of received signal level would be:

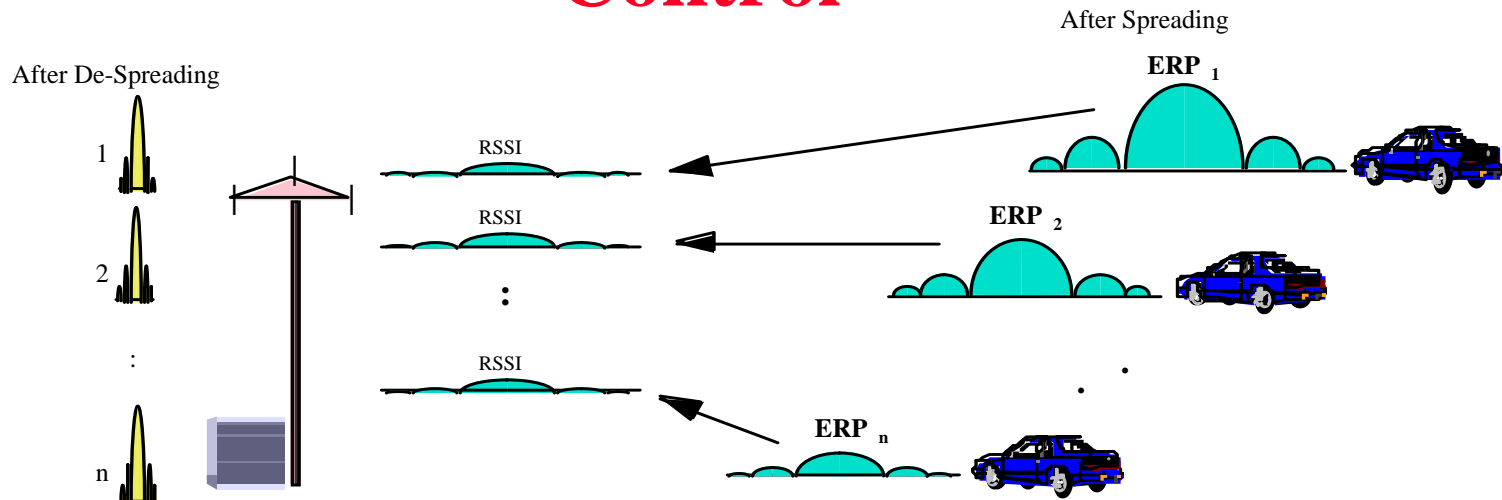
$$\frac{RSSI_1}{RSSI_2} = \left( \frac{d_2}{d_1} \right)^\gamma$$

- This implies that if  $d_1=4d_2$ , and prop. constant=4 (Urban), RSSI from  $d_2$  will be 256 times stronger and the base station will be unable to recover the other mobile
- Therefore power control is necessary so that RSSI's are same

## Power Control Process in CDMA

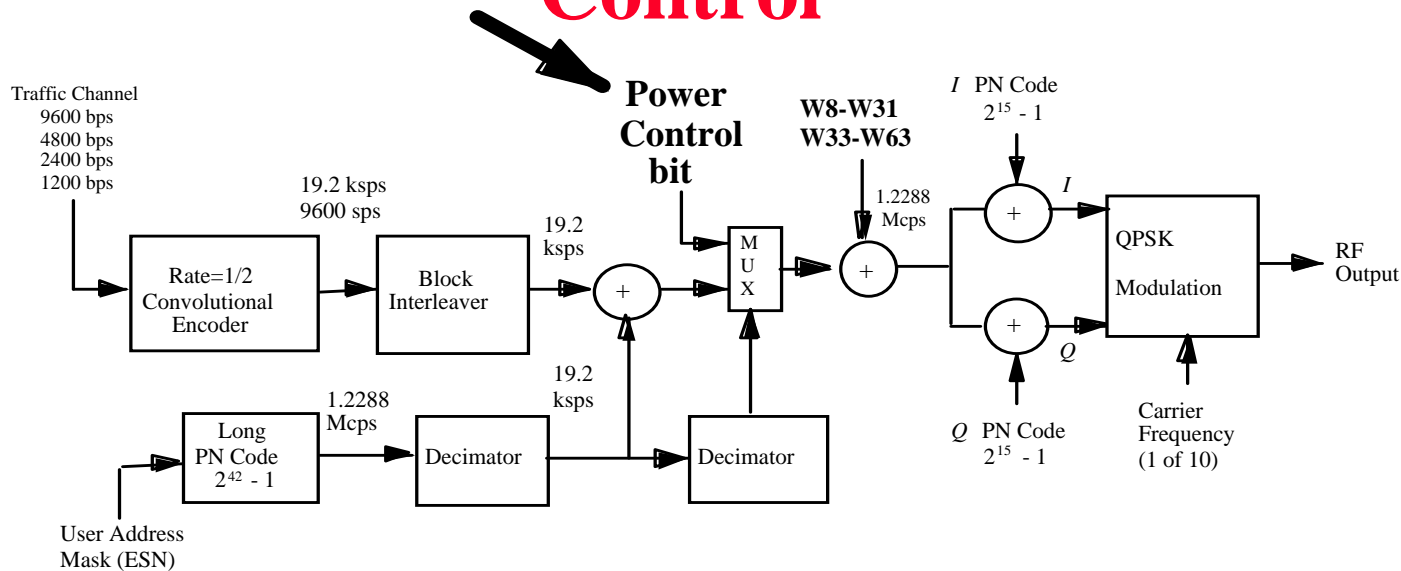
- **Three step process:**
  - **Reverse link open loop power control**
  - **Reverse link closed loop power control**
  - **Forward link power control**

## Reverse Link Open Loop Power Control



- Each mobile computes differential path loss and adjusts its transmit power so that the received signal level at the base is constant
- The total received power from all mobiles determines the system capacity

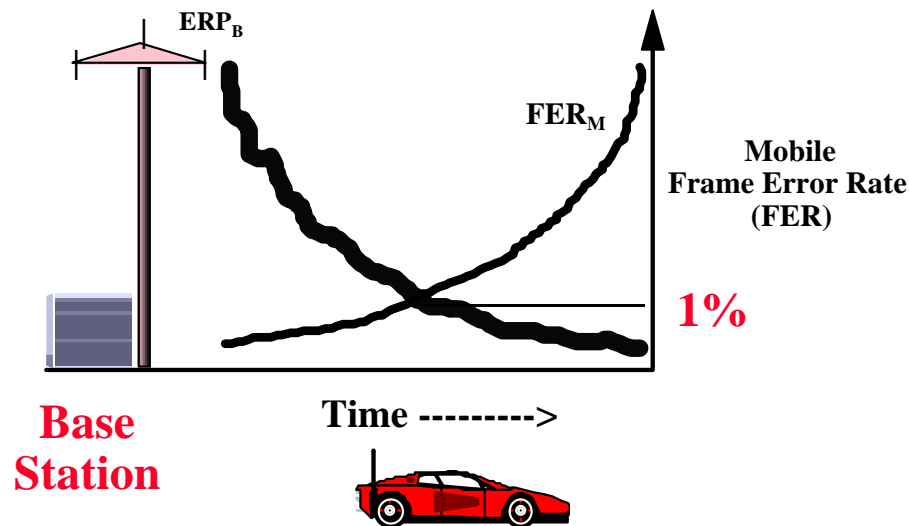
## Reverse Link Closed Loop Power Control



- Reverse link closed loop power control is accomplished by a power up/down command from the base station
- A single power control bit is inserted in the data stream, rate=1.25ms:
  - 1 for 0.5dB up, and 0 for 0.5dB down



## Forward Link Power Control



- Base station reduces ERP
- Mobile computes FERRO
- if FER=1%, stop reducing ERP

- The base station reduces its power while the mobile computes Frame Error Rate (FER)
- When the mobile detects 1% FER, it sends a request to stop the power reduction
- The process occurs every 15 to 20 ms
- This adjustment process is limited to only 6dB in 0.5dB steps because all

## CDMA Hand-Off

### ■ CDMA to CDMA Soft Hand-Off

- Mobile is directed to adjacent cell or sector, same frequency, different code, without dropping the serving code
- Mobile keeps both channels during the process
- Once the new link is well established, the original link is dropped
- This process is known as “Soft-Hand Off” or “Make before Break”

### ■ CDMA to CDMA Hard Hand-Off

- Mobile is directed to a different CDMA frequency
- Voice is muted momentarily during this process

### ■ CDMA to AMPS Hard Hand-Off

- A dual mode mobile is directed to an AMPS channel
- Voice is muted momentarily during this process

**End of Lecture 23**