

RSA

Don't forget Proj topics due WED 28th
(sorry for typo)

TODAY

- ▶ multiplicative inverses, Bezout
- ▶ totient, exponentiation
- ▶ RSA

} pari examples

\mathbb{Z}_n vs \mathbb{Z}_n^*

$$\mathbb{Z}_{15} = \{0, 1, 2, \dots, 14\}$$

↓ throw away
mults of 3 & 5

$$\mathbb{Z}_{15}^* = \{1, 2, 4, 7, 8, 11, 13, 14\}$$

$$\mathbb{Z}_{11} = \{0, \dots, 10\}$$

$$\mathbb{Z}_{11}^* = \{1, \dots, 10\}$$

Claim: 7 has mult
inverse mod 15

Bezout's theorem gives

(integers)

$$(-2) \cdot 7 + 1 \cdot 15 = 1$$

↓ mod 15

$$\underbrace{-2 \cdot 7}_{\text{mult inverse of 7}} \equiv_{15} 1$$

$$(-2 \equiv_{15} 13)$$

Does 9 have
inverse mod 15?

(shouldn't since $9 \notin \mathbb{Z}_{15}^*$)

$$2 \cdot 9 + (-1) \cdot 15 = 3 = \gcd(9, 15)$$

↓ mod 15

$$2 \cdot 9 \equiv_{15} 3$$

all true, but doesn't give you inverses

Totient

$$\phi(n) = |\mathbb{Z}_n^*| = \# \text{ of elements (in } \mathbb{Z}_n) \text{ relatively prime to } n$$

$$\phi(15) = \# \text{ elements in } \{1, 2, 4, 7, 8, 11, 13, 14\} \\ = 8$$

$$\phi(p) = p-1 \quad \text{when } p \text{ is prime}$$

Euler's theorem:

$$x^{\phi(n)} \equiv_n 1 \quad \text{for all } x \in \mathbb{Z}_n^*$$

$$\text{try: } n=15, \quad \phi(n)=8$$

$$1^8 \longrightarrow 1$$

$$2^8 \longrightarrow 256 \bmod 15 = 1$$

if $x^t \equiv_n 1$
then x^{t-1} is
an inverse of x

Note: if operations are mod n , then
all exponents can be reduced mod $\phi(n)$

$$3^{141592} \bmod 11 ?$$

$$3^{141592} \equiv 3^{141592 \bmod 10} \equiv 3^2 \equiv 9$$

Exponentiation with big numbers:

$$x^{p-1} \bmod p ;$$

bad way: first compute x^{p-1} too big !!! over \mathbb{Z}
then reduce mod p

Ex:

Computing 13^{1024}

Bad way: for $i = 1$ to 1024

$$x = 13 \cdot x$$

good way: for $i = 1$ to 10

$$x = x^2$$

computing x^{37}

$$\left(\overbrace{100101}^9 \right)_{18}$$

$$x^{37} = x \cdot x^{36}$$

$$= x \cdot (x^{18})^2$$

$$= x \cdot ((x^9)^2)^2$$

$$= x \cdot ((x \cdot x^8)^2)^2$$

$$= x \cdot ((x((x^2)^2)^2)^2)^2$$

► good way:

$$X^{37} \bmod n$$

$$X^{37} = X \cdot ((X((X^2)^2)^2)^2) \bmod n$$

$$\dots ((X \cdot ((X^2 \bmod n)^2 \bmod n)^2 \bmod n) \bmod n) \dots$$

Idea: If final result is needed mod n
then can reduce intermediate values mod n

RSA:

$$N = pq \quad \Rightarrow \quad \varphi(N) = (p-1)(q-1)$$

$$ed \equiv_{\varphi(N)} 1$$

$$m \longrightarrow m^e \longrightarrow (m^e)^d = m$$

all operations mod N