

RSA stuff

!! ★ Proj topics due today ★ !!

Roots of unity WTF? (what's the formula?)

roots of unity mod 15?

												-4	-3	-2	-1	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14		$: \mathbb{Z}_{15}$
square ↓	↓	↓	↓	↓	...						↓	↓	↓	↓		
0	1	4	9	1							1	9	4	1		

look at this from CRT perspective

0,0	1,1	2,2	0,3	1,4							2,1		2,4			$\mathbb{Z}_3 \times \mathbb{Z}_5$
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14		\mathbb{Z}_{15}
↓	↓	↓	↓	↓	...						↓	↓	↓	↓		
0	1	4	9	1							1	9	4	1		

roots of unity are

$(1, 1)$	$=$	$(1, 1)$
$(1, 4)$	$=$	$(1, -1)$
$(2, 1)$	$=$	$(-1, 1)$
$(2, 4)$	$=$	$(-1, -1)$

CRT

Something is true mod pq

\Leftrightarrow same thing true mod p & mod q

Ex: $x^2 \equiv_{pq} 1 \Leftrightarrow \begin{matrix} x^2 \equiv_p 1 & \& \\ x^2 \equiv_q 1 \end{matrix}$

Claim: If you have **nontrivial** **Sqrt** unity
then you can factor

$$x^2 \equiv_{pq} 1$$

$$\Rightarrow x^2 - 1 \equiv_{pq} 0$$

$$\Rightarrow (x+1)(x-1) \equiv_{pq} 0$$

$\Rightarrow (x+1)(x-1)$ is multiple
of pq

prime factorization of $(x+1)(x-1)$
contains p and q

$$x \not\equiv_{pq} 1$$

$$\Rightarrow x-1 \not\equiv_{pq} 0$$

$\Rightarrow x-1$ is not multiple
of pq

prime fact of $x-1$
doesn't contain both p & q

$$x \not\equiv_{pq} -1$$

$$\Rightarrow x+1 \not\equiv_{pq} 0$$

$\Rightarrow x+1$ is not mult of pq

prime fact of $x+1$
doesn't contain both p & q

Note: say $N = p \cdot q$. For any x ,

$\gcd(x, N)$ is either 1, p , q , or N

Strategy: identify some \underline{x} that is mult of p
but not q } q

Computing SQRT UNITY given e & d

given N, e, d : $ed \equiv_{\varphi(N)} 1$

can factor N

by first computing nontriv. $\text{sqrt}(1)$
(then do gcd to factor)

Idea:

$ed-1$ is mult of $\varphi(N)$

$$\text{So } w^{ed-1} \equiv_N w^0 \equiv 1$$

what if $ed-1$ is even: $ed-1 = 2k$

then

$$1 \equiv_N w^{ed-1} \equiv \underbrace{(w^k)^2}_{\text{sqrt of unity}}$$

If $ed-1$ is divisible by 4: $ed-1 = 4k$

$$1 \equiv w^{ed-1} \equiv ((w^k)^2)^2$$

In general: $ed-1 = 2^s \cdot r$

$$w^r \rightarrow w^{2r} \rightarrow w^{4r} \rightarrow w^{8r} \rightarrow \dots \rightarrow w^{2^s r}$$

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Claim: given N, e, m^e
hard to guess MSB of m

(because IF you could figure out $\text{MSB}(m)$ from
this information, THEN you could factor N)

Consequence: public key encryption!

public key: N, e

private key: d

to Encrypt $b \in \{0, 1\}$ using public key:

▶ choose random $m \leftarrow \mathbb{Z}_N$

▶ compute $c \equiv m^e$

$$x = \boxed{\text{MSB}(m)} \oplus b$$

▶ output (c, x)

looks random
to eavesdropper,
given N, e, c

to Decrypt:

$$b = x \oplus \text{MSB}(c^d)$$