USCS303 – OS: Practical – 03:RR Scheduling Algorithm

Practical – 03: Round-Robin (RR) Scheduling Algorithm	
Practical Date: 27 th July, 2021	
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Sample Output - 03	

Round Robin (RR) Scheduling Algorithm

Round-robin (RR) scheduling Algorithm is mainly designed for time-sharing systems.

This algorithm is similar to FCF scheduling, but in Round-robin (RR) scheduling, preemption is added which enables the system to switch between processes.

Round-robin scheduling algorithm is used to schedule process fairly each job a time slot or quantum and the interrupting the job if it is not completed by then the job come after the other job which is arrived in the quantum time that makes these scheduling fairly

Algorithm

Step 1: Input the number of processes and time quanta or time slice required to be scheduled using RR, burst time for each process.

Step 2: Choose the first process in the ready queue, set a timer to interrupt it after time quantum and dispatches it. Check if any other process request has arrived. If a process request arrives during the quantum time in which another process is executing, then add the new process to the Ready queue.

Step 3: After the quantum time has passed, check for any processes in the Ready queue. If the ready queue is empty and the current process is not complete, then add the current process to the end of the ready queue.

Step 4: Take the first process from the Ready queue and start executing it. Calculate the **Turn Around Time** and **Waiting Time** for each process using RR.

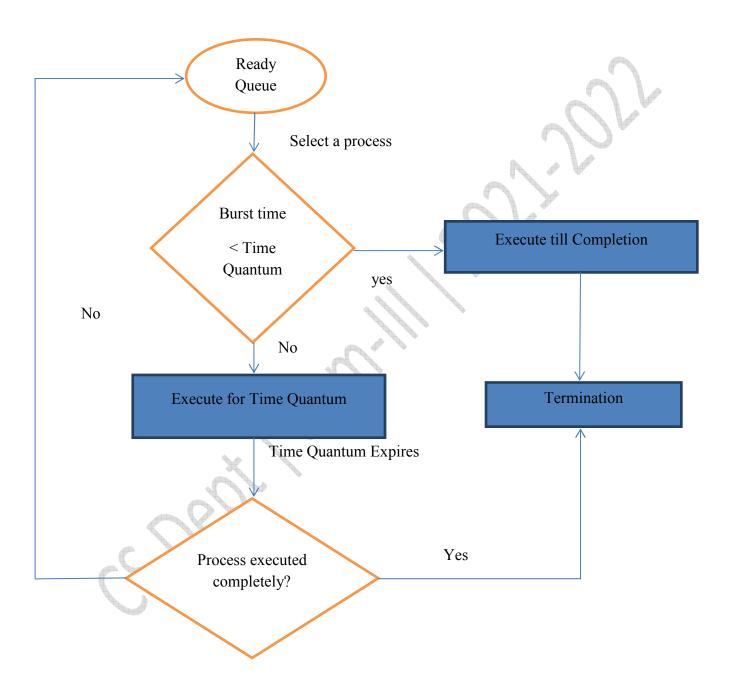
Step 5: Repeat all steps above from Step 2 to Step 4.

Step 6: If the process is complete and the ready queue is empty then the task is complete.

Step 7: Calculate the Average waiting time and Average turn around time.

Step 8: Stop.

RR Flowchart:



Solved Example of the Round-robins Algorithm

Process ID	Burst Time
P0	24
P1	3
P2	3

Assume Time Quanta: 4ms

Step 1: Consider the time quanta/ time slice = 4ms.

Step 2: Following shows the scheduling and execution of processes.

Step2.1: P0 process arrives at 0 with 24ms as the burst time which is greater than time quanta = 4ms. So P0 executes for 4ms and goes in waiting queue.

System Time	0
Process Scheduled	P0
Remaining Time	24-4 = 20
Waiting Time	0 - 0 = 0
Turn Around Time	0+4 = 4

Step 2.2: Next P1 process executes for 3ms which is greater than quanta time. So P1 executes and gets terminated.

System Time	4
Process Scheduled	P0,P1
Remaining Time	3-4 = 0
Waiting Time	4 - 0 = 4
Turn Around Time	4+3 = 7

Step 2.3: Next P2 process executes for 3ms which is greater than quanta time. So P2 executes and gets terminated

System Time	7
Process Scheduled	P0,P1,P2
Remaining Time	3-4 = 0
Waiting Time	7 - 0 = 7
Turn Around Time	7+3 = 10

Step 2.4: Now P0 turns comes again and it's the only process for execution so for 4ms of quanta it gets executed.

System Time	10
Process Scheduled	P0,P1,P2,P0
Remaining Time	20-4 = 16
Waiting Time	0
Turn Around Time	10+4 = 14

Step 2.5: Again, P0 continues to execute for next 4ms. Waiting for P0 will be zero.

System Time	14
Process Scheduled	P0,P1,P2,P0,P0
Remaining Time	16-4=12
Waiting Time	0
Turn Around Time	14+4=18

Step3: Calculate Average Waiting Time and Average Turn Around Time.

Average Waiting Time =
$$(6+4+7)/3 = 17/3 = 5.666667$$

Average Turn Around Time = (30+7+10)/3 = 47/3 = 16

Step 4: After scheduling of all provided processes:

Process ID	Burst Time	Turn Around Time	Waiting Time (Turn	
		(Completion Time –	Around Time – Burst	
		Arrival Time)	Time)	
P0	24	30-0=30	30-24=6	
P1	3	4+3=7	7-3=4	
P2	3	7+3=10	10-3=7	
Average		15.666667	5.666667	

Gnatt Chart:

P0	P1	P2	P0	P0	P0	P0	P0
0	4	7	10		30		

Example 2

Sample Output - 2

Process ID	Burst time
P0	2
Pl	1
P2	6

Assume Time Slice = Ims.

Process ID	Burst	Turn around time	Waiting time	
	time	(Finish time – Arrival time)	(Turn Around time – Burst Time)	
P0	2	(4-0) = 4	(4-2)=2	
P1	1	(2-0)=2	(2-1)=1	
P2	6	(9-0) = 9	(9-6)=3	
Average		2,000000	5.000000	

Gnatt Chart :-

	P0	P1	P2	P0	P2	P2	P2	P2	P2	
0)	1	2	3	4	5	6	7	8	9

Sample Output - 3

Process ID	Burst time
P0	7
P1	3
P2	2
P3	10
P4	8

Assume Time Slice = 3ms.

Process ID	Burst	Turn around time	Waiting time (Turn Around time – Burst Time)			
	time	(Finish time – Arrival time)				
P0	7	(24-0) = 24	(24 – 7) = 17			
P1	3	(6-0)=6	(6-3)=3			
P2	2	(8-0) = 8	(8-2)=6			
Р3	10	(30-0)=30	(30-10)=20			
P4	8	(29-0)=29	(29-8)=21			
Average		13.400000	5.000000			

Gnatt Chart

P	20	P1	P2	P3	P4	P0	P3	P4	P0	P3	P4	P3
	0	3	6	8	11	14	17	20	23 24	27	29	30

Implementation Source code:

```
\\NAME: Shraddha Sawant
\\Batch: B1
\\PRN: 2020016400773862
\\Date: 27th July, 2021
\\Prac-03: RR Scheduling Algorithm
import java.util.Scanner;
class P3 RR SS{
     public static void main(String args[]){
        Scanner input = new Scanner(System.in);
        int i,j,k,q,sum = 0;
        System.out.print("Enter number of process: ");
        int n= input.nextInt();
        int burstTime[] = new int[n];
        int waitingTime[] = new int[n];
        int turnAroundTime[] = new int[n];
        int a[] = new int[n];
System.out.println("Enter Burst Time of each process: ");
for(i = 0; i < n; i++){
        System.out.print("Enter brust Time for Process -P" + (i+1) + ":");
        burstTime[i] = input.nextInt();
        a[i] = burstTime[i];
}
System.out.print("Enter Time quantum: ");
```

```
q= intput.nextInt();
for(i = 0; i < n; i++)
         waitingTime[i] = 0;
int timer = 0;//current time
//Keep traversing processes in round robin manner
//until all of them are not done.
do{
      for(i=0;i< n; i++)
          //If burst time of a process is greater than 0 then only need to process further
          if(burstTime[i] > q){
             //Increase the value of t i.e shows how much time a process has been processed
             timer +=q;
             //Decrease the burst time of current process by quantum
             burstTime[i] -=q;
             for(j=0;j< n;j++)
                 if((j!=i) && (burstTime[j]!=0))
                     waitingTime[j]+=q;
}//if ends
//if burst time is smaller than or equal to quantum.Last cycle for this process
else{
       //increase the value of t i.e.shows how much time a process has been processed
       timer+=burstTime[i];
       for(j=0;j< n;j++)
          if((j!=i) \&\& (burstTime[j]!=0))
              waitingTime[i]+=burstTime[i];
```

```
}
       //as the process gets fully executed make its remaining burst time =0
       burstTime[i]=0;
}//else ends
sum = 0;
for(k=0;k< n;k++)
      sum+=burstTime[k];
while(sum!=0);
//calculating turnaround time by adding waitingTime+burstTime
for(i=0;i<n;i++)
   turnAroundTime[i]= waitingTime[i]+a[i];
float total=0;
for(int n= waitingTime){
     total += n;
}
float averageWaitingTime= total/n
total=0;
for(int n= turnAroundTime){
      total+=n;
float averageTurnAroundTime= total/n;
System.out.println("RR Algorithm: ")
System.out.format("%20s%20s%20s%20s\n", "ProcessId", "BurstTime". "WaitingTime",
"TurnAroundTime");
for (i=0; i< n; i++)
       System.out.format("%20s%20d%20d%20d\n",waitingTime[i],turnAroundTime[i]);
```

```
System.out.format("%40s%20f%20f\n", "Average", averageWaitingTime,averageTurnAroundTime);
}

INPUT:

Enter number of process:
Enter Brust Time of each process:
Enter brust Time for Process - P1: 24
Enter brust Time for Process - P2: 3
Enter brust Time for Process - P3: 3
Enter brust Time for Process - P3:
```

OUTPUT:



Sample output 01

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Sample output 02

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Sample output 03

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Enter Burst Time of each process:
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