

Prove that:

Intersection of two CFL  
or may not be CFL

1) May be a CFL  $L_1 \cap L_2$

language  
contained  
in another

$L_1$ : Palindrome [odd + even]

$L_2$ :  $\{a^n \mid n=1,2,\dots\}$

$= \{a, aa, aaa, aaaa, \dots\}$

$L_2 \subseteq L_1$

$L_1 \cap L_2 = L_2$

$$\begin{array}{l}
 \vdots \\
 \downarrow \\
 L_2: \{0^i 1^n 2^n \mid n \geq 1, i \geq 1\} \Rightarrow \left. \begin{array}{l} S \rightarrow S_1 S_2 \\ S_1 \rightarrow 0 S_1 \\ S_2 \rightarrow 1 S_2 2 \end{array} \right\} \\
 \searrow L_2 = \{012, 0012, 001122, \dots, 0001122\}
 \end{array}$$

$$\underline{\{0^n 1^n 2^n \mid n \geq 1\}}$$

$\rightarrow$  Using pumping lemma we have proved that it is not

Assume  $L_1$  and  $L_2$  to be two CFLs

$\Rightarrow L_1'$  and  $L_2'$  would also be CFL

$\Rightarrow L_1' \cup L_2'$  would also be a CFL (CFL is closed under union)

$\Rightarrow (L_1' \cup L_2')'$  would also be CFL

But we know,  $(L_1' \cup L_2')' = L_1 \cap L_2$  which is not a CFL

i)  $\{a^n b^m c^{(n+m)}, n \geq 1, m \geq 1\}$

and

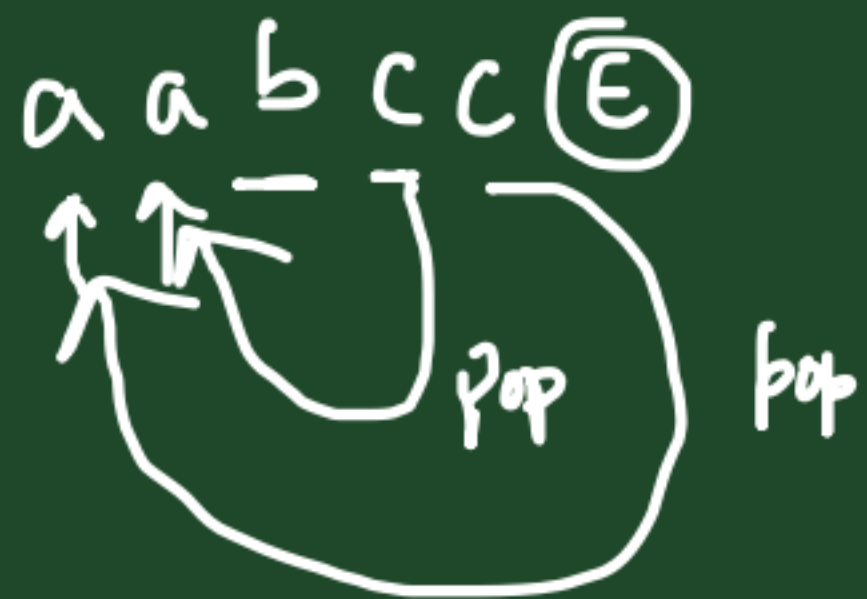
ii)  $\{a^n b^{(n+m)} c^m, n \geq 1, m \geq 1\}$

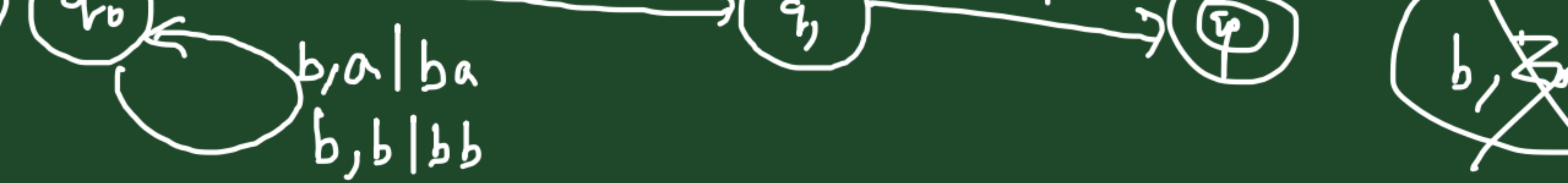
iii)  $\{a^{n+m} b^m c^n, n \geq 1, m \geq 1\}$

iv)  $\{a^n b^m a^m b^n, n \geq 1, m \geq 1\}$

→ PDA

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$b, a \in$



$b, z_0 / bz_0$

$b, b / bb$

$\in \{z_0 | z_1\}$

v)  $a^{n+m} b^m c^n$ ,  $n, m \geq 1$

(HW)

v)  $a^n b^m a^m b^n$ ,  $n, m \geq 1$

(HW)

aaaaa bb c



