

UNIT-IV

ELEMENTARY NUMBER THEORY AND CRYPTOGRAPHY

DIVISIBILITY

If a and b are any two integers such that $b \neq 0$ then we say that " b divides a " if there exists an integer, k such that $a = kb$. And it is written as $b|a$.

Note

If " b divides a ", then we say that " b is factor of a " or " a is multiple of b ".

DIVISION ALGORITHM

If a and b are any two integers such that $b > 0$ then there exist unique integer q and r such that,

$$a = bq + r$$

where q is called quotient

r is called remainder

CONGRUENCE RELATION

Let m be a positive integer. Then an integer ' a ' is said to be congruent to an integer ' b ', under modulo m , if " m divides $(a-b)$ ". ($m|(a-b)$) symbolically it is written as $a \equiv b \pmod{m}$ or $a \equiv b \pmod{m}$

It is read as "a is congruent to b modulo m".

Note

b is called remainder or residue of $a \pmod{m}$

OR b is remainder when m divides a

Properties of Congruence Relation

- i) If $a \equiv b \pmod{m}$, then $m \mid (a-b)$
- ii) If $a \equiv b \pmod{m}$ then $b \equiv a \pmod{m}$
- iii) If $a \equiv b \pmod{m}$ & $b \equiv c \pmod{m}$ then $a \equiv c \pmod{m}$

MODULAR ARITHMETIC OPERATION

If $a \equiv b \pmod{m}$, then for $k \neq 0 \in \mathbb{Z}$

- i) $a+k \equiv b+k \pmod{m}$
- ii) $a-k \equiv b-k \pmod{m}$
- iii) $a \cdot k \equiv b \cdot k \pmod{m}$
- iv) $a^k \equiv b^k \pmod{m}$

Properties of Modular Arithmetic

- Residue system modulo m

Define the set \mathbb{Z}_m as a set of non-negative integers less than m ,

$$\mathbb{Z}_m = \{0, 1, 2, \dots, (m-1)\}$$

It is called residue system modulo m .

- Residue classes

Each integer in \mathbb{Z}_m represents a residue class and it is divided by $[a]$ and defined by

$$[r] = \{x : x \equiv r \pmod{m}\}$$

For Example:

(i) Residue system modulo 3 is,

$$\mathbb{Z}_3 = \{0, 1, 2\}$$

∴ Residue classes of elements of \mathbb{Z}_3 are

$$[0] = \{x : x \equiv 0 \pmod{3}\}$$

$$[0] = \{\dots, -9, -6, -3, 0, 3, 6, 9, \dots\}$$

$$[1] = \{\dots, -8, -5, -2, 1, 4, 7, 10, \dots\}$$

$$[2] = \{\dots, -7, -4, -1, 2, 5, 8, 11, \dots\}$$

Theorem

Let m be a positive integer and $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$. Then prove that $a+c \equiv b+d \pmod{m}$. Then prove that $a+c \equiv b+d \pmod{m}$ and $ac \equiv bd \pmod{m}$.

Proof: Given $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$

$$m \mid (a-b) \text{ and } m \mid (c-d) \quad \text{--- (1)}$$

$$\Rightarrow m \mid (a-b)+(c-d)$$

$$\Rightarrow m \mid (a-b+c-d)$$

$$\Rightarrow m \mid (a+c)-(b+d)$$

$$\Rightarrow [a+c \equiv b+d \pmod{m}]$$

By (1) we get a solution for $a+b+c-d$

$$a-b = K_1 m \text{ and } c-d = K_2 m$$

$$c(a-b) = CK_1 m \text{ and } b(c-d) = BK_2 m$$

$$ac-bc = (K_1 c) m \text{ and } bc-bd = (K_2 b) m$$

$$ac-bc+bc-bd = (K_1 c) m + (K_2 b) m$$

$$ac-bd = K'_1 m + K''_1 m \quad \text{where } K'_1 = K_1 c, K''_1 = K_2 b$$

$$ac-bd = (K'_1 + K''_1) m$$

$$ac-bd = K_1 m \quad \text{where } K_1 = K'_1 + K''_1 \in \mathbb{Z}$$

$$\Rightarrow m \mid (ac-bd)$$

$$\Rightarrow [ac \equiv bd \pmod{m}]$$

PRIME NUMBER

An integer $p \geq 2$ is called prime number if it is divisible by 1 and itself. Otherwise a number is called composite number.

Ex : ① Prime numbers are $2, 3, 5, 7, 11, 13, 17, 19, 23, \dots$

② Composite numbers are $4, 6, 8, 9, 10, \dots$

Note

Every composite number can be expressed as product of prime integers:

$$\text{Ex: } \textcircled{1} 10 = 2 \times 5 \quad \textcircled{2} 20 = 2 \times 10 = 2 \times 2 \times 5 \quad \textcircled{3} 35 = 5 \times 7$$

RELATIVELY PRIME NUMBER (CO-PRIME NUMBERS)

Two numbers a and b are said to be relatively prime if they have no common divisors other than 1.
 $(\text{GCD}(a, b) = 1)$

Ex: $\textcircled{1}$ 10 and 21 are relatively prime as $\text{GCD}(10, 21) = 1$

EULER'S ϕ -FUNCTION / EULER'S TOTIENT FUNCTION

REDUCED RESIDUE SYSTEM MODULO m

The reduced residue system modulo m is the set of all elements from residue system modulo m . $\mathbb{Z}_m = \{0, 1, 2, \dots, m-1\}$ which are relatively prime to m .

i.e $S = \{x : \text{GCD}(x, m) = 1\}$

EULER'S ϕ -FUNCTION / EULER'S TOTIENT FUNCTION

The Euler's ϕ function of an integer $n \geq 1$ is denoted by $\phi(n)$ and defined by the number of non-zero positive integers less than n that are relatively prime to n .

$$\text{Ex: } \phi(1) = 0, \phi(2) = n(\{1, 2\}) = 1, \phi(3) = n(\{1, 2\}) = 2$$

$$\phi(4) = n(\{1, 3\}) = 2, \phi(5) = n(\{1, 2, 3, 4\}) = 4$$

$$\phi(6) = n(\{1, 3\}) = 2, \phi(7) = 6$$

Note

1 If $n = p$ is a prime number then Euler ϕ function of p is $\phi(p) = p-1$

2 If n is number that can be expressed as a product of relatively prime numbers (a, b) Euler's ϕ function of n is,

$$\text{Ex:- } \begin{aligned} \textcircled{1} \quad \phi(20) &= \phi(4 \cdot 5) = \phi(4) \cdot \phi(5) = 2 \times 4 = 8 \\ \textcircled{2} \quad \phi(15) &= \phi(5 \cdot 3) = \phi(5) \cdot \phi(3) = 4 \times 2 = 6 \\ \textcircled{3} \quad \phi(10) &= \phi(2 \cdot 5) = \phi(2) \cdot \phi(5) = 1 \times 4 = 4 \end{aligned}$$

Euler's Theorem.

Statement: Let n and a be positive integers which are relatively prime ($\text{GCD}(n, a) = 1$). Then

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

where $\phi(n)$ is Euler ϕ -function.

Proof:

Given that $n, a \in \mathbb{Z}_{>0}$ such that $\text{GCD}(n, a) = 1$.

Euler ϕ -function of n is,

$$\phi(n) = k$$

Let us take

Consider a reduced system modulo n ,

$$S = \{a_1, a_2, a_3, \dots, a_{\phi(n)}\}$$

$$\text{OR } S = \{a_1, a_2, a_3, \dots, a_k\}$$

Take $a \neq 0 \in \mathbb{Z}$ such that $\text{GCD}(n, a) = 1$

$$aS = \{aa_1, aa_2, aa_3, \dots, aa_k\}$$

$$aa_1 \equiv a_1 \pmod{n}; aa_2 \equiv a_2 \pmod{n}; \dots; aa_k \equiv a_k \pmod{n}$$

Now, we know that,

$$aa_1 \cdot aa_2 \cdot aa_3 \cdot \dots \cdot aa_k \equiv a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_k \pmod{n}$$

$$a^k (a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_k) \equiv (a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_k) \pmod{n}$$

$$\Rightarrow a^k \equiv 1 \pmod{n}$$

$$\Rightarrow a^{\phi(n)} \equiv 1 \pmod{n}$$

FERMAT'S THEOREM / FERMAT'S LITTLE THEOREM

Statement Let P be a prime number and $P \nmid a$. Then

$$a^{P-1} \equiv 1 \pmod{P}$$

$$\text{OR } a^P \equiv a \pmod{P}$$

Proof

Given that, P is a prime number and $P \nmid a$.

$$\Rightarrow \text{GCD}(P, a) = 1$$

Euler ϕ -function of prime no P is.

$$\phi(P) = P - 1$$

Consider a residue system modulo P ,

$$S = \{a_1, a_2, a_3, \dots, a_{P-1}\}$$

Take $a \neq 0 \in \mathbb{Z}$ such that $P \nmid a$

$$as = \{aa_1, aa_2, aa_3, \dots, aa_{P-1}\}$$

\in

$$aa_1 \equiv a_1 \pmod{P}; aa_2 \equiv a_2 \pmod{P}, \dots, aa_{P-1} \equiv a_{P-1} \pmod{P}$$

Next we know that,

$$aa_1 \cdot aa_2 \cdot aa_3 \cdot \dots \cdot aa_{P-1} \equiv a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_{P-1} \pmod{P}$$

$$\Rightarrow a^{P-1} (a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_{P-1}) \equiv (a_1 \cdot a_2 \cdot a_3 \cdot \dots \cdot a_{P-1}) \pmod{P}$$

$$\Rightarrow a^{P-1} \equiv 1 \pmod{P}$$

$$\Rightarrow a^P \equiv a \pmod{P}$$

For example

$$\textcircled{1} \quad 4^{\phi(7)} \equiv 1 \pmod{7}.$$

$$\phi(7) = 6$$

$$4^6 \equiv 1 \pmod{7}$$

$$4^2 \cdot 4^2 \cdot 4^2 \equiv 1 \pmod{7}$$

$$2 \cdot 2 \cdot 2 \equiv 1 \pmod{7}$$

$$8 \equiv 1 \pmod{7}$$

$$\textcircled{2} \quad 5^{\phi(6)} \equiv 1 \pmod{6}$$

$$\phi(6) = \varphi(\{1, 3\}) = 2$$

$$5^2 \equiv 1 \pmod{6}$$

CHINESE REMAINDER THEOREM

Statement: If $m_1, m_2, m_3, \dots, m_k$ are pairwise relatively prime numbers and $a_1, a_2, a_3, \dots, a_k$ are any integers, then, the simultaneous congruence relations,

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

$$x \equiv a_3 \pmod{m_3}$$

$$\vdots \quad \vdots$$

$$x \equiv a_k \pmod{m_k}$$

has a unique solution under modulo M ,

where $M = m_1 \cdot m_2 \cdot m_3 \cdots \cdot m_k$

STEPS TO SOLVE SIMULTANEOUS CONGURENCE RELATION BY CHINESE THEOREM.

Step 1: Check $\text{GCD}(m_i, m_j) = 1$ for $i \neq j$

Step 2: $x \equiv (M_1 x_1 a_1 + M_2 x_2 a_2 + \dots + M_k x_k a_k) \pmod{M}$

where $M = m_1 \cdot m_2 \cdot \dots \cdot m_k$

and $M_i = \frac{M}{m_i} = m_1 \cdot m_2 \cdots m_{i-1} \cdot m_{i+1} \cdots m_k$

Next, to calculate x_i^o . x_i^o is a multiplicative inverse of M_i^o under modulo m_i^o
i.e. $M_i^o x_i^o \equiv 1 \pmod{m_i^o}$

Step 3: Using values of M_i^o and x_i^o in ① we can find the value of x

Examples

① Solve the following system of congruences by using Chinese remainder theorem.

$$x \equiv 2 \pmod{3}, \quad x \equiv 1 \pmod{4}, \quad x \equiv 3 \pmod{5}$$

→ Given,

$$x \equiv 2 \pmod{3}$$

$$x \equiv 1 \pmod{4}$$

$$x \equiv 3 \pmod{5}$$

$$a_1 = 2, \quad a_2 = 1, \quad a_3 = 3$$

$$m_1 = 3, \quad m_2 = 4, \quad m_3 = 5$$

$$\text{GCD}(3, 4) = \text{GCD}(4, 5) = \text{GCD}(3, 5) = 1$$

∴ 3, 4, 5 are pairwise relatively prime

$$\text{Next, } x = (M_1 x_1 a_1 + M_2 x_2 a_2 + M_3 x_3 a_3) \pmod{M} \quad \text{--- (1)}$$

$$\text{where } M = m_1 m_2 m_3 = 3 \cdot 4 \cdot 5 = 60$$

$$M = 60$$

$$M_1 = \frac{M}{m_1} = \frac{60}{3} = 20$$

$$M_2 = \frac{M}{m_2} = \frac{60}{4} = 15$$

$$M_3 = \frac{M}{m_3} = \frac{60}{5} = 12$$

Calculate x_1

$$M_1 x_1 \equiv 1 \pmod{m_1}$$

$$M_1 x_1 \equiv 1 \pmod{m_1}$$

$$\Rightarrow 20x_1 \equiv 1 \pmod{3}$$

$$\Rightarrow 2x_1 \equiv 1 \pmod{3}$$

②

$$\therefore x_1 = 2 \text{ as } 3 \nmid 2 - 1$$

$$M_2 x_2 \equiv 1 \pmod{m_2}$$

$$15x_2 \equiv 1 \pmod{5}$$

$$3x_2 \equiv 1 \pmod{5}$$

$$x_2 = 3 \text{ as } 5 \nmid 3 - 1$$

$$M_3 x_3 \equiv 1 \pmod{m_3}$$

$$12x_3 \equiv 1 \pmod{5}$$

$$2x_3 \equiv 1 \pmod{5}$$

$$x_3 = 3 \text{ as } 5 \nmid 3 - 1$$

∴ Relation ① becomes,

$$x \equiv [20 \cdot 2 \cdot 2 + 15 \cdot 3 \cdot 1 + 12 \cdot 3 \cdot 3] \pmod{60}$$

$$x \equiv [80 + 45 + 108] \pmod{60}$$

$$x \equiv 233 \pmod{60}$$

$$x \equiv 53 \pmod{60}$$

$$\boxed{x \equiv 53}$$

$$② x \equiv 2 \pmod{3}, x \equiv 3 \pmod{5}, x \equiv 2 \pmod{7}$$

$$\rightarrow a_1 = 2, a_2 = 3, a_3 = 2$$

$$m_1 = 3, m_2 = 5, m_3 = 7$$

$$\text{GCD}(3, 5) = \text{GCD}(5, 7) = \text{GCD}(3, 7) = 1$$

3, 5, 7 are pair wise relatively prime.

$$(1) \text{ Next, about } (x_1, x_2, x_3) \text{ such that } (0, 0, 0), (0, 0, 1), (0, 1, 0), (0, 1, 1), (1, 0, 0), (1, 0, 1), (1, 1, 0), (1, 1, 1)$$

$$x \equiv (M_1 x_1 a_1 + M_2 x_2 a_2 + M_3 x_3 a_3) \pmod{M} \quad \text{--- (1)}$$

$$\text{where } M = m_1 m_2 m_3 = 3 \cdot 5 \cdot 7 = 105$$

$$M = 105$$

$$M_1 = \frac{M}{m_1} = 5 \cdot 7 = 35$$

$$M_2 = \frac{M}{m_2} = 3 \cdot 7 = 21$$

$$M_3 = \frac{M}{m_3} = 3 \cdot 5 = 15$$

(calculate x_i)

$$M_1 x_1 = 1 \pmod{m_1}$$

$$35x_1 = 1 \pmod{3}$$

$$2x_1 = 1 \pmod{3}$$

$$x_1 = 2 \quad 3|2(2)-1$$

$$M_2 x_2 = 1 \pmod{m_2}$$

$$21x_2 = 1 \pmod{5}$$

$$1x_2 = 1 \pmod{5}$$

$$x_2 = 1 \quad 5|1(1)-1$$

$$M_3 x_3 = 1 \pmod{m_3}$$

$$15x_3 = 1 \pmod{7}$$

$$1x_3 = 1 \pmod{7}$$

$$x_3 = 1 \quad 7|1(1)-1$$

\therefore Relation ① becomes

$$x = 35 \cdot 2 \cdot 2 + 21 \cdot 1 \cdot 3 + 15 \cdot 1 \cdot 2 \pmod{105}$$

$$x = 140 + 63 + 30$$

$$x \equiv 233 \pmod{105}$$

$$x \equiv 23 \pmod{105}$$

$$\boxed{x = 23}$$

GREATEST COMMON DIVISION BETWEEN THE POSITIVE INTEGER

Find GCD between following pair of integers.

$$(i) 100, 37$$

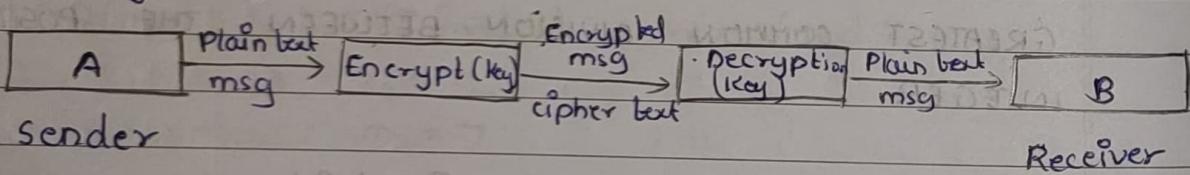
$$\begin{aligned} \rightarrow \text{GCD}(37, 100) &= \text{GCD}(37, 300 \pmod{37}) \\ &= \text{GCD}(37, 26) \\ &= \text{GCD}(26, 37 \pmod{26}) \\ &= \text{GCD}(26, 11) \\ &= \text{GCD}(11, 26 \pmod{11}) \\ &= \text{GCD}(11, 4) \\ &= \text{GCD}(4, 11 \pmod{4}) \\ &= \text{GCD}(4, 3) \\ &= \text{GCD}(3, 4 \pmod{3}) \\ &= \text{GCD}(3, 1) \\ &= \text{GCD}(1, 3 \pmod{1}) \\ &= 1 \end{aligned}$$

(ii) 252, 52

$$\begin{aligned} \rightarrow \text{GCD}(52, 252) &= \text{GCD}(52, 252 \text{ mod } 52) \\ &= \text{GCD}(52, 44) \\ &= \text{GCD}(44, 52 \text{ mod } 44) \\ &= \text{GCD}(44, 8) \\ &= \text{GCD}(8, 44 \text{ mod } 8) \\ &= \text{GCD}(8, 4) \\ &= \text{GCD}(4, 8 \text{ mod } 4) \end{aligned}$$

CRYPTOGRAPHY AND NETWORK SECURITY

Cryptography: A method of protecting information by transforming it into an unreadable format is called cryptography.



case

- 1) If Key are same then it is symmetric cryptography
- 2) If Key are different then it is asymmetric cryptography

Encryption: The process of transforming information from readable format to unreadable format.

Decryption: The process of transforming information from unreadable format to readable format.

Key: String of bits used by cryptography algorithm convert plain text to cipher text and vice-versa.

RSA ALGORITHM (Rivest, Shamir)

RSA Algorithm is asymmetric cryptographic algorithm it means that 2 keys public key and private key are different.

Public Key: It is known to all users in the network

Private Key: It is kept secret and not shareable by all.

If a public key of a user A is used for encryption then we have to use the private key of the same user A for decryption.

RSA is a block cipher in which the plain text and cipher text are integers between zero to $n-1$ for some value n .

Algorithm

STEP 1: Key Generation.

[i] Select two large prime numbers p and q if p and q are more and more large then security is higher.

[ii] Calculate $n = p * q$.

[iii] Calculate Euler's ϕ function.

$$\phi(n) = \phi(pq) = \phi(p)\phi(q) = (p-1)(q-1)$$

[iv] Choose the value e such that,

$$1 < e < \phi(n) \text{ and } \text{GCD}(\phi(n), e) = 1$$

[v] calculate $ed \equiv 1 \pmod{\phi(n)}$

$$\text{OR } d \equiv e^{-1} \pmod{\phi(n)}$$

It means that e is multiplicative inverse of d under modulo $\phi(n)$

[vi] Public Key : { e, n }

[vii] Private Key : { d, n }

STEP 2: RSA Encryption.

$$C \equiv M^e \pmod{n}$$

STEP 3: RSA Decryption

$$M \equiv C^d \pmod{n}$$

Where M is a plain text.

C is a cipher text.

i) Perform encryption and decryption using RSA algorithm

for $P=7$, $q=11$ and $M=6$

Given that $p=7$, $q=11$

(i) $n = pq = 7 \cdot 11 = 77$.

$$\boxed{n=77}$$

(ii) Euler's ϕ function is

$$\phi(n) = \phi(pq) = (p-1)(q-1)$$

$$= (7-1)(11-1)$$

$$= 6 \times 10$$

$$\boxed{\phi(n) = 60}$$

(iii) Choose $e=7$ such that $1 < e < 60$ and $\text{GCD}(60, 7)=1$

(iv) $ed \equiv 1 \pmod{\phi(n)}$

$$7d \equiv 1 \pmod{60}$$

$$d = 43 \quad \text{as} \quad 60 \mid 7(43-1)$$

$$60k + 1 = 7d$$

3633
77 67

Public key : $\{e, n\} = \{7, 77\}$

Private Key : $\{d, e\} = \{43, 77\}$

STEP 2 RSA Encryption

$$C \equiv M^e \pmod{n}$$

$$C \equiv 6^7 \pmod{77}$$

$$C \equiv 41 \pmod{77}$$

$$C \equiv 1$$

Cipher text - $C = 41$

STEP 3 :

$$M \equiv C^d \pmod{n}$$

$$M \equiv 41^{43} \pmod{77}$$

$$M \equiv 6 \pmod{77}$$

Plain text - $M = 6$.

② Perform encryption and decryption using RSA algorithm for $P=13$, $q=17$ and $n=18$

Given that $P=13$, $q=17$

$$n = P \cdot q$$

$$= 13 \cdot 17$$

$$n = 221$$

Euler's ϕ function

$$\phi(n) = \phi(P \cdot q)$$

$$= (P-1)(q-1)$$

$$= 12 \cdot 16$$

$$= 192$$

Choose $e = 17$ such that $1 < e < 192$ and $\text{GCD}(192, 17) = 1$

$$ed \equiv 1 \pmod{\phi(n)}$$

$$75d \equiv 1 \pmod{192}$$

$$d = 55 \text{ for } k=2$$

$$d = \frac{192}{75} = 192 \times \frac{1}{75} = 192 \times (-1)$$

$$75d - 1 = 192$$

$$d = 192 \times (+) = 192 \times \frac{1}{75}$$

$$\text{Public Key } \rightarrow \{e, n\} = \{5, 221\}$$

$$\text{Private Key } \rightarrow \{d, n\} = \{55, 221\}$$

$$(55 \text{ mod } 221)^{-1} = 1286290$$

RSA encryption

$$\begin{aligned} C &\equiv M^e \pmod{n} \\ &= 18^5 \pmod{221} \\ &= 18^5 \pmod{221} \\ &= 86 \pmod{221} \end{aligned}$$

Cipher text is $c = 86$

RSA decryption.

$$\begin{aligned} M &\equiv c^d \pmod{n} \\ &= 86^{55} \pmod{221} \end{aligned}$$

$$M = 18 \pmod{221}$$

Plain text is $M = 18$

GROUPS

A non-empty set G with binary operation $*$ denoted by $(G, *)$ is said to be group if satisfies following axioms (Laws)

(1) Closure Laws: $\forall a, b \in G : (a * b) \in G$

(2) Associative law: $a * (b * c) = (a * b) * c$

(3) Existence of Identity Element: $\exists e \in G \ni e * a = a * e = a$
 $\forall a \in G$

(4) Existence of inverse element: $\forall a \in G \exists a^{-1} \in G \ni a * a^{-1} = a^{-1} * a = e$

If $(G, *)$ is said to be abelian group if it satisfies

- (5) Commutative law: $a * b = b * a \quad \forall a, b \in G$.

RINGS

A set R with two binary operations $+$ and \times is denoted by $(R, +, \times)$ is said to be ring if it satisfies following axioms.

- 1 Closure under addition: $\forall a, b \in R, a+b \in R$
- 2 Associative under addition: $a+(b+c) = (a+b)+c$
- 3 Additive identity element: $\exists a \in R, \exists a+0 = 0+a = a \quad \forall a \in R$
- 4 Additive inverse element: $\forall a \in R, \exists -a \in R, \exists -a+a = 0$
- 5 Commutative under addition: $a+b = b+a, \forall a, b \in R$
- 6 Closure under multiplication: $\forall a, b \in R, ab \in R$
- 7 Associative under multiplication: $a(bc) = (ab)c$
- 8 Distributive Law: $a(b+c) = ab+ac \quad \forall a, b, c \in R$.

A ring $(R, +, \times)$ is said to be commutative ring if it satisfies.

- 9 Commutative under multiplication: $ab = ba \quad \forall a, b \in R$

A commutative ring $(R, +, \times)$ is called integral domain if it satisfies

- 10 Multiplicative identity $\exists 1 \in R, \exists 1 \cdot a = a \cdot 1 = a, \forall a \in R$
- 11 No zero Divisions: $\forall a, b \in R, \text{ if } ab = 0, \text{ then } a=0 \text{ or } b=0$.

FIELDS

A set F under the operation $+$ and \times is denoted by

$(F, +, \times)$ and said to be field if $(F; +, \times)$ is integral domain and satisfies

- 12) Multiplicative inverse: $\forall a \in F \setminus \{0\}, \exists a^{-1} \in F$
 $\Rightarrow aa^{-1} = a^{-1}a = 1$

Finite Field.

For a given prime number P we define the finite field of order P denoted by $GF(P)$ as the set \mathbb{Z}_P of all integers $\{0, 1, 2, \dots, (P-1)\}$ together with arithmetic operation modulo P .

Finite field of the form $GF(P)$
 GF stands for Galois Field in honor of great, mathematician Evariste Galois (25th October 1811 - 31st 1832) who first studied finite fields. We have finite field of GF . This finite field is studied for prime number P .

Example

①

Show that $GF(3)$ is a finite field.

Given that $P=3$ is a prime number.

$$\mathbb{Z}_3 = \{0, 1, 2\}$$

Addition under modulo 3

| \oplus | 0 | 1 | 2 |
|----------|---|---|---|
| 0 | 0 | 1 | 2 |
| 1 | 1 | 2 | 0 |
| 2 | 2 | 0 | 1 |

Multiplication under modulo 3.

| \times | 0 | 1 | 2 |
|----------|---|---|---|
| 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 2 |
| 2 | 0 | 2 | 1 |

Additive inverse under modulo 3

| a | 0 | 1 | 2 |
|----|---|---|---|
| -a | 0 | 2 | 1 |

Multiplicative inverse under modulo 3

| a | 0 | 1 | 2 |
|----------|---|---|---|
| a^{-1} | - | 1 | 2 |

$\therefore GF(3)$ is a finite field.

② Show that $GF(5)$ is a finite field.

Given that $P=5$ is a prime number.

$$\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$$

Addition under modulo 5

| \oplus | 0 | 1 | 2 | 3 | 4 |
|----------|---|---|---|---|---|
| 0 | 0 | 1 | 2 | 3 | 4 |
| 1 | 1 | 2 | 3 | 4 | 0 |
| 2 | 2 | 3 | 4 | 0 | 1 |
| 3 | 3 | 4 | 0 | 1 | 2 |
| 4 | 4 | 0 | 1 | 2 | 3 |

Multiplication under modulo 5

| \otimes | 0 | 1 | 2 | 3 | 4 |
|-----------|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 2 | 3 | 4 |
| 2 | 0 | 2 | 4 | 1 | 3 |
| 3 | 0 | 3 | 1 | 4 | 2 |
| 4 | 0 | 4 | 3 | 2 | 1 |

Additive Inverse under modulo 5

| a | 0 | 1 | 2 | 3 | 4 | 5 | 1 | 0 | 0 |
|----|---|---|---|---|---|---|---|---|---|
| -a | 0 | 4 | 3 | 2 | 1 | 1 | 4 | 0 | 0 |

Multiplicative inverse under modulo 5.

| a | 0 | 1 | 2 | 3 | 4 |
|----|---|---|---|---|---|
| -a | - | 1 | 3 | 2 | 4 |

$\therefore GF(5)$ is a finite field.