

Termwork 7

Implement All-pairs Shortest Paths Problem using  
Floyd's Algorithm.

### Objective of the Experiment :

- To introduce the concept of dynamic programming
- Present the working of Floyd's Algorithm
- To find the shortest path from all nodes to all other nodes in a given graph
- Analyze the Algorithm Complexity.

### Theory:

#### Dynamic programming:

- It is a technique for solving problems with overlapping sub problems
- It suggests solving each of the smaller sub-problems only once and recording the results in a table from which a solution to the original problem can then be obtained

#### Transitive Closure:

If there exists an edge between  $V_i$  and  $V_k$  and also between  $V_k$  to  $V_j$  then there is a path between  $V_i$  and  $V_j$

#### Graphs:

Is a non-linear data structure consisting of nodes and edges. Consists of finite set of vertices and set of edges which connect a pair of nodes.

#### Weight Matrix:

$$W(i, j) = 0 \quad \text{if } i = j$$

$$W(i, j) = \infty \quad \text{if no edge between } i \text{ and } j$$

$$W(i, j) = \text{"weight of edge"}$$



Algorithm:

input: weight Matrix  $W$  of a graph with no-negative-length cycle

output: The distance matrix of the shortest path's length

begin

$D \leftarrow W$ ,  $P[i, j] \leftarrow 0$

for  $k \leftarrow 1$  to  $n$ , do

for  $i \leftarrow 1$  to  $n$ , do

for  $j \leftarrow 1$  to  $n$ , do

$D[i, j] \leftarrow \min \{ D[i, j], D[i, k] + D[k, j] \}$ ,  ~~$D[i, j] \leftarrow \min \{ D[i, j], D[i, k] + D[k, j] \}$~~

return  $D$

end.



Program:

```
#include <stdio.h>
```

```
void printMatrix (int D[10][10], int n) {
```

```
    int i, j;
```

```
    for (i=0; i<n; i++) {
```

```
        for (j=0; j<n; j++)
```

```
            printf ("%d ", D[i][j]);
```

```
        printf ("\n");
```

```
    }
```

```
}
```

```
void floyds (int D[10][10], int n) {
```

```
    int i, j, k;
```

```
    for (k=1; k<=n; k++) {
```

```
        printf ("With %d as intermediate vertex\n", k);
```

```
        printf ("Cost Matrix now: \n");
```

```
        for (i=1; i<=n; i++)
```

```
            for (j=1; j<=n; j++)
```

```
                if (i==j)
```

```
                    D[i][j] = 0;
```

```
            else
```

```
                D[i][j] = min (D[i][j], D[i][k] + D[k][j]);
```

```
        printMatrix (D, n);
```

```
    }
```

```
}
```

```
int minimum (int a, int b) {
```

```
    return (a<b) ? a : b;
```

```
}
```



```

int main() {
    int D[10][10], w, n, e, u, v, i, j;
    clrscr();
    printf("\n Enter the number of vertices: ");
    scanf("%d", &n);

    printf("\n Enter the Cost Matrix: \n");
    for (i=1; i<=n; i++)
        for (j=1; j<=n; j++)
            scanf("%d", &D[i][j]);
    printf("\n The initial cost matrix: \n");
    printMatrix(D, n);
    printf("\n The shortest paths are: \n");
    for (i=1; i<=n; i++)
        for (j=1; j<=n; j++) {
            if (i != j)
                printf("\n <%d, %d> ==> %d", i, j, D[i][j]);
        }
    }
}

```



### References :

- Kenneth Berman, Jerome Paul, Algorithms, Cengage learning
- Thomas H Cormen, Charles E. Leiserson, Ronal L Rivest, Clifford Stein, Introduction to Algorithms, MIT, 2nd Edition & Onwards.

### Conclusion:

- In this Termwork, we learnt about Dynamic Programming, Graphs, Floyd's Algorithm and applied to solve this problem.
- We also learnt ~~comparing~~ the basic problem solving techniques.