TERMWORK 9.

PROBLEM STATEMENT

Find a subset of a given set 5-951, 52. -. 5 n 3 of n positive integers whose sum is equal to a give positive Integer d.

For example

7/ 5= 11,2,5,683 and d=9, there are two soluboons

A suitable message to do displayed if given problem

OBJECTIVE

· To inbroduce the concept of backtracking

· Present the work of subset -sum problem.

· To find the subset as a solution for a given positive integer. · Analyze the Algorithm completity.

THEORY

In Subset-Sum problem ; we find a subset of a given set s= {s1....sn} of n positive integer where sum is equal to a given positive integer d.

Example.

5={1,2,5,6,83 and d=9 solution aver \$1,83 and \$1,2,63 Some of the instances may have no solutions.

Il is convincent to soft element in increasing order. The state space true can be constructed as a benavy true. A path from the root to a node on the level of tree Indicates which of first i numbers have been included in subsets represented by that node. We record the value obs, the sum of there numbers, in node of s is equal to d, we have the solution to problem.

ALGORITHM Il gives a template on genné backtracking algorithm. 11 Input : X[1. --- ?] 11 Output: All buples respresenting solutions. 2) × [1....]) & a solution weste × [1....] else for each element or E Siti consistent with x[1...i] and constraints do · [x(i+1) < x) Backbracking (x[1...?+1])

```
PROGRAM
# include (stdio h)
# include (stillib b)
# Defore MAX 50
# define TRUE 1
# define FAISE O
Put inc[HAX], w[HAX], sum, n;
int promising (int i, int wt, int total)!
   return (((wt+bobal)) = sum) && ((wt=sum))
              11 (wt +w[i+1] <= sum )));
 void sumset (int i int wt, int botal) {
   of (promising (i wt, total))
        ( (wt = = sum)
            print (4 10 & 12");
             for (j=0; j <= 1; j++)
            ig (inc [j] print ( " ".d", w[j]);
print ("\n3 \t").
           incliti] = TRUE,
           simset (i+i) wt twt [i+i] , bolal w (i+i) jasmie
           inc (i+ i)= FALSE;
           sum set (i+1, wt, bobal-w(i+1);
```

```
int mais (int arg, that & argv[]) {
   int i, j, n, temp, botal=0;
   prints (" in Enter how many numbers: in");
    scanf (" ", d", &n);
    prints ("In Enter of d numbers to set: In", n).
    for (i=0; i<n; i++)
        for (j=0; j<n-1; j++)
           of (wij) > wijti)
               · bemp = w[]]:
                 w[j]= w[j+1].
                w[j+1] = bamp;
     prints ('In The given '1.d numbers in ascending order: In', n)
     for (i=0; i<n; i++)
         printly ("%d", w(I)).
     if ((bokal (sum))
        print ("In Subset construction is not possible");
     else
         for (i=0; i<n; i++)
              Pro [i]=0;
        printly ("In The solution using backbracking is: In");
         subset (-1,0, tobal);
    system ("PAUSE");
    rebour 0;
```

REFERENCES K. Berman, J. Paul, "Algoribhne" (engage learning. · Thomas, H. (Charles Terison: Introduction to Algorithm. CONCLUSION In this tournwork we leavent the concept of Backteracking and also learned how to apply back braking in the sum

TERMWORK 09

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OUTPUT:

Enter how many numbers: 6

Enter 6 numbers to the set:

1 3 5 6 4 2

Input the sum value to create sub set: 7

The given 6 numbers in ascending order:

1 2 3 4 5 6

The solution using backtracking is:

SOLUTION:1 { 1 2 4 }

SOLUTION:2 { 1 6 }

SOLUTION: 3 { 2 5 }

SOLUTION:4 { 3 4 }

Press any key to continue . . .

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TERM WORK - LO

PROBLEM

Implement N queen's problem using Back Fracking.

OBJECTIVE · To implement solution to place or queens on oby or chess board as that no 2 queens attack each other THEORY In this method main aim is to place in queens on ixi chesis board such that no two queens attack each other by being in same now or in same column or on same diagonal. For n=1 there is a trivial solution and for n=2 and n=3 there is no solution. So in they bermwork we consider n=4.

```
ALGORITHM.
 Procedure quen (i, index, n)
   var i < index
     promising (i) then
      if P=n bhen
       write (collid through collin)
       else
          for jel to ndo.
            col [i+1] ej
         queens (itin)
         end
fundion promising (i=index): boolean;
      var Keindex
     begin.
         KCI
         promising & brue;
         while ICC's and promising do
            of collid=collic) or abs (collid-collic)=i then
               promising < false.
           KE K+1
```

```
PROGRAM
#include (sbdio.h)
#include (stellib.h)
# include (mabh.h)
int a (30) , count = 0;
int place (int pas)
      for (i=1; i <pos; i++)
          % ((a[i]== a[pos]) | ((abs (a[i]-a[pos])== abs (i-pos))))
                return 0;
       return 1:
 void printsol (Pat in)
     count + =
     printly ("In In Solution #1d In In", count);
      for (i=1; i<=n; i++) {
           forG=1:j<=n;j++){

{| (a(i)=j)|
                    print ("Q\t");
                     brint (" * XF").
```

```
void queen (int n)
     int K=1;
     alleJ=0;
     while (12)=0)
        alk] =alk]+1
         while (a(K) (= n &8! place (K))
           alk)++;
         g (a(K) (=n)
             il (K==D)
              printsol(n);
             else f
               K++
               a[K]=0:
            K--
I Cham pion
   int n;
   printly ("Enter the number of queens: \n"):
  scanf ("%d " &n):
  queen(n):
 printly (" In Tobal Number of solubions = "1.d", count);
```

REFERENCES · K Berman J Paul ! Algorithm ' lingage learning.
· Thomas H (, Char leism "Introduction to Algorithm" PHL 2nd edition CONCLUSIONS In this termwork we have successfully solved the or queen problem and implemented with code.

TERMWORK 10 N QUEENS

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