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- We introduce a buying strategy for players in a data market including a dataset allocation and price prediction (Section 3).
- We provide a novel data market simulation to be publicly available upon acceptance (Section 4.1) based on a model defined Section 2.
- Using the simulation, we demonstrate a proof-of-concept of our approach, showing its superiority over several baseline methodologies (Section 4.3).

A *data market*  $\mathcal{M}$  is a composition of three sets of entities, namely *buyers*  $\mathcal{B}$ , *sellers*  $\mathcal{S}$ , and *mediator*  $m$ . We shall refer to the first two entities also as the *players*  $\mathcal{A} = \mathcal{B} \cup \mathcal{S}$  in the market. Players may be individuals or groups (e.g., companies or organizations) interested in trading datasets. We use  $\mathcal{D}$  to denote the set of authorized datasets

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Recall that a player  $a$  has a  $\mathcal{D} = \langle D_1, D_2, \dots, D_n \rangle$  (set of relevant datasets),  $L$  (budget), and a valuation value  $v_i$  for each relevant dataset  $D_i$ . At timestamp 0, the set of datasets  $\mathcal{D}$ , its corresponding valuations  $V$  and an initial  $L$  are set. While the valuation values

of datasets stays constant throughout the horizon, the  $\mathcal{D}$  and  $L$  change with respect to the interaction in the market. We denote the datasets set available at time  $t$  and the current budget at time  $t$  as  $\mathcal{D}^t$  and  $L^t$ , respectively. We assume that the only thing that changes the datasets is that a player has purchased a dataset and thus,  $\mathcal{D}^t \subseteq \mathcal{D}^{t'}$ ,  $t' < t$  and  $\mathcal{D}^0 = \mathcal{D}$ . The budget changes with respect to the revenue from each purchased product and its cost.

At time  $t$  the player has to select a subset of datasets  $\bar{\mathcal{D}} \subseteq \mathcal{D}^{t-1}$  that maximizes her future revenues (recall Section 2.1). To simplify the notation, we denote that number of datasets available at time  $t$  as  $m \leq n$ . In what follows, let  $C = \langle c_1, c_2, \dots, c_m \rangle$ ,  $c_i \in \mathbb{R}$  and  $W = \langle w_1, w_2, \dots, w_m \rangle$ ,  $w_m \in \{0, 1\}$  represent a realization of costs and win indicators of the relevant datasets after timestamp  $t$  has completed, respectively.

In practice, a player has to allocate a subset of products at the beginning of time  $t$ . Thus, the player does not know the actual cost of datasets when the dataset allocation takes place. Accordingly, the player has to estimate the costs  $\hat{C} = \langle \hat{c}_1, \hat{c}_2, \dots, \hat{c}_m \rangle$  and wining indicators  $\hat{W} = \langle \hat{w}_1, \hat{w}_2, \dots, \hat{w}_m \rangle$ . Using these estimations, the *datasets allocation* problem can be formalized as follows:

$$\begin{aligned} & \underset{X}{\text{minimize}} && \sum_{i=1}^m (v_i - \hat{c}_i) \cdot \hat{w}_i \cdot X_i \\ & \text{subject to} && \sum_{i=1}^m \hat{c}_i \cdot \hat{w}_i \cdot X_i \leq L \\ & && X_i \in \{0, 1\} \quad i = 1, \dots, m. \end{aligned} \quad (1)$$

Recalling that  $\hat{w}_i \in \{0, 1\}$  the player can actually decrease the size of  $\mathcal{D}^{t-1}$  to the set of product she estimates she would win, i.e.,  $\text{win}(\mathcal{D}^{t-1}) = \{p_i \in \mathcal{D}^{t-1} | \hat{w}_i = 1\}$ . We denote the size of  $\text{win}(\mathcal{D}^{t-1})$  as  $m_{\text{win}}$

$$\begin{aligned} & \underset{X}{\text{minimize}} && \sum_{i=1}^{m_{\text{win}}} (v_i - \hat{c}_i) \cdot X_i \\ & \text{subject to} && \sum_{i=1}^{m_{\text{win}}} \hat{c}_i \cdot X_i \leq L \\ & && X_i \in \{0, 1\} \quad i = 1, \dots, m_{\text{win}}. \end{aligned} \quad (2)$$

Solving Eq. 2 is [RS: hard? can we prove it?](#). In addition, simultaneous auctions combining infinite supply with multiple buyers and sellers is far from easy [elaborate](#). Thus, in this paper we focus on an action-free market as described next.

### 3.2 An Auction-Free Market

we begin with problem definition assuming auctions and than relax the auction such that each buyer is allocated to a seller and in case there is a match (the price the buyer is willing to pay is higher than the price limit set by the seller), the buyer gets the product.

relax the problem, without  $w$

### 3.3 Price Prediction and Bidding

estimating costs

biding strategies

## 4 EVALUATION

bla bla bla

### 4.1 A Data Market Simulation

### 4.2 Empirical Settings

### 4.3 Results

## 5 RELATED WORK

complete

## 6 CONCLUSIONS

complete

## REFERENCES

- [1] Raul Castro Fernandez, Pranav Subramaniam, and Michael J Franklin. 2020. Data Market Platforms: Trading Data Assets to Solve Data Problems [Vision Paper]. (2020).