

# Lecture 2

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# Plan for the day

- ▶ Revise binary Heaps.
- ▶ Binary Search Trees.

# Abstract Data Type - Heap

## Heap

A binary max-heap:

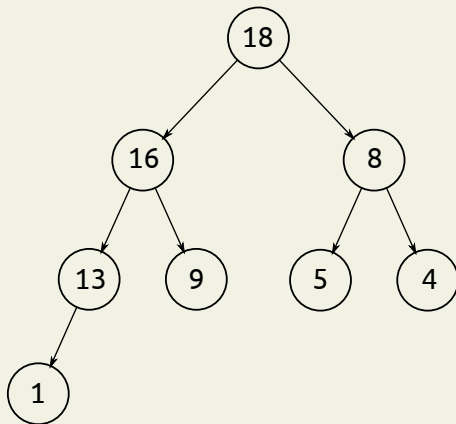
1. Is a complete binary tree except possibly for the lowest level.
2. The value of a node is greater than that of both its children.

A max-heap supports the following functions:

- ▶ `INSERT(val)` – Inserts *val* into the heap.
- ▶ `EXTRACTMAX()` – Returns and removes the maximum element from the heap.

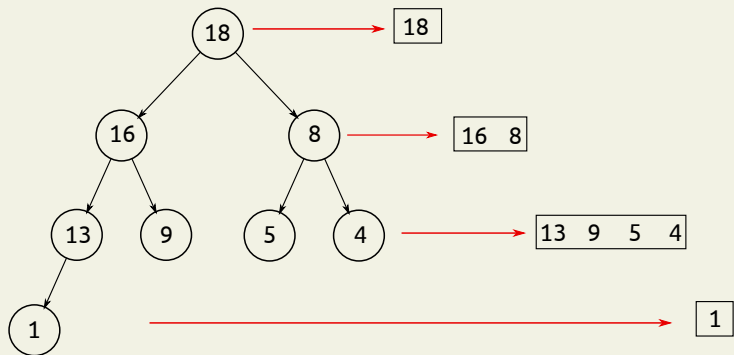
# Data Structure

Heaps are usually implemented using arrays.



# Data Structure

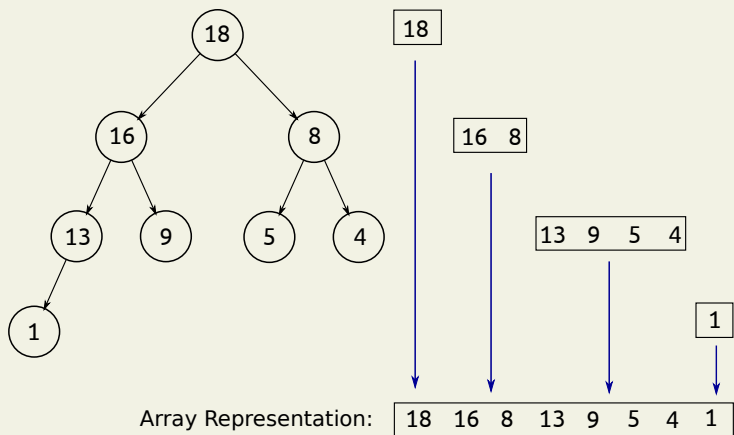
Read off from top to bottom, left to right.



Array Representation:

# Data Structure

Read off from top to bottom, left to right.



# Questions

About Heaps:

1. How many nodes does a height  $h$  heap have? (both bounds)
2. What is the maximum height of a heap with  $n$  nodes?

About the array implementation:

1. What is the array index of the children of the node at  $A[i]$ ?
2. What is the array index of the right sibling of the node at  $A[i]$ ?

# Heaps using arrays

Typically, a heap is built starting with an arbitrary array:

- ▶ Procedure BUILDHEAP(Array  $A$ ) – Takes an array and rearranges the elements to form a heap.

In Object Oriented languages, BUILDHEAP is essentially the *Constructor* of class Heap.

The procedure BUILDHEAP works by using a method called HEAPIFY(*node*).



# Heapify

The  $\text{HEAPIFY}(node)$  procedure:

- ▶ If  $node$  violates the heap property:
  1. Swap value of  $node$  with the largest of its two children.
  2. Call  $\text{HEAPIFY}$  on the child replaced.
- ▶ Else, do nothing and return.

# Heapify

The HEAPIFY(*node*) procedure:

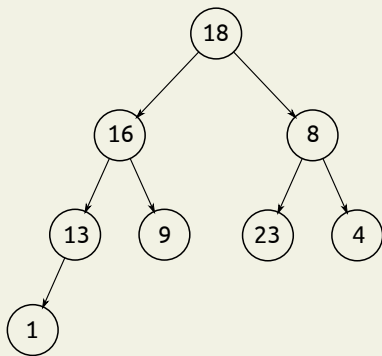
- ▶ If *node* violates the heap property:
  1. Swap value of *node* with the largest of its two children.
  2. Call HEAPIFY on the child replaced.
- ▶ Else, do nothing and return.

Note:

- ▶ The Heapify procedure assumes that both the subtrees under *node* are already heaps.
- ▶ It merely resolves the possible conflict between the value at *node* and its children and recurses.

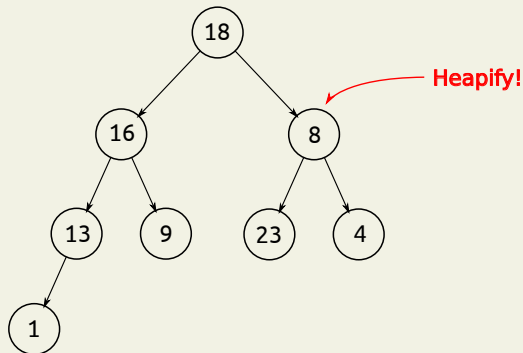
# Heapify

Example:



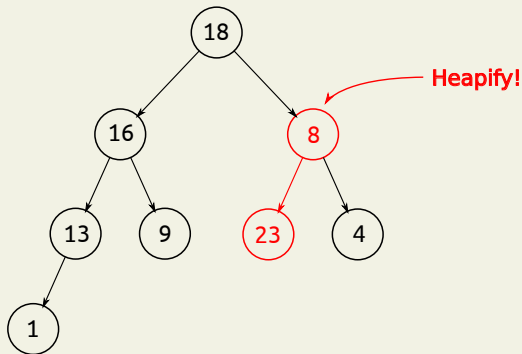
# Heapify

Example:



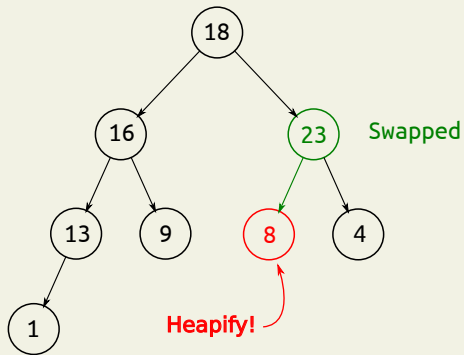
# Heapify

Example:



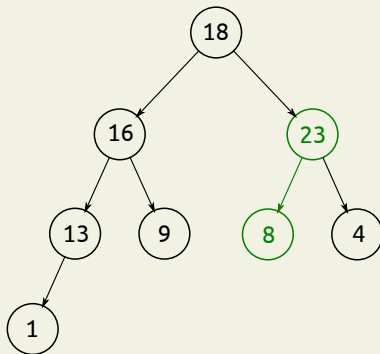
# Heapify

Example:



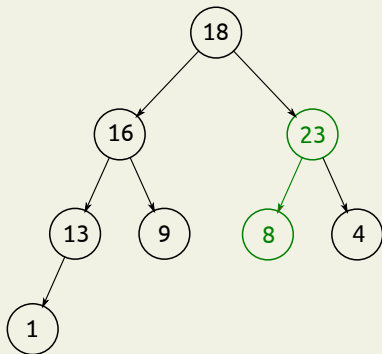
# Heapify

Example:



# Heapify

Note that Heapify only resolves conflicts downwards.





# Building a Heap

Two ways:

- ▶ William's method: Take each element and use INSERT procedure.
- ▶ Floyd's method: Take all elements in an arbitrary array. Heapify repeatedly.

# Building a Heap

The procedure  $\text{BUILDHEAP}(A)$  by Floyd is the following:

- ▶ For  $i$  from  $n$  to 1:
  - ▶  $\text{HEAPIFY}(i)$

# Building a Heap

The procedure  $\text{BUILDHEAP}(A)$  by Floyd is the following:

- ▶ For  $i$  from  $n$  to 1:
  - ▶  $\text{HEAPIFY}(i)$

Note: Indices  $n/2$  to  $n$  form leaves of the heap.

The leaves are already heaps (trivially).

Hence it suffices to run the above loop from  $n/2$  to 1.

# Exercises

Write the following procedures:

- ▶ **INSERT(*val*):**
  - ▶ Insert new value as the last element in the array.
  - ▶ Repeatedly Heapify *upwards* from the new element.
- ▶ **EXTRACTMAX():** Swap positions of root with last leaf. Heapfiy at new root.

## Binary Search Tree

A Binary Search Tree (BST) is a tree that satisfies the following:

- For every node  $X$  in the BST, we have:

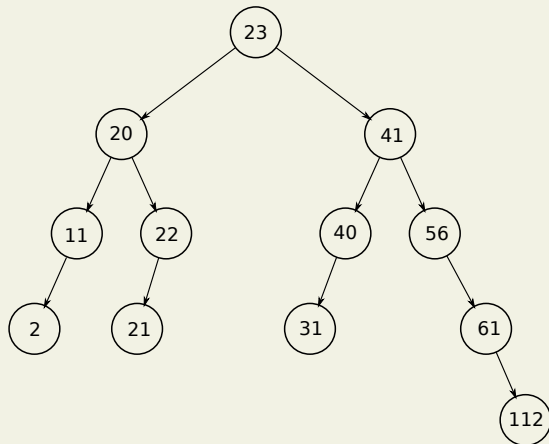
Values in left subtree  $\leq$  value( $X$ )  $\leq$  Value in right subtree

# Abstract Data Type - BST

A BST supports the following functions:

- ▶  $\text{INSERT}(node, val)$  – Inserts  $val$  into the BST rooted at  $node$ .
- ▶  $\text{SEARCH}(node, val)$  – Returns True if  $val$  exists in the BST rooted at  $node$ . False otherwise.
- ▶  $\text{SUCC}(val)$  – Returns the smallest element greater than  $val$  in the BST.
- ▶  $\text{PRED}(val)$  – Returns the largest element lesser than  $val$  in the BST.
- ▶  $\text{DELETE}(val)$  – Deletes  $val$  from the BST.

## Example BST



## Example BST

The order in which elements are inserted makes a difference!  
Consider two different sequences of values:

- ▶ **Sequence A:** 23, 11, 20, 21, 2, 56, 40, 41
- ▶ **Sequence B:** 2, 11, 20, 21, 23, 40, 41

(See whiteboard).



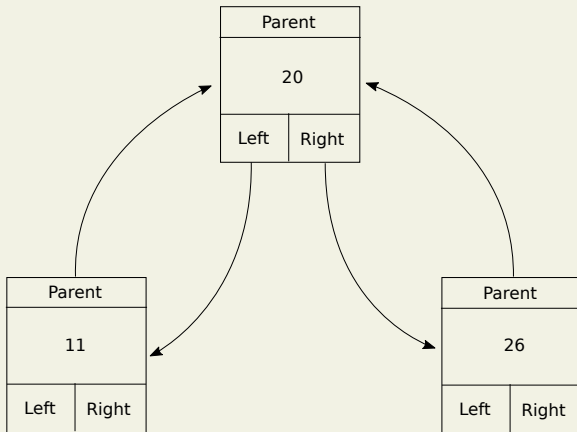
# Data Structure

BSTs are implemented using the Tree datastructure.

Similar to a node in a linked list, each node in the Tree has the following:

- ▶ `int val` – holds the data/value of the node.
- ▶ Left child pointer.
- ▶ Right child pointer.
- ▶ Parent node pointer.

# Data Structure



# Questions

1. What is the maximum height of a BST with  $n$  nodes?
2. What is the minimum height of a BST with  $n$  nodes?

# INSERT procedure

The INSERT(*node*, *x*) procedure:

- ▶ If *node* = NULL, create new node with *x* and attach to parent.
- ▶ Else If  $x < \text{value}(\textit{node})$ ,
  - ▶ INSERT(*node*  $\rightarrow$  *left*, *x*)
- ▶ Else If  $x > \text{value}(\textit{node})$  Then,
  - ▶ INSERT(*node*  $\rightarrow$  *right*, *x*)