(10) Plaintext : 10000001

Key : 11111 00000

P_{10} : 1110101000

P_{10}

3 5 2 7 4 10 1 9 7 6

P_8

6 3 7 4 8 5 10 9

IP_4

2 6 3 1 4 8 5 7

IP^{-1}

4 1 3 5 7 2 8 6

E_P

4 1 2 3 2 3 4 1

P_4 :

2 4 3 1

LS-1 LS-1

11011 10000

P_8 : 10010100

K_1 : 10010100

→ 1101110000

LS-2 LS-2

01111 00010

P_8 : 01010101

K_2 : 01010101

P : 10000001

IP : 00010100

L : 0001 R : 0100

E_P : 00101000

K_1 : 10010100

⊕ 10111100

low = 11 → 3 low = 10 → 2

Col = 01 → 1 Col → 10 → 2

S_0 : 01 S_2 : 01

0101

P_4 : 1100

R : 0001

⊕ 1101

L: 0100 R: 1101

Ep: 11101011

K_2: 01010101

⊕ 10111110

L: 1011 R: 1110

δ = 11 → 3 δ = 10

C: 01 → 1 C = 11

S_0: 01 S_1: 00

→ 0100

P_4: 1000

L: 0100

⊕ 1100

11001101

$\bar{E}p^{-1}$: 01010111

Ciphertext: 01010111

(ii)

(a) Plaintext = 1001

K_1: 01

K_2: 11

$P_1 = E(P, K_1) = E(1001, 01)$

$= 1001$

$E(P_1, K_2) = E(1001, 11)$

$C_1 = 1000$

$D(C, K_1) = (1000, 01) = 1111$

$D(C, K_2) = (1000, 11) = 1001$

$K_0(1001) = 1111$

$K_1(1001) = 1001$

$K_2(1001) = 0001$

$K_3(1001) = 1000$

+ here decryption for $D(C, K_2) = 1001$ and it is matched with $K_1(1001) = 1001$ which is K_1 (plain text)

This is called meet in the middle attack.

b) (i) Time Complexity or possible keys for DES $= 2^{56}$

for 2DES is $2 \times 2^{56} = 2^{57}$

(ii) Time Complexity of brute force attack on double encryption is $2^{56+56} = 2^{112}$.