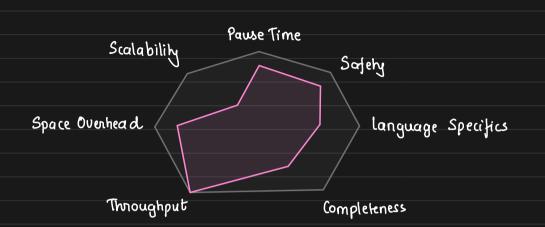


How to pick a Garbage Collector



BY ARPIT BHAYANI

Metrics of a Garbage Collector



There is no best garbage collector

From a Study done on Garbage Collectors in 2000s

There are many garbage
Collectors and each one is
superior than others in atleast

* Most GCs give you a lot of knobs to have the performance for the load

After a certain scale you will be spending

you are handling

More fime tuning these params than coding.

one use-case by atleast 15%

-xx:UseSerial&C -xx: Use Panallel&C

- xx: MaxacPawe Millis

- XX: Min Keap Free Ratio

- xx: GCTime Ratio

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1. Safety collector must never neclaim the storage of live objects.



obj: and obj: are pointing to the same location in memony. GC should NOT be deleting f

2. Throughput The fime spent in garbage collection should be as low as possible.

as possible.

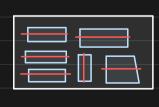
stelaiming the space

User wank "program" + "garbage collection" to execute in as little time

Most GC trades its performance and execution time in favour of higher program exec. throughput

eg. Once a while GC will nun an expensive defrogmentation phase So as to improve on program's memory allocation performance.

3. Complekness eventually, all garbage in the heap should be neclaimed



Brogram Execution

Complete cleanup in one shot is not desirable now always possible

Hence the word Eventual

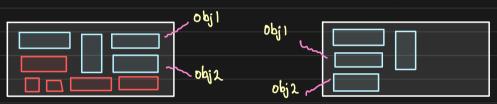
One cycle u/s many cycles

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4. Pause Time Many garbage collectors pause the program execution during cleanup; this pause should be as low as possible

* One of the most important and impactful metric

Most execution of GC do not "Stop the world" but some times to maintain the connectness of references it needs to stop the world



Garbage collectors do their best to reduce the Pause Time but it comes with its own set of challenges & complications

5. Space Overhead GC may require auxiliany data structures to brack objects and decide efficiently

but it puls additional load on memory and consumption

Bitmap Tables → To keep a check on objects abready considered Graphs → To manage & maintain object dependency

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Defragmentation

Object Shuffling

- 6. language Specific Optimizations A GC may provide language specific optimizations to gain that extra ounce of performance
 - eg: some languages are pure functional

 Some languages have only heap allocation

Some languages may have explicit de-allocation

some languages have only persisknt data structures

A GC explois its understanding of the language and its constructs to optimize its execution.

- * Some Crc runs in constant time because of how objects are laid out by the memory manager
- T. Scalability GC needs to leverage the modern hardware capabilities to make its execution faster.

Servers are growing ... los and loos GB of heap and hence GC will have a lot of work to do going through this massive heap

This would increase the GC time and hence GC that always

Stop the world become inefficient of Favoured

Some CCs have evolved & become Pause free

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