

DBMS - 3

1 - Database funda.

2 - ERD + Data integ.

3 - Data Normalisation

- Anomalies
- Functional Dps
- Normal Forms

Transaction

- Intro
- ACID properties



1	John	1	Sherlock
2	Mary	1	Sherlock
3	Kilush	2	Shakti
4	NULL	3	Rebellion
5	Tontia	NULL	NULL

How can I add a student
w/o a batch X

① Cannot create batch w/o student

② Cannot create student w/o batch



- ① missing data (batch w/o student)
- ② FK constraints (student w/o batch)

Deletion anomaly

ID	Name	Batch
1	John	Shivam
-	Max	Shivam



Shobti

None | PSO | Agen org

Deletion anomaly

↳ bad design

Updation Anomaly

→ ID | Name | BID | BNone
 | John | |
 | | | Star lock |



Update Sherlock to Sherlock

Season 6

Failure

- manual
- automated

Inconsistent data

Common issues in BB

- bad design

Insertion

- ↳ missing data
- ↳ FK constraint

Deletion

- ↳ Tightly couple

Update

- ↳ errors or partial failures

Good design

- Normalisation

Functional dependencies

CAR

VIN

Engine Capacity

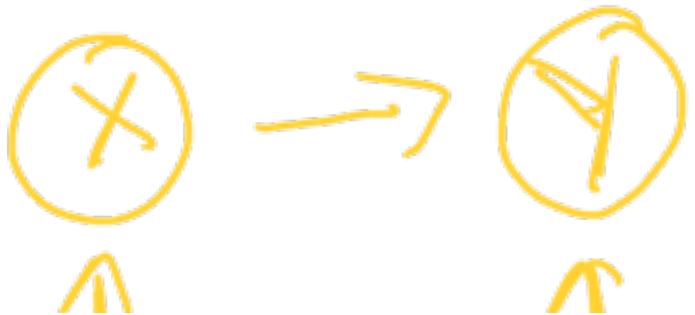
Model

VIN \rightarrow Capacity

VIN \rightarrow Engine Capacity

VIN \rightarrow Model

Engine Capacity \rightarrow VIN



Functional
dependency

ID	NAME	PHONE	BATCH-ID	B-NAME
----	------	-------	----------	--------

→

(ID) → Name

(ID) → Phone

ID → ↗ ↘ ↙

BATCH-ID → BATCH-NAME

NAME → (ID)



Not Unique

Phone \rightarrow ID

ID \rightarrow Phone

Select phone from
stu chirity

ID \rightarrow batch ID

Rating, feedback

Session

Queney

Student ID	Batch ID	Session ID	Rating	Feedback
1	1	1		



PK - Student, Mentor, SessionID

(FD)

1 Student can only have 1 mentor

Student ID \rightarrow Mentor ID

2 Student, Mentor ID, SID

\rightarrow noting
 \rightarrow feedback

Normalization

Multiple tables

↳ reduce data redundancy

↳ remove anomalies

→ reduce redundancy

→ improve data integrity

Normal Forms (Relational)

DB

- | N.F. Strictness

- > N.C. |



6NF

Denormalisation

2NF to 3NF

Set of rules to normalise your database

1NF

2NF - 1NF + new rules

1NF

Rule

→ All your values should be
atomic

ER model

— multivalued

— composite



Mys & 75.7

1NF

→ collection

ID	NAME	PHONE-NOS	
1	John	9876543210, 9876543211	1NF

T	Sherlock	[1234, ...]	Lists
2	Mortany	[9876, ...]	

① Multiple columns

ID	NAME	Phone #1	Phone #2	...	#5
1	Sherlock	1234	9876 ...		

1NF ✓

③

1.m.c =

..... -

① We don't know the upper bound

② Efficient

→ Find a student by phone

→ All phone no.

③ Sparse

X

② Multiple rows

1NF ✓

ID	Name	phone
1	Sherlock	12345...
1	Sherlock	98765...
2	Son	123456...

2 | John | 98765...

① redundant data (n-1)
2 5 x

② PK - ID + Phone

Phone CSV a, b, c

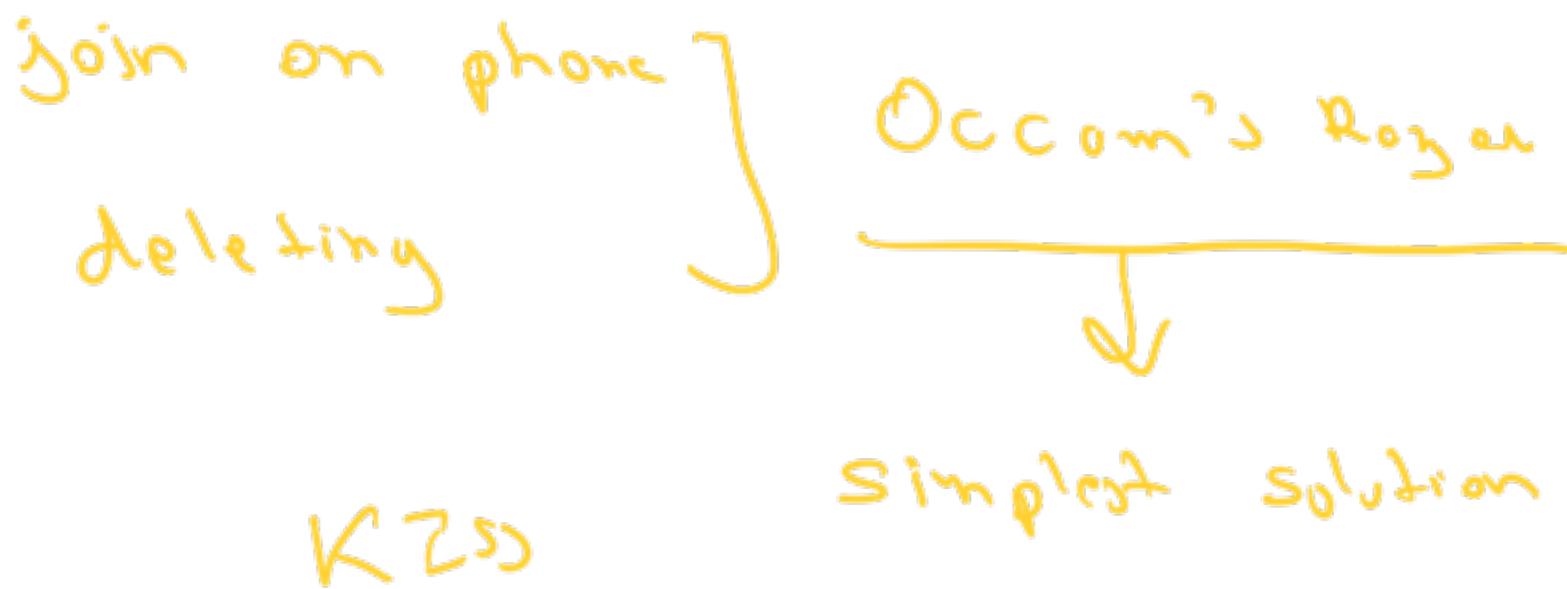
"12345;98765;"

{"123,987,..."} - von ihm



Search user after no = 123

Complicated Solution



Simplest - Create a new table

New Table

STID	Name	Email	Student

Phone

S ID	Phone number
1	12345
1	987654..

Normalize

Normal forms

User

ID	NAME	INTERESTS
1	Tony	EDBMS

[1NF]

→ [atomic values]

→ Multiple columns

— multiple Rows

→ Phone number

ID Phone No

1 123

1 987

→ Create a new table

Students

Phone number

MC

User ID

Name

Interest 1

Interest 2

Separate table

Student ID	Interest
1	DBMS
1	Books
1	Movies

User interest

- └ 1 table (User)
- └ 4 tables
 - performance issue
 - optimisations

6:12

10:42

6:17

10:47

4 sessions

6 sessions

4 sessions

— Theory

5th → Queries

5th → Databases

6th → Transactions

→ Joins

→ Views

→ Triggers

→ Procedures

⋮ ⋮ ⋮

2: 30

Inters. (S3D Int)

→ window
→ Sliding Optimisation



2NF

tanmay. kaicker

SQL Rule

2NF

① 1NF

② No partial dependency

A \rightarrow B

$A \rightarrow X$

ID	NAME	BATCH_ID	b_NAME	psp
----	------	----------	--------	-----

ID → NAME

One user can be part of multiple batches

ID	None	bID	bName	PSP
1	Shenlock	1	Season1	and
1	Shenlock	2	Season2	loop

PK - $ID + bID$

$(ID, bID) \rightarrow PSP$ ✓

Functional dep.

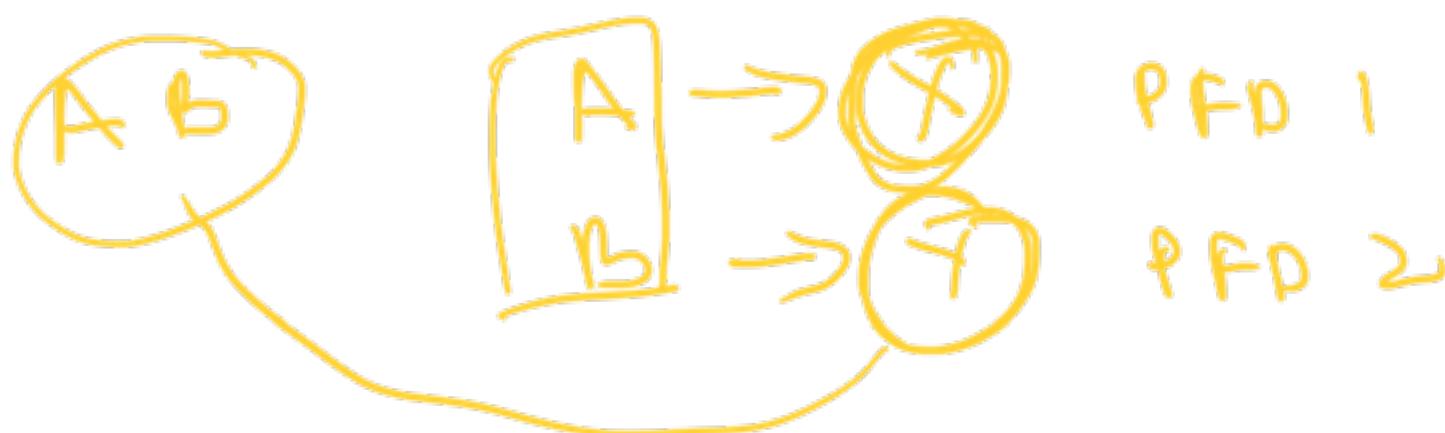
bID \rightarrow batch Name }
job

ID \rightarrow Name

\rightarrow part of a composite
key

parts of CPK on LHS

partial dependency



2NF

→ Not have partial FD

→ Every non-candidate key attribute should depend on the whole candidate key
→ Just not a part of it

2NF - Composite

1NF + single PK = 2NF

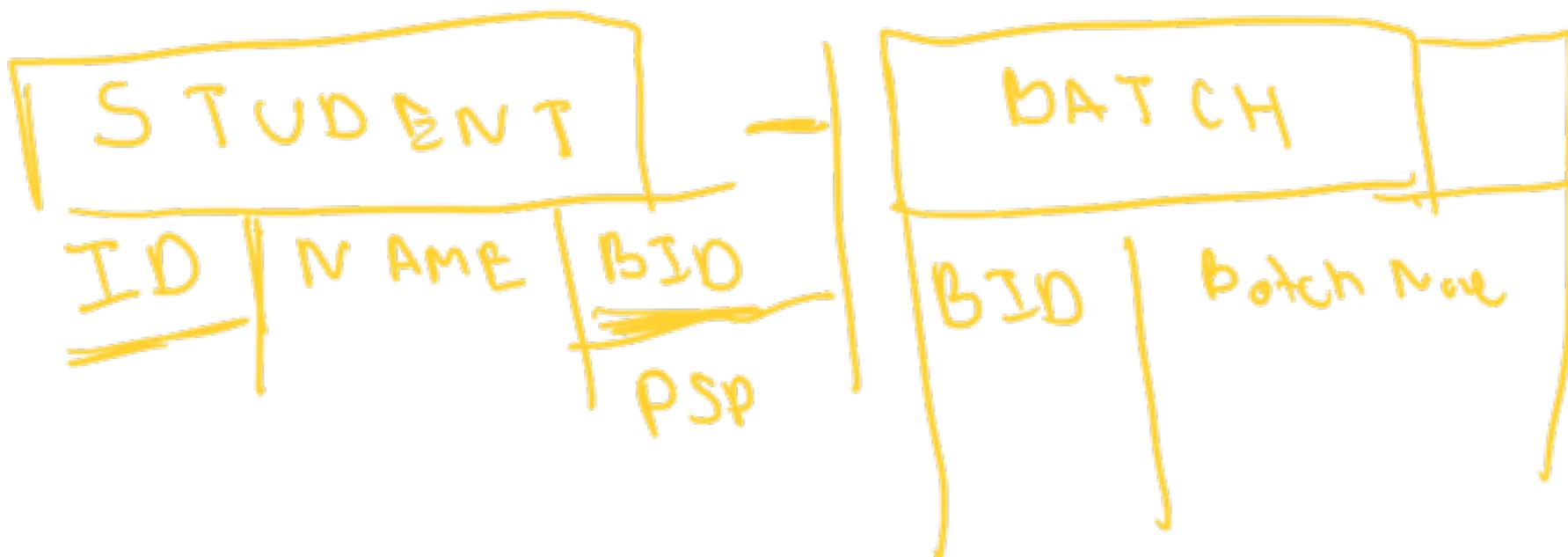
→ composite = PFD

ID	NAME	B3D	BN ANE	P3P
1	Tony	✓	✗	✓

UML Diagram $\vdash 2 \rightarrow 1$

$BID \rightarrow BNAME$ — PFD

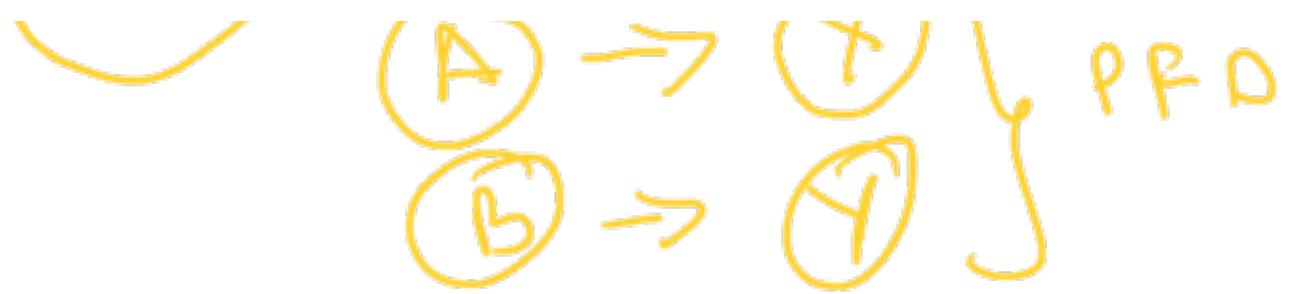
PK \Rightarrow ID, BID



2NF — composite key

Partial Functional Dependency

$(A, B) \rightarrow C$



To \sim to fix)

Create a new table
with duplicate rows

- ① Student can be part of multiple batches

ID BAL

STUDENT	ID	BID	PSD
	1	2	90
	1	1	100

MIN \rightarrow BNAME \rightarrow

ID \rightarrow NAME

\hookrightarrow User info

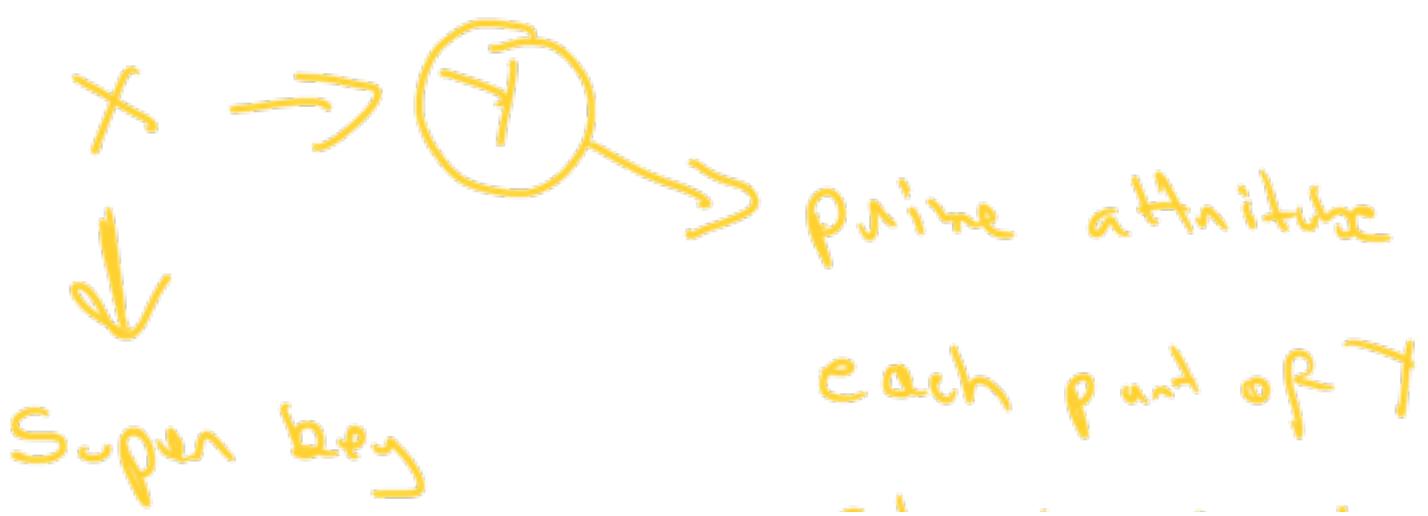
\exists 3NF

- 2NF
- It should not contain transitive dependencies

In sensitivity

$$a = b \quad b = c$$

$$\boxed{a = c}$$



Should be in a
Candidate

ID	NAME	PHONE	BID	BNAME
----	------	-------	-----	-------

ID → NAME ✓
Phone → ID }

BID → BNAME ✗

3NF

1 No transitive depen den circ

$ID \rightarrow B \Sigma D \rightarrow BNAME$

2 $X \rightarrow Y \rightarrow \{A, B\}$
prime attribute

Super key

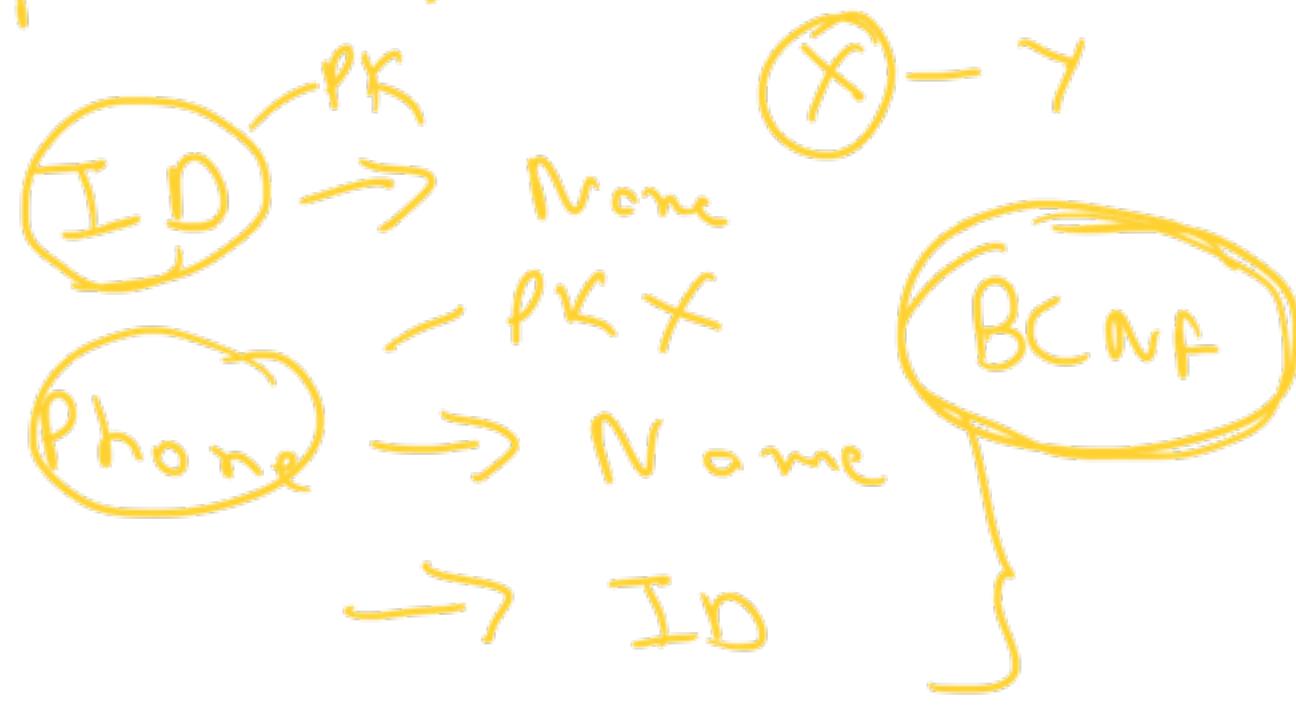
3NF \rightarrow No transitive depen den circ





X is not PK

ID	Phone	Name	Batch ID



Phone = SK, CK, PK

FD

BCNF

3NF

$nNF \Rightarrow (n-1)NF$

$\Rightarrow n^{th}$ rule \hookrightarrow other PK int

ID	NAME	Phone	BID	BNAME
----	------	-------	-----	-------

Student

ID	NAME	BID	STD	Phone No.
1	Tony	1	1	123

① → ③

BID	BNAME
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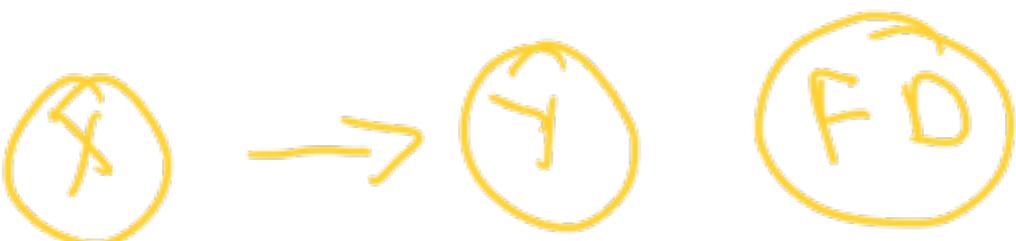
Normalisation

- Check if rules are met
 - If not, create new rule
-

- Data anomalies

- Insertion
 - Deletion
 - Update
- loss of data
errors
inconsistent

- Functional Dependency



→ Partial Popon decay

Composite box \textcircled{A}^B

$$\begin{array}{l} A \rightarrow X \\ B \rightarrow Y \end{array}$$

→ Transitions

$$a = b = c$$

$$a = c$$

$$X \rightarrow Y \quad Y \rightarrow Z$$

$$X \rightarrow Z$$

$\gamma \rightarrow \text{BID} \rightarrow \text{brane}$

1NF - Atomic

2NF - No PFD

3NF - No TFD

BCNF = Only PK on LHS of FD



Create a separate table

- Transactions + ACID
- Indexes

NF - redundancy

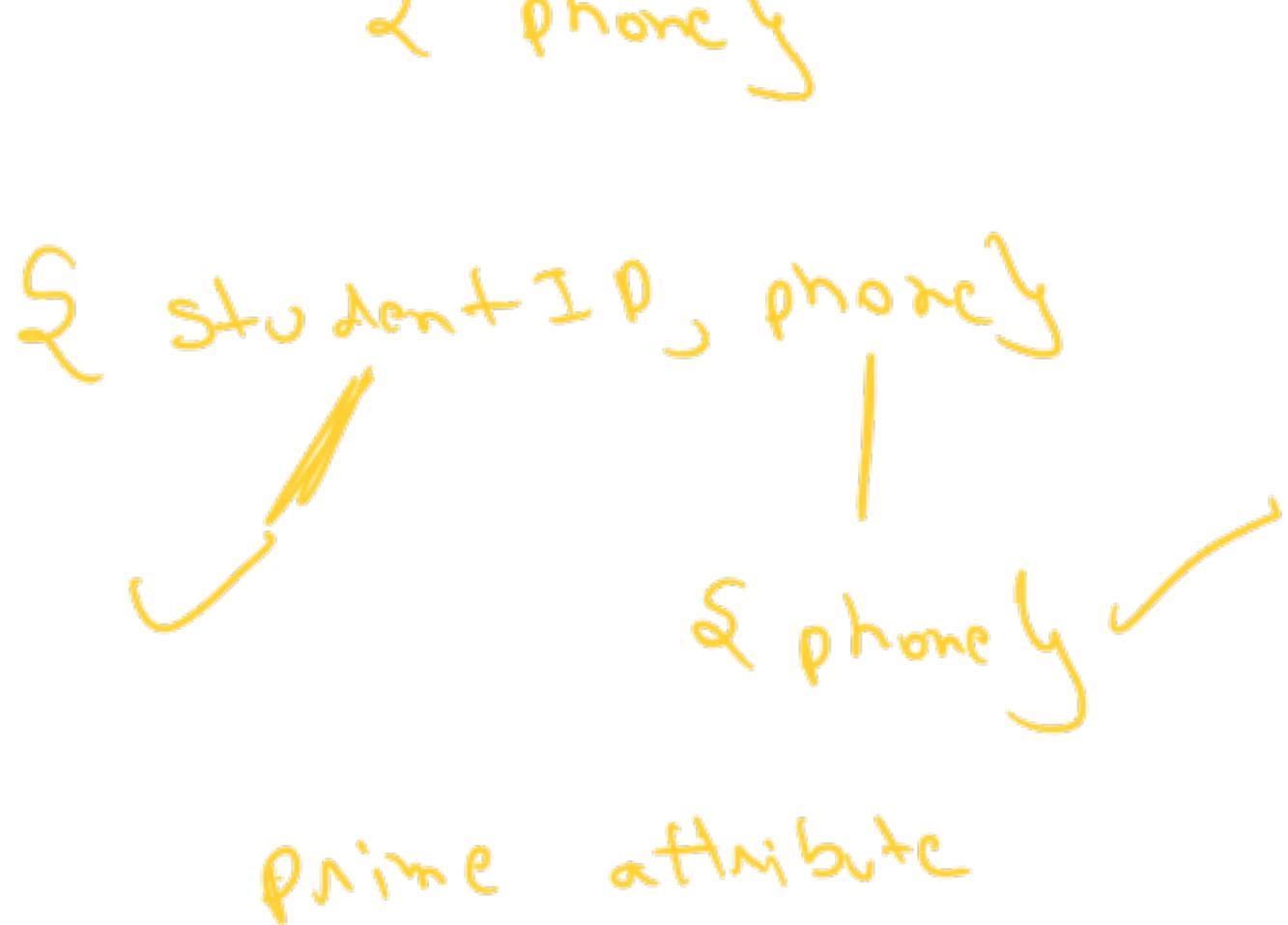
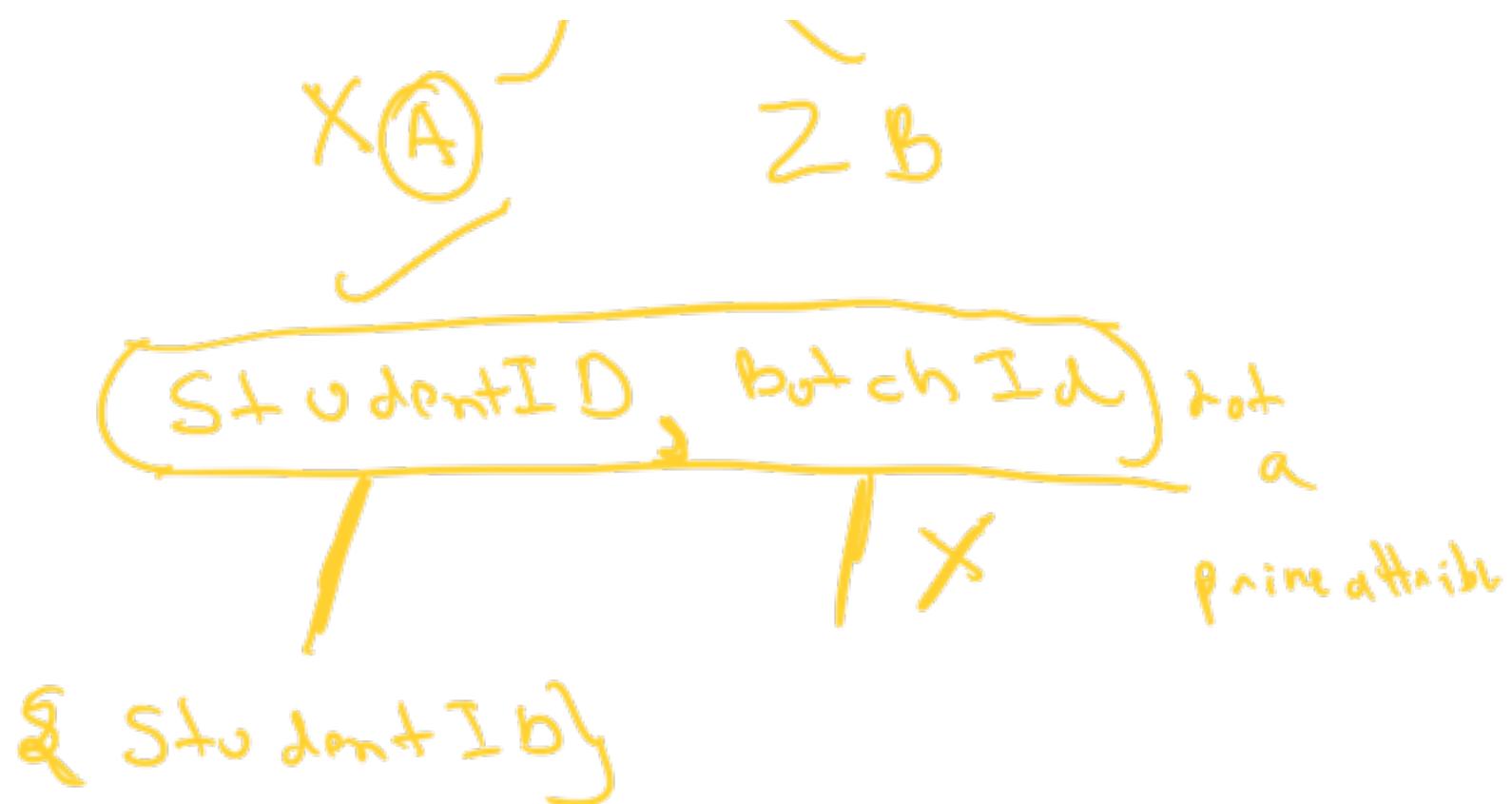
— inconsistency



3NF \rightarrow Transitive FD

- ① X should be a superkey
- ② Y should be a prime attribute

$Y = A B$



3NF $\cancel{X} \rightarrow SK$

$Y \rightarrow \cancel{PA}$

2 Student ID, Phone

Present in a candidate key

2 S ID, Name $\cancel{+ PA}$

S ID \rightarrow PK

$\hookrightarrow CK \cancel{X}$

\rightarrow CK

S ST n PA

CK \rightarrow CK?

CK = Prime Attribute

3NF

$\begin{cases} X = SK \\ Y = PA \end{cases}$

SD - $\begin{cases} PK \\ PA \end{cases}$

PA - $\begin{cases} \text{None} \end{cases}$

CK \subset Y

id, None

CK - $\Sigma_{\{F, H\}} 2x$

$\Sigma_{\{A, B, C\}}$

$\Sigma_{\{A, B\}} \Sigma_{\{A\}}$

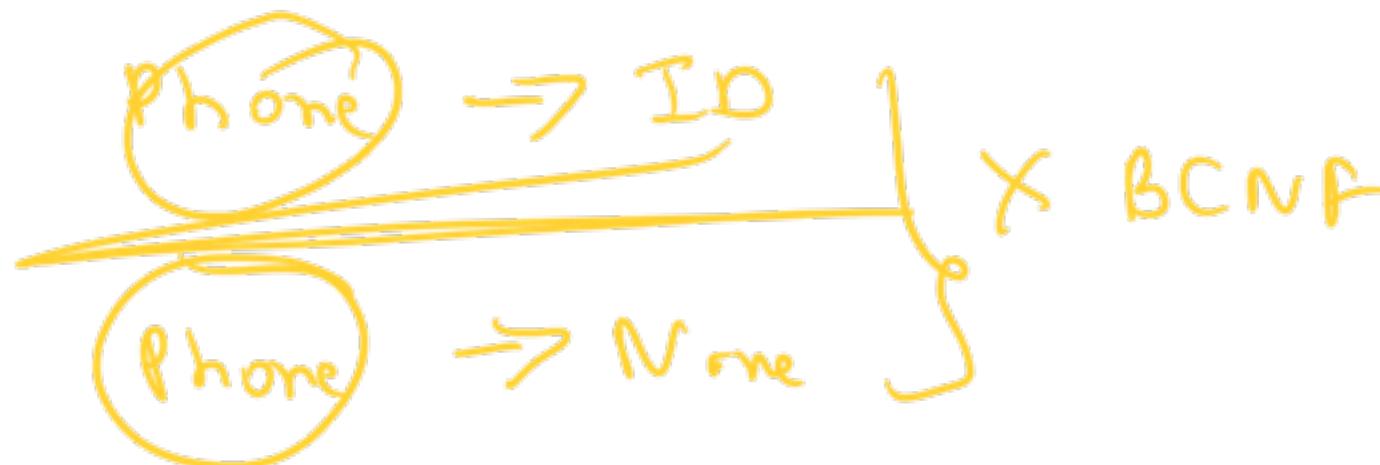
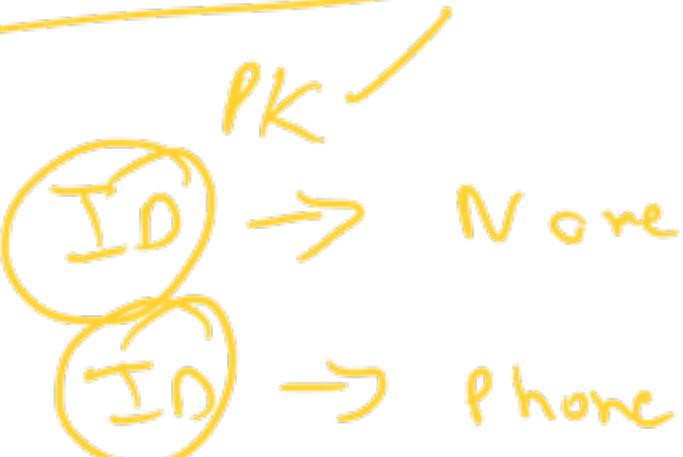
Σ_{NF}

\Rightarrow  TFD

$\Sigma_{CNF} - \Sigma_{NF}$

 \Rightarrow X
 \Rightarrow  PK

ID	NAME	Phone
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ID	NAME		ID	Phone