CS3500 : Operating Systems Lab 3 : System Calls

21th August 2020

In this lab you will add some new system calls to xv6, which will help you understand how they work and will expose you to some of the internals of the xv6 kernel. You will add more system calls in later labs.

The goals of this assignment are:

- Understand the system call interface
- Understand how user programs send parameters to the kernel, and receive values back

Resources:

Before you start coding, read Chapter 2 of the xv6 book, and Sections 4.3 and 4.4 of Chapter 4, and related source files:

- 1. user-space code for systems calls is in user/user.h and user/usys.pl
- 2. kernel-space code is kernel/syscall.h, kernel/syscall.c
- 3. The process-related code is kernel/proc.h and kernel/proc.c

Problem 1: 10 points

In this problem you will create a new system call: $echo_simple$, which should receive one argument, a string, and print it to stdout. As part of this problem, you should write a user program $test_problem_1$ that invokes the $echo_simple$ system call.

```
$ test_problem_1 Hello
Hello
```

Your solution is correct if your program behaves as shown above (though the inputs may be different).

Some hints:

- Add $U/\text{_test_problem_1}$ to UPROGS in Makefile
- Run make qemu and you will see that the compiler cannot compile user/test_problem_1.c, because the user-space stubs for the system call don't exist yet: add a prototype for the system call to user/user.h, a stub to user/usys.pl, and a syscall number to kernel/syscall.h. The Makefile invokes the perl script user/usys.pl, which produces user/usys.S, the actual system call stubs, which use the RISC-V ecall instruction to transition to the kernel. Once you fix the compilation issues, run test_problem_1; it will fail because you haven't implemented the system call in the kernel yet.

• Add a $sys_echo_simple()$ function in kernel/sysproc.c that implements the new system call. The functions to retrieve system call arguments from user space are in kernel/syscall.c, and you can see examples of their use in kernel/sysproc.c.

Problem 2: 20 points

echo is a built-in user command that writes its arguments to standard output. In this problem you will create a new system call: echo_kernel, that is similar to echo user command in its functionality, but executes in Kernel space rather than user space. As part of this problem, you should write a user program test_problem_2 that invokes the echo_kernel system call.

```
$ test_problem_2 Hello World is passed
Hello World is passed
```

Your solution is correct if your program behaves as shown above (though the inputs may be different).

Some hints:

- Add $U/\text{_test_problem_2}$ to UPROGS in Makefile
- Run make qemu and you will see that the compiler cannot compile user/test_problem_2.c. Add the system call echo_kernel following the same steps as in previous problem. Once you fix the compilation issues, run test_problem_2; it will fail because you haven't implemented the system call in the kernel yet.
- Add a $sys_echo_kernel()$ function in kernel/sysproc.c that implements the new system call. It should have the same functionality as echo() function. (see user/echo.c)

Problem 3: 20 points

In this problem you will create a new *get_process_info* system call that returns the Process ID, Process Name and the Size of process memory. As part of this problem, you should write a user program *test_problem_3* that invokes the *get_process_info* system call. The system call takes one argument: a pointer to a *struct processinfo* (see *processinfo.h*). The kernel should fill out the fields of this struct.

```
$ test_problem_3
Process ID -> 23
Process Name -> test_problem_3
Memory Size -> 2405 Bytes
```

Your solution is correct if your program behaves as shown above (though the output may be different).

Some hints:

- Add \$U/_test_problem_3 to UPROGS in Makefile
- Run make qemu; user/test_problem_3.c will fail to compile. Add the system call get_process_info, following the same steps as in the previous problem. To declare the prototype for get_process_info() in user/user.h you need predeclare the existence of struct processinfo:

```
struct processinfo;
int get_process_info(struct processinfo *);
```

Once you fix the compilation issues, run *test_problem_3*; it will fail because you haven't implemented the system call in the kernel yet.

- get_process_info needs to copy a struct processinfo back to user space; see sys_fstat() (ker-nel/sysfile.c) and filestat() (kernel/file.c) for examples of how to do that using copyout().
- To fill the processinfo structure, make use of the proc structure (see kernel/proc.h).

Problem 4: 20 points

In this problem you will add a system call tracing feature that may help you when debugging later labs. You'll create a new system call: trace, that will control tracing. It should take one argument, an integer "mask", whose bits specify which system calls to trace. For example, to trace the read system call, a program calls trace ($1 << SYS_read$), where SYS_read is a syscall number from kernel/syscall.h. You have to modify the xv6 kernel to print out a line when each system call is about to return, if the system call's number is set in the mask. The line should contain the process id, the name of the system call and the return value; you don't need to print the system call arguments. The trace system call should enable tracing for the process that calls it, but should not affect other processes.

We provide a trace user-level program that runs another program with tracing enabled (see user/trace.c). When you're done, you should see output like this:

```
$ trace 32 grep hello README
3: syscall read -> 1023
3: syscall read -> 959
3: syscall read -> 0
$
$ trace 2147483647 grep hello README
4: syscall trace -> 0
4: syscall exec -> 3
4: syscall open -> 3
4: syscall read -> 1023
4: syscall read -> 959
4: syscall read -> 0
4: syscall close -> 0
$
$ grep hello README
$
```

In the first example above, *trace* invokes *grep* tracing just the *read* system call. The 32 is 1<*SYS_read*. In the second example, *trace* runs *grep* while tracing all system calls; the 2147583647 has all 31 low bits set. In the third example, the program isn't traced, so no trace output is printed. Your solution is correct if your program behaves as shown above (though the process IDs may be different).

Some hints:

- Add \$U/_trace to UPROGS in Makefile
- Run make qemu and you will see that the compiler cannot compile user/trace.c. Add the system call trace following the same steps as in the previous problem. Once you fix the compilation issues, run trace 32 grep hello README; it will fail because you haven't implemented the system call in the kernel yet.

- Add a $sys_trace()$ function in kernel/sysproc.c that implements the new system call by remembering its argument in a new variable in the proc structure (see kernel/proc.h).
- Modify the syscall() function in kernel/syscall.c to print the trace output.

Challenge Problem 1: 30 points (Optional)

Extend the trace system call by printing the system call arguments for the traced system calls.

Lab Workflow:

- 1. Copy the *trace.c* file (provided as part of the question) to *user*/ directory and *processinfo.h* file to *kernel*/ directory
- 2. Solve the problems
- 3. Answer the Lab_3_Questionnaire (provided as part of the question).
- 4. Once your ready to submit, run *make clean* and zip the answered Lab_3_Questionnaire along with the entire *xv6-riscv* repo, containing your solutions and upload it on Moodle. Follow the naming convention as *<ROLLNO>_<LABNO>.<EXTENSION>*.

\$ tar -cvzf XXX-XXX.tar.gz xv6-riscv/* Lab_3_Questionnaire.docx

Note: You need not submit any Writeup. Ensure that you have zipped the filled Questionnaire along with the solution repo.