§1

An example of backtracking. Given a list of m-letter words and another list of n-letter words, we find all $m \times n$ matrices whose rows and columns are all listed. This program was written a week or so after I wrote BACK-MXN-WORDS, because I realized that I ought to try a scheme that fills in the cells of the matrix one by one. (That program fills entire columns at each level.)

I'm thinking m=5 and n=6 as an interesting case to try in TAOCP, but of course the problem makes sense in general.

The word list files are named on the command line. You can also restrict the list length to, say, at most 500 words, by appending ':500' to the file name.

```
#define maxm = 7
                        /* largest permissible value of m */
#define maxn = 10
                         /* largest permissible value of n */
#define maxmwds 30000
                                 /* largest permissible number of m-letter words */
#define maxtriesize
                       1000000
                                    /* largest permissible number of prefixes */
#define o mems++
#define oo mems += 2
#define bufsize maxm + maxn
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
  unsigned long long mems;
                                     /* memory references */
  unsigned long long thresh = 100000000000;
                                                     /* reporting time */
  typedef struct {
               /* a character (the other seven bytes are zero) */
    int c;
    int link;
                  /* a pointer */
               /* a "tri element" */
  } trielt:
  int maxmm = maxmwds, maxnn = maxtriesize;
  char mword[maxmwds][maxm + 1];
  int trie[maxtriesize][27];
                  /* this many nodes in the trie of n-letter words */
  int trieptr;
  trielt tri[maxtriesize];
  int triptr;
                 /* this many elements in the tri of m-letter words */
  char buf[bufsize];
  unsigned int count;
                             /* this many solutions found */
  FILE *mfile, *nfile;
  int a[maxn + 1][maxn + 1];
  int x[maxm * maxn + 1];
  long long profile [maxm * maxn + 2];
  \langle \text{Subroutines 8} \rangle;
  main(\mathbf{int} \ argc, \mathbf{char} * argv[])
    register int i, j, k, l, m, n, p, q, mm, nn, xl, li, lj, lnk, chr;
    register char *w;
    \langle \text{Process the command line } 3 \rangle;
     \langle Input the m-words and make the tri 4\rangle:
    \langle \text{Input the } n\text{-words and make the trie 5} \rangle;
    fprintf(stderr, "(\%1lu_mems_ito_iinitialize_ithe_idata_istructures) \n", mems);
    (Backtrack thru all solutions 9);
    fprintf(stderr, "Altogether_\", "u_\solutions_\", (%llu\", mems). \", count, mems);
    \langle \text{ Print the profile 2} \rangle;
```

```
2. \langle \text{Print the profile 2} \rangle \equiv
  fprintf(stderr, "Profile:____1\n");
  for (k = 2; k \le m * n + 1; k \leftrightarrow) fprintf (stderr, "\%1911d\n", profile[k]);
This code is used in section 1.
3. \langle \text{Process the command line 3} \rangle \equiv
  if (argc \neq 3) {
     exit(-1);
  w = strchr(arqv[1], ':');
  if (w) {
                 /* colon in filename */
     if (sscanf(w+1, "%d", \& maxmm) \neq 1) {
        fprintf(stderr, "I_{\sqcup}can't_{\sqcup}parse_{\sqcup}the_{\sqcup}m-file_{\sqcup}spec_{\sqcup}'%s'!\n", argv[1]);
        exit(-20);
     *w = 0;
  if (\neg(mfile = fopen(argv[1], "r")))  {
     fprintf(stderr, "I_{\sqcup}can't_{\sqcup}open_{\sqcup}file_{\sqcup}'%s'_{\sqcup}for_{\sqcup}reading_{\sqcup}m-words!\n", argv[1]);
     exit(-2);
  w = strchr(argv[2], ':');
  if (w) { /* colon in filename */
     if (sscanf(w+1, "%d", \& maxnn) \neq 1) {
        fprintf(stderr, "I_{\sqcup}can't_{\sqcup}parse_{\sqcup}the_{\sqcup}n-file_{\sqcup}spec_{\sqcup}'%s'!\n", argv[1]);
        exit(-22);
     }
     *w = 0;
  if (\neg(nfile = fopen(argv[2], "r"))) {
     fprintf(stderr, "I_{\sqcup}can't_{\sqcup}open_{\sqcup}file_{\sqcup}'%s'_{\sqcup}for_{\sqcup}reading_{\sqcup}n-words!\n", argv[2]);
     exit(-3);
This code is used in section 1.
```

```
4. (Input the m-words and make the tri 4) \equiv
  m = mm = 0;
  while (1) {
     if (mm \equiv maxmm) break;
    if (\neg fgets(buf, bufsize, mfile)) break;
     for (k = 0; o, buf[k] \ge \texttt{'a'} \land buf[k] \le \texttt{'z'}; k++) o, mword[mm][k] = buf[k];
     if (buf[k] \neq '\n') {
       fprintf(stderr, "Illegal_m-word:_u%s", buf);
       exit(-10);
     if (m \equiv 0) {
       m=k;
       if (m > maxm) {
          fprintf(stderr, "Sorry, \_m\_should\_be\_at\_most\_%d! \n", maxm);
          exit(-16);
     } else if (k \neq m) {
       fprintf(stderr, "The_m-file_has_words_of_lengths_%d_and_%d!\n", m, k);
       exit(-4);
     mm ++;
  \langle \text{Build the tri } 7 \rangle;
  fprintf(stderr, "OK, LI've\_successfully\_read\_%d\_words\_of\_length\_m=%d.\n", mm, m);
  fprintf(stderr, "(The_{\sqcup}tri_{\sqcup}has_{\sqcup}%d_{\sqcup}elements.)\n", triptr);
This code is used in section 1.
```

5. For simplicity, I make a sparse trie with 27 branches at every node. An *n*-letter word $w_1
ldots w_n$ leads to entries $trie[p_{k-1}][[w_k] = p_k$ for $1 \le k \le n$, where $p_0 = 0$ and $p_k > 0$. Here $1 \le w_k \le 26$; I am reserving slot 0 for later enhancements.

Mems are counted as if trie[x][y] is array[27 * x + y]. (I mean, 'trie[x]' is not a pointer that must be fetched, it's a pointer that the program can compute without fetching.)

```
#define trunc(c) ((c) & #1f)
                                         /* convert 'a' to 1, ..., 'z' to 26 */
\langle \text{Input the } n\text{-words and make the trie 5} \rangle \equiv
  n = nn = 0, trieptr = 1;
  while (1)
     if (nn \equiv maxnn) break;
     if (\neg fgets(buf, bufsize, nfile)) break;
     for (k = p = 0; o, buf[k] \ge 'a' \land buf[k] \le 'z'; k++, p = q) {
       o, q = trie[p][trunc(buf[k])];
       if (q \equiv 0) break;
     for (j = k; o, buf[j] \ge \texttt{'a'} \land buf[j] \le \texttt{'z'}; j \leftrightarrow ) {
       if (j < n - 1 \lor n \equiv 0) {
          if (trieptr \equiv maxtriesize) {
             fprintf(stderr, "Overflow, (maxtriesize=%d)!\n", maxtriesize);
             exit(-66);
          o, trie[p][trunc(buf[j])] = trieptr;
          p = trieptr ++;
        }
     if (buf[j] \neq `\n') {
       fprintf(stderr, "Illegal_n-word: _%s", buf);
        exit(-11);
     \langle Check the length of the new line _{6}\rangle;
     o, trie[p][trunc(buf[n-1])] = nn + 1;
                                                     /* remember index of the word */
                        /* we knew trie[p] when p = 0 and when q = 0; buf[j] when j = k */
     mems = 3;
     nn ++;
  fprintf(stderr, "Plus_{\square}%d_{\square}words_{\square}of_{\square}length_{\square}n=%d.\n", nn, n);
  fprintf(stderr, "(The_itrie_ihas_i), d_inodes.) \n", trieptr);
This code is used in section 1.
```

```
6. \langle Check the length of the new line _{6}\rangle \equiv
  if (n \equiv 0) {
    n = j;
    if (n > maxn) {
       fprintf(stderr, "Sorry, \_n\_should\_be\_at\_most\_%d! \n", maxn);
                          /* we allocated an unnecessary node, since n wasn't known */
     p--, trieptr--;
  } else {
     if (n \neq j) {
       fprintf(stderr, "The_n-file_has_words_of_lengths_%d_and_%d!\n", n, j);
       exit(-5);
     if (k \equiv n) {
       buf[j] = 0;
       fprintf(stderr, "The \ n-file \ has \ the \ duplicate \ word \ '%s'! \ h", buf);
       exit(-6);
  }
```

This code is used in section 5.

7. In this program I build what's called here (only) a "tri," by which I mean a trie that has been compressed into the following simple format: When node k has c children, the **trielt** entries tri[k], tri[k+1], ..., tri[k+c-1] will contain the next character and a pointer to the relevant child node. The following entry, tri[k+c], will be zero. (More precisely, its link part will be zero.)

It's easiest to build a normal trie first, and to compress it afterwards. So the following code—which is actually performed *before* the *n*-letter words are input—uses the *trie* array to do this.

```
 \begin{cases} \text{Build the tri 7} \rangle \equiv \\ \text{for } (i=0, trieptr=1; \ i < mm; \ i++) \ \{ \\ \text{for } (o, k=p=0; \ k < m; \ k++, p=q) \ \{ \\ o, q=trie[p][trunc(mword[i][k])]; \\ \text{if } (q\equiv 0) \ \text{break}; \\ \} \\ \text{for } (j=k; \ j < m; \ j++) \ \{ \\ \text{if } (j < m-1) \ \{ \\ \text{if } (trieptr\equiv maxtriesize) \ \{ \\ fprintf(stderr, "Overflow_{\sqcup}(maxtriesize=\%d)! \ ", maxtriesize); \\ exit(-67); \\ \} \\ o, trie[p][trunc(mword[i][j])] = trieptr; \\ p=trieptr++; \\ \} \\ \} \\ o, trie[p][trunc(mword[i][m-1])] = i+1; \ /* \ \text{remember the word } */ \\ \} \\ compress(0,m);
```

This code is used in section 4.

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8. With a small change (namely to cache the value of a compressed node) this program would actually compress a trie that has overlapping subtries. But our trie doesn't have that property, so I don't worry about it here.

```
\langle \text{Subroutines 8} \rangle \equiv
  int compress(int p, int l)
     \mathbf{register} \ \mathbf{int} \ a, c, k;
                 /* subroutine call overhead */
     00,00;
     if (l \equiv 0) return p;
                               /* prefix has its maximum length */
     for (c = 0, k = 1; k < 27; k++)
       if (o, trie[p][k]) c \leftrightarrow;
     a = triptr;
     triptr += c + 1;
     if (triptr \ge maxtriesize) {
       fprintf(stderr, "Tri_overflow_(maxtriesize=%d)!\n", maxtriesize);
        exit(-67);
     for (c = 0, k = 1; k < 27; k++)
       if (o, trie[p][k]) {
          o, tri[a+c].link = compress(trie[p][k], l-1);
          tri[a+c++].c=k;
          o, trie[p][k] = 0; /* clear the trie for the next user */
     \mathbf{return}\ a;
This code is used in section 1.
```

```
9. Here I follow Algorithm 7.2.2B.
\langle Backtrack thru all solutions 9\rangle \equiv
b1: l = 1;
  for (k = 0; k < m; k++) o, a[k][0] = 0;
                                                 /* root of trie at left of each row */
b2: profile[l]++;
  \langle \text{Report the current state, if } mems > thresh 11 \rangle;
  if (l > m * n) (Print a solution and goto b5 10);
  li = (l-1) \% m;
  lj = ((l-1)/m) + 1;
                             /* at this level we work on row li and column lj */
  if (li) xl = lnk;
  else xl = 0;
                    /* root of tri at top of each column */
  o, lnk = tri[xl].link;
b3: chr = tri[xl].c;
                         /* no mem cost, this word has already been fetched */
  oo, q = trie[a[li][lj-1]][chr];
  if (\neg q) goto b4;
  o, x[l] = xl;
  o, a[li][lj] = q;
  l++;
  goto b2;
b4: o, lnk = tri[++xl].link;
  if (lnk) goto b3;
b5: l--;
  if (l) {
     o, xl = x[l];
     li = (l-1) \% m;
     lj = ((l-1)/m) + 1;
     goto b4;
This code is used in section 1.
10. (Print a solution and goto b5 \ 10) \equiv
     count ++; printf("%d:", count);
     for (k = 1; k \le n; k++) printf("\ ", mword[tri[x[m*k]].link-1]);
     for (p = 0, k = 1; k \le n; k++)
       if (tri[x[m*k]].link > p) p = tri[x[m*k]].link;
     for (q = 0, j = 1; j \le m; j++)
       if (a[j-1][n] > q) q = a[j-1][n];
     printf(" (\%06d,\%06d; sum_\%07d, prod_\%012d) \n", p, q, p + q, p * q);
     goto b5;
This code is used in section 9.
11. \langle Report the current state, if mems \geq thresh |11\rangle \equiv
  if (mems \ge thresh) {
     thresh += 100000000000;
     fprintf(stderr, "After_\%lld_\mems:", mems);
     for (k = 2; k < l; k \leftrightarrow) fprintf (stderr, "_{\perp}) flid", profile [k]);
     fprintf(stderr, "\n");
This code is used in section 9.
```

12. Index.

```
a: 1, 8.
argc: \underline{1}, 3.
argv: \underline{1}, 3.
array: 5.
buf: \underline{1}, 4, 5, 6.
bufsize: \underline{1}, 4, 5.
b1: <u>9</u>.
b2: \underline{9}.
b3: \underline{9}.
b4: <u>9</u>.
b5: 9, 10.
c: \underline{1}, \underline{8}.
chr: \underline{1}, \underline{9}.
compress: 7, \underline{8}.
count: \underline{1}, \underline{10}.
exit: 3, 4, 5, 6, 7, 8.
fgets: 4, 5.
fopen: 3.
fprintf: 1, 2, 3, 4, 5, 6, 7, 8, 11.
i: \underline{1}.
j: \underline{1}.
k: \underline{1}, \underline{8}.
l: \underline{1}, \underline{8}.
li: \underline{1}, \underline{9}.
link: 1, 7, 8, 9, 10.
lj: \underline{1}, \underline{9}.
lnk: \underline{1}, \underline{9}.
m: 1.
main: \underline{1}.
maxm: \underline{1}, \underline{4}.
maxmm: 1, 3, 4.
maxmwds: \underline{1}.
maxn: \underline{1}, \underline{6}.
maxnn: \underline{1}, \underline{3}, \underline{5}.
maxtriesize: 1, 5, 7, 8.
mems: \underline{1}, 5, 11.
mfile: \underline{1}, 3, 4.
mm: \underline{1}, 4, 7.
mword: \underline{1}, 4, 7, 10.
n: \underline{1}.
nfile: \underline{1}, 3, 5.
nn: \underline{1}, \underline{5}.
o: \underline{1}.
oo: 1, 8, 9.
p: \underline{1}, \underline{8}.
printf: 10.
profile: \underline{1}, \underline{2}, \underline{9}, \underline{11}.
q: \underline{\mathbf{1}}.
sscanf: 3.
stderr: 1, 2, 3, 4, 5, 6, 7, 8, 11.
strchr: 3.
```

 $\begin{array}{llll} \textit{thresh}\colon \ \underline{1}, \ 11. \\ \textit{tri}\colon \ \underline{1}, \ 7, \ 8, \ 9, \ 10. \\ \textit{trie}\colon \ \underline{1}, \ 5, \ 7, \ 8, \ 9. \\ \textbf{trielt}\colon \ \underline{1}, \ 7. \\ \textit{trieptr}\colon \ \underline{1}, \ 5, \ 6, \ 7. \\ \textit{triptr}\colon \ \underline{1}, \ 4, \ 8. \\ \textit{trunc}\colon \ \underline{5}, \ 7. \\ w\colon \ \underline{1}. \\ x\colon \ \underline{1}. \\ xl\colon \ \underline{1}, \ 9. \end{array}$

BACK-MXN-WORDS-MXN NAMES OF THE SECTIONS 9

```
\langle Backtrack thru all solutions 9\rangle Used in section 1.

\langle Build the tri 7\rangle Used in section 4.

\langle Check the length of the new line 6\rangle Used in section 5.

\langle Input the m-words and make the tri 4\rangle Used in section 1.

\langle Input the n-words and make the trie 5\rangle Used in section 1.

\langle Print a solution and goto b5 10\rangle Used in section 9.

\langle Print the profile 2\rangle Used in section 1.

\langle Process the command line 3\rangle Used in section 1.

\langle Report the current state, if mems \geq thresh 11\rangle Used in section 9.

\langle Subroutines 8\rangle Used in section 1.
```

BACK-MXN-WORDS-MXN

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