(See https://cs.stanford.edu/~knuth/programs.html for date.)

1. Primitive sorting networks at random. This program is a quick-and-dirty implementation of the random process studied in exercise 5.3.4–40: Start with the permutation $n \dots 21$ and randomly interchange adjacent elements that are out of order, until reaching $12 \dots n$. I want to know if the upper bound of $4n^2$ steps, proved in that exercise, is optimum.

This Monte Carlo program computes a number c such that c(n-1) random adjacent comparators would have sufficed to complete the sorting. This number is the sum of $1/t_k$ during the $\binom{n}{2}$ steps of sorting, where t is the number of adjacent out-of-order pairs before the kth step. If c is consistently less than 4n, the exercise's upper bound is too high.

In fact, ten experiments with n = 10000 all gave 19904 < c < 20017; hence it is extremely likely that the true asymptotic behavior is $\sim 2n^2$.

```
#include <stdio.h>
#include <math.h>
#include "gb_flip.h"
  int *perm;
  int *list;
  int seed;
                   /* random number seed */
               /* this many elements */
  int n;
  main(argc, argv)
        int argc;
        \mathbf{char} * argv[];
     register int i, j, k, t, x, y;
     register double s;
     \langle Scan \text{ the command line } 2 \rangle;
     \langle \text{Initialize everything 3} \rangle;
     while (t) \langle Move 4 \rangle;
     \langle \text{ Print the results 5} \rangle;
   \langle Scan the command line 2\rangle \equiv
  if (argc \neq 3 \lor sscanf(argv[1], \text{"%d"}, \&n) \neq 1 \lor sscanf(argv[2], \text{"%d"}, \&seed) \neq 1) {
     fprintf(stderr, "Usage: \_\%s\_n\_seed n", argv[0]);
     exit(-1);
This code is used in section 1.
```

3. We maintain the following invariants: the indices i where perm[i] > perm[i+1] are list[j] for $0 \le j < t$.

```
 \begin{split} &\langle \text{Initialize everything 3} \rangle \equiv \\ & gb\_init\_rand (seed); \\ & perm = (\textbf{int} *) \ malloc (4 * (n+2)); \\ & list = (\textbf{int} *) \ malloc (4 * (n-1)); \\ & \textbf{for} \ (k=1; \ k \leq n; \ k++) \ perm[k] = n+1-k; \\ & perm[0] = 0; \ perm[n+1] = n+1; \\ & \textbf{for} \ (k=1; \ k < n; \ k++) \ list[k-1] = k; \\ & t = n-1; \\ & s = 0.0; \end{split}
```

This code is used in section 1.

2

This code is used in section 1.

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6. Index.

 $argc: \underline{1}, \underline{2}.$ $argv: \quad \underline{1}, \quad 2.$ $exit: \quad 2.$ fprintf: 2. gb_init_rand : 3. gb_unif_rand : 4. i: $\underline{1}$. j: $\underline{\underline{1}}$. $k: \underline{1}.$ $list: \underline{1}, 3, 4.$ $\begin{array}{ll} \textit{main} \colon \ \underline{1}. \\ \textit{malloc} \colon \ \underline{3}. \end{array}$ $n: \underline{1}.$ $perm: \underline{1}, \underline{3}, \underline{4}.$ printf: 5. $s: \underline{1}.$ seed: $\underline{1}$, $\underline{2}$, $\underline{3}$. sscanf: 2.stderr: 2.t: $\underline{1}$. $x: \underline{1}$.

y: <u>1</u>.

4 NAMES OF THE SECTIONS

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