1

(See https://cs.stanford.edu/~knuth/programs.html for date.)

1.* Intro. This program generates clauses for the transition relation from time 0 to time r in Conway's Game of Life (thus simulating r steps), on an $m \times n$ torus, given m, n, and r.

The Boolean variable for cell (x, y) at time t is named by its so-called "xty code," namely by the decimal value of x, followed by a code letter for t, followed by the decimal value of y. For example, if x = 10 and y = 11 and t = 0, the variable that indicates liveness of the cell is 10a11; and the corresponding variable for t = 1 is 10b11.

Up to 19 auxiliary variables are used together with each xty code, in order to construct clauses that define the successor state. The names of these variables are obtained by appending one of the following two-character combinations to the xty code: A2, A3, A4, B1, B2, B3, B4, C1, C2, C3, C4, D1, D2, E1, E2, F1, F2, G1, G2. These variables are derived from the Bailleux-Boufkhad method of encoding cardinality constraints: The auxiliary variable Ak stands for the condition "at least k of the eight neighbors are alive." Similarly, Bk stands for "at least k of the first four neighbors are alive," and Ck accounts for the other four neighbors. Codes D, E, F, and C refer to pairs of neighbors. Thus, for instance, 10a11C2 means that at least two of the last four neighbors of cell (10,11) are alive.

Those auxiliary variables receive values by means of up to 77 clauses per cell. For example, if u and v are the neighbors of cell z that correspond to a pairing of type D, there are six clauses

$$\bar{u}d_1$$
, $\bar{v}d_1$, $\bar{u}\bar{v}d_2$, $uv\bar{d}_1$, $u\bar{d}_2$, $v\bar{d}_2$.

The sixteen clauses

define b variables from d's and e's; and another sixteen define c's from f's and g's in the same fashion. A similar set of 21 clauses will define the a's from the b's and c's.

Once the a's are defined, thus essentially counting the live neighbors of cell z, the next state z' is defined by five further clauses

$$\bar{a}_4\bar{z}', \quad a_2\bar{z}', \quad a_3z\bar{z}', \quad \bar{a}_3a_4z', \quad \bar{a}_2a_4\bar{z}z'.$$

For example, the last of these states that z' will be true (i.e., that cell z will be alive at time t+1) if z is alive at time t and has ≥ 2 live neighbors but not ≥ 4 .

Nearby cells can share auxiliary variables, according to a tricky scheme that is worked out below. In consequence, the actual number of auxiliary variables and clauses per cell is reduced from 19 and 77 + 5 to 13 and 57 + 5, respectively, except at the boundaries.

2 INTRO SAT-LIFE-TORUS §2

So here's the overall outline of the program. #define maxx 50/* maximum number of lines in the pattern supplied by stdin */ #define maxy 50 /* maximum number of columns per line in stdin */ #include <stdio.h> #include <stdlib.h> char p[maxx + 2][maxy + 2];/* is cell (x, y) potentially alive? */ char $have_b[maxx + 2][maxy + 2];$ /* did we already generate b(x,y)? */ **char** $have_d[maxx + 2][maxy + 2];$ /* did we already generate d(x,y)? */ **char** $have_{-}e[maxx + 2][maxy + 4];$ /* did we already generate e(x,y)? */ char $have_{-}f[maxx + 4][maxy + 2];$ /* did we already generate f(x-2,y)? */ int tt; /* the time being considered */ /* the command-line parameters */ int mm, nn, r; int xmax, ymax; /* the number of rows and columns in the input pattern */ int xmin = maxx, ymin = maxy;/* limits in the other direction */ $char \ timecode[] = "abcdefghijklmnopqrstuvwxyz"$ "ABCDEFGHIJKLMNOPQRSTUVWXYZ" "!\"#\$%&'()*+,-./:;<=>?@[\\]^_'{|}~"; /* codes for $0 \le t \le 83 */$ **char** buf[maxy + 2]; /* input buffer */ /* clauses are assembled here */ unsigned int clause [4]; /* this many literals are in the current clause */ int clauseptr; ⟨Subroutines 6*⟩ $main(\mathbf{int} \ argc, \mathbf{char} * argv[])$ register int j, k, x, y; $\langle \text{Process the command line } 3^* \rangle$; for (tt = 0; tt < r; tt ++) { ymax = nn, ymin = 1;xmax = mm, xmin = 1;for $(x = xmin; x \le xmax; x++)$ for $(y = ymin; y \leq ymax; y++)$ { a(x,y); zprime(x, y);3* $\langle \text{Process the command line } 3^* \rangle \equiv$ if $(argc \neq 4 \lor sscanf(argv[1], "%d", \&mm) \neq 1 \lor sscanf(argv[2], "%d", \&nn) \neq 1 \lor sscanf(argv[3], "%d", \&nn) \Rightarrow 1 \lor sscanf(argv[3],$ $\&r) \neq 1$) { $fprintf(stderr, "Usage: _\%s _m _n _r \n", argv[0]);$ exit(-1); This code is used in section 2*.

§4 SAT-LIFE-TORUS INTRO 3

```
4. \langle \text{Input the pattern 4} \rangle \equiv
  for (x = 1; ; x++) {
    if (\neg fgets(buf, maxy + 2, stdin)) break;
    if (x > maxx) {
       fprintf(stderr, "Sorry, \_the\_pattern\_should\_have\_at\_most\_%d\_rows! \n", maxx);
    for (y = 1; buf[y - 1] \neq '\n'; y++)  {
       if (y > maxy) {
         fprintf(stderr, "Sorry, \bot the \_pattern \_should \_have \_at \_most \_%d \_columns! \n", maxy);
         exit(-4);
       if (buf[y-1] \equiv "") {
         p[x][y] = 1;
         if (y > ymax) ymax = y;
         if (y < ymin) ymin = y;
         if (x > xmax) xmax = x;
         if (x < xmin) xmin = x;
       } else if (buf[y-1] \neq '.') {
         fprintf(stderr, "Unexpected_character_'%c'_found_in_the_pattern!\\n", buf[y-1]);
         exit(-5);
       }
    }
5. #define pp(xx, yy) ((xx) \ge 0 \land (yy) \ge 0 ? p[xx][yy] : 0)
\langle \text{ If cell } (x,y) \text{ is obviously dead at time } t+1, \text{ continue } 5 \rangle \equiv
  if (pp(x-1,y-1) + pp(x-1,y) + pp(x-1,y+1) + pp(x,y-1) + p[x][y] + p[x][y+1] + pp(x+1,y+1)
         y-1)+p[x+1][y]+p[x+1][y+1]<3 continue;
```

4 Intro sat-life-torus §6

6.* Clauses are assembled in the *clause* array (surprise), where we put encoded literals.

The code for a literal is an unsigned 32-bit quantity, where the leading bit is 1 if the literal should be complemented. The next three bits specify the type of the literal (0 thru 7 for plain and A-G); the next three bits specify an integer k; and the next bit is zero. That leaves room for two 12-bit fields, which specify x and y.

Type 0 literals have k=0 for the ordinary xty code. However, the value k=1 indicates that the time code should be for t+1 instead of t. And k=2 denotes a special "tautology" literal, which is always true. If the tautology literal is complemented, we omit it from the clause; otherwise we omit the entire clause. Finally, k=7 denotes an auxiliary literal, used to avoid clauses of length 4.

Here's a subroutine that outputs the current clause and resets the *clause* array.

```
#define taut (2 \ll 25)
#define sign (1_{\rm U} \ll 31)
\langle \text{ Subroutines } 6^* \rangle \equiv
  void outclause(void)
     register int c, k, x, y, p;
     for (p = 0; p < clauseptr; p++)
       if (clause[p] \equiv taut) goto done;
     for (p = 0; p < clauseptr; p \leftrightarrow)
       if (clause[p] \neq taut + sign) {
          if (clause[p] \gg 31) printf("\square"); else printf("\square");
          c = (clause[p] \gg 28) \& #7;
          k = (clause[p] \gg 25) \& #7;
          x = (clause[p] \gg 12) \& #fff;
          y = clause[p] \& #fff;
          if (c)
            printf("%d%c%d%c%d", 1 + ((x + mm - 1) \% mm), timecode[tt], 1 + ((y + nn - 1) \% nn), c + '@', k);
          else if (k \equiv 7) printf ("%d%c%dx", 1 + ((x + mm - 1) \% mm), timecode [tt], 1 + ((y + nn - 1) \% nn));
          else printf("%d%c%d", 1 + ((x + mm - 1) \% mm), timecode[tt + k], 1 + ((y + nn - 1) \% nn));
     printf("\n");
  done: clauseptr = 0;
  }
See also sections 7*, 8*, 9*, 10, 11*, 12*, 14, and 15.
This code is used in section 2*.
7.* And here's another, which puts a type-0 literal into clause.
\langle \text{Subroutines } 6^* \rangle + \equiv
  void applit(int x, int y, int bar, int k)
     clause[clauseptr++] = (bar ? sign : 0) + (k \ll 25) + (x \ll 12) + y;
```

§8 SAT-LIFE-TORUS INTRO 5

8.* The d and e subroutines are called for only one-fourth of all cell addresses (x, y). Indeed, one can show that x is always odd, and that $y \mod 4 < 2$.

```
Therefore we remember if we've seen (x, y) before.
```

```
Slight trick: If yy is not in range, we avoid generating the clause \bar{d}_k twice.
```

```
#define newlit(x, y, c, k) clause[clauseptr ++] = ((c) \ll 28) + ((k) \ll 25) + ((x) \ll 12) + (y)
#define newcomplit(x, y, c, k) clause[clauseptr ++] = sign + ((c) \ll 28) + ((k) \ll 25) + ((x) \ll 12) + (y)
\langle \text{Subroutines } 6^* \rangle + \equiv
  void d(\mathbf{int} \ x, \mathbf{int} \ y)
     register x1 = x - 1, x2 = x, yy = y + 1;
    if (have_{-}d[x \% mm][y \% nn] \neq tt + 1)  {
       applit(x1, yy, 1, 0), newlit(x, y, 4, 1), outclause();
       applit(x2, yy, 1, 0), newlit(x, y, 4, 1), outclause();
       applit(x1, yy, 1, 0), applit(x2, yy, 1, 0), newlit(x, y, 4, 2), outclause();
       applit(x1, yy, 0, 0), applit(x2, yy, 0, 0), newcomplit(x, y, 4, 1), outclause();
       applit(x1, yy, 0, 0), newcomplit(x, y, 4, 2), outclause();
       applit(x2, yy, 0, 0), newcomplit(x, y, 4, 2), outclause();
       have_{-}d[x \% mm][y \% nn] = tt + 1;
     }
  }
  void e(\mathbf{int} \ x, \mathbf{int} \ y)
     register x1 = x - 1, x2 = x, yy = y - 1;
    if (have_{-}e[x \% mm][y \% nn] \neq tt + 1)  {
       applit(x1, yy, 1, 0), newlit(x, y, 5, 1), outclause();
       applit(x2, yy, 1, 0), newlit(x, y, 5, 1), outclause();
       applit(x1, yy, 1, 0), applit(x2, yy, 1, 0), newlit(x, y, 5, 2), outclause();
       applit(x1, yy, 0, 0), applit(x2, yy, 0, 0), newcomplit(x, y, 5, 1), outclause();
       applit(x1, yy, 0, 0), newcomplit(x, y, 5, 2), outclause();
       applit(x2, yy, 0, 0), newcomplit(x, y, 5, 2), outclause();
       have_{-}e[x \% mm][y \% nn] = tt + 1;
  }
```

9.* The f subroutine can't be shared quite so often. But we do save a factor of 2, because x + y is always even.

```
 \begin{array}{l} \text{Void } f(\textbf{int } x, \textbf{int } y) \\ \{ \\ \textbf{register } xx = x - 1, \ y1 = y, \ y2 = y + 1; \\ \textbf{if } (have\_f[x \% mm][y \% nn] \neq tt + 1) \ \{ \\ applit(xx, y1, 1, 0), newlit(x, y, 6, 1), outclause(); \\ applit(xx, y2, 1, 0), newlit(x, y, 6, 1), outclause(); \\ applit(xx, y1, 1, 0), applit(xx, y2, 1, 0), newlit(x, y, 6, 2), outclause(); \\ applit(xx, y1, 0, 0), applit(xx, y2, 1, 0), newcomplit(x, y, 6, 2), outclause(); \\ applit(xx, y1, 0, 0), newcomplit(x, y, 6, 2), outclause(); \\ applit(xx, y2, 0, 0), newcomplit(x, y, 6, 2), outclause(); \\ have\_f[x \% mm][y \% nn] = tt + 1; \\ \} \\ \} \end{array}
```

6 Intro sat-life-torus §10

10. The g subroutine cleans up the dregs, by somewhat tediously locating the two neighbors that weren't handled by d, e, or f. No sharing is possible here.

```
 \begin{array}{l} & \text{ void } g(\textbf{int } x, \textbf{int } y) \\ \{ & \text{ register } x1, \ x2, \ y1, \ y2; \\ & \text{ if } (x \& 1) \ x1 = x - 1, y1 = y, x2 = x + 1, y2 = y \oplus 1; \\ & \text{ else } \ x1 = x + 1, y1 = y, x2 = x - 1, y2 = y - 1 + ((y \& 1) \ll 1); \\ & applit(x1, y1, 1, 0), newlit(x, y, 7, 1), outclause(); \\ & applit(x2, y2, 1, 0), newlit(x, y, 7, 1), outclause(); \\ & applit(x1, y1, 1, 0), applit(x2, y2, 1, 0), newlit(x, y, 7, 2), outclause(); \\ & applit(x1, y1, 0, 0), applit(x2, y2, 0, 0), newcomplit(x, y, 7, 1), outclause(); \\ & applit(x2, y2, 0, 0), newcomplit(x, y, 7, 2), outclause(); \\ & applit(x2, y2, 0, 0), newcomplit(x, y, 7, 2), outclause(); \\ \} \end{array}
```

11.* Fortunately the b subroutine can be shared (since x is always odd), thus saving half of the sixteen clauses generated.

```
 \begin{array}{l} \text{ void } b(\textbf{int } x, \textbf{int } y) \\ \{ & \textbf{register } j, \ k, \ xx = x, \ y1 = y - (y \& 2), \ y2 = y + (y \& 2); \\ \textbf{if } (have\_b[x \% mm][y \% nn] \neq tt + 1) \ \{ \\ d(xx,y1); \\ e(xx,y2); \\ \textbf{for } (j=0; \ j < 3; \ j++) \\ \textbf{for } (k=0; \ k < 3; \ k++) \\ \textbf{if } (j) \ newcomplit(xx,y1,4,j); \ \ /* \ \bar{d}_j \ */ \\ \textbf{if } (k) \ newcomplit(xx,y2,5,k); \ \ /* \ \bar{e}_k \ */ \\ newlit(x,y,2,j+k); \ \ /* \ b_{j+k} \ */ \\ outclause(); \\ \textbf{if } (j) \ newlit(xx,y1,4,3-j); \ \ \ /* \ d_{3-j} \ */ \\ \textbf{if } (k) \ newlit(xx,y2,5,3-k); \ \ \ /* \ e_{3-k} \ */ \\ newcomplit(x,y,2,5-j-k); \ \ \ /* \ \bar{b}_{5-j-k} \ */ \\ outclause(); \\ \} \\ have\_b[x \% mm][y \% nn] = tt + 1; \\ \} \\ \} \end{array}
```

§12 SAT-LIFE-TORUS INTRO 7

12.* The (unshared) c subroutine handles the other four neighbors, by working with f and g instead of d and e.

If y = 0, the overlap rules set y1 = -1, which can be problematic. I've decided to avoid this case by omitting f when it is guaranteed to be zero.

```
\langle Subroutines 6*\rangle + \equiv
  void c(\mathbf{int} \ x, \mathbf{int} \ y)
  {
     register j, k, x1, y1;
     if (x \& 1) x1 = x + 2, y1 = (y - 1) | 1;
     else x1 = x, y1 = y \& -2;
     g(x,y);
     if (0) \langle \text{Set } c \text{ equal to } g \text{ 13} \rangle
     else {
        f(x1, y1);
        for (j = 0; j < 3; j ++)
          for (k = 0; k < 3; k ++)
             if (j+k) {
                if (j) newcomplit (x1, y1, 6, j); /* \bar{f}_j */
                if (k) newcomplit(x, y, 7, k); /* \bar{g}_k */
                newlit(x, y, 3, j + k); /* c_{j+k} */
                outclause();
                if (j) newlit(x1, y1, 6, 3 - j); /* f_{3-j} */
                if (k) newlit(x, y, 7, 3 - k);
                                                      /* g_{3-k} */
                newcomplit(x, y, 3, 5 - j - k); /* \bar{c}_{5-j-k} */
                outclause();
             }
  }
13. \langle \text{Set } c \text{ equal to } g \text{ 13} \rangle \equiv
     for (k = 1; k < 3; k ++) {
        newcomplit(x, y, 7, k), newlit(x, y, 3, k), outclause();
                                                                           /* \bar{g}_k \vee c_k */
        newlit(x, y, 7, k), newcomplit(x, y, 3, k), outclause();
                                                                            /* g_k \vee \bar{c}_k */
     newcomplit(x, y, 3, 3), outclause();
                                                    /* \bar{c}_3 */
     newcomplit(x, y, 3, 4), outclause();
                                                    /* \bar{c}_4 */
This code is used in section 12*.
```

8 INTRO SAT-LIFE-TORUS §14

14. Totals over all eight neighbors are then deduced by the a subroutine.

```
\langle Subroutines 6*\rangle + \equiv
   void a(int x, int y)
       register j, k, xx = x \mid 1;
       b(xx,y);
       c(x,y);
       for (j = 0; j < 5; j++)
          for (k = 0; k < 5; k++)
              if (j + k > 1 \land j + k < 5) {
                 \begin{array}{lll} \textbf{if} \ (j) \ \ newcomplit(xx,y,2,j); & /* \ \bar{b}_j \ */ \\ \textbf{if} \ (k) \ \ newcomplit(x,y,3,k); & /* \ \bar{c}_k \ */ \end{array}
                  newlit(x, y, 1, j + k); /* a_{j+k} */
                  outclause();
       for (j = 0; j < 5; j++)
          for (k = 0; k < 5; k++)
              if (j+k > 2 \land j+k < 6 \land j*k) {
                 \begin{array}{lll} \textbf{if} \ (j) \ \ newlit(xx,y,2,j); & /* \ b_j \ */ \\ \textbf{if} \ (k) \ \ newlit(x,y,3,k); & /* \ c_k \ */ \end{array}
                  newcomplit(x, y, 1, j + k - 1); /* \bar{a}_{j+k-1} */
                  outclause();
              }
   }
```

15. Finally, as mentioned at the beginning, z' is determined from z, a_2 , a_3 , and a_4 . I actually generate six clauses, not five, in order to stick to 3SAT.

```
 \begin{array}{l} \langle \, \text{Subroutines} \,\, 6^* \rangle \, + \equiv \\ \mathbf{void} \,\, zprime(\mathbf{int} \,\, x, \mathbf{int} \,\, y) \\ \{ \\ \quad newcomplit(x,y,1,4), applit(x,y,1,1), outclause(\,); \qquad / * \,\, \bar{a}_4 \bar{z}' \,\, */\\ \quad newlit(x,y,1,2), applit(x,y,1,1), outclause(\,); \qquad / * \,\, a_2 \bar{z}' \,\, */\\ \quad newlit(x,y,1,3), applit(x,y,0,0), applit(x,y,1,1), outclause(\,); \qquad / * \,\, \bar{a}_3 a_4 z' \,\, */\\ \quad newcomplit(x,y,1,3), newlit(x,y,1,4), applit(x,y,0,1), outclause(\,); \qquad / * \,\, \bar{a}_3 a_4 z' \,\, */\\ \quad applit(x,y,0,7), newcomplit(x,y,1,2), newlit(x,y,1,4), outclause(\,); \qquad / * \,\, \bar{x}\bar{z}z' \,\, */\\ \quad applit(x,y,1,7), applit(x,y,1,0), applit(x,y,0,1), outclause(\,); \qquad / * \,\, \bar{x}\bar{z}z' \,\, */\\ \} \end{array}
```

§16 SAT-LIFE-TORUS INDEX 9

16* Index.

```
The following sections were changed by the change file: 1, 2, 3, 6, 7, 8, 9, 11, 12, 16.
                                                                                       ymax: 2^*, 4.
applit: \ \ \underline{7}, \ 8, \ 9, \ 10, \ 15.
                                                                                       ymin: \underline{2}^*, 4.
argc: 2* 3*
                                                                                       yy: 5, 8*
argv: \underline{2}^*, \underline{3}^*
                                                                                       y1: \quad \underline{9}^*, \, \underline{10}, \, \, \underline{11}^*, \, \underline{12}^*.
b: <u>11</u>*
                                                                                       y2: \ \underline{9}^*, \ \underline{10}, \ \underline{11}^*
bar: \underline{7}^*
                                                                                      zprime: 2^*, \underline{15}.
buf: 2^*, 4.
c: <u>6</u>*, <u>12</u>*
clause: 2*, 6*, 7*, 8*
clauseptr: 2, 6, 7, 8.
d: <u>8</u>*
done: \underline{6}^*
e: <u>8</u>*
exit: 3^*, 4.
f: <u>9</u>*
fgets: 4.
fprintf: 3, 4.
g: \underline{10}.
have_b: 2*, 11*
have_d: 2* 8*
have_{-}e: \ \ \underline{2}^{*}, \ 8.^{*}
have_f: 2^* 9*
j: \quad \underline{2}^*, \ \underline{11}^*, \ \underline{12}^*, \ \underline{14}.
k: 2,* 6,* 7,* 11,* 12,* 14.
main: \underline{2}^*
maxx: 2^*, 4.
maxy: 2^*, 4.
mm: 2,* 3,* 6,* 8,* 9,* 11.*
newcomplit: 8,* 9,* 10, 11,* 12,* 13, 14, 15.
newlit: 8, 9, 10, 11, 12, 13, 14, 15.
nn: 2* 3* 6* 8* 9* 11*
outclause: 6, 8, 9, 10, 11, 12, 13, 14, 15.
p: 2* 6*
pp: \underline{5}.
printf: 3,* 6.*
r: \underline{2}^*
sign: 6* 7* 8*
sscanf: 3*
stderr: 3* 4.
stdin: 2, 4.
taut: 6*
timecode: \underline{2}, \underline{6}, \underline{6}
tt: 2* 6* 8* 9* 11*
x: \quad 2, 6, 7, 8, 9, 10, 11, 12, 14, 15.
xmax: 2^*, 4.
xmin: 2^*, 4.
xx: 5, \underline{9}^*, \underline{11}^*, \underline{14}.
x1: \ \ \underline{8}^*, \ \underline{10}, \ \underline{12}^*
x2: 8^*, 10.
y: \quad 2, 6, 7, 8, 9, 10, 11, 12, 14, 15.
```

10 NAMES OF THE SECTIONS SAT-LIFE-TORUS

```
\langle If cell (x,y) is obviously dead at time t+1, continue 5\rangle \langle Input the pattern 4\rangle \langle Process the command line 3*\rangle Used in section 2*. \langle Set c equal to g 13\rangle Used in section 12*. \langle Subroutines 6*, 7*, 8*, 9*, 10, 11*, 12*, 14, 15\rangle Used in section 2*.
```

SAT-LIFE-TORUS

	Section	Page
Intro	 1	1
Index	16	O