

1. Intro. This program is an iterative implementation of an interesting recursive algorithm due to Willard L. Eastman, *IEEE Trans. IT-11* (1965), 263–267: Given a sequence of nonnegative integers $x = x_0x_1 \dots x_{n-1}$ of odd length n , where x is not equal to any of its cyclic shifts $x_k \dots x_{n-1}x_0 \dots x_{k-1}$ for $1 \leq k < n$, we output a cyclic shift σx such that the set of all such σx forms a commafree code of block length n (over an infinite alphabet).

The integers are given as command-line arguments.

The simplest nontrivial example occurs when $n = 3$. If $x = abc$, where a , b , and c aren't all equal, then exactly one of the cyclic shifts $y_0y_1y_2 = abc, bca, cab$ will satisfy $y_0 > y_1 \leq y_2$, and we choose that one. It's easy to check that the triples chosen in this way are commafree.

Similar constructions are possible when $n = 5$ or $n = 7$. But the case $n = 9$ already gets a bit dicey, and when n is really large it's not at all clear that commafreeness is possible. Eastman's paper resolved a conjecture made by Golomb, Gordon, and Welch in their pioneering paper about comma-free codes (1958).

(Of course, it's not at all clear that we would want to actually *use* a commafree code when n is large; but that's another story, and beside the point. The point is that Eastman discovered a really interesting algorithm.)

```
#define maxn 105
#include <stdio.h>
#include <stdlib.h>
int x[maxn + maxn + maxn];
int b[maxn + maxn + maxn];
int bb[maxn];
<Subroutines 5>;
main(int argc, char *argv[])
{
    register int i, j, k, n, p, q, t, tt;
    <Process the command line 2>;
    <Do Eastman's algorithm 3>;
}

2. <Process the command line 2> ≡
if (argc < 4) {
    fprintf(stderr, "Usage: %s x1 x2 . . . xn\n", argv[0]);
    exit(-1);
}
n = argc - 1;
if ((n & 1) == 0) {
    fprintf(stderr, "The number of items, n, should be odd, not %d!\n", n);
    exit(-2);
}
for (j = 1; j < argc; j++) {
    if (sscanf(argv[j], "%d", &x[j - 1]) != 1 ∨ x[j - 1] < 0) {
        fprintf(stderr, "Argument %d should be a nonnegative integer, not '%s'!\n", j, argv[j]);
        exit(-3);
    }
}
```

This code is used in section 1.

3. The algorithm. We think of x as written cyclically, with $x_{n+j} = x_j$ for all $j \geq 0$. The basic idea in the algorithm below is to also think of x as partitioned into $t \leq n$ subwords by boundary markers b_j where $0 \leq b_0 < b_1 < \dots < b_{t-1} < n$; then subword y_j is $x_{b_j}x_{b_j+1} \dots x_{b_{j+1}-1}$, for $0 \leq j < t$, where $b_t = b_0$. If $t = 1$, there's just one subword, and it's a cyclic shift of x . The number t of subwords during each phase will be odd.

Eastman's algorithm essentially begins with $b_j = j$ for $0 \leq j < n$, so that x is partitioned into n subwords of length 1. It successively *removes* boundary points until only one subword is left; that subword is the answer. It operates in phases, so that all subwords during the j th phase have length 3^{j-1} or more; thus at most $\lceil \log_3 n \rceil$ phases are needed. (For example, the case $n = 9$ is “dicey” because it might require two phases.)

The algorithm is based on comparison of adjacent subwords y_{j-1} and y_j . If those subwords have the same length, we use lexicographic comparison; otherwise we declare that the longer subword is bigger.

(After the first phase, all subwords not only have length ≥ 3 , they also always begin with a nonzero entry; in other words, $x_{b_j} > 0$ for every boundary marker b_j . However, we won't need to use that fact explicitly.)

The algorithm can be described with terminology based on the topography of Nevada: Say that i is a *basin* if the subwords satisfy $y_{i-1} > y_i \leq y_{i+1}$. There must be at least one basin; otherwise all the y_j would be equal, and x would equal one of its cyclic shifts. We look at consecutive basins, i and j ; this means that $i < j$ and that i and j are basins, and that $i+1$ through $j-1$ are *not* basins. If there's only one basin we have $j = i + t$. The indices between consecutive basins are called *ranges*.

Since t is odd, there's an odd number of consecutive basins for which $j-i$ is odd. Each phase of Eastman's algorithm retains exactly one boundary point in the range between such basins, and deletes all the others. The retained point is the smallest $k = i + 2l$ such that $y_k > y_{k+1}$.

(For example, suppose $i = 2$ and $j = 9$ are consecutive basins. Then we have $y_1 > y_2 \leq y_3 \leq \dots \leq y_q > y_{q+1} > \dots > y_9 \leq y_{10}$, for some range element $2 < q < 9$. We choose $k = 4$ if $q = 3$ or $q = 4$, $k = 6$ if $q = 5$ or $q = 6$, and $k = 8$ if $q = 7$ or $q = 8$.)

\langle Do Eastman's algorithm 3 $\rangle \equiv$

\langle Initialize 4 \rangle ;

for ($p = 1, t = n; t > 1; t = tt, p++$)

\langle Do one phase of Eastman's algorithm, putting tt boundary points into bb 6 \rangle ;

This code is used in section 1.

4. \langle Initialize 4 $\rangle \equiv$

for ($j = n; j < n + n + n; j++$) $x[j] = x[j - n]$;

for ($j = 0; j < n + n + n; j++$) $b[j] = j$;

$t = n$;

This code is used in section 3.

5. Here's a basic subroutine that returns 1 if subword y_{i-1} exceeds subword y_i , otherwise it returns 0.

⟨Subroutines 5⟩ \equiv

```

int compare(register int i)
{
    register int j;
    if ( $b[i] - b[i-1] \equiv b[i+1] - b[i]$ ) {
        for ( $j = 0; b[i] + j < b[i+1]; j++$ ) {
            if ( $x[b[i-1] + j] \equiv x[b[i] + j]$ ) continue;
            return ( $x[b[i-1] + j] > x[b[i] + j]$ );
        }
        return 0;    /*  $y_{i-1} = y_i$  */
    }
    return ( $b[i] - b[i-1] > b[i+1] - b[i]$ );
}

```

This code is used in section 1.

6. ⟨Do one phase of Eastman's algorithm, putting tt boundary points into bb 6⟩ \equiv

```

{
    for ( $tt = 0, i = 1; i \leq t; i++$ )
        if (compare(i)) break;
    if ( $i > t$ ) {
        fprintf(stderr, "The input is cyclic!\n");
        exit(-666);
    }
    for ( ; compare( $i+1$ );  $i++$ ) ;    /* advance to the first basin */
    for ( ;  $i \leq t; i = j$ ) {
        for ( $q = i+1; compare(q+1) \equiv 0; q++$ ) ;    /* climb the range */
        for ( $j = q+1; compare(j+1); j++$ ) ;    /* advance to the next basin */
        if ( $((j-i) \& 1)$ ) ⟨Choose a boundary point to retain 7⟩;
    }
    printf("Phase %d leaves", p);
    for ( $k = 0; k < tt; k++$ )  $b[k] = bb[k]$ , printf(" %d",  $bb[k]$ );
    printf(" \n");
    for ( ;  $b[k-tt] < n+n; k++$ )  $b[k] = b[k-tt] + n$ ;
}

```

This code is used in section 3.

7. ⟨Choose a boundary point to retain 7⟩ \equiv

```

{
    if ( $(q-i) \& 1$ )  $q++$ ;
    if ( $q < t$ )  $bb[tt++] = b[q]$ ;
    else {
        for ( $k = tt++; k > 0; k--$ )  $bb[k] = bb[k-1]$ ;
         $bb[0] = b[q-t]$ ;
    }
}

```

This code is used in section 6.

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COMMAFREE-EASTMAN

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