

Processor State (x86-64, Partial)

Information about currently executing program

- Temporary data (`%rax, ...`)
- Location of runtime stack (`%rsp`)
- Location of current code control point (`%rip, ...`)
- Status of recent tests (`CF, ZF, SF, OF`)

Registers

<code>%rax</code>	<code>%r8</code>
<code>%rbx</code>	<code>%r9</code>
<code>%rcx</code>	<code>%r10</code>
<code>%rdx</code>	<code>%r11</code>
<code>%rsi</code>	<code>%r12</code>
<code>%rdi</code>	<code>%r13</code>
<code>%rsp</code>	<code>%r14</code>
<code>%rbp</code>	<code>%r15</code>

`%rip` Instruction pointer

`CF` `ZF` `SF` `OF` Condition codes

Current stack top

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Condition Codes (Explicit Setting: Test)

Explicit Setting by Test instruction

- `testq Src2, Src1`
 - `testq b, a` like computing `a & b` without setting destination

- Sets condition codes based on value of `Src1` & `Src2`
- Useful to have one of the operands be a mask

ZF set when `a & b == 0`

SF set when `a & b < 0`

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Today

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Condition Codes (Explicit Setting: Compare)

Explicit Setting by Compare Instruction

- `cmpq Src2, Src1`
 - `cmpq b, a` like computing `a - b` without setting destination

CF set if carry out from most significant bit (used for unsigned comparisons)

ZF set if `a == b`

SF set if `(a - b) < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a > 0 && b < 0 && (a - b) < 0) || (a < 0 && b > 0 && (a - b) > 0)`

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Machine-Level Programming II: Control

Condition Codes (Implicit Setting)

Single bit registers

- CF** Carry Flag (for unsigned)
- SF** Sign Flag (for signed)
- ZF** Zero Flag
- OF** Overflow Flag (for signed)

Implicitly set (think of it as side effect) by arithmetic operations

Example: `addq Src, Dest` \leftrightarrow `t = a + b`

CF set if carry out from most significant bit (unsigned overflow)

ZF set if `t == 0`

SF set if `t < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a > 0 && b > 0 && t < 0) || (a < 0 && b < 0 && t >= 0)`

Not set by `leaq` instruction

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Reading Condition Codes (Cont.)

SetX Instructions:

- Set single byte based on combination of condition codes
- One of addressable byte registers
 - Does not alter remaining bytes
 - Typically use `movzbl` to finish job
 - 32-bit instructions also set upper 32 bits to 0

```
int gt (long x, long y)
{
    return x > y;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
cmpq %rsi, %rdi # Compare x:y
setg %al # Set when >
movzbl %al, %eax # Zero rest of %rax
ret
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Conditional Branch Example (Old Style)

Generation

```
shark> gcc -Og -S -fno-if-conversion control.c
```

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq %rsi, %rdi # x:y
    jle .L4
    movq %rdi, %rax
    subq %rsi, %rax
    ret
.L4:
    # x <= y
    movq %rsi, %rax
    subq %rdi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



x86-64 Integer Registers

%rax	%a1	%r8	%r8b
%rbx	%b1	%r9	%r9b
%rcx	%c1	%r10	%r10b
%rdx	%d1	%r11	%r11b
%rsi	%s11	%r12	%r12b
%rdi	%d11	%r13	%r13b
%rsp	%sp1	%r14	%r14b
%rbp	%bp1	%r15	%r15b

- Can reference low-order byte

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Jumping

jX Instructions

- Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
jz	ZF	Equal / Zero
jnz	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF) & ~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Reading Condition Codes

SetX Instructions

- Set low-order byte of destination to 0 or 1 based on combinations of condition codes
- Does not alter remaining 7 bytes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~(SF^OF) & ~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF & ~ZF	Above (unsigned)
setb	CF	Below (unsigned)

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Today

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Using Conditional Moves

Conditional Move Instructions

- Instruction supports:
if (Test) Dest ← Src
- Supported in post-1995 x86 processors
- GCC tries to use them
 - But, only when known to be safe

Why?

- Branches are very disruptive to instruction flow through pipelines
- Conditional moves do not require control transfer

C Code

```
val = Test
    ? Then_Expr
    : Else_Expr;
```

Goto Version

```
result = Then_Expr;
eval = Else_Expr;
nt = !Test;
if (nt) result = eval;
return result;
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

15

Today

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

16

General Conditional Expression Translation (Using Branches)

C Code

```
val = Test ? Then_Expr : Else_Expr;
```

```
val = x > y ? x - y : y - x;
```

Goto Version

```
ntest = !Test;
if (ntest) goto Else;
val = Then_Expr;
goto Done;
Else:
    val = Else_Expr;
Done:
    . . .
```

- Create separate code regions for then & else expressions
- Execute appropriate one

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

14

Bad Cases for Conditional Move

Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

- Both values get computed
- Only makes sense when computations are very simple

Risky Computations

```
val = p ? *p : 0;
```

- Both values get computed
- May have undesirable effects

Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- Must be side-effect free

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

17

Expressing with Goto Code

- C allows goto statement
- Jump to position designated by label

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x - y;
    else
        result = y - x;
    return result;
}
```

```
long absdiff_j
(long x, long y)
{
    long result;
    int ntest = x <= y;
    if (ntest) goto Else;
    result = x - y;
    goto Done;
Else:
    result = y - x;
Done:
    return result;
}
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

13

Conditional Move Example

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x - y;
    else
        result = y - x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax    # x
    subq    %rsi, %rax    # result = x - y
    movq    %rsi, %rdx
    subq    %rdi, %rdx    # eval = y - x
    cmpq    %rsi, %rdi    # x:y
    cmovle  %rdx, %rax    # if <=, result = eval
    ret
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

18

General “Do-While” Translation

C Code

```
do
    Body
while (Test);
```

Body: {
 Statement₁;
 Statement₂;
 ...
 Statement_n;
}

Goto Version

```
loop:
    Body
    if (Test)
        goto loop
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



General “While” Translation #2

While version

```
while (Test)
    Body
```

- “Do-while” conversion
- Used with -O1

Do-While Version

```
if (!Test)
    goto done;
do
    Body
while (Test);
done:
```

Goto Version

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
        goto loop;
done:
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



“Do-While” Loop Compilation

Goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rax	result

```
movl $0, %eax # result = 0
.L2: # loop:
movq %rdi, %rdx
andl $1, %edx # t = x & 0x1
addq %rdx, %rax # result += t
shrq %rdi # x >>= 1
jne .L2 # if (x) goto loop
rep; ret
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



While Loop Example #1

C Code

```
long pcount_while
(unsigned long x) {
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Jump to Middle

```
long pcount_goto_jtm
(unsigned long x) {
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x) goto loop;
    return result;
}
```

- Compare to do-while version of function
- Initial goto starts loop at test

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



“Do-While” Loop Example

C Code

```
long pcount_do
(unsigned long x) {
    long result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

Goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

- Count number of 1's in argument x (“popcount”)
- Use conditional branch to either continue looping or to exit loop

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



General “While” Translation #1

- “Jump-to-middle” translation
- Used with -Og

While version

```
while (Test)
    Body
```

Goto Version

```
goto test;
loop:
    Body
test:
    if (Test)
        goto loop;
done:
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



“For” Loop → While Loop

For Version

```
for (Init; Test; Update )
    Body
```

While Version

```
Init;
while (Test) {
    Body
    Update;
}
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Today

- Control: Condition codes
- Conditional branches
- Loops
- Switch Statements

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



“For” Loop Form

General Form

```
for (Init; Test; Update )
    Body
```

```
#define WSIZE 8*sizeof(int)
long pcount_for
(unsigned long x)
{
    size_t i;
    long result = 0;
    for (i = 0; i < WSIZE; i++)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
    unsigned bit =
        (x >> i) & 0x1;
    result += bit;
}
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



“For” Loop Do-While Conversion

C Code Goto Version

```
long pcount_for
(unsigned long x)
{
    size_t i;
    long result = 0;
    for (i = 0; i < WSIZE; i++)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

- Initial test can be optimized away

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

```
long pcount_for_goto_dw
(unsigned long x) {
    size_t i;
    long result = 0;
    i = 0;
    if (! (i < WSIZE))
        goto done;
    loop:
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
    }
    i++;
    if (i < WSIZE)
        goto loop;
done:
    return result;
}
```



While Loop Example #2

C Code

```
long pcount_while
(unsigned long x) {
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Do-While Version

```
long pcount_goto_dw
(unsigned long x) {
    long result = 0;
    if (!x) goto done;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
done:
    return result;
}
```

- Compare to do-while version of function
- Initial conditional guards entrance to loop

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



For-While Conversion

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

Body

```
{
    unsigned bit =
        (x >> i) & 0x1;
    result += bit;
}
```

```
long pcount_for_while
(unsigned long x)
{
    size_t i;
    long result = 0;
    i = 0;
    while (i < WSIZE)
    {
        unsigned bit =
            (x >> i) & 0x1;
        result += bit;
        i++;
    }
    return result;
}
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition



Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi    # x:6
    ja      .L8
    jmp     *.L4(, %rdi, 8)
```

What range of values
takes default?

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Note that w not
initialized here

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Jump Table

Jump table

```
.section .rodata
.align 8
.L4:
.quad .L8 # x = 0
.quad .L3 # x = 1
.quad .L5 # x = 2
.quad .L9 # x = 3
.quad .L8 # x = 4
.quad .L7 # x = 5
.quad .L7 # x = 6
```

```
switch(x) {
    case 1: // .L3
        w = y*z;
        break;
    case 2: // .L5
        w = y/z;
        /* Fall Through */
    case 3: // .L9
        w += z;
        break;
    case 5:
    case 6: // .L7
        w -= z;
        break;
    default: // .L8
        w = 2;
}
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Jump Table Structure

Switch Form

```
switch(x) {
    case val_0:
        Block 0
    case val_1:
        Block 1
    . . .
    case val_n-1:
        Block n-1
}
```

Jump Table

```
jtab:
    Targ0
    Targ1
    .
    .
    Targn-1
```

Jump Targets

```
Targ0: Code Block 0
Targ1: Code Block 1
Targ2: Code Block 2
.
.
Targn-1: Code Block n-1
```

Translation (Extended C)

```
goto *JTab[x];
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Assembly Setup Explanation

Table Structure

- Each target requires 8 bytes
- Base address at .L4

Jumping

- Direct: `jmp .L8`
- Jump target is denoted by label .L8
- Indirect: `jmp *.L4(, %rdi, 8)`
- Start of jump table: .L4
- Must scale by factor of 8 (addresses are 8 bytes)
- Fetch target from effective Address `.L4 + x*8`
- Only for $0 \leq x \leq 6$

Jump table

```
.section .rodata
.align 8
.L4:
.quad .L8 # x = 0
.quad .L3 # x = 1
.quad .L5 # x = 2
.quad .L9 # x = 3
.quad .L8 # x = 4
.quad .L7 # x = 5
.quad .L7 # x = 6
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

```
long switch_eg
(long x, long y, long z)
{
    long w = 1;
    switch(x) {
    case 1:
        w = y*z;
        break;
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
    case 5:
    case 6:
        w -= z;
        break;
    default:
        w = 2;
    }
    return w;
}
```

Switch Statement Example

Multiple case labels

- Here: 5 & 6

Fall through cases

- Here: 2

Missing cases

- Here: 4

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Switch Statement Example

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    movq    %rdx, %rcx
    cmpq    $6, %rdi    # x:6
    ja      .L8          # Use default
    jmp     *.L4(, %rdi, 8) # goto *JTab[x]
```

Indirect
jump

Jump table

```
.section .rodata
.align 8
.L4:
.quad .L8 # x = 0
.quad .L3 # x = 1
.quad .L5 # x = 2
.quad .L9 # x = 3
.quad .L8 # x = 4
.quad .L7 # x = 5
.quad .L7 # x = 6
```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Code Blocks (x == 2, x == 3)

```

long w = 1;
...
switch(x) {
...
case 2:
    w = y/z;
    /* Fall Through */
case 3:
    w += z;
    break;
...
}

```

```

.L5:                # Case 2
    movq    %rsi, %rax
    cqto
    idivq   %rcx    # y/z
    jmp     .L6     # goto merge
.L9:                # Case 3
    movl    $1, %eax # w = 1
.L6:                # merge:
    addq    %rcx, %rax # w += z
    ret

```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Summary

- Today**
 - Control: Condition codes
 - Conditional branches & conditional moves
 - Loops
 - Switch statements
- Next Time**
 - Stack
 - Call / return
 - Procedure call discipline

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Handling Fall-Through

```

long w = 1;
...
switch(x) {
...
case 2:
    w = y/z;
    /* Fall Through */
case 3:
    w += z;
    break;
...
}

```

```

case 2:
    w = y/z;
    goto merge;

```

```

case 3:
    w = 1;

```

```

merge:
    w += z;

```

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Summarizing

- C Control**
 - if-then-else
 - do-while
 - while, for
 - switch
- Assembler Control**
 - Conditional jump
 - Conditional move
 - Indirect jump (via jump tables)
 - Compiler generates code sequence to implement more complex control
- Standard Techniques**
 - Loops converted to do-while or jump-to-middle form
 - Large switch statements use jump tables

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Code Blocks (x == 1)

```

switch(x) {
case 1: // .L3
    w = y*z;
    break;
...
}

```

```

.L3:
    movq    %rsi, %rax # y
    imulq   %rdx, %rax # y*z
    ret

```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition

Code Blocks (x == 5, x == 6, default)

```

switch(x) {
...
case 5: // .L7
case 6: // .L7
    w -= z;
    break;
default: // .L8
    w = 2;
}

```

```

.L7:
    movl    $1, %eax # Case 5,6
    subq    %rdx, %rax # w -= z
    ret
.L8:
    movl    $2, %eax # Default:
    ret

```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Notes adapted from Bryant and O'Hallaron, Computer Systems: A Programmer's Perspective, Third Edition