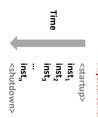
Control Flow

- Processors do only one thing:
- From startup to shutdown, a CPU simply reads and executes (interprets) a sequence of instructions, one at a time This sequence is the CPU's control flow (or flow of control)

Physical control flow





Altering the Control Flow

- Up to now: two mechanisms for changing control flow:
- Jumps and branches
- Call and return
- React to changes in program state

Insufficient for a useful system: Difficult to react to changes in system state

- Data arrives from a disk or a network adapter
- Instruction divides by zero
- User hits Ctrl-C at the keyboard
- System timer expires
- System needs mechanisms for "exceptional control flow"



1118

Exceptional Control Flow

- Exists at all levels of a computer system
- Low level mechanisms
- 1. Exceptions
- Change in control flow in response to a system event (i.e., change in system state)

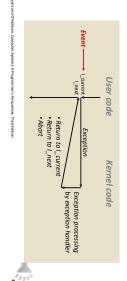
Implemented using combination of hardware and OS software

- Higher level mechanisms
- 2. Process context switch Implemented by OS software and hardware timer
- 3. Signals
- Implemented by OS software
- 4. Nonlocal jumps: setjmp() and longjmp()
 Implemented by C runtime library

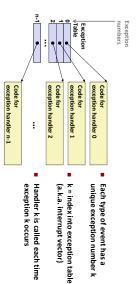


Exceptions

- An exception is a transfer of control to the OS kernel in response to some event (i.e., change in processor state)
 Kernel is the memory-resident part of the OS
 Examples of events: Divide by 0, arithmetic overflow, page fault, I/O request completes, typing Ctrl-C



Exception Tables



Asynchronous Exceptions (Interrupts)

- Caused by events external to the processor
- Indicated by setting the processor's interrupt pin Handler returns to "next" instruction

Examples:

- Timer interrupt
- Every few ms, an external timer chip triggers an interrupt
- Used by the kernel to take back control from user programs I/O interrupt from external device
 Hitting Ctrl-C at the keyboard
 Arrival of a packet from a network
 Arrival of data from a disk







Synchronous Exceptions

- Caused by events that occur as a result of executing an instruction:
- Intentional
- Examples: system calls, breakpoint traps, special instructions
 Returns control to "next" instruction
- Faults

- Aborts
- Unintentional and unrecoverable
 Examples: Illegal instruction, parity error, machine check
 Aborts current program



System Calls

- Each x86-64 system call has a unique ID number
- Examples:

62	60	59	57	4	ω	2	ь	0	Number
kill	_exit	execve	fork	stat	close	open	write	read	Name
Send signal to process	Terminate process	Execute a program	Create process	Get info about file	Close file	Open file	Write file	Read file	Description



System Call Example: Opening File

- User calls: open (filename, options)
 Calls open function, which invokes system call instruction syscall

e5d79: b8 02 00 00 00 mov \$0x2,%eax # open is syscall #2 e5d7e: 0f 05 syscall # Return value in %rax e5d80: 48 3d 01 f0 ff ff cmp \$0xfffffffffff001,%rax 00000000000e5d70 <__open>: ೭



- %rax contains syscall number
- Other arguments in %rdi, %rsi, %rdx, %r10, %r8, %r9 Return value in %rax

Negative value is an error
corresponding to negative
errno

- Unintentional but possibly recoverable
 Examples: page faults (recoverable), protection faults (unrecoverable), floating point exceptions
 Either re-executes faulting ("current") instruction or aborts

Fault Example: Page Fault

- User writes to memory location
- That portion (page) of user's memory is currently on disk

80483ъ7:



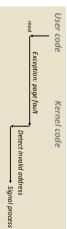
c7 05 10 9d 04 08 0d movl \$0xd,0x8049d10





Fault Example: Invalid Memory Reference

int a[1000]; main () 8048357: User code a[5000] = 13; 05 60 e3 04 08 0d Kernel code \$0xd,0x804e360



- Sends **SIGSEGV** signal to user process User process exits with "segmentation fault"

