



UC Berkeley EECS  
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## Lecture 6: Environment Diagrams, Recursion Review, Midterm Review

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<http://inst.eecs.berkeley.edu/~cs88>



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### Recursion Key concepts – by example



1. Test for simple “base” case

2. Solution in simple “base” case

```
def sum_of_squares(n):  
    if n < 1:  
        return 0  
    else:  
        return sum_of_squares(n-1) + n**2
```

3. Assume recursive solution  
to simpler problem

4. Transform soln of simpler  
problem into full soln

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## Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
4         return 1
5     else:
6         return n_squared + sum_of_squares(n-1)
7
8 sum_of_squares(3)

```

Edit code

Frames

Global frame

sum\_of\_squares

Objects

func sum\_of\_squares(n) [parent=Global]

python3.3

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## How does it work?

- Each recursive call gets its own local variables
  - Just like any other function call
- Computes its result (possibly using additional calls)
  - Just like any other function call
- Returns its result and returns control to its caller
  - Just like any other function call
- The function that is called happens to be itself
  - Called on a simpler problem
  - Eventually bottoms out on the simple base case
- Reason about correctness “by induction”
  - Solve a base case
  - Assuming a solution to a smaller problem, extend it



## Environments Example

Python 3.3

```

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8 sum_of_squares(3)

```

Edit code

Frames

Global frame

sum\_of\_squares

Objects

func sum\_of\_squares(n) [parent=Global]

f1: sum\_of\_squares [parent=Global]

n 3

n\_squared 9

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## Local variables

```

def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 0
    else:
        return n_squared + sum_of_squares(n-1)

```

- Each call has its own “frame” of local variables
- What about globals?
- Let’s see the environment diagrams

<https://goo.gl/CiFaUJ>



## Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
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Edit code

Frames

Global frame

sum\_of\_squares

Objects

func sum\_of\_squares(n) [parent=Global]

f1: sum\_of\_squares [parent=Global]

n 3

n\_squared 9

f2: sum\_of\_squares [parent=Global]

n 2

n\_squared 4


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## Environments Example



```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
4         return 1
5     else:
6         return n_squared + sum_of_squares(n-1)
7
8 sum_of_squares(3)

```

Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)

n 2  
n\_squared 4

```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
4         return 1
5     else:
6         return n_squared + sum_of_squares(n-1)
7
8 sum_of_squares(3)

```

Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)


n 2  
n\_squared 4

f3: sum\_of\_squares (parent=Global)

n 1

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## Environments Example



```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
4         return 1
5     else:
6         return n_squared + sum_of_squares(n-1)
7
8 sum_of_squares(3)

```

Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)


n 2  
n\_squared 4  
Return value 5

f3: sum\_of\_squares (parent=Global)

n 1  
n\_squared 1  
Return value 1

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## Environments Example



```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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7
8 sum_of_squares(3)

```

Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)

n 2  
n\_squared 4

f3: sum\_of\_squares (parent=Global)

n 1  
n\_squared 1

```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)


n 2  
n\_squared 4

f3: sum\_of\_squares (parent=Global)

n 1  
n\_squared 1

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## Environments Example



```

Python 3.3
1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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8 sum_of_squares(3)

```

Global frame

sum\_of\_squares → func sum\_of\_squares(n) (parent=Global)

f1: sum\_of\_squares (parent=Global)

n 3  
n\_squared 9

f2: sum\_of\_squares (parent=Global)


n 2  
n\_squared 4

f3: sum\_of\_squares (parent=Global)

n 1  
n\_squared 1  
Return value 1

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## How much ???



- Time is required to compute `sum_of_squares(n)` ?
  - Recursively?
  - Iteratively?
- Space is required to compute `sum_of_squares(n)` ?
  - Recursively?
  - Iteratively?
- Count the frames...
- Recursive is linear, iterative is constant!

Linear proportional to  $cn$  for some  $c$

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## Tail Recursion

- All the work happens on the way down the recursion
- On the way back up, just return

```
def sum_up_squares(i, n, accum):
    """Sum the squares from i to n in incr. order"""
    if i > n:
        # Base Case
    else:
        # Tail Recursive Case

>>> sum_up_squares(1,3,0)
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```

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## Tree Recursion with HOF

```
def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""

    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split_fun(less_maker(pivot), rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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## Tree Recursion

- Break the problem into multiple smaller sub-problems, and Solve them recursively

```
def split(x, s):
    return [i for i in s if i <= x], [i for i in s if i > x]

def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""
    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split(pivot, rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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## Computational Concepts Toolbox

- Data type: values, literals, operations,
  - e.g., int, float, string
- Expressions, Call expression
- Variables
- Assignment Statement
- Sequences: tuple, list
  - indexing
- Data structures
- Tuple assignment
- Call Expressions
- Function Definition Statement
- Conditional Statement
- Iteration:
  - data-driven (list comprehension)
  - control-driven (for statement)
  - while statement
- Higher Order Functions
  - Functions as Values
  - Functions with functions as argument
  - Assignment of function values
- Recursion
- Environment Diagrams

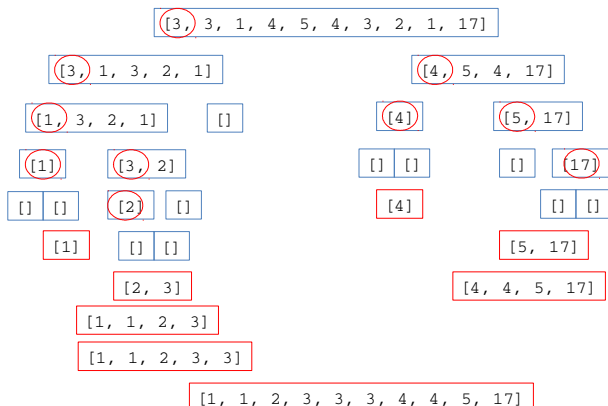


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## QuickSort Example



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## Answers for the Wandering Mind

The computer chooses a random element  $x$  of the list generated by `range(0,n)`. What is the smallest amount of iteration/recursion steps the best algorithm needs to guess  $x$ ?

$\log_2 n$

How would the algorithm look like?

Guess the binary digits of  $x$  starting with the highest significant digit. This is, ask questions of the form "smaller than  $2^{n-1}$ ?" (yes => 0...), "smaller than  $2^{n-2}$ ?" (no => 0 1...), "smaller than  $2^{n-2}+2^{n-3}$ ?", ...

This method is also called: binary search

Quantum physics: Allow less than  $\log_2 n$  guesses

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