

Lab: vacuum world cleaner

function REFLIX - VACUUM - AGENT
([location, status]) return

any action of { up, down, left, right }

if status = Dirty then return suck

else if location = A then return right

else if location = B then return left

(A, B) (A, B) (left, right)

Algorithm for two quadrants,

Step 1: Initialization.

- Input current room (either A or B)
- Input the status of the room clean or dirty.
- Initialize a cost to 0.

Step 2: Display initial room status

Step 3: Cleaning loop.

- While either of the room is dirty.
- 1) If current room is A and is dirty then clean and increase cost by one
 - 2) If current room is B and is dirty then clean and increase cost by one.

Step 4: Display cost
Step 5: Stop.

Algorithm for four quadrants.

Step 1: Start

Accept a current room input from user (either A, B, C or D).

Step 2: Display the initial status.

Clean rooms

While all rooms are clean,

1) If first room is dirty clean and go to next room

2) If second room is dirty clean and go to next room

3) If third room is dirty clean and go to next room

4) If fourth room is dirty clean and go to next.

5) Increase cost when cleaning & decrease count.

Step 4: Show the last status.
Step 5: End.

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same solving

Labs: 8. Puzzle Problem -

(BFS) & (DFS)

Algorithm for BFS:

1) Initialization :

- Create a node data structure to represent the current state of the puzzle
- Initialize the node with the initial state of the puzzle
- Create a queue to store nodes to be explored
- Enqueue the initial node.

2) Loop until goal state is reached or queue is empty

3) Return solution path.

Algorithm for DFS:

1) Initialization :

- Initialize the node with the initial state of puzzle
- Push the initial node onto the stack

2) Loop until goal state is reached or stack is empty.

3) Backtrack : if the stack is empty and the goal state is not reached, return failure.

Output.

Enter the initial state (space-separated):

1 2 3 0 4 6 7 5 8

Move 0:	Move 1:	Move 2:	Move 3:
1 2 3	1 2 3	1 2 3	1 2 3
0 4 6	4 0 6	4 5 6	4 5 6
7 5 8	7 5 8	7 0 8	7 8 0

No. of moves = 3.

Lab: for 8 puzzle problem using A* implementation to calculate $f(n)$ in

(a) $g(n)$ = depth of node

$h(n)$ = heuristic value \rightarrow no. of misplaced tiles

$$f(n) = g(n) + h(n)$$

(b) $g(n)$ = depth of node

$h(n)$ = heuristic value \rightarrow manhattan distance

$$f(n) = g(n) + h(n)$$

Algorithm for No. of misplaced tiles:

Step 1: Initialize the initial state and goal state

Step 2: Based on the position of '-' tile calculated possible movements and calculate $g(n)$ depth of puzzle and $h(n)$:

Step 3: calculate $f(n) = h(n) + g(n)$ for all possibilities with minimum $f(n)$.

Step 4: When $f(n) = 0$, the goal is achieved

Algorithm for manhattan distance \rightarrow

Step 1: Initialize the initial stage and goal state.

Step 2: Loop until not become a i.e. goal state is achieved.

Step 3: Perform $f(n) = g(n) + h(n)$ and move forward by selecting minimum value.

Step 4: Once the goal is achieved point $g(n)$ and No. of movements.

Lesson 3: Iterative deepening search Algo

Algorithm:

- 1) For each child of the current node
- 2) If it is the target node, return
- 3) If current maximum depth is reached, return
- 4) Set the current node to this node and go back to 1.
- 5) After having gone through all children, go to the next child of the parent (the next sibling).
- 6) After having gone through all children of the start node, increase the maximum depth and go back to 1.
- 7) If we have reached all leaf (bottom) nodes, the goal node doesn't exist.

function ~~ITERATIVE - DEEPENING - SEARCH~~ (problem) return
a solution, or failure

for depth = 0 to ∞ do

result \leftarrow Depth-limited-Search
(problem, depth,
if result \neq cutoff then
return result)

Labs: Hill climbing Algo for N-queens

function HILL-CLIMBING (problem)
returns a state that
is a local maximum

current ← MAKE-NODE (problem,
INITIAL-STATE)

loop do

neighbour ← a highest-valued
successor of current

if neighbor.VALUE ≤ current
VALUE then return

current.STATE

current ← neighbor.

- State: 4 queens on the board
 - variables: n_0, n_1, n_2, n_3 where n_i is the row position of the queen in column i .
Assume there is one queen per column
 - domain for each variable n_i : $\in \{0, 1, 2, 3\}, \forall i$.
- Initial state: a random state
- Goal state: 4 queens on the board.
No pair of queens attacking each other.

Neighbour relation : Swap the row position of two queens.

- Cost function : The no. of pairs of queens attacks on each other, directly and indirectly

Lab: Annealing N-queens
function simulated-Annealing (problem, schedule) return a solution state
inputs : problem , a problem schedule a mapping from time to temperature

current \leftarrow Make - Node (problem, Initiate - state).
from $t \leftarrow 1$ to ∞ do
 $T \leftarrow$ schedule (t)
if $T=0$ then return current
if $\Delta E > 0$ then current \leftarrow next
else current \leftarrow next only with probability $e^{\Delta E/T}$

OR

- 1 Start a random point x
- 2 Choose a new point x_j on a neighbourhood $N(n)$
- 3 Decide whether or not to move to the new point x_j
The decision will be made based on probability function $P(n, x_j, T)$.

$$P(x, x_j, T) = \begin{cases} 1 & \text{if } f(x_j) \geq f(n) \\ e^{\frac{f(n_j) - f(n)}{T}} & \text{if } f(n_j) < f(n) \end{cases}$$

4. Reduce T

Output it with move

1) for 8 queen problem

The best position found is
[6 3 1 7 8 0 2 4]

Lab: Propositional Entailment

function TT-ENTAILS? (KB, α) returns true or false
inputs: KB, the knowledge base,
a sentence in propositional logic.
 α , query, a sentence in propositional logic.

Symbols \leftarrow a list of propositional symbols in KB and α

return TT-CHECK-ALL (KB, α , symbols, § 3)

function TT-CHECK-ALL (KB, α , symbols, model) return true or false

if EMPTY? (symbols) then
if PL-TRUE? (KB, model) then
return

PL-True? (α , model)

else do ,

~~P \leftarrow FIRST (symbols)~~
rest \leftarrow REST (symbols)

return (TT-CHECK-ALL (KB, α , rest,
 \rightarrow model $\cup \{ P = \text{true} \} \rightarrow$

and

(TT-CHECK-ALL (KB, α , rest,
 \rightarrow model $\cup \{ P = \text{false} \}).$

Lab: Unification in FOL

Algorithm: Unify (Ψ_1, Ψ_2)

Step 1: If Ψ_1 or Ψ_2 is variable or constant, then:

- (a) If Ψ_1 or Ψ_2 are identical, then return nil.
- (b) Else if Ψ_1 is variable then if Ψ_1 occurs in Ψ_2 , then return failure.
- (c) Else if Ψ_2 is variable If Ψ_2 occurs in Ψ_1 , then return failure, else return (Ψ_1/Ψ_2) .
- (d) Else return failure.

Step 2: If the initial predicate symbol in Ψ_1 and Ψ_2 are not same, then return failure.

Step 3: Set substitution set (SUBST) to NIL.

Step 4: For $i=1$ to the number of elements in Ψ_1 ,

- (a) Call unify function with i^{th} element of Ψ_1 and i^{th} element of Ψ_2 , and put result into S.

- (b) If $S \neq \text{NIL}$ then do,
Apply S to the remainder

of both.

(b) SUBST = APPEND (s, SUBST)
Step 6: Return SUBST

eg : $P(n, F(y)) \xrightarrow{①}$
 $P(a, F(g(n))) \xrightarrow{②}$

in ① & ② predicate are identical
and no. of argument are equal.

in ① replace n with a

$$P(a, F(y)) \xrightarrow{a/x} ①$$

in ① and ② are same
replace y with $g(n)$.

$$P(a, F(g(x))) \xrightarrow{y/g(n)} ①$$

Now ① and ② are same.

Let's create KB consisting of for statement & process the query

→ {FOL - Forward + Deagreen}

function FOL-FC-ASK(KB, α) returns
a substitution or false

inputs : KB, the knowledge base,
a set of first-order definite clauses.

α , the query, an atomic sentence.

local variables : new, the new sentence inferred on each iteration.

repeat until new is empty

new $\leftarrow \{ \}$

for each rule in KB do
 $(P_1 \wedge \dots \wedge P_n \Rightarrow q) \leftarrow \text{STANDARDIZE-VARIABLES}(\text{rule})$

for each θ such that subset
 $(\theta, P_1 \wedge \dots \wedge P_n) = \text{SUBSET}(\theta, P'_1 \wedge \dots \wedge P'_n)$
for some $P'_1 \dots P'_n$ in KB

$q' \leftarrow \text{SUBST}(\theta, q)$

if q' does not unify with some sentences already in KB or

new then add q' to new.
 $\phi \leftarrow \text{UNIFY } (q', \alpha)$
if ϕ is not fail then
return ϕ
add new to KB
return false.

Lab: FOL statement to CNF: (Resolution)

Basic steps for proving a conclusion S given premises premise₁, ..., premise_n.

(all expressed in FOL):

- 1) Convert all sentences to CNF
- 2) Negate conclusion S and convert in CNF.
- 3) Add negated premise to clauses.
- 4) Repeat until contradiction or no progress is made.
 - (a) Select two clauses
 - (b) Resolve them together, performing all required unifications.
 - (c) If resolvent is the empty clause, a contradiction has been found (ie, S follows from the premise).
 - (d) If not, add resolvent to the premises.

Lab: Alpha Beta Pruning

function **ALPHA-BETA-SEARCH** (state)
return an action

$v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$

return the action in **ACTIONS**(state) with value v

function **MAX-VALUE** (state, α, β)
returns a utility value

if **TERMINAL-TEST**(state) then return
UTILITY(state)

$v \leftarrow -\infty$

for each a in **ACTIONS**(state) do

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(\text{RESULT}(s, a), \alpha, \beta))$

if $v \geq \beta$ then return v

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return v

function **MIN-VALUE** (state, α, β)

return a utility value if
TERMINAL-TEST(state) then
return **UTILITY**(state)

$v \leftarrow \infty$

for each a in **ACTIONS**(state) do

$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a), \alpha, \beta))$

if $v \leq \alpha$ then return v

$\beta \leftarrow \text{MIN}(\beta, v)$

return v

Con

Max α

6

Min
Max β

N9A-72 3

6

S

Max α

5

3

6

7

5

8

Min β

5

4

3

6

7

6

7

5

8

6

(a) 5 6 7 4 5 3 6 1 6 - X M 7 5 9 8 6

value utility a min

where want (state) T>T - I A U M 9 7 T ji

Output : (state) PTI J TU

Enter the leaf node value
separated by space:-

-1 8 -3 -1 2 1 < -3 | 4

Optimal value calculated using
Minimax : 2

91

2 3 12

10 14 V - U M

position a min

want (state) T>T - I A U M 9 7 T

(state) PTI J TU a min