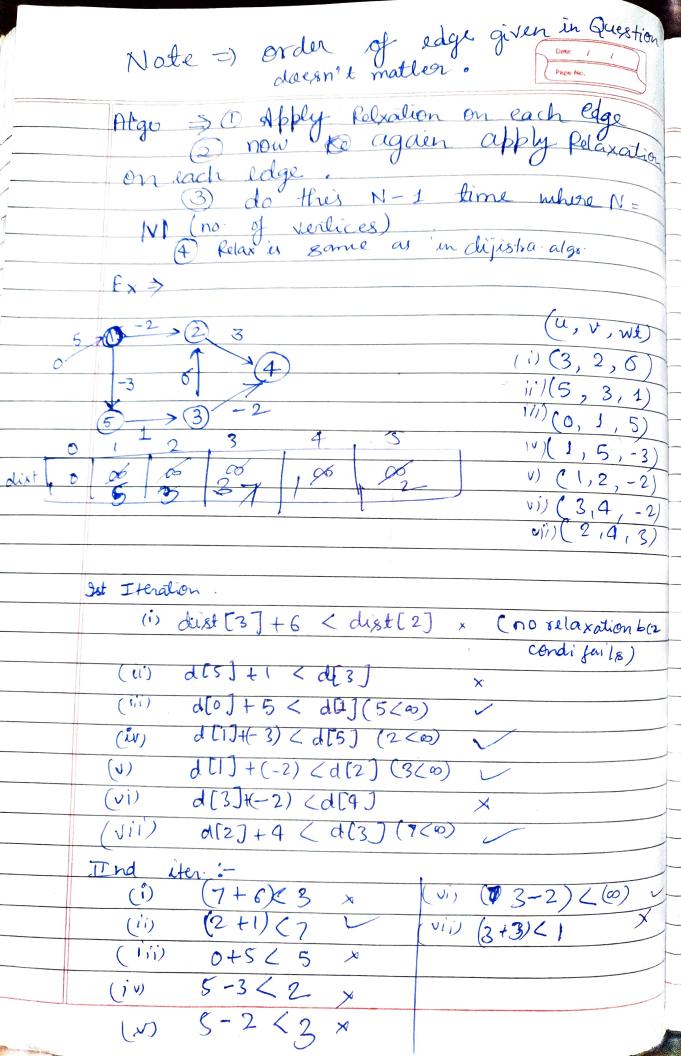
Bellmon Ford Alg:—
(Sperdest fath finding)

(Donly applicable for directed graph.

(9) given undirected graph then change it
by side edges). edges and - ve mighted y le , it will sook fine (3) Ot helps to & dollet no - ue cycle.





total 5 iterations will be completed. Ans: - Only (n-1) iteration original 3.

Ontution: 6 -> 1 -> 2 -> 3 -> 4 (u, v, wt) St 9ter. 2nd (3,4,1)  $\omega+\kappa\omega(x)$   $\times$  (2,3,1)  $\omega+1<\omega(x)$   $\times$ 0+160 (x) 1+160 (1,2,1)0+1 L &(V) X (0, 1, 1)0 0 0 0 0 dist We can clearly observe that at max upation of starting node can reach to last node in (n-1) no of time. O More togical Armen - Since in a graph of N nodes, will take n-1 edges to neach from 1st node to last, seeme (n-1) at max extention required

How to udelict cycle in graph?

Apply Bellman Ford Algo;  $\S(N-1)$  time itteration. Of at this iteration some updation in distance corrary are observed then we can sure sort say this will be a sign a -ve weight ayole. Ans = Ex =) 18+ 9te. 2nd 9tex. fu, v, w? 3rd Oter  $(0-2)<\infty(v)$  -1+(-2)<(-2) -2-2<-3(v)(-2-1) < 00 () -3+(-1) < -3(v) -4-1 < -4(v) (-3+2)(0)(1) -4+2<<-1(V) -5+2<-2(V) Bellman Cycle Ford Algo Detection at 3rd itention ( Nth Iteration)